Telelogic **Statemate® MicroC Code Generator**

Statemate®

MicroC Code Generator

Before using the information in this manual, be sure to read the "Notices" section of the Help or the PDF file available from **Help > List of Books**.

This edition applies to Telelogic Statemate 4.5 and to all subsequent releases and modifications until otherwise indicated in new editions.

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Contents

Statemate's MicroC Code Generator

Statemate's MicroC Code Generator can be used to develop embedded real-time software for micro-controllers. In addition to the generation of the code, tools are provided for debugging and testing the software.

The Generator Overview

The MicroC Code Generator generates compact, readable ANSI C code, based on the model you have designed using Statemate's graphical tools. The graphical elements can be supplemented by linking in user-supplied C and or Assembly code.

The code generator allows you to define a wide range of settings in order to tailor the generated code to your target operating systems and hardware.

This guide contains detailed information regarding:

- The mapping of the various model elements to C code
- Configuring the MicroC Code Generator to output code tailored to your target operating systems and hardware
- Using the MicroC Code Generator to generate code and to then run, debug, and test the compiled code

Basic and Advanced Generator Information

In order to get you up and running quickly, this guide is divided into two sections:

The Basics

This section describes the basics of using the MicroC Code Generator. For many users, this information will be sufficient for their work.

Advanced Topics

This section covers more advanced topics, allowing you to further fine-tune the generated code to your target system.

MISRA Compliance

Statemate generates code that is compliant with MISRA Rule 60. This is accomplished when the Check-box **Options > Settings > General > Advanced Options > Generate** are selected to enable the "else" clause after "if...else if..." construct in the generated code that is required for this compliance.

Statemate Block in Rhapsody

MicroC code-generator supports integration of a Statemate model (block) into a Rhapsody model. When generating code for a Statemate block, the Statemate MicroC code generator performs the following operations:

- Generates all code in one single file that includes other required nongenerated files.
- Generates unique code to allow using few Statemate Blocks in a single Rhapsody model.

For more detailed information about this feature, refer to the *Statemate User Guide*.

Selecting an OSI for your Statemate Project

When you create a new project in Statemate, you select an *OSI (Operating System Implementation)* for the project. The OSI selected customizes the code generator to produce code that is appropriate for the target operating system.

Statemate contains a standard set of OSIs that can be selected. In addition, these basic OSIs can be modified and saved as a new OSI using the OSDT tool that can optionally be installed with Statemate (see **[Advanced: Creating Customized OSIs](#page-82-2)**).

The list of available OSIs is generated from the content of the CTD directory under Statemate. If the CTD directory contains "customized" OSIs, you will see these additional OSIs in addition to the standard set.

Once you have selected an OSI for a project, you cannot change the OSI used. However, the OSI list also contains an option **None**. This value can be selected temporarily, and you can then select an OSI at a later time. (Note that once you have selected an OSI, you cannot return to **None**.)

Working with Profiles

While the OSDT allows you to configure generated code for an operating system at a project level, profiles allow you to further configure the generated code for specific hardware targets. Because of the highly-variable hardware configurations available for embedded systems, this ability to finetune the code for the planned target hardware is essential. A number of different profiles can be defined and used for generating code in the same project.

Creating a Profile

A profile consists of the charts for which you would like to generate code, as well as the code generation options you would like to use when the code is generated.

Profiles are defined using the Code Generation Profile Editor. To open the profile editor, select **File > New/Open > Profile > Statemate MicroC Code Generator** from the Statemate main menu. The New/Open Compilation dialog box displays.

Setting the Target Configuration

The MicroC Code Generator allows you to modify a large number of settings that affect the generated code.

In addition, the code generator provides you with a list of configurations that have been set up for various hardware targets. When you select one of these configurations, the appropriate values are automatically set for the various code generation options. In many cases, this may be sufficient, and you may not have to make any more changes to the code generation options.

In some cases, however, you may want to further refine the code generation settings by manually changing the settings for certain options. In such cases, it is recommended that you choose one of the target configurations as a starting point, and then modify the specific options that must be changed.

To select a target configuration, select **Options > Set Target Configuration** from the MicroC Code Generator's main menu.

After you have selected a target configuration, when you open one of the code generation option tabs, you will see the appropriate values for that target.

Code Generation Options

The code generation options allow you to customize the generated code for your target hardware.

To modify the settings for these code generation options, select **Options > Settings** from the menu in the MicroC Code Generator Window.

The code generation options are categorized as follows:

- **[Target Properties](#page-16-0)**
- **[Code Instrumentation](#page-20-1)**
- **[Application Configuration](#page-24-0)**
- \bullet <u>[OS](#page-30-0)</u>
- **[General](#page-32-0)**
- **[Test Driver](#page-37-0)**
- **[Optimization](#page-39-0)**

Each of these categories is described in detail in the following sections.

Note

Any changes made to these code generation options affects all of the modules included in the profile.

Target Properties

The following section describes the options available on the **Target Properties** tab.

Memory Structure

- **Word Size**
	- Determines the word size used for bit buffers, such as conditions and events.
	- You can select from among the following values: 8, 16, 32.
	- If a buffer size smaller than the selected word size is sufficient, then the smaller buffer is used.
	- Examples:
		- If there are 3 conditions in the model and the word size selected is 16, then an 8-bit buffer is allocated to hold the 3 conditions.
		- If there are 20 conditions in the model and the word size selected is 16, then two buffers are allocated to hold the 20 conditions—one with 16 bits (for the first 16 conditions) and one with 8 bits (for the remaining 4 conditions).
	- **Note:** The State variable that holds the current state of a control activity does not follow this rule, and will allocate a buffer with sufficient size to hold the State's topology (up to 32 bits).

Bit Orientation

- Controls the orientation of the bits inside a single byte.
- You can select one of the following values:
	- LSbit First
	- MSBit First

Byte Orientation

- Controls the orientation of the bytes inside allocated data larger than a single byte.
- You can select one of the following values:
	- LSByte First
	- MSByte First

In addition, you can select the **Use Instrumentation** check box to control the generation of byte orientation directives in the code (#ifdef LSBYTE_FIRST directives). If this option is selected, #ifdef directives will be used in the code to accommodate the two byte orientations. This adds flexibility by making the code easier to change manually.

Memory Initialization

Reset Global (Internal) Data

Selecting this check box enables the global data reset options:

Compile Time, Static

Enables initialization of the data in the model using static initialization at the data allocation location.

Run Time, Dynamic

Enables initialization of all the data in the model through a call to the macro RESET DATA in the TASKINIT function.

The RESET DATA macro uses the function memset which should be defined in the environment. If this function is not defined, you can define the macro AVOID_MEMSET and use the function rimc_mem_set which is defined in the file <profile-name>.c.

Reset User Model Data

Selecting this check box enables the user model data reset options:

Compile Time, Static

Enables initialization of the data in the model using static initialization at the data allocation location.

Run Time, Dynamic

Enables initialization of all the data in the model through a call to the macro RESET DATA in the TASKINIT function.

The RESET DATA macro uses the function memset which should be defined in the environment. If this function is not defined, you can define the macro AVOID_MEMSET and use the function rimc_mem_set which is defined in the file

Static Initialization of "Frequency-of-Activation" Data

The generated code for static initialization of data related to Frequency of Activation complies with the profile settings.

When setting the Statemate MicroC profile options:

- Reset Global (Internal) Data, to be: RunTime, Dynamic, and
- Reset User Model Data, to be Run Time, Dynamic, the static initialization of the global data: <task_name>_COUNTER_FREQ in the file glob_dat.c will be omitted.

Use "const" Keyword to define Constant Values

If this option is selected, constant elements will be generated with the const modifier in the files glob data.c and glob dat.h, rather than being generated as pre-processor macros in the file macro_def.h.

The attribute for Constant elements (Conditions and Data-items) controls whether the specific element should be generated using the "const" keyword. The name of the new design-attribute is "Use 'const' Keyword" with these possible values:

- **no** The constant will be generated according to the settings of the MicroC code generation option: **Options > Settings... > Memory > Use "const" Keyword** to define Constant Values
- **yes** The constant will be generated using the "const" keyword, regardless of the settings of the MicroC code generation option: **Options > Settings... >Memory > Use "const" Keyword to define Constant Values**

The "Use 'const' Keyword" design attribute is available with all predefined OSIs. This attribute has the following parameters in the OSDT APIs in Memory Management page: Variable Declaration() and Extern Variable Declaration().

- IsConstant has the value "yes" if the element is defined as Constant in the model
- InitValue has the initial value related with the element (for constants it will be the element's definition)

Timeout Variable Type

In this text box, enter the data type for the Timeout Variable which holds the expiration time of a pending timeout.

Default Data Types

These text boxes are used to specify the default data types that should be used for each of the following:

- Signed Integer
- Unsigned Integer
- Bit Field
- Floating Point

The data types specified will be used when declaring data, and, in the case of signed integers and unsigned integers, will be used when using bit-wise shift operators.

