

IBM Storage Protect
8.1.27

Optimizing Performance



Note:

Before you use this information and the product it supports, read the information in [“Notices” on page 251](#).

Edition notice

This edition applies to version 8, release 1, modification 27 of IBM Storage Protect products (product numbers 5725-W98, 5725-W99, 5725-X15), and to all subsequent releases and modifications until otherwise indicated in new editions.

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About this publication

This information helps you to optimize the performance of IBM Storage Protect servers and clients, and identify and solve performance problems.

Standard subscription and support services from IBM do not include extensive performance analysis and tuning. The comprehensive analysis of a performance problem is a billable service that is offered to IBM Storage Protect customers. For more information, see the IBM® Software Support Handbook at <http://www.ibm.com/support/customer/sas/f/handbook/home.html>.

Who should read this guide

The guide is intended for administrators who want to improve the performance of IBM Storage Protect servers and clients.

Before using this information, ensure that you are familiar with your IBM Storage Protect solution:

- How the IBM Storage Protect servers and clients are used and monitored
- The operating systems on which your IBM Storage Protect servers and clients run
- The networks that are in use for IBM Storage Protect server and client operations
- The storage devices that are used for IBM Storage Protect operations

Publications

The IBM Storage Protect product family includes IBM Storage Protect Plus, IBM Storage Protect for Virtual Environments, IBM Storage Protect for Databases, and several other storage management products from IBM.

To view IBM product documentation, see [IBM Documentation](#).

Part 1. Where to start

This information helps you to optimize the performance of IBM Storage Protect servers and clients, and identify and solve performance problems.

Where you start in this information depends on what your goal is:

- If you are installing or upgrading a new server and clients, start with [Part 2, “Configuration best practices,” on page 7.](#)
- If you need to investigate performance degradation, start with [Part 3, “Solving performance problems,” on page 57.](#)

Before using this information, ensure that you are familiar with your IBM Storage Protect solution:

- How the IBM Storage Protect servers and clients are used and monitored
- The operating systems on which your IBM Storage Protect servers and clients run
- The networks that are in use for IBM Storage Protect server and client operations
- The storage devices that are used for IBM Storage Protect operations

Standard subscription and support services from IBM do not include extensive performance analysis and tuning. The comprehensive analysis of a performance problem is a billable service that is offered to IBM Storage Protect customers. For more information, see the IBM Software Support Handbook at <http://www.ibm.com/support/customercare/sas/f/handbook/home.html>.

Chapter 1. Operating system and performance information cross-reference

Most of the performance information applies to any client or server, on any operating system. For certain operating systems, specific information about configuration of clients and servers for performance is available.

Table 1. Topics for IBM Storage Protect servers, by operating system		
Server operating system	Key topics	Topics specifically for the operating system
AIX®	Chapter 2, “Configuring the server for optimal performance,” on page 9 Chapter 4, “Monitoring and maintaining the environment for performance,” on page 49 Chapter 8, “Identifying performance bottlenecks,” on page 73 Chapter 11, “Tuning server performance,” on page 127 Chapter 12, “Tuning disk storage for the server,” on page 177 Chapter 14, “Tuning network performance,” on page 237	“Tuning AIX systems for IBM Storage Protect server performance” on page 173 “Configuring AIX systems for disk performance” on page 187 “Monitoring performance with operating system tools” on page 50 “Setting network options for IBM Storage Protect on AIX systems” on page 238
Linux®	Chapter 2, “Configuring the server for optimal performance,” on page 9 Chapter 4, “Monitoring and maintaining the environment for performance,” on page 49 Chapter 8, “Identifying performance bottlenecks,” on page 73 Chapter 11, “Tuning server performance,” on page 127 Chapter 12, “Tuning disk storage for the server,” on page 177 Chapter 14, “Tuning network performance,” on page 237	“Tuning Linux systems for IBM Storage Protect server performance” on page 174 “Tuning Linux on System z systems for IBM Storage Protect server performance” on page 175 “Configuring Linux systems for disk performance” on page 188 “Monitoring performance with operating system tools” on page 50

Table 1. Topics for IBM Storage Protect servers, by operating system (continued)

Server operating system	Key topics	Topics specifically for the operating system
Windows	<p>Chapter 2, “Configuring the server for optimal performance,” on page 9</p> <p>Chapter 4, “Monitoring and maintaining the environment for performance,” on page 49</p> <p>Chapter 8, “Identifying performance bottlenecks,” on page 73</p> <p>Chapter 11, “Tuning server performance,” on page 127</p> <p>Chapter 12, “Tuning disk storage for the server,” on page 177</p> <p>Chapter 14, “Tuning network performance,” on page 237</p>	<p>“Tuning Windows systems for IBM Storage Protect server performance” on page 175</p> <p>“Monitoring performance with operating system tools” on page 50</p>

Table 2. Topics for IBM Storage Protect clients, by operating system

Client operating system or environment	Key topics	Topics specifically for the operating system
AIX	<p>Chapter 3, “Configuring clients for optimal performance,” on page 47</p> <p>Chapter 8, “Identifying performance bottlenecks,” on page 73</p> <p>Chapter 13, “Tuning client performance,” on page 189</p> <p>Chapter 14, “Tuning network performance,” on page 237</p>	<p>“Tuning journal-based backups” on page 221</p> <p>“File space tuning” on page 228</p> <p>“Tuning for IBM Storage Protect for Space Management” on page 243</p>
Linux	<p>Chapter 3, “Configuring clients for optimal performance,” on page 47</p> <p>Chapter 8, “Identifying performance bottlenecks,” on page 73</p> <p>Chapter 13, “Tuning client performance,” on page 189</p> <p>Chapter 14, “Tuning network performance,” on page 237</p>	<p>“Tuning journal-based backups” on page 221</p> <p>“File space tuning” on page 228</p> <p>“Tuning for IBM Storage Protect for Space Management” on page 243</p>

<i>Table 2. Topics for IBM Storage Protect clients, by operating system (continued)</i>		
Client operating system or environment	Key topics	Topics specifically for the operating system
Mac OS X	Chapter 3, “Configuring clients for optimal performance,” on page 47 Chapter 8, “Identifying performance bottlenecks,” on page 73 Chapter 13, “Tuning client performance,” on page 189 Chapter 14, “Tuning network performance,” on page 237	
Oracle Solaris	Chapter 3, “Configuring clients for optimal performance,” on page 47 Chapter 8, “Identifying performance bottlenecks,” on page 73 Chapter 13, “Tuning client performance,” on page 189 Chapter 14, “Tuning network performance,” on page 237	“File space tuning” on page 228 “Tuning for IBM Storage Protect for Space Management” on page 243
VMware	“Tuning virtual machine backup operations” on page 230 “Resolving common performance problems with virtual machine backup operations” on page 203	
Windows	Chapter 3, “Configuring clients for optimal performance,” on page 47 Chapter 8, “Identifying performance bottlenecks,” on page 73 Chapter 13, “Tuning client performance,” on page 189 Chapter 14, “Tuning network performance,” on page 237	“Tuning journal-based backups” on page 221 “Windows system state backups” on page 229

Resources for operating system information

User groups and other sites can be good sources of information for tuning and solving problems for your operating system. The following list gives some examples.

AIX

Search for performance management and tuning information for AIX in the [AIX product information](#).

Windows

Search for performance information for Windows hardware at <http://msdn.microsoft.com/windows/hardware>.

Part 2. Configuration best practices

Typically, hardware configuration and selection have the most significant effect on the performance of an IBM Storage Protect solution. Other factors that affect performance are the operating system selection and configuration, and the configuration of IBM Storage Protect.

Procedure

- The following best practices are the most important for optimal performance and problem prevention.
- Review the table to determine the best practices that apply to your environment.

Best practice	More information
Use fast disks for the server database. Enterprise-grade solid-state disks (SSD), with Fibre Channel or SAS interface, offer the best performance.	Use fast, low-latency disks for the database. Using SSD is essential if you are using data deduplication and data replication. Avoid Serial Advanced Technology Attachment (SATA) and Parallel Advanced Technology Attachment (PATA) disks. For details and more tips, see the following topics: <ul style="list-style-type: none">– “Checklist for server database disks” on page 14– Choosing the correct type of storage technology
Ensure that the server system has enough memory.	Review operating system requirements in technote 84861 . Heavier workloads require more than the minimum requirements. Advanced features such as data deduplication and data replication can require more than the minimum memory that is specified in the system requirements document. If you plan to run multiple instances, each instance requires the memory that is listed for one server. Multiply the memory for one server by the number of instances that are planned for the system.
Separate the server database, the active log, the archive log, and disk storage pools from each other.	Keep all IBM Storage Protect storage resources on separate disks. Keep storage pool disks separate from the disks for the server database and logs. Storage pool operations can interfere with database operations when both are on the same disks. Ideally, the server database and logs are also separated from each other. For details and more tips, see the following topics: <ul style="list-style-type: none">– “Checklist for server database disks” on page 14– “Checklist for server recovery log disks” on page 16– “Checklist for storage pools on DISK or FILE” on page 27
Use at least four directories for the server database. For larger servers or servers that use advanced features, use eight directories.	Place each directory on a LUN that is isolated from other LUNs and from other applications. A server is considered to be large if its database is larger than 2 TB or is expected to grow to that size. Use eight directories for such servers. See “Checklist for server database disks” on page 14 .

Best practice	More information
If you are using data deduplication, data replication, or both, follow the guidelines for database configuration and other items.	<p>Configure the server database according to the guidelines, because the database is extremely important to how well the server runs when these features are being used. For details and more tips, see the following topics:</p> <ul style="list-style-type: none"> – “Checklist for container storage pools” on page 18 – “Checklist for data deduplication” on page 35 – “Checklist for implementing data replication” on page 43
For storage pools that use FILE type device classes, follow the guidelines for the size of storage pool volumes. Typically, 50 GB volumes are best.	<p>Review the information in “Optimal number and size of volumes for storage pools that use disk” on page 142 to help you to determine volume size.</p> <p>Configure storage pool devices and file systems based on throughput requirements, not only on capacity requirements.</p> <p>Isolate the storage devices that are used by IBM Storage Protect from other applications that have high I/O, and ensure that there is enough throughput to that storage.</p> <p>For more details, see “Checklist for storage pools on DISK or FILE” on page 27.</p>
Schedule IBM Storage Protect client operations and server maintenance activities to avoid or minimize overlap of operations.	<p>For more details, see the following topics:</p> <ul style="list-style-type: none"> – “Tuning the schedule for daily operations” on page 149 – “Checklist for IBM Storage Protect server configuration” on page 31
Monitor operations constantly.	<p>By monitoring, you can find problems early and more easily identify causes. Keep records of monitoring reports for up to a year to help you identify trends and plan for growth. See Chapter 4, “Monitoring and maintaining the environment for performance,” on page 49.</p>