The data type specified for **Bit Field** is used for conditions and bit data items. This text box is only enabled when the option (**Use Macros For**) **Conditions and Bit Data Items** on the **General** tab is not selected.

Use Fixed Point variables for "Real"

If this check box is selected, the code generated will use fixed point variables for data items of type "real."

If you have chosen to use fixed point variables in the generated code, you can use the **Word Size** radio buttons and **LSB^-2** box to select a default word size and LSB. These default settings will be used for individual variables where "*" has been selected as the value for word size and LSB in the variable properties. (The value in the LSB box represents the negative exponent to use.)

Code Instrumentation

The following section describes the options available on the **Code Instrumentation** tab.

Graphical Back Animation (GBA)

Graphical Back Animation (GBA) allows you to view an animation of the activities/states/ flowchart actions in your model. This is not a simulation, but rather an animation that runs in parallel to the running of your executable.

The GBA data is passed from the application to the Statemate model via the GBA server, which processes the data received from the application and performs the actual painting of the charts in the model.

If the **Graphical Back Animation** check box is selected, the code generator will generate the code required for this feature.

Animate Activities

If this check box is selected, the code generator will generate the flags required for displaying activity animation when the animation is run.

Animate States and Flowchart Actions

If this check box is selected, the code generator will generate the flags required for displaying animated statecharts and flowcharts when the animation is run.

If the **Graphical Back Animation** check box was selected, you can then choose one of the following animation methods:

- **Indirect, using target debugger**
	- Enables usage of GBA with a target debugger. When this method is used, the animation data is passed from the running application to the GBA Server indirectly, using a 3rd party debugger.
	- The text box next to the radio button contains the name of the source code file for the animation functions. If you prefer to use animation code that you have modified, type in the name of the appropriate source code file, or click **...** to browse for the file.
- **Direct, using Sync. TCP/IP**

Enables usage of GBA synchronized with the running application. When this method is used, the animation data is read directly from the running application, and is passed immediately to the GBA server, using the TCP/IP protocol.

- **Direct, using A-Sync., Buffered, TCP/IP**
	- Enables usage of GBA where the animation is not synchronized with the running application. When this method is used, the animation data is read directly from the running application, but rather than being passed immediately to the GBA server, the animation data is stored in a buffer which sends the data to the GBA server when the GBA task is running, using the TCP/IP protocol.
	- This method allows you to have the animation run as a separate task.

When the buffered method is selected, you can define the following settings for the GBA task:

GBA Buffer Size

The size of the buffer to use to store the animation data.

GBA Task Priority (OSEK OS only)

Allows you to specify the priority of the GBA task.

More... (OSEK OS only)

Opens the **Task Specific Attributes** dialog box.

Trace

The Trace option allows you to include code that will print to the screen information regarding the current status of tasks and/or ISRs.

Trace Tasks

When selected, code will be included for tracing tasks.

Trace ISRs

When selected, code will be included for tracing ISRs.

- **Trace Implementation File**
	- The text box next to the check boxes contains the name of the source code file for the trace functions. If you prefer to use trace code that you have modified, type in the name of the appropriate source code file, or click **...** to browse for the file.
	- The trace functions receive two parameters—one is a value which identifies the task/ISR, and the other is a character that identifies the status of the task/ISR (started, terminated, entering/exiting wait for event).

Debug

The Debug option allows you to include code that will print to the screen the state the application is in for each level of detail in the statechart. (Level 1 means that each super-step is reported, while Level 2 means that each step is reported.)

When the Debug check box is selected, you can define the following debug settings:

Debug Level 1

The state the application is in will be reported at the end of every superstep of the state machine (stable state).

Debug Level 2

The state the application is in will be reported at the end of every step of the state machine.

State Dumper File

The text box next to the radio buttons contains the name of the source code file for the implementation of the debug functions. If you prefer to use debug code that you have modified, type in the name of the appropriate source code file, or click **...** to browse for the file.

Application Configuration

The following section describes the options available on the Application Configuration tab.

Application Files...

In order to create the code for your project, the MicroC Code Generator requires certain files. You can define the names of the files and their include lists, as well as other source files and object files by clicking **Application Files...** on the Application Configuration tab.

The types of files are as follows:

OS

The following section describes the options available on the OS tab.

General

The following section describes the options available on the General tab.

Test Driver

The following section describes the options available on the Test Driver tab.

Optimization

The following section describes the options available on the Optimization tab.

Working with Profiles

Setting the Time Expression Scale

The **Options** menu in the MicroC Code Generator Window also contains an item called **Time Expression Scale**.

This option allows you to set the time scale for all expressions in the model. For a detailed explanation of this feature, see the MicroC Code Generator section in the Statemate User Guide.

Modules, Adding Charts

As mentioned earlier, defining a profile consists of defining what code should be generated, and then defining options to determine how it will be generated. You define what code should be generated by creating modules in a profile and then adding charts to the modules created.

Create Modules

Modules allow you to organize the code that is to be generated. Once you have created a module, select the charts that you would like to add to the module. When the code is generated, the code for all charts in a module will be included in a single file.

Add Charts to Modules

To add charts to modules:

- **1.** Select the relevant module
- **2.** Select **Edit > Add with Descendants** from the main menu or click the corresponding button in the toolbar.

Direct Editing of Profile Files

While you will most likely use the Profile GUI for making modifications to the profile, it is also possible to directly edit the file that contains the profile information.

To edit the file:

- **1.** Select **File> Profile Management** from the MicroC Code Generator menu.
- **2.** Select the profile.
- **3.** Click **Show**.

Checking for Errors in Profiles

After a profile has been defined, before generation of the code, you can check the profile for errors.

To check the profile, select **Compile > Check Profile** from the MicroC Code Generator menu.

Generating and Running MicroC Code

Once a profile has been defined, you can use the MicroC Code Generator to generate code from your model. After the code has been generated, you can:

- Edit the code
- Compile the code
- Run the code
- Run the code with animation

All of these options are selected from the Compile menu in the MicroC Code Generator window.

Checking Profile Before Generating Code

Since the content of the generated code is determined by the profile you have defined, it is recommended that you use the Check Profile feature before generating code from the model.

This verifies that the profile complies with the scoping rules. For example, the profile settings will be checked to make sure that they are legal and that they do not conflict with each other.

To check the current profile, follow these steps:

- **1.** Select **Compile > Check Profile...** from the MicroC Code Generator window. Any error, warning, or information messages will be displayed in the Check Profile dialog.
- **2.** Click **Dismiss** to close the Check Profile dialog.

Generating Code

To generate code, follow these steps:

- **1.** Select **Compile** > **Generate Code**... from the MicroC Code Generator window.
- **2.** When the directory tree is displayed, select the directory where the generated code files should be stored, and click **OK**.
- **3.** The Generate Code dialog will display relevant messages regarding the progress and results of the code generation process. Select **Dismiss** to close the Generate Code dialog.

Editing Code

To edit one of the generated files, follow these steps:

- **1.** Select **Compile** > **Edit Code...** from the MicroC Code Generator window.
- **2.** The contents of the output directory you selected will be displayed. Select the file that you would like to edit.
- **3.** The contents of the selected file will be displayed in the default text editor that has been defined for Statemate.
	- **Note:** To change the editor that is launched, select **Project** > **General Preferences** from the main Statemate menu, and enter the relevant command line for the parameter Editor Command Line. (You may have to close and reopen the MicroC Code Generator window for the change to take effect.)

Compiling Code

To compile the code, follow these steps:

- **1.** Select **Compile** > **Make Code...** from the MicroC Code Generator window.
- **2.** The contents of the output directory you defined will be displayed. Select the makefile to use, and click **Open**. (The name of the makefile defined in the profile's properties is offered by default.)

Running Code

To run the compiled code, follow these steps:

- **1.** Select **Compile** > **Run Code** from the MicroC Code Generator window. The Run Command dialog is displayed.
- **2.** Locate the file to run, and click **Open**.

Running Code with Animation

If you have defined an interface panel for your application, you can run the code with animation to follow the transition between different states, the activation/deactivation of activities, or the execution of a flowchart.

To run your code with animation, follow these steps:

- **1.** Select **Options** > **Settings** from the MicroC Code Generator menu.
- **2.** On the Code Instrumentation tab, select the **Graphical Back Animation (GBA)** check box.
- **3.** Save the modified profile (**OK** to close the Settings dialog, and then **File** > **Save** from the MicroC Code Generator menu.)
- **4.** Regenerate the code (**Compile** > **Generate Code...** from the MicroC Code Generator menu).
- **5.** Recompile the code (**Compile** > **Make Code...** from the MicroC Code Generator menu).
- **6.** Launch the GBA server (**Tools** > **Open GBA** from the MicroC Code Generator menu).
- **7.** Run the application (**Compile** > **Run Code...** from the MicroC Code Generator menu).
	- **Note:** If you selected the GBA option in the profile settings, then the generated application will automatically try to connect with the GBA server in order to run the animation. If you have not launched the GBA server, a message will be displayed indicating that the connection could not be established, and the generated application will continue running without the animation.

Designing Your Model: Model-Code Correspondence

The code that is generated is based on the model's graphical elements, textual elements, and the design attribute settings within these elements.

Note

This section contains many code examples. Most of these were taken from OSEK projects. Keep in mind that these are examples only. The same principles apply to non-OSEK projects as well.

To design the model, Statemate provides graphical tools for the following:

- Activity charts
- **+** Statecharts
- Flowcharts

In addition, Statemate allows you to define the following:

- Truth tables
- Lookup tables

Activity Charts

When designing an activity chart, activities are broken down into sub-activities, which are further broken down into their sub-activities, and so on, until no further decomposition is possible. Activities that cannot be broken down further are considered "basic" activities.

The code generated for an activity is a function (or a C preprocessor macro). For a non-basic activity, the function calls each of the activity's subactivity functions. For a basic activity, the function contains the implementation code.

Activities can represent functions, tasks, or ISRs (Interrupt Service Routines).

Task Activities

In the OSEK 2.0 operating system, Tasks can be divided into *basic tasks* and *extended tasks*:

- A *basic Task* runs once, upon activation, and then terminates.
- An *extended Task* runs once, upon activation, and then suspends itself, calling the API function "WaitEvent".

Basic Task - Generated Code

The code for a basic Task that contains activities A11 and A12 will resemble the following:

```
TASK (TASK1)
{
       cgActivity_A11();
       cgActivity_A12();
       TerminateTask();
}
```
If the Task is periodic, with a period of 10 ticks, the code will resemble the following:

```
TASK (TASK1)
{
       if ((cgGlobalFlags & ALARM SET TASK1) == 0) {
          cgGlobalFlags |= ALARM_SET_TASK1;
          SetRelAlarm(TASK1_ALARM, 10, 10);
       }
       cgActivity_A11();
       cgActivity_A12();
       TerminateTask();
```

```
}
```
The code for a periodic Task, containing activities A11 and A12 with CTRL1 as a controller, will resemble the following:

```
TASK (TASK1)
{
       if ((cgGlobalFlags & ALARM_SET_TASK1) == 0){
          cgGlobalFlags |= ALARM_SET_TASK1;
          SetAbsAlarm(TASK1_ALARM, 10, 10);
       }
      do {
          cgGlobalFlags &= ~BITSUPERSTEP_TASK3;
          cgActivity_A11();
          cgActivity A12();
          cgActivity_CTRL1cnt1();
       } while ( (cgGlobalFlags & BITSUPERSTEP_TASK1) != 0);
       TerminateTask();
}
```
Extended Task - Generated Code

The code for an extended Task that contains activities A21 and A22 will resemble the following:

```
TASK (TASK2)
{
       cgSingleBuffer_TASK2.eventMask = 0xff;
       start_activity_A21;
       start_activity_A22;
       while(1) {
              cgActivity_A21();
              cgActivity_A22();
              WaitEvent(cgSingleBuffer TASK2.eventMask);
              ClearEvent(cgSingleBuffer_TASK2.eventMask);
       }
       /* TerminateTask(); */
```
}

If a statechart is added beneath the Task, but not as a direct descendant, the code will resemble the following:

```
TASK (TASK2)
{
       cgSingleBuffer_TASK2.eventMask = 0xff;
       start_activity_A21;
       start_activity_A22;
       while(1) {
          do {
              cgGlobalFlags &= ~BITSUPERSTEP_TASK2;
             cgActivity_A21();
             cgActivity A22();
              if(cgDoubleBufferNew_TASK2.cg_Events)
              cgGlobalFlags |= BITSUPERSTEP_TASK2;
              cgDoubleBufferOld_TASK2 = cgDoubleBufferNew_TASK2;
              cgDoubleBufferNew_TASK2.cg_Events = 0;
       } while ( (cgGlobalFlags & BITSUPERSTEP_TASK2) != 0);
```

```
WaitEvent(cgSingleBuffer TASK2.eventMask);
   GetEvent(TASK2, &cgSingleBuffer TASK2.eventsBuff);
   ClearEvent(cgSingleBuffer_TASK2.eventMask);
}
/* TerminateTask(); */
```
If the Task is periodic, with a period of 10 ticks, the code will resemble the following:

```
TASK (TASK2)
{
       SetRelAlarm(TASK2_ALARM, 1, 10);
       cgSingleBuffer_TASK2.eventMask = 0xff;
       start_activity_A21;
       start activity A22;
       while(1) {
          do {
             cgGlobalFlags &= ~BITSUPERSTEP_TASK2;
             cgActivity_A21();
             cgActivity A22();
              if(cgDoubleBufferNew_TASK2.cg_Events)
              cgGlobalFlags |= BITSUPERSTEP_TASK2;
              cgDoubleBufferOld_TASK2 = cgDoubleBufferNew_TASK2;
              cgDoubleBufferNew_TASK2.cg_Events = 0;
          } while ( (cgGlobalFlags & BITSUPERSTEP_TASK2) != 0);
       WaitEvent(cgSingleBuffer TASK2.eventMask);
       GetEvent(TASK2, &cgSingleBuffer TASK2.eventsBuff);
       ClearEvent(cgSingleBuffer_TASK2.eventMask);
       if(cgSingleBuffer_TASK2.eventsBuff & 0x01)
          GEN_IN_CURRENT(TASK2_EV);
       }
       /* TerminateTask(); */
}
```
}

ISR (Interrupt Service Routine) Activities

Interrupt Service Routines (ISRs) run once, upon activation, and then end.

ISR Categories

For OSEK 2.0, three ISR categories can be used: 1, 2, and 3.

The decision of which ISR category to use depends on the content of the functions it runs. According to the OSEK/OS specification, an OS API function cannot be called from a category 1 ISR. For categories 2 and 3, some OS API functions can be called, but only within code sections marked by EnterISR()/LeaveISR() calls.

The following are some code examples for different types of ISRs:

ISR - Examples of Generated Code

Code for a category 1 or 2 ISR, named ISR0, containing activities I01 and I02:

```
ISR (ISR0)
{
       cgActivity_I01();
       cgActivity I02();
}
```
Code for a category 3 ISR function names ISR0, containing activities I01 and I02:

```
ISR (ISR0)
{
      EnterISR();
      cgActivity_I01();
      cgActivity_I02();
      LeaveISR();
}
```
Code for a category 3 ISR function named ISR1, containing activities I11 and I12, and a controller named CTRL1:

```
ISR (ISR1)
{
      EnterISR();
      do {
          cgGlobalFlags &= ~BITSUPERSTEP_ISR1; MicroC 41
          TASK/ISR Run Modes
          cgActivity_I11();
          cgActivity_I12();
          cgActivity_CTRL1cnt1();
      } while ( (cgGlobalFlags & BITSUPERSTEP_ISR1) != 0);
      LeaveISR();
}
```
Task/ISR Run Mode

A Task/ISR can have one of the following run modes:

- Single Step—the Task/ISR always runs a single step, then returns handling to the operating system.
- Super Step—the Task/ISR runs the necessary number of Tasks before returning handling to the operating system.

Decomposition of Non-basic Activities

When a non-basic activity does not contain an immediate descendant that is a control activity, all of the activity's subactivities are considered active when the activity is active. For such a non-basic activity, the generated code will resemble the following:

```
void
cgActivity_A11acy1(void)
{
      cgActivity_A111();
      cgActivity_A112();
}
```
Execution Order (for Subactivities)

The order in which the subactivities are called within the A11 activity body is determined by the subactivity design attribute *Execution Order*, as defined for A111, A112, A113. In the previous example, the value of this attribute was "1" for subactivity A111 and "2" for subactivity A112.

If the *Execution Order* attribute is not set, the calling order is not defined.

Code for Basic Subactivities

Basic activities can be defined in one of three activation modes:

- Reactive controlled
- Reactive self
- Procedure-like

For reactive controlled and reactive self modes, the code for the basic activity will resemble the following:

```
void
cgActivity_A111(void)
{
       … Body implementation
}
```
For the procedure-like mode, the code for the basic activity will resemble the following:

```
void
cgActivity_A112(void)
{
       if ((cgActiveActivities1 & BITAC_A112) != 0) {
       … Body implementation
       stop_activity(A112);
       }
}
```
Adding a controller A11_CTRL to A11 will make the code look like:

```
void
cgActivity_A11acy1(void)
{
       cgActivity_A111();
       cgActivity_A112();
       cgActivity_A11_CTRLcnt1();
}
```
with the controller function, *cgActivity_A11_CTRLcnt1()*, looking like:

```
void
cgActivity_A11_CTRLcnt1(void)
{
       cgDo_A11_CTRLcnt1();
}
```
The implementation of *cgDo_A11_CTRLcnt1()* depends on whether A11_CTRL is implemented as a statechart or as a flowchart.