Related concepts

[Samples of data protection solutions](#)

Samples of data protection solutions that use IBM Storage Protect for selected scenarios are available on the Service Management Connect wiki. The samples describe specific hardware and software configurations and give performance measurements that were obtained in IBM test labs.

Chapter 2. Configuring the server for optimal performance

Evaluate the characteristics and configuration of the system where the server is installed to ensure that the server is set up for good performance.

Before you begin

First, review the most recent, [basic requirements for a server](#). Then, review the following information for more details.

Procedure

1. Review the [“Checklist for the server hardware and the operating system”](#) on page 10. Correct items as needed.
2. Review the [“Checklist for server database disks”](#) on page 14. Correct items as needed.
3. Review the [“Checklist for server recovery log disks”](#) on page 16. This checklist covers the active log, archive log, and other logs. Correct items as needed.
4. Review the [“Checklist for container storage pools”](#) on page 18. Correct items as needed.
5. Review the [“Checklist for storage pools on DISK or FILE”](#) on page 27. Correct items as needed.
6. If you are getting new storage, test the storage system before you implement it. You can use tools to evaluate the characteristics of storage systems before you use them for the IBM Storage Protect database or storage pools. For more information, see [“Analyzing the basic performance of disk systems”](#) on page 82.
7. Review the tips for disk systems on specific operating systems. Operating systems can require different techniques for optimizing disk operations. For details, see [“Configuring the operating system for disk performance”](#) on page 187.
8. Review the [“Checklist for IBM Storage Protect server configuration”](#) on page 31 for tips about configuring schedules and other operations.
9. If you use data deduplication, review the [“Checklist for data deduplication”](#) on page 35.
10. If you use data replication, review the [“Checklist for implementing data replication”](#) on page 43.

Related tasks

[Grouping data by using collocation in server storage pools](#)

[Use collocation to improve IBM Storage Protect performance and maintain optimal data organization.](#)

[Tuning the schedule for daily operations](#)

Typically, backup operations must be completed every day for all clients. Certain server maintenance processes must also run every day. Ensuring that resources for these critical operations are available when needed requires planning and tuning.

Checklist for the server hardware and the operating system

Use the checklist to verify that the system where the server is installed meets requirements for hardware and software configuration.

Question	Tasks, characteristics, options, or settings	More information
<p>Does the operating system and hardware meet or exceed requirements?</p> <ul style="list-style-type: none"> • Number and speed of processors • System memory • Supported operating system level 	<p>If you are using the minimum required amount of memory, you can support a minimal workload.</p> <p>You can experiment by adding more system memory to determine whether the performance is improved. Then, decide whether you want to keep the system memory dedicated to the server. Test the memory variations by using the entire daily cycle of the server workload.</p> <p>If you run multiple servers on the system, add the requirements for each server to get the requirements for the system.</p> <p>Restriction: AIX Do not use Active Memory Expansion (AME). When you use AME, IBM Db2 software uses 4 KB pages instead of 64 KB pages. Each 4 KB page must be decompressed when accessed, and compressed when not needed. When the compression or decompression occurs, Db2 and the server wait for access to the page, which degrades the server performance.</p>	<p>Review operating system requirements at technote 84861.</p> <p>Additionally, review the guidance in Tuning tasks for operating systems and other applications.</p> <p>For more information about requirements when these features are in use, see the following topics:</p> <ul style="list-style-type: none"> • Checklist for data deduplication • Checklist for node replication <p>To help you determine whether processor or memory characteristics are the cause of performance problems, see Identifying server performance problems.</p> <p>For more information about sizing requirements for the server and storage, see the IBM Storage Protect Blueprint.</p>

Question	Tasks, characteristics, options, or settings	More information
Are disks configured for optimal performance?	The amount of tuning that can be done for different disk systems varies. Ensure that the appropriate queue depths and other disk system options are set.	<p>For more information, see the following topics:</p> <ul style="list-style-type: none"> • Checklist for server database disks • Checklist for server recovery log disks • Checklist for storage pools on DISK or FILE
Does the server have enough memory?	<p>Heavier workloads and advanced features such as data deduplication and data replication require more than the minimum system memory that is specified in the system requirements document.</p> <p>For databases that are not enabled for data deduplication, use the following guidelines to specify memory requirements:</p> <ul style="list-style-type: none"> • For daily ingest of data of less than 1 TB, you need 24 GB of memory. • For daily ingest of data of 1 TB - 10 TB, you need 64 GB of memory. • For daily ingest of data of 10 TB - 30 TB, you need 192 GB of memory. • For daily ingest of data up to 100 TB, you need 384 GB of memory. <p>Ensure that you allocate extra space for the active log and the archive log for replication processing.</p>	<p>For more information about requirements when these features are in use, see the following topics:</p> <ul style="list-style-type: none"> • Checklist for data deduplication • Checklist for node replication • Memory requirements <p>To help you determine whether processor or memory characteristics are the cause of performance problems, see Identifying server performance problems.</p>