For a statechart implementation:

```
void cgDo_A11_CTRLcnt1(void)
{
       StateInfo_A11_CTRLcnt1 nextState_A11_CTRLcnt1 = 0;
       if (currentState A11 CTRLcnt1 == 0) {
          nextState_A11_CTRLcnt1 = FS_A11_CTRLst2;
       }
       else
       {
          … Rest of the Statechart logic
       }
       if (nextState_A11_CTRLcnt1 != 0) {
          if (currentState_A11_CTRLcnt1 !=
          nextState_A11_CTRLcnt1)
          cgGlobalFlags |= BITSUPERSTEP_TASK1;
          currentState_A11_CTRLcnt1 = nextState_A11_CTRLcnt1;
       }
}
```
For a flowchart implementation:

```
void
cgDo_A11_CTRLcnt1(void)
{
… The Flowchart logic
}
```
Communication and Synchronization Services between Activities

Communication and synchronization services between activities, including those not residing in the same task/ISR, consist of the following:

- Non-queued messages
- Queued messages
- \bullet Signals
- Global Data
- Semaphores

Non-queued Messages

Non-queued messages uses a message identifier (i.e., the message name) to share data between various tasks in the application. The sender and/or receiver task for such a message can be running in the same ECU, share the same memory address space, or run across and ECU network on a remote MCU. The user of the message need not be aware of the concrete implementation. Thus, use of this mechanism ensures that the resulting design is correct, flexible, and efficient.

Queued Messages

The queued messages mechanism is similar to that of non-queued messages. The difference is that queued messages do not contain values but rather signal the occurrence of some event.

Signals

Signals indicate the occurrence of some event. However, since they are not queued, there is no information regarding how many such events occurred, until they are processed.

Also, these Task Event signals require that a specific task is addresses with a specific event, thus requiring knowledge of the application structure. The Task Event implementation is more efficient than ordinary or queued messages, however the task must be of type extended, which is not always possible or efficient.

The downside of requiring knowledge of the application is balanced by the improved performance. The weight assigned to these two issues will depend on the problem at hand.

Global Data

As always with real time applications, when using global data, caution should be taken regarding the validity of the data when running in a preemptive environment with multiple tasks and ISRs. The protection mechanism supported is the OSEK RESOURCE mechanism, which is similar to a binary semaphore.

Semaphores

It is common for data to arrive through the bus or board ports, in some predefined messages and addresses, and must be reproduced to the bus or board in the form of other predefined messages and addresses.

In such situations, the designer simply uses the defined interface for his application. However, the discussion above is relevant when one tries to build an implementation that will use the appropriate interfaces but will also be easy to maintain, modify, and ported to various other environments, usually unknown at design time.

Statecharts

Statecharts define the behavior of activities defined in activity diagrams, and are linked to an activity with a control activity. Statecharts can contain sub-charts (nested statecharts).

This section provides details regarding the code that is generated for statecharts in your Statemate model.

Functions Generated for Statecharts

For a control activity All_CTRL, the following two functions will be generated:

```
void cgActivity_A11_CTRLcnt1(void)
void cgDo_A11_CTRLcnt1(void)
```
The code generated for these functions will be as follows:

```
void cgDo_A11_CTRLcnt1(void)
{
      StateInfo_A11_CTRLcnt1 nextState_A11_CTRLcnt1 = 0;
      if (currentState A11 CTRLcnt1 == 0) {
          nextState A11 CTRLcnt1 = FS A11 CTRLst2;
      }
      else
```

```
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```

```
{
          … The rest of the Statechart logic
       }
       if (nextState_A11_CTRLcnt1 != 0) {
          if (currentState_A11_CTRLcnt1 !=
          nextState_A11_CTRLcnt1)
          cgGlobalFlags |= BITSUPERSTEP_TASK1;
          currentState A11 CTRLcnt1 = nextState A11 CTRLcnt1;
       }
}
void cgActivity_A11_CTRLcnt1(void)
{
       cgDo_A11_CTRLcnt1();
}
```
Statechart - Data Usage

When a statechart is created, a stateInfo data type is defined and a few variables of that type are declared.

For the previous example, the stateInfo data type would be named stateInfo All CTRLcntl, and would be defined as an unsigned type of 8, 16, or 32 bits (e.g., typedef int8 StateInfo A11 CTRLcnt1)

The StateInfo variables will be currentState, nextState, and staySame:

```
StateInfo_A11_CTRLcnt1 currentState_A11_CTRLcnt1;
```
(global variable)

```
StateInfo_A11_CTRLcnt1 nextState_A11_CTRLcnt1;
```
(automatic variable)

StateInfo_A11_CTRLcnt1 staySame_A11_CTRLcnt1;

(automatic variable)

The currentState and nextState variables will always be allocated. The staySame variable will be allocated only if the entering or exiting reaction function is required.

currentState is allocated as a global variable, while nextState and staySame are allocated as local, automatic variables to the statechart function cgDo_....

Statechart - Generated Functions

In general, functions generated from statecharts will resemble the following (the provided line numbers are used in explanations of the code below):

```
1 void
2 cgDo_A11_CTRLcnt1(void)
3 {
4 StateInfo A11 CTRLcnt1 nextState A11 CTRLcnt1 = 0;
5 if (currentState A11 CTRLcnt1 == 0) {
6 nextState_A11_CTRLcnt1 = FS_A11_CTRLst2;
7 }
8
9 else
10 {
11 … The rest of the Statechart logic
12 }
13 if (nextState_A11_CTRLcnt1 != 0) {
14 if (currentState A11 CTRLcnt1 !=
             nextState_A11_CTRLcnt1)
15 cgGlobalFlags |= BITSUPERSTEP TASK1;
16 currentState A11 CTRLcnt1 =
             nextState_A11_CTRLcnt1;
17 }
18 }
```
In line 4, the next state variable is reset. This variable will be set only if a transition has been made, and will hold the statechart's new state configuration.

Lines 13 and 14 check the next State variable to determine whether a transition was made, and whether to enforce another step in the task holding the statechart.

Line 16 advances the statechart configuration a step to hold the configuration of the next step.

In your statechart, lines 5 to 12 will be replaced with specific code resulting from the specified statechart logic. For example, in many cases, two additional functions will be generated here entry actions and exit actions. If the statechart logic requires entering/exiting reactions, the functions will resemble the following.

```
void
cgEnterActions_A11_CTRLcnt1(void)
{
… entering reactions code
}
void
cgExitActions_A11_CTRLcnt1(void)
{
… exiting reactions code
}
```
When either of these function are needed, the following changes to cgDo_... will also be made:

```
void cgDo_A11_CTRLcnt1(void)
{
       StateInfo_A11_CTRLcnt1 nextState_A11_CTRLcnt1 = 0;
       staySame_A11_CTRLcnt1 = 0;
       if (currentState A11 CTRLcnt1 == 0) {
          nextState_A11_CTRLcnt1 =
          FS_DefaultOf_Chart_A11_CTRL;
       }
       else
       {
```

```
… The rest of the Statechart logic
      }
      if (nextState_A11_CTRLcnt1 != 0) {
          cgExitActions A11 CTRLcnt1();
          cgEnterActions_A11_CTRLcnt1();
          if (currentState_A11_CTRLcnt1 !=
          nextState_A11_CTRLcnt1)
          cgGlobalFlags |= BITSUPERSTEP_TASK1;
          currentState_A11_CTRLcnt1 = nextState_A11_CTRLcnt1;
      }
}
```
Order of Function Execution Rules

The following rules determine the order in which exiting actions, entering actions, transition actions, and static reactions are carried out for a state transition:

- **1.** Where the state has not changed, static reactions are carried out in descending order down the state hierarchy.
- **2.** When a transition is detected, the transition action is carried out immediately.
- **3.** Exiting actions are then carried out, from the innermost state exited to the outermost state.
- **4.** Finally, entering actions are carried out, from the outermost state entered to the innermost state.
	- **Note:** In order to optimize code, you may sometimes want to inline entering/exiting reactions. For more details, see **[Optimization](#page-39-0)**.

State Variable Validation Macro

The IS IN VALID STATE <ctrl-activity-name> macro is defined in the generated file *macro_def.h* for each Control-Activity Statechart hierarchy. This macro includes code for validating that the state variable has a valid value. The validation is accomplished using the inLeafState()(inState() for And States macros against all possible leaf states in the hierarchy. In addition, there is a test against the valid "0" value.

For example, for the Control Activity (**CTRL)** with two leaf states (**S2** and **S2),** the generated macro is as follows:

```
#define IS IN VALID STATE CTRL ( \backslashinLeafState(currentState CTRL, S1, StateInfo CTRL) || \ \rangleinLeafState(currentState_CTRL, S2, StateInfo_CTRL) || \
currentState CTRL == 0 \setminus)
```
Timeout Implementation

Software counters are used as the basis for the implementation of timeouts. When a timeout or delay is set, the current value of the relevant software counter will be added to the requested delay time and stored in a variable, using a defined macro, INSTALL_TIMEOUT. By default, MicroC relates to the primary software counter defined in the compilation profile.

Other software counters can be referenced using an optional third argument in the timeout operator. The name of the counter is as written in the model, using the syntax:

```
tm(en(S1), 12, myCounter)
```
where myCounter is the name of the counter. Each counter receives an index value defined as <counter_name>_INDEX. The index value identifies that specific counter in the application.

INSTALL_TIMEOUT Macro

The INSTALL_TIMEOUT macro has three arguments:

- The name of the event.
- The requested delay.
- The index of the counter on which it is pending

This allows the code to reuse the same timeout variable with different counters. The first argument is concatenated to the INSTALL macro, as shown here. In the code, a call like the following will be used:

INSTALL_TM(tm_999999962, 10, SYS_TIMER)

This call will set a timeout to expire 10 ticks from the current time of SYS TIMER. The macro itself will be defined as follows:

```
#define INSTALL TM tm 999999962(D, C) \setminuscgTimeoutsMask |= tm 999999962 TM MASK;
tm 999999962 TIME = currentTick + (D);
```
This call will assign to $tm = 999999962$ TIME which is a variable of type Timeout Variable Type the current counter value, help in currentTick plus the requested delay time help in D. In addition, the bit tm_999999962_TM_MASK is set to flag that this timeout is pending.

A test for timeout expiration is carried out in the function:

```
genTmEvent_<CTRL_CHART_NAME>(<Timeout Variable Type>
\overline{\text{currentTickVar}}, \overline{\text{c} \text{buffer}} * buff, uint8 counterIndex)
```
The third parameter, uint 8 counterIndex, holds the index of the counter that is referred to in the current call to this function. Before each call to this function, the correct counter would be read into the currentTick global variable.

For each Timeout Variable, there are three options for code generation inside the genTmEvent ... function:

- **1.** When there is only one counter in the model, no check will be made for the counter.
- **2.** When there is only one counter that the timeout.variable can be installed for, then the code will look.like:

```
if(counterIndex == <ITS_COUNTER_NAME>_INDEX &&
cgTimeoutsMask & tm_999999993_TM_MASK &&
currentTickVar >= t\overline{m} 999999993 TIME) {
    GEN IN BUFF(tm 999999993, \overline{\text{buffer}});
cgT\overline{1}me\overline{o}utsMask\overline{\&}="x=tm_999999993\overline{1}m_MASK;
```
- **3.** If there is more than one counter that the Timeout Variable can be installed for, then the code will include the following provisions:
	- **a.** In the file, *glob* dat.c a uint8 variable tm 999999993 counter; is generated, holding the index of the current relevant counter.
	- **b.** In the file *macro* $def.h$, along with the previous code that was generated for the INSTALL_TIMEOUT macro, there is one more statement that keeps the INDEX of the counter for which the timeout was installed.

The index that is passed to the function is compared with the index of the counter that was used when the timeout was installed. This enables the application to identify the counter on which the timeout is pending.

Special Requirements for OSEK-targeted Applications

OSEK-targeted applications have special requirements:

- **1.** For each counter, an overflow task named <counter_name>_OVERFLOW is generated. This includes the task declaration (found in $os_dec.l.h$) and body code (found in glob func.c).
- **2.** In each task, there is overflow management provided *only* for the Timeout variables that refer to the specific counter.
- **3.** For each counter, an alarm named <counter_name>_ALARM is generated. This includes the alarm declaration (found in *os_decl.h*) and installation (found in *macro_def.h*). In the *macro def.h* file, a new macro is generated:

```
#define SET ADDITIONAL OVERFLOW ALARMS() \{\setminusSetAbsAlarm(<counte_name>_ALARM, 0,
    OSMAXALLOWEDVALUE); \overline{\setminus}}
```
This macro installs all the overflow alarms that activates the overflow tasks. A call to this macro is in the file <profile-name>.c after the installation of the SYS_TIMER_ALARM (formerly known as SYS_TIME_OVERFLOW).

Compare this to non-OSEK implementations:

- **1.** For each counter, an overflow function named on<counter_name>_OVERFLOW is generated. In each task, overflow management is provided *only* for the Timeout variables that refer to that specific counter.
- **2.** IMPORTANT there is no call to these functions in the generated code. Therefore, in order to use them, additional code should be added by the developer that decides when to call these functions (on overflow), possibly in *usercode.c*.
	- **Note:** Set from within the Code Generation Profile Editor. Use **Options > Settings >** The goal is to have a variable that is bigger then the counter, thus avoiding the "value overflow" problem.

History and Deep History Implementation

History and deep history implementation require a *StateInfo* variable for each state with a history connector or deep history connector.

The state configuration is stored in this *StateInfo* variable, such that when a transition occurs into the history/deep history, this configuration is assigned to the *nextState* variable, causing an entrance to the stored state configuration.

The operators history clear and deep clear assign the corresponding default state configuration to the corresponding *StateInfo* variable.

Optimization of Statechart Code

There are a number of optimization options available for state chart code. For more information regarding these optimization options, see **[Optimization](#page-39-0)** in **[Working with Profiles](#page-14-0)**.

Recommendations for Efficient Code

To increase the efficiency of your code, it is recommended that you:

- Avoid redundant intermediate states (i.e., not persistent states).
- Avoid duplication of code segments use functions or defined actions instead of hardcoded duplicates.
- For a simple single state with self-transition scheduling some operation, use a static reaction or an ISR.
- Use state hierarchy to represent priorities.

Flowcharts

A flowchart is another type of diagram that can be used to define the behavior of an activity.

For the purpose of code generation, a flowchart is considered to be the flowchart directly connected to a control activity, as well as all of its sub-charts and the generics instantiated within them.

Functions Generated for Flowcharts

For a control activity, A12_CTRL, the following two C functions will be generated:

```
void cgActivity_A12_CTRLcnt1(void)
void cgDo_A12_CTRLcnt1(void)
```
The body of these functions will resemble the following:

```
void
cgDo_A12_CTRLcnt1(void)
{
       … The flowchart logic
}
void
cgActivity_A12_CTRLcnt1(void)
{
       cgDo_A12_CTRLcnt1();
}
```
The function cgActivity_A12_CTRLcnt1 simply calls cgDo_A12_CTRLcnt1.

Flowchart Implementation

The flowchart language that is used graphically describes a structured C program.

While both flowcharts and statecharts define the behavior of activities, the graphics and semantics used in flowcharts are very different from those used in statecharts. In some cases, you may prefer this approach to the statechart approach.

The code of a flowchart runs from beginning to end, without stopping and without explicitly maintaining its internal state. Flowcharts do not have a notion of state or internal state.

While flowcharts allow the creation of highly visible, graphical algorithms, there is no inherent overhead in the code that is generated from the chart. The MicroC code generator produces optimized structured code, just as it does for statecharts.

If a flowchart is properly constructed, it will result in the generation of highly optimized code. However, it is the responsibility of the designer to build appropriate charts with proper syntax, logic, and association with a valid control activity. Otherwise, the results could be non-structured code.

Flowchart Elements

The elements found in flowcharts can be divided into two categories:

- Boxes
- Arrows

Compound boxes—boxes containing other boxes—represent code blocks.

Labels

As with statecharts, the graphical elements in statecharts can be assigned labels for purposes of identification and describing associated logic or value assignments. Labels on arrows are considered to be literal constants and are allowed only for arrows exiting either decision or switch elements.

Decision Expressions

The following expressions are permitted for decisions:

- Event (such as ON_POWERUP)
- Condition (such as [POWER_ON])
- Expressions (such as $[TEMP > 27]$)

The following expressions are permitted on arrows exiting decisions:

- Yes
- \bullet No
- \bullet True
- False

Switch Expressions

The following expressions are permitted for switches:

 \blacktriangleright Value-type expressions (such as F1(3) + 5)

The following expressions are permitted on arrows exiting switches:

- Literal constants
- \bullet else
- \bullet default

Minimization of Goto Statements

The MicroC Code Generator tries to minimize the number of goto statements in the code. This results in more structured and readable code. However, this is not always possible, and in some cases goto statements may appear in the generated code.

Restructuring the flowchart or using statecharts instead of flowcharts may eliminate generated goto code.

Code Structure

The code is generated in C blocks.

Compound (non-basic) boxes in the flowchart are translated into blocks.

Basic boxes are interpreted as control positions between executable statements.

Begin/End Points

The start point for each block (the point at which the non-basic box is entered) is marked using a Start arrow in that box. The end point is marked using an End connector in that box.

The start point for the entire flowchart is marked using a Start arrow at the highest level. The end point for the entire flowchart is marked using an End connector at the highest level.

Flowchart execution stops when it can make no more progress. This may be due to reaching an End connector, or it may be due to reaching a box for which all of the outgoing arrows evaluate to false.

Arrows and Labels

In the case of nested boxes, all arrows on the inside boxes are tried first. If none of them can be taken, then higher-level arrows are tried. This continues until the highest level is reached. If no arrows can be taken at that level, the code finishes executing, i.e., the function returns.

Flowchart Examples

In the examples below, the flowchart is followed by the code generated for the flowchart.

Simple Flowchart