Question	Tasks, characteristics, options, or settings	More information
Does the system have enough host bus adapters (HBAs) to handle the data operations that the IBM Storage Protect server must run simultaneously?	<p>Understand what operations require use of HBAs at the same time.</p> <p>For example, a server must store 1 GB/sec of backup data while also doing storage pool migration that requires 0.5 GB/sec capacity to complete. The HBAs must be able to handle all of the data at the speed required.</p>	See Tuning HBA capacity .
Is network bandwidth greater than the planned maximum throughput for backups?	<p>Network bandwidth must allow the system to complete operations such as backups in the time that is allowed or that meets service level commitments.</p> <p>For data replication, network bandwidth must be greater than the planned maximum throughput.</p>	<p>For more information, see the following topics:</p> <ul style="list-style-type: none"> • Tuning network performance • Checklist for node replication
Are you using a preferred file system for IBM Storage Protect server files?	Use a file system that ensures optimal performance and data availability. The server uses direct I/O with file systems that support the feature. Using direct I/O can improve throughput and reduce processor use. For more information about the preferred file system for your operating system, see IBM Storage Protect server-supported file systems .	For more information, see Configuring the operating system for disk performance .

Question	Tasks, characteristics, options, or settings	More information
Did you configure enough paging space?	<p>Paging space, or swap space, extends the memory that is available for processing. When the amount of free RAM in the system is low, programs or data that is not in use are moved from memory to paging space. This action releases memory for other activities, such as database operations.</p> <p>Restriction: Do not use paging space to add memory to your system. Paging space is intended to provide only a limited and temporary extension of space. If your system uses paging space, system memory is full and must be extended.</p> <p>AIX Linux</p> <p>Use a minimum of 32 GB of paging space or 50% of your RAM, whichever value is larger.</p> <p>Windows Paging space is automatically configured.</p>	
<p>Linux</p> <p>Linux Did you tune kernel parameters after installation of the server?</p>	You must tune kernel parameters.	See the information about tuning kernel parameters: Tuning kernel parameters for Linux systems .

Checklist for server database disks

Use the checklist to verify that the system where the server is installed meets requirements for hardware and software configuration.

Question	Tasks, characteristics, options, or settings	More information
Is the database on fast, low-latency disks?	<p>Do not use the following drives for the IBM Storage Protect database:</p> <ul style="list-style-type: none">• Nearline SAS (NL-SAS)• Serial Advanced Technology Attachment (SATA)• Parallel Advanced Technology Attachment (PATA) <p>Do not use internal disks that are included by default in most server hardware.</p> <p>Enterprise-grade solid-state disks (SSD), with Fibre Channel or SAS interface, offer the best performance.</p> <p>If you plan to use the data deduplication functions of IBM Storage Protect, focus on disk performance in terms of I/O operations per second (IOPS).</p>	<p>For more information, see Checklist for data deduplicationChecklist for data deduplication.</p>
Is the database stored on disks or LUNs that are separate from disks or LUNs that are used for the active log, archive log, and storage pool volumes?	<p>Separation of the server database from other server components helps reduce contention for the same resources by different operations that must run at the same time.</p> <p>Tip: The database and the active log can share an array when you use solid-state drive (SSD) technology.</p>	
If you are using RAID, did you select the optimal RAID level for your system? Did you define all LUNs with the same size and type of RAID?	<p>When a system must do large numbers of writes, RAID 10 outperforms RAID 5. However, RAID 10 requires more disks than RAID 5 for the same amount of usable storage.</p> <p>If your disk system is RAID, define all your LUNs with the same size and type of RAID. For example, do not mix 4+1 RAID 5 with 4+2 RAID 6.</p>	

Question	Tasks, characteristics, options, or settings	More information
<p>If an option to set the strip size or segment size is available, did you optimize the size when you configured the disk system?</p>	<p>If you can set the strip size or segment size, use 64 KB or 128 KB sizes on disk systems for the database.</p>	<p>The block size that is used for the database varies depending on the table space. Most table spaces use 8 KB blocks, but some use 32 KB blocks.</p>
<p>Did you create at least four directories, also called storage paths, on four separate LUNs for the database?</p> <p>Create one directory per distinct array on the subsystem. If you have fewer than three arrays, create a separate LUN volume within the array.</p>	<p>Heavier workloads and use of some features require more database storage paths than the minimum requirements.</p> <p>Server operations such as data deduplication drive a high number of input/output operations per second (IOPS) for the database. Such operations perform better when the database has more directories.</p> <p>Use the following guidelines to create directories in the server database:</p> <ul style="list-style-type: none"> • For server databases less than 2 TB, you need 4 directories. • For server databases with a size of 2 - 4 TB, you need 8 directories. • For server databases greater than 4 TB, you need 12 directories. <p>Consider planned growth of the system when you determine how many storage paths to create. The server uses the higher number of storage paths more effectively if the storage paths are present when the server is first created.</p> <p>Use the <i>DB2_PARALLEL_IO</i> variable to force parallel I/O to occur on table spaces that have one container, or on table spaces that have containers on more than one physical disk. If you do not set the <i>DB2_PARALLEL_IO</i> variable, I/O parallelism is equal to the number of containers that are used by the table space. For example, if a table space spans four containers, the level of I/O parallelism that is used is 4.</p>	<p>For more information, see the following topics:</p> <ul style="list-style-type: none"> • Checklist for data deduplication • Checklist for node replication <p>For help with forecasting growth when the server deduplicates data, see technote 1596944.</p> <p>For the most recent information about database size, database reorganization, and performance considerations for IBM Storage Protect servers, see technote 1683633.</p> <p>For information about setting the <i>DB2_PARALLEL_IO</i> variable, see Recommended settings for IBM Db2® registry variables.</p>