```
void cgDo_FL_CH_TEST_3()
{
         \mathtt{DI}\text{=} \mathtt{FUNC1} ( ) \mathtt{;}if (DI > 5) {
               ACT_2();
          }
          else {
               ACT_1();
          }
}
```
Find/Merge Logic


```
void cgDo_FL_CH_FIND_MERGE_BOX()
```

```
{
      DI = 1;if ((DI == 1)) {
         if ((DI == 3)) {
            DI = 4;}
         DI = 5;}
      else {
         DI = 3;}
      DI = 2;
}
```
Switch Control


```
void cgDo_USE_SWITCH_CTRL() 
{ 
         switch(DI + 1) {
            case 3: 
               if ((DI < 3)) {
                   switch(DI * 2) {
                    case 4: 
                    DI = 43; break; 
                    default: 
                     DI = 87; 
                     break; 
                    } 
                   DI = 4; }
```

```
 else { 
                   DI = 455; } 
             break; 
             case 5: 
             { 
                 switch(COLOR) { 
                     case BLACK: 
                    DI = 5; break; 
                     case BLUE: 
                    DI = 65; break; 
                     case RED: 
                     DI = 99; 
                     break; 
                     default: 
                     break; 
                 }
           } 
                DI = 34; 
                break; 
             default: 
                 if (EV) { 
                    GENERATE_EVENT(EV) ;
                 } 
                 else { 
                     SetRelAlarm(EV_ALARM, 11, 0); 
                 } 
                break; 
         } 
}
```
Truth Table Implementation

The code implementation of truth tables is demonstrated below using the following sample truth table implementing function *F*, using data items *DI1* and *DI2* as input.

For this truth table, the following code would be generated:

```
void f(void)
{
  if (DI1 == 1) {
    if (DI2 == 1) {
       A1();
     }
     else {
      if (DI2 == 2) {
          A2();
       }
     }
   }
   else {
    if (DI1 == 2 & & DI2 == 3) {
       A3();
     }
   }
}
```
Lookup Table Implementation

Statemate allows the definition of lookup tables to represent the type of non-linear $Y=F(X)$ functions that are so common in the world of microcontrollers. The data for a lookup table can be defined manually in Statemate, or imported from any ASCII data file. You can elect to have linear interpolation between defined points, or a histogram-like mode. Upper and lower bounds can be defined, as can the search order to use (low-to-high, high-to-low).

In the sample lookup table below, the input is defined as *Integer*, and the return value of the function is defined as *Real*.

Using the settings *linear interpolation*, *high-to-low* search order, *lower bound* = 0, *upper bound* = 4, the following code will be generated:

```
double LOOKUP1(int IN1)
{
       /*
      Interpolation Function:
      if(In < X2 && In >= X1)
      Out = (Y2-Y1)/(X2-X1)*(In-X1)+Y1*/
      double LOOKUP1_retval;
      if(IN1 < 1)LOOKUP1_retval = (0);
      else if(IN1 >= 1000)
          LOOKUP1 retval = (4);
      else if(IN1 >= 100)
          LOOKUP1_retval = (4 - 3)/((double)1000 - 100)*(IN1 -
          100) + 3;
      else if(IN1 >= 10)
          LOOKUP1 retval = (3 - 2)/((\text{double})100 - 10)*(IN1 -
```

```
10) + 2;else if(IN1 > = 1)LOOKUP1_retval = (2 - 1)/((double)10 - 1)*(IN1 - 1)+ 1;return(LOOKUP1_retval);
```
Fixed-Point Variable Support

}

This section describes the MicroC Code Generator's fixed-point support for integer arithmetic, which scales integer variables so that they can represent non-integral values (fractions). This feature allows you to perform calculations involving fractions without requiring floating-point support from the target.

Statemate's MicroC Code Generator supports fixed-point arithmetic at the model level, as well as in the generated code.

Fixed-Point Variable Implementation Method

Statemate's MicroC Code Generator uses the "2 factorials" implementation method—redefining the least significant bit (LSB) to represent zero, or the negative power of 2. This implementation is not the most accurate method but it provides reasonable code size and runtime performance.

For example, take the binary 8-bit value 0b00010001. Usually, the value represented here is "17":

- The LSB (1st bit) corresponds to 2^0 (1).
- The 5th bit corresponds to 2^4 (16).

Rescaling this value to begin at 2^{-3} gives: 2.125 = $1*2^{-3}$ (or 0.125) + $1*2^{1}$ (or 2)

The parameter required here is the power (of 2) represented by the LSB. This is also the resolution.

Supported Operators

You can use the following operators with fixed-point variables:

- \bullet Arithmetic $(+, -, *, \wedge)$
- \triangleleft Assignment (=)
- Comparison $(<,>,<,>=, ==, !=)$
- Functions (return value, parameters, local variables)

Evaluating the wordSize and shift

The wordsize and shift of an object are defined by its design attributes (specified in the element properties). The MicroC Code Generator determines the wordsize and shift of expressions made of objects and operators by using the formulas listed in the macro definition table below.

The conventions used in the table are as follows:

- **WS—**The wordSize of the object
- **SH—**The shift of the object
- **RG—**The range (wordSize shift)
- \blacklozenge **MAX(A, B)**—A>B:A:B
- \bullet **SUM(A, B)**—A+B
- **SUB(A, B)—**A–B:

If the evaluated wordSize is greater than 32 bits, MicroC displays the following messages:

- wrn_err.inf Warning: Fixed-Point Overflow in Expression:<Expression>
- generated code /* Warning Fixed- Point Overflow in Expression.*

This message is located right after the expression.

When you use fixed-point variables in integer arithmetic, the special functions (or C macros) provided in the FXP package are used to perform the calculations. The following table lists these macros

Unsupported Functionality

The following functionality is not supported:

FXP parameter passed by reference

The MicroC Code Generator generates the following error message:

Error: Unsupported usage of Fixed-Point parameter used by reference.

In function: <FUNC_NAME> Parameter number: <PARAM_NUM>.

MicroC ignores the remainder in division operations that result in remainders

For example:

```
FXP1(WS=8, SH=2) = 5FXP2(WS=8, SH=2) = 2
FXP1/FXP2 = 2 (not 2.5)
```
Specifying Fixed-Point Variables

To specify fixed-point variables in the Code Generator, follow these steps:

- **1.** Select **Options** > **Settings** > **Target Properties** from the MicroC Code Generator window.
- **2.** Select the option **Use Fixed Point variables for "Real"**.
- **3.** Select the default **Word Size** (8/[16]/32) and **LSB=** 2**^-** (0,1,2,..n).

The Generated Code

Fixed-point variables are implemented using uint variables (sints, sint16, sint32), with hardcoded shift values. Data is allocated according to the wordsize of the variable:

All calls to functions or expressions requiring integer values are done through an FXP-to-int cast, including the test driver / panel driver. Specifically, the operators "ROUND" and "TRUNC" are called with an FXP-to-int cast.

For example, given a fixed-point variable f_{xp} var, an integer variable int var, and the following actions:

```
INT_VAR := FXP_VAR + 4;
FXP_VAR := INT_VAR/5;
```
The generated code is as follows, if you specify fixed-point mode:

```
INT_VAR = RS_FXP2FXP16(FXP_VAR + LS_FXP2FXP16(0x4,
0, FXP_VAR_FXP_SHIFT), FXP_VAR_FXP_SHIFT, 0);
FXP_VAR = LS_FXP2FXP16(INT_VAR / 0x5, 0, FXP_VAR_FXP_SHIFT);
```
Usage of Upper Case / Lower Case in Statemate

When you create a model element in Statemate, the element name that is saved in the Statemate database reflects your exact case usage (lower case, upper case, mixed).

However, Statemate also provides a preference called *Exact Case Mode* (under **General Preferences**), which allows you to specify how the element name should appear in both the graphical editors and in the generated code. If this preference is set to *On*, then the element names will appear just as you typed them, in both the graphical editors and in the generated code. If this preference is set to *Off*, then the element names will appear in all upper case.

The mode that you select does not affect the way the names are saved in the database. There, they always remain exact case. So to restore exact case usage in both the graphical editors and the generated code, all you have to do is change the setting of the *Exact Case Mode* preference.

Statemate will not let you define two elements with the same name, differing only in the case used, for example, an event called *aB* and an event called *Ab*. If you enter a different-case variation of a name, Statemate automatically converts it to the exact case usage of the original element.

If you want to change the case used for an element, you must use the *Rename* option.

In addition to the code customization that can be achieved through the use of profiles (described above), you can use the OSDT (Operating System Definition Tool) to further customize code generation in order to create code appropriate for your target operating system. The result of this process is an OSI (Operating System Implementation).

Since it is assumed that this level of customization will be performed at the project level (and not at the individual user level), the OSDT is only installed with Statemate if it is selected during a custom installation.

Using the OSDT to Customize OSIs

The OSDT includes predefined OSIs (operating system implementations) for a number of operating systems. If you are using one of these systems, you do not need to perform any customization.

If, however, customization is required, it is recommended that you select the OSI that is closest to the system you will be using and use this as the base to which you add the required customizations.

To use an existing OSI as a template, select the desired OSI from the **OS Implementation** dropdown list. Then, select **File > Save As** from the OSDT menu, and provide the name that you would like to use for the new OSI you are creating.

After making any changes to individual OSI settings on the various tabs, select **File > Save** from the OSDT menu to save the changes to your OSI. If you try to close the OSDT before changes have been saved, you will be asked if you want to save the changes.

Statemate will not allow you to save changes to the predefined OSIs. If you make any changes to these OSIs, you will be asked to save the modified profile under a different name.

Static OS Configuration

If you are designing software for a system that requires a static configuration file, such as an OSEK OIL file, select the check box **Static OS Configuration**.

When this option is selected, the **Static OS Configuration...** button is enabled, allowing you to define the relevant APIs. This will also result in the static OS options being displayed on the OS tab of the profile settings.

Memory Management

If you select the **Memory Management** check box, the **Memory Management...** button will be enabled, allowing you to define various code options such as:

- directives specifying how data should be stored in memory
- directives such as *#ifdef* to include/exclude parts of the code

OSEK API

If you are designing an OSEK-based system, select the **Use OSEK API** check box, and select one of the OSEK implementations from the drop-down list below the check box.

If you want to customize the OSEK API, select the **Allow API Overriding** check box. When this options is selected, the **API Definitions...** button is enabled.

Types of Customization Available

The code customization that the OSDT allows can be categorized as follows:

[Customizing Design Attributes](#page-84-0)

These attributes allow you to specify additional information for Statemate model elements. They can also be used as building blocks when defining APIs (see below).

[Customizing API Definitions](#page-91-0)

These allow you to define the code that should be generated for a wide variety of basic actions. In addition, you can customize the code for static OS configuration and memory management. You can also specify formats for issues such as variable and function naming or file headers/footers.

[Specifying Related Files](#page-217-0)

This allows you to select the files that you will want to have available in your workarea, for example, makefiles, .h files, and .oil files.

Customizing Design Attributes

Elements in a Statemate model are further defined using design attributes. Every element has attributes that are relevant to that type of element. For example, the design attributes for activities may include Type, Task Run Mode, Generate Function (yes/no). These attributes appear on the Design Attributes tab of the Properties dialog.

For most of the provided OSIs, a default set of design attributes are defined for the various Statemate elements. You can also add new design attributes for an element, using the Attributes Editor, which is part of the OSDT. In addition to allowing you to define new attributes, this editor allows you to modify the various field values for the default design attributes.

The design attributes for an element, both the default attributes and any new ones, can be used as tokens when defining the different APIs with the OSDT, allowing you to include the values of element attributes in the generated code.

For each design attribute, the editor displays:

- The basic attribute information, such as name, type, default value, key name.
- The attribute's dependencies
- Information that should be displayed when Statemate's Info tool is used.

Note

The *key name* is the string that is used when the attribute is used as a token. Since both other attributes and API's may reference an attribute's key value, you will be asked to confirm any changes to an attribute's key name. When such a change is made, the name change will be propagated to any attributes for the same element that have a reference to the attribute.

To edit attribute information, follow these steps:

- **1.** Click **Edit Attributes...** on the main OSDT screen.
- **2.** The Design Attributes Definition dialog is displayed. This screen is not used for editing attribute information, but simply for turning on/off the use of design attributes for specific model elements. If you clear the check box next to an element name, the design attributes will not appear when you display the Properties dialog for elements of this type.

3. In the Design Attributes Definition dialog, click **Edit Attributes...** . The Attribute Editor screen will be displayed. All editing of information is done on the right side of this dialog. The left side serves as a browser, and contains controls that can be used to modify the order in which the attributes appear in the Properties dialog.

To save any changes/additions you have made to the design attributes, select **File > Save All** from the menu, and after you have closed the Attribute Editor dialog, select **File > Save** from the OSDT menu. The changes will not be saved to the OSI if you have only saved them in the Attribute Editor dialog.

Design Attribute Fields

General

Dependency

The Dependency tab allows you to make the existence of a design attribute dependent upon other attributes such that the design attribute will be used only under certain conditions. If these conditions are not met, the attribute is not included in the generated code and not visible in the Design Attributes tab in the Properties dialog.

For example, the attribute Task Run Mode will be displayed only if the activity is of type *Task*.

The drop-down lists can be used to compose conditions involving relevant attributes.

You can then specify whether the design attribute is available only if all of the listed conditions are met or if any of the listed conditions are met.

Info

The Info tab allows you to specify whether the design attribute should be included in Statemate's Info dialog when the user displays this dialog for a given element, and, if so, how it should be displayed.

Customizing API Definitions

The API definitions are divided into the following categories:

- **[General API Definitions](#page-96-0)**
- **[Customizing Code Style](#page-149-0)**
- **[Customizing Memory Management](#page-171-0)**
- **[Customizing the Static OS Configuration](#page-193-0)**

The OSDT contains a number of features that facilitate API definition. These features apply to each of the above categories, and will be described in the following section. Afterward, each of the API definition categories will be discussed in detail.

Features that Facilitate API Definition

The following features can be used when defining the APIs:

- **[Browse Properties from OSDT](#page-91-1)**
- **[Using Parameters for the Generated Code](#page-92-0)**
- **[Conditional Expressions in API Definitions](#page-94-0)**

Browse Properties from OSDT

When defining APIs in the OSDT, you can easily access tokens available for use in the definitions.

To see the available tokens, type *\$.* The list of available values will be displayed.

The list of available tokens includes the API's formal arguments, and all the design attributes defined for the various elements.

Note

Design attributes have element-level scope (i.e, they are relevant only for the element for which they are defined). Therefore, if a design attribute from a different element is used in an API definition, its value will be an empty string.

Using Parameters for the Generated Code

To represent parameters in the API definitions you provide, you add the prefix "\$<" and the suffix ">" to the parameter name, for example:

API name:

Terminate Task(*nameid*)

API Definition in OSDT:

TerminateThread (t \$<nameid>.hndl , 0);

Code that will be generated for a task named T1:

TerminateThread (t_T1.hndl , 0);

A second way to use parameters in the API definition is to use the design attribute value of the element itself. For example, suppose the element has a design attribute named *Create Mode* that uses the attribute key word *CK_createdMode*, which then evaluates to:

CREATE_SUSPENDED API Name: Create Task(*nameid*) API Definition in OSDT:

```
t $<nameid>. hndl = CreateThread ( NULL ,
0 , ( LPTHREAD_START_ROUTINE )$<nameid> , NULL ,
$<CK_createdMode> , &t_$<nameid>.tid );
```
Code that will be generated for a task named T1:

t T1. hndl = CreateThread (NULL , 0 , (LPTHREAD_START_ROUTINE) T1, NULL , CREATE SUSPENDED, &t T1.tid);

A third way to use parameters in the API definition is to use the property value of the element as the API definition. For example, suppose the element has a design attribute, possibly hidden, that uses the attribute key word CK_sendMessagesAPI

This evaluates to:

mySendMessage(\$<*nameid*>,...)

For the following API definition:

API Name:

Send Message(*nameid*)

API Definition in OSDT:

\$<<CK_sendMessagesAPI>>

and design attribute definition:

mySendMessage(\$<*nameid*>, \$<CK_MessagePriority>);

The resulting generated code, for a data item named *DI1*, will be:

mySendMessage(DI1, 1);

assuming that the CK_MessagePriority property evaluates to 1.

Conditional Expressions in API Definitions

The OSDT also allows the use of conditional expressions in API definitions. This feature allows the inclusion of a certain string if the condition is met and the inclusion of an alternative string if the condition is not met. Basically, the feature mimics the C conditional expression, "? :", although the syntax is slightly different.

The basic syntax is as follows:

```
? <br/>begin> expression 1 ? <?> expression 2 ? <<<<<<<<>>
expression 3 ?<<r/>ead>
```
If expression 1 evaluates to true, then expression 2 will be used in the API definition; otherwise, expression 3 will be used.

Example of using conditional expressions:

```
?<begin> $<prop1> ?<==> prop1val ?<?> expression when yes ?<:>
expression when no ?<end>
```
In the above example, the string used in the API definition will be expression when yes if **\$<prop1>** evaluates to prop1val. Otherwise, the string used will be expression when no.

When defining the condition (expression 1), the following symbols can be used:

 $? \leq == \leq$ (equal strings) $? \leq ! \leq$ (not equal strings) $? < \&\&>$ (logical AND) $?$ < $|$ | > (logical OR)

Expression 2 and expression 3 referred to above can consist of any expression that is legal in the API definition, including additional conditional expressions.

The following is a more complex example, which uses nested conditional expressions.