Question	Tasks, characteristics, options, or settings	More information
Are all directories for the database the same size?	Directories that are all the same size ensure a consistent degree of parallelism for database operations. If one or more directories for the database are smaller than the others, they reduce the potential for optimized parallel prefetching. This guideline also applies if you must add storage paths after the initial configuration of the server.	
Did you raise the queue depth of the database LUNs on AIX systems?	The default queue depth is often too low.	See “Configuring AIX systems for disk performance” on page 187 Configuring AIX systems for disk performance .

Related tasks

[Choosing the correct type of storage technology for IBM Storage Protect](#)

Storage devices have different capacity and performance characteristics. These characteristics affect which devices are better for use with IBM Storage Protect.

Checklist for server recovery log disks

The recovery log for the server consists of the active log, the archive log, and optional logs for mirroring and failover. Use the checklist to verify that disk systems that are being used for the logs have the characteristics and configuration that are key to good performance.

Question	Tasks, characteristics, options, or settings	More information
Are the active log and archive log stored on disks or LUNs that are separate from what is used for the database and storage pool volumes?	Ensure that the disks where you place the active log are not used for other server or system purposes. Do not place the active log on disks that contain the server database, the archive log, or system files such as page or swap space. Note: If capacity limitations forcing the need to use fewer disks or file systems in the environment and workloads cannot be separated, then the following can be combined on the same physical disks or file systems: <ul style="list-style-type: none"> • Db2 Database active log and Database disks together • Db2 Archive log and Storage Protect storage pool disks 	Separation of the server database, active log, and archive log helps to reduce contention for the same resources by different operations that must run at the same time.

Question	Tasks, characteristics, options, or settings	More information
Are the logs on disks that have nonvolatile write cache?	Nonvolatile write cache allows data to be written to the logs as fast as possible. Faster write operations for the logs can improve performance for server operations.	
Did you set the logs to a size that adequately supports the workload?	<p>If you are not sure about the workload, use the largest size that you can.</p> <p>Active log</p> <p>Configure the server to have a maximum active log size that is appropriate for the size of your deployment. For example, for a small to medium sized deployment, you can configure a maximum active log size of 128 GB by setting the ACTIVELOGSIZE server option to a value of 131072. For more information about active log size, go to the IBM Storage Protect Blueprints and locate the Blueprint for your operating system.</p> <p>Ensure that there is at least 8 GB of free space on the active log file system after the fixed size active logs are created.</p> <p>Archive log</p> <p>The size of the archive log is limited by the size of the file system on which it is located, and not by a server option. Make the archive log at least as large as the active log.</p>	<ul style="list-style-type: none"> For information about sizing when you use data deduplication, see Checklist for data deduplication.
Did you define an archive failover log? Did you place this log on a disk that is separate from the archive log?	The archive failover log is for emergency use by the server when the archive log becomes full. Slower disks can be used for the archive failover log.	<p>Use the ARCHFAILOVERLOGDIRECTORY server option to specify the location of the archive failover log.</p> <p>Monitor the usage of the directory for the archive failover log. If the archive failover log must be used by the server, the space for the archive log might not be large enough.</p>

Question	Tasks, characteristics, options, or settings	More information
If you are mirroring the active log, are you using only one type of mirroring?	<p>You can mirror the log by using one of the following methods. Use only one type of mirroring for the log.</p> <ul style="list-style-type: none"> • Use the MIRRORLOGDIRECTORY option that is available for the IBM Storage Protect server to specify a mirror location. • Use software mirroring, such as Logical Volume Manager (LVM) on AIX. • Use mirroring in the disk system hardware. 	<p>If you mirror the active log, ensure that the disks for both the active log and the mirror copy have equal speed and reliability.</p> <p>For more information, see Configuring the recovery log.</p>

Related tasks

Choosing the correct type of storage technology for IBM Storage Protect

Storage devices have different capacity and performance characteristics. These characteristics affect which devices are better for use with IBM Storage Protect.

Checklist for container storage pools

Review how your server is set up to ensure optimal performance when using directory-container and cloud-container storage pools.

Question	Tasks, characteristics, options, or settings	More information
Measured in terms of input/output operations per second (IOPS), are you using fast disk storage for the IBM Storage Protect database?	<p>Use a high-performance disk for the database. Use solid-state drive technology for data deduplication processing.</p> <p>Ensure that the database has a minimum capability of 3000 IOPS. For each TB of data that is backed up daily (before data deduplication and compression), add 1000 IOPS to this minimum.</p> <p>For example:</p> <ul style="list-style-type: none"> • For daily ingest of data of 1 TB, the server needs 4000 IOPS. • For daily ingest of data of 10 TB, the server needs 13000 IOPS. • For daily ingest of data of 10 TB - 30 TB, the server needs 13000 IOPS - 33000 IOPS. • For daily ingest of data of 30 TB - 100 TB, the server needs 33000 IOPS - 100300 IOPS. <div> 3000 IOPS minimum + 30000 (30 TB x 1000 IOPS) = 33000 IOPS </div>	<p>For recommendations about disk selection, see Checklist for server database disks.</p> <p>For more information about IOPS, see the IBM Storage Protect Blueprints and locate the Blueprint for your operating system.</p>

Question	Tasks, characteristics, options, or settings	More information
Do you have enough memory for the size of your database?	<p>Use a minimum of 40 GB of system memory for IBM Storage Protect servers, with a database size of 100 GB, that are deduplicating data. If the retained capacity of backup data grows, the memory requirement might need to be higher.</p> <p>Monitor memory usage regularly to determine whether more memory is required.</p> <p>Use more system memory to improve caching of database pages. The following memory size guidelines are based on the daily amount of new data that you back up:</p> <ul style="list-style-type: none"> • For daily consumption of data up to 1 TB, you need 16 - 24 GB of memory. • For daily ingest of data up to 10 TB, you need 64 GB of memory. • For daily ingest of data up to 10 - 30 TB, you need 192 GB of memory. • For daily ingest of data up to 100 TB, you need 384 GB of memory. 	Memory requirements
Have you properly sized the storage capacity for the database active log and archive log?	<p>Configure the server to have a minimum active log size of 128 GB by setting the ACTIVELOGSIZE server option to a value of 131072.</p> <p>The suggested starting size for the archive log is 1 TB. The size of the archive log is limited by the size of the file system on which it is located, and not by a server option. Ensure that there is at least 10% extra disk space for the file system than the size of the archive log.</p> <p>Use a directory for the database archive logs with an initial free capacity of at least 1 TB. Specify the directory by using the ARCHLOGDIRECTORY server option.</p> <p>Define space for the archive failover log by using the ARCHFAILOVERLOGDIRECTORY server option.</p>	For more information about sizing for your system, see the IBM Storage Protect Blueprints and locate the Blueprint for your operating system.

Question	Tasks, characteristics, options, or settings	More information
Is compression enabled for the archive log and database backups?	<p>Enable the ARCHLOGCOMPRESS server option to save storage space.</p> <p>This compression option is different from inline compression. Inline compression is enabled by default with IBM Storage Protect 7.1.5 and later.</p> <p>Restriction: Do not use this option if the amount of backed up data exceeds 6 TB per day.</p>	For more information about compression for your system, see the IBM Storage Protect Blueprints and locate the Blueprint for your operating system.
<p>Are the IBM Storage Protect database and logs on separate disk volumes (LUNs)?</p> <p>Is the disk that is used for the database configured according to best practices for a transactional database?</p>	The database must not share disk volumes with IBM Storage Protect database logs or storage pools, or with any other application or file system.	For more information about server database and recovery log configuration, see Server database and recovery log configuration and tuning .
Are you using a minimum of eight (2.2 GHz or equivalent) processor cores for each IBM Storage Protect server that you plan to use with data deduplication?	If you are planning to use client-side data deduplication, verify that client systems have adequate resources available during a backup operation to complete data deduplication processing. Use a processor that is at least the minimum equivalent of one 2.2 GHz processor core per backup process with client-side data deduplication.	<ul style="list-style-type: none"> • Data Deduplication FAQ • IBM Storage Protect Blueprints
Did you allocate enough storage space for the database?	<p>For a rough estimate, plan for 100 GB of database storage for every 25 TB of data that is to be protected in deduplicated storage pools. <i>Protected data</i> is the amount of data before data deduplication, including all versions of objects stored.</p> <p>For database backup operations with a large number of small files, where the average size of the file is less than 512 KB, you need more database space. For smaller object sizes, plan on 100 GB of database space for every 10 TB stored.</p> <p>As a best practice, define a new container storage pool exclusively for data deduplication. Data deduplication occurs at the storage-pool level, and all data within a storage pool, except encrypted data, is deduplicated.</p>	The optimal IBM Storage Protect environment is set up by using the IBM Storage Protect Blueprints .

Question	Tasks, characteristics, options, or settings	More information
<p>Have you estimated storage pool capacity to configure enough space for the size of your environment?</p>	<p>You can estimate capacity requirements for a deduplicated storage pool by using the following technique:</p> <ol style="list-style-type: none"> 1. Estimate the base size of the source data. 2. Estimate the daily backup size by using an estimated change and growth rate. 3. Determine retention requirements. 4. Estimate the total amount of source data by factoring in the base size, daily backup size, and retention requirements. 5. Apply the deduplication ratio factor. 6. Apply the compression ratio factor. 7. Round up the estimate to consider transient storage pool usage. 	<p>For an example of using this technique, see Data Deduplication FAQ.</p>

Question	Tasks, characteristics, options, or settings	More information
Have you distributed disk I/O over many disk devices and controllers?	<p>Use arrays that consist of as many disks as possible, which is sometimes referred to as wide striping. Ensure that you use one database directory per distinct array on the subsystem.</p> <p>Set the <i>DB2_PARALLEL_IO</i> registry variable to enable parallel I/O for each table space used if the containers in the table space span multiple physical disks.</p> <p>When I/O bandwidth is available and the files are large, for example 1 MB, the process of finding duplicates can occupy the resources of an entire processor. When files are smaller, other bottlenecks can occur.</p> <p>Use the following guidelines to create file systems:</p> <ul style="list-style-type: none"> • For daily ingest of data less than 10 TB, you need 8 or more file systems. • For daily ingest of data of 10 TB - 30 TB, you need 12 or more file systems. • For daily ingest of data unto 100 TB, you need 32 or more file systems. 	<p>For guidelines about setting up storage pools, see Checklist for storage pools on DISK or FILE</p> <p>For information about setting the <i>DB2_PARALLEL_IO</i> variable, see Recommended settings for IBM DB2® registry variables.</p>
Have you scheduled daily operations based on your backup strategy?	<p>The best practice sequence of operations is in the following order:</p> <ol style="list-style-type: none"> 1. Client backup 2. Storage pool protection 3. Data replication 4. Database backup 5. Expire inventory 	<ul style="list-style-type: none"> • Scheduling data deduplication and node replication processes • Daily operations for directory-container storage pools
Have you scheduled audit operations to identify corrupted files in storage pools?	<p>To schedule audit operations, use the DEFINE STGRULE command and specify the ACTIONTYPE=AUDIT parameter.</p> <p>As a best practice, to ensure that audit operations run continuously, do not specify the DELAY parameter.</p>	

Question	Tasks, characteristics, options, or settings	More information
Do you have enough storage to manage the IBM Db2 lock list?	<p>If you deduplicate data that includes large files or large numbers of files concurrently, the process can result in insufficient storage space. When the lock list storage is insufficient, backup failures, data management process failures, or server outages can occur.</p> <p>File sizes greater than 500 GB that are processed by data deduplication are most likely to deplete storage space. However, if many backup operations use client-side data deduplication, this problem can also occur with smaller-sized files.</p>	For information about tuning the Db2 LOCKLIST parameter, see Tuning server-side data deduplication .
Is sufficient bandwidth available to transfer data to an IBM Storage Protect server?	To transfer data to an IBM Storage Protect server, use client-side or server-side data deduplication and compression to reduce the bandwidth that is required.	For more information, see the enablededup client option.
Have you determined how many storage pool directories to assign to each storage pool?	<p>Assign directories to a storage pool by using the DEFINE STGPOOLDIRECTORY command.</p> <p>Create multiple storage pool directories and ensure that each directory is backed up to a separate disk volume (LUN).</p>	
Did you allocate enough disk space in the cloud-container storage pool?	<p>To prevent backup failures, ensure that the local directory has enough space. Use the following list as a guide for optimal disk space:</p> <ul style="list-style-type: none"> • For serial-attached SCSI (SAS) and spinning disk, calculate the amount of new data that is expected after daily data reduction (compression and data deduplication). Allocate up to 100 percent of that amount, in terabytes, for disk space. • For Flash or SSD based systems with fast network connections to high-performance cloud systems refer to the IBM Storage Protect Blueprints for the latest guidance on disk recommendations. 	

Question	Tasks, characteristics, options, or settings	More information
Have you bench marked the performance of your cloud container pool cloud cache?	To prevent disk hot spots, ensure that the local directory has a single storage pool directory and file system for the cloud cache.	For more information about optimizing backup operations, refer to Sizing a cloud cache to optimize backup operations.
Did you select the appropriate type of local storage?	<p>Ensure that data transfers from local storage to cloud finish before the next backup cycle starts.</p> <p>Tip: Data is removed from local storage soon after it moves to the cloud.</p> <p>Use the following guidelines:</p> <ul style="list-style-type: none"> • Use flash or SSD for large systems that have high-performing cloud systems. Ensure that you have a dedicated 10 GB wide area network (WAN) link with a high-speed connection to the object storage. For example, use flash or SSD if you have a dedicated 10 GB WAN link plus a high-speed connection to either an IBM Cloud Object Storage location or to an Amazon Simple Storage Service (Amazon S3) data center. • Use larger capacity 15000 rpm SAS disks for these scenarios: <ul style="list-style-type: none"> – Medium-sized systems – Slower cloud connections, for example, 1 GB – When you use IBM Cloud Object Storage as your service provider across several regions • For SAS or spinning disk, calculate the amount of new data that is expected after daily data reduction (compression and data deduplication). Allocate up to 100 percent of that amount for disk space, in terabytes. 	

Question	Tasks, characteristics, options, or settings	More information
<p>For cloud-container storage pools, have you specified the total maximum number of parallel processes for the storage tiering rule and each of its subrules?</p>	<p>To specify the maximum number of parallel processes, issue the DEFINE STGRULE command and specify the MAXPROCESS parameter. The default value is 8. For example, if the default value of 8 is specified, and the storage rule has four subrules, the storage rule can run eight parallel processes and each of its subrules can run eight parallel processes.</p> <p>For optimal throughput, use the following maximum number of parallel processes for small, medium, and large Blueprint systems:</p> <ul style="list-style-type: none"> • Small system: 10 processes • Medium system: 25 processes • Large system: 35-50 processes 	
<p>For cloud-container storage pools, have you defined multiple Accesser endpoints if you are using an on-premises IBM Cloud Object Storage system with IBM Storage Protect?</p>	<p>To optimize performance, define access for the following number of Accessers for small, medium, and large blueprint systems, depending on your data ingestion requirements:</p> <ul style="list-style-type: none"> • Small system: 1 Accesser • Medium system: 2 Accessers • Large system: 3-4 Accessers 	<p>For more information, see the IBM Storage Protect Cloud Blueprints.</p>

Question	Tasks, characteristics, options, or settings	More information
<p>For cloud-container storage pools, have you defined multiple Accesser endpoints if you are using an on-premises IBM Cloud Object Storage system with IBM Storage Protect?</p>	<p>Generally, the following Ethernet capability is required to connect to private IBM Cloud Object Storage endpoints for small, medium, and large Blueprint systems:</p> <ul style="list-style-type: none"> • Small system: 1 Gbit • Medium system: 5 Gbit • Large system: 10 Gbit <p>Tip: Depending on client data ingestion and simultaneous data transfer to object storage, you might require more than one 10 Gbit Ethernet network.</p> <p>When you configure the Ethernet connection, work with a network administrator and consider the following factors:</p> <ul style="list-style-type: none"> • The Ethernet capability of the server • The nature of the network between the server and the IBM Cloud Object Storage endpoint • The final ingestion point on object storage via a cloud-container storage pool 	

Checklist for storage pools on DISK or FILE

Use the checklist to review how your disk storage pools are set up. This checklist includes tips for storage pools that use DISK or FILE device classes.

Question	Tasks, characteristics, options, or settings	More information
Can the storage pool LUNs sustain throughput rates for 256 KB sequential reads and writes to adequately handle the workload within the time constraints?	<p>When you are planning for peak loads, consider all the data that you want the server to read or write to the disk storage pools simultaneously. For example, consider the peak flow of data from client backup operations and server data-movement operations such as migration that run at the same time.</p> <p>The IBM Storage Protect server reads and writes to storage pools predominantly in 256 KB blocks.</p> <p>If the disk system includes the capability, configure the disk system for optimal performance with sequential read/write operations rather than random read/write operations.</p>	<p>For more information, see Analyzing the basic performance of disk systems.</p>

Question	Tasks, characteristics, options, or settings	More information
Did you allocate enough storage space for the database?	<p>For a rough estimate, the following database size guidelines are based on the small, medium, and large blueprint systems to allow for database growth:</p> <ul style="list-style-type: none"> • Small system: At least 2 TB • Medium system: At least 4 TB • Large system: At least 8 TB <p>Tip: You might need more memory based on the amount of data that must be protected, the number of files that are stored, and whether you use data deduplication. With data deduplication, the load on the database becomes greater because there are frequent queries to the database to determine what deduplicated extents are on the server.</p> <p>For a rough estimate, plan for 100 GB of database storage for every 50 TB of data that is to be protected in deduplicated storage pools. Protected data is the amount of data before data deduplication, including all versions of objects stored.</p> <p>If you have several hundred TB of protected data, or if you are backing up multiple TBs of data daily, the starting size for the database must be at least 1 TB. Use the IBM Storage Protect to size the database for your system.</p>	<p>The optimal IBM Storage Protect environment is set up by using the IBM Storage Protect Blueprints.</p> <p>For information about the minimum amount of memory you must allocate on the server to complete operations, based on the database size, see Memory requirements</p>
Is the disk configured to use read and write cache?	Use more cache for better performance.	
Do you need to backup the IBM Storage Protect database to cloud object storage?	<p>You can back up a database to, and restore a database from, cloud object storage for disaster recovery purposes.</p> <p>You can tune object storage endpoints, IBM Cloud Object Storage Accessers, network bandwidth, and data streams to ensure that database backup operations run efficiently.</p>	<p>“Tuning database backups to cloud object storage” on page 162.</p>

Question	Tasks, characteristics, options, or settings	More information
For storage pools that use FILE device classes, have you determined a good size to use for the storage pool volumes?	Review the information in “Optimal number and size of volumes for storage pools that use disk” on page 142. If you do not have the information to estimate a size for FILE device class volumes, start with volumes that are 50 GB.	Typically, problems arise more frequently when the volumes are too small. Few problems are reported when volumes are larger than needed. When you determine the volume size to use, as a precaution choose a size that might be larger than necessary.
For storage pools that use FILE device classes, are you using preallocated volumes?	Scratch volumes can cause file system fragmentation. To ensure that a storage pool does not run out of volumes, set the MAXSCRATCH parameter to a value greater than zero.	Use the DEFINE VOLUME server command to preallocate volumes in the storage pool. Use the DEFINE STGPOOL or UPDATE STGPOOL server command to set the MAXSCRATCH parameter.
For storage