```
Some prefix code ?<br/>begin> $<prop1> ?<==> prop1val
?<&&> $<prop1.1> ?<==> prop1.1val ?<?> ?<begin>
$<p>prop2> ?<=&gt; prop2val ?<|&gt; $<prop2.1> ?<=&gt;prop2.1val ?<?> exp 1.1 when yes ?<:> exp 1.2 when no
?<end> ?<:> exp 2 when no ?<end> Some postfix code,
then another conditional expression ?<br/>begin> $<prop3>
?<==> prop3val ?<?> \exp 3.1 when yes ?<:> \exp 3.2 when
no ?<end>
```
Start with the inner conditional expression:

```
?<begin> $<prop2> ?<==> prop2val ?<||> $<prop2.1> ?<==> prop2.1val 
?<?> exp 1.1 when yes ?<:> exp 1.2 when no ?<end>
```
This expression will evaluate to exp 1.1 when yes if either γ prop²> evaluates to "prop2val" or **\$<prop2.1>** evaluates to "prop2.1val". If neither of these conditions are met, then the expression will evaluate to exp 1.2 when no.

Now look at the outer conditional expression, replacing the result of the inner expression with the string "result of inner conditional expression":

```
?<br/>begin> $<prop1> ?<==> prop1val ?<&&> $<prop1.1< ?<==>
prop1.1val ?<?> result of inner conditional expression ?<:> exp 2 when no 
?<end>
```
This expression will evaluate to the result of the inner conditional expression if $\mathcal{S}\leq \mathbf{prop1}$ evaluates to prop1val and $\overline{\ast}$ **corp1.1**> evaluates to prop1.1val. Otherwise, the expression will evaluate to exp 2 when no.

So, assuming that:

 S-prop1 = prop1val $\text{S
1.1>}= \text{prop1.1val}$ \S <prop2> \le prop2val $\text{S-prop2.1}> \text{prop2.1val}$ \$<prop3> <> prop3val

The API result will be:

Some prefix code exp 1.2 when no Some postfix code, then another conditional expression exp 3.2 when no

General API Definitions

The OSDT allows you to define how code should be generated for the various model elements. These APIs represent the code that will be generated when you use specific elements in your charts.

For example, if you add an activity to an activity diagram, and select "Task" as the type, the code that is generated for activities related to this task will be determined by the APIs you have defined.

In the case of a Task, you can define the code for activities such as:

- creating the Task
- activating the Task
- destroying the Task

Using the OSDT, you can define the code to generate in connection with the following categories:

- OS data types
- \bullet timeouts
- \bullet Tasks
- \bullet events
- software counters
- \bullet timers
- synchronization
- critical sections
- messages
- interrupts
- get-set functions
- \bullet queues
- scheduler

In the tables containing the API details, the generated code is based on the definition that appears in the column *Sample Definition*, unless specified otherwise.

OS Data Type APIs

Timeout APIs

Task APIs

Event APIs

Software Counter APIs

Timer APIs

Synchronization APIs

Critical Section APIs

Message APIs

Interrupt APIs

Scheduler Definition APIs

A scheduler file is a file that includes the Tasks defined in the application.

If this option is selected, the code generator will look for the specified scheduler file, and insert the list of Tasks to be performed.

File Name

By default, the tasks are added to the file <Profile Name>.c. If you want the generator to use a different file, enter the name of the file in the text box or use the **...** button to select the file.

- **Scheduler Key Words**
	- These are the words that demarcate the beginning and end of the task list. The code generator requires the end keyword as well because it removes the tasks that were previously included in the list.
	- The keywords used can include tokens enclosed with "\$<" and ">" (API definition notation).
	- When using "\$<token>" as part of the keyword, the name of the keyword can vary between groups of tasks, depending on the data of the task. The tokens that can be used are:
		- \$<nameid> The name of the task
		- \$<Design-Attribute-Name> Any design attribute defined for the task.
	- For example:

If a task has a design attribute named "CK_timeSlice" with three possible values ("10ms", "40ms" and "100ms"), and the begin and end keywords are defined as */* User \$<CK_timeSlice> Tasks Begin */*, and */* User \$<CK_timeSlice> Tasks End */respectively, then tasks with CK timeSlice = 40ms will be put in the* scheduler file between the keywords */* User 40msTasks Begin */* and the keyword */* User 40ms Tasks End */*

Task Separator

The delimiter to use to separate the individual tasks in the list. When using a custom delimiter, you can use " \sqrt{n} " to specify a new line.

Get-Set Function APIs

Queue APIs

Internal Data Types APIs

Customizing Code Style

The OSDT allows you to modify a large number of settings that affect the code style. These settings are accessed by clicking **Code Style...** in the main OSDT window, and they are categorized as follows:

- ◆ Code Style
- Types Naming Style
- Variables Naming Style
- Model Data Naming Style
- Functions Naming Style
- File Header/Footer

Code Style

This category includes the following settings, which affect the visual appearance of the code:

- \bullet Indent Size
- New Line and Brace Style

These settings determine whether the opening brace appears on the same line as the preceding code or on a new line. There are separate settings for:

- Function Blocks
- Code Blocks
- Type Blocks

Types Naming Style

Variables Naming Style

Model Data Naming Style

Functions Naming Style

File Header/Footer

Customizing Memory Management

The OSDT allows you to modify settings that affect memory management. These settings are accessed by clicking **Memory Management...** in the main OSDT window. (This option will be grayed out unless the **Memory Management** check box is selected.) These settings are categorized as follows:

- Data—Variable Declaration
- Data—Declaration Section
- Code—Task/ISR and Related Activities
- Code—Activities Definition Section
- ◆ Code—Per-User Function
- Code—User Functions Definition Section

Advanced: Creating Customized OSIs

Code—Activities Definition Section

Code—Per-User Function

Advanced: Creating Customized OSIs

Code—User Functions Definition Section

Customizing the Static OS Configuration

The OSDT allows you to modify a large number of static OS configuration settings. These settings are accessed by clicking **Static OS Configuration...** in the main OSDT window. (This option will be grayed out unless the **Static OS Configuration** check box is selected.) These settings are categorized as follows:

- Task Definition
- Event Definition
- **+** Timer Definition
- Synchronization Definition
- Critical Section Definition
- Message Definition
- ISR Definition

Where Definition is Used, Code Generated

For OSIs that use a static OS configuration file, you provide the names of the following two files, when defining the profile to use for code generation:

- *OS CFG Input*
	- The template file to use for the creation of the OS configuration file.
	- The keywords used in the template file will be replaced with concrete data from the model to create the OS configuration file that reflects the OS objects in the model.
- *OS CFG Output*

The name to use for the generated static OS configuration file.

The task, event, timer, synchronization, critical section, message, and ISR definitions listed in the tables in this section are used for building the file used for generation of the static OS configuration file.

The following code will be generated for an activity named T1, defined to be a task:, and the input file given below.

Input File:

… task_definition_keyword_1 … task_definition_keyword_6 semaphore_definition_keyword_9 … event_definition_keyword_4

```
...
```
Generated Code:

```
…
task_definition_1
…
task definition 6
semaphore_definition_9
…
event_definition_4
...
```
Task Definition

Event Definition

Timer Definition

Synchronization Definition

Critical Section Definition

Message Definition

ISR Definition

OS Definition

Specifying Related Files

The main OSDT screen contains a button labeled **Related Files...**. When this button is pressed, a list of files is displayed. These files are files that you will want to have available in your workarea, for example, makefiles, .h files, and .oil files. The specific files listed will depend upon the OSI you are using.

When you create a new workarea, these files are copied to the workarea's *prt* directory.

To remove a file from the files that will be copied to this directory, clear the check box next to the file. To add a file, click the **Add File** button.

Upgrading an OSI

The OSDT is capable of upgrading OSIs that were created with previous versions of the tool.

To upgrade an existing OSI, select **File** > **Update from OSI...** from the main menu.

The upgrade operation carries out the following actions:

- Checks for API definitions that exist in the reference OSI but not in the OSI being upgraded. You will be prompted to approve these additions.
- Checks for obsolete API definitions—those that exist in the OSI being upgraded but not in the reference OSI. You will be prompted to approve the removal of these obsolete definitions.
- Checks the list of Related Files for new, modified, or obsolete files.
- Checks for design attributes that exist in the reference OSI but not in the OSI being upgraded. You will be prompted to approve these additions.
- Checks for obsolete design attributes—those that exist in the OSI being upgraded but not in the reference OSI. You will be prompted to approve the removal of these obsolete design attributes.

All changes made during the upgrade operation are recorded in a log file located in the ctd directory (file is named <OSI NAME>DD/MM/YY.txt).

When the upgrade operation is completed, the OSI will be marked as "modified." The changes will only take effect when the OSI is saved.

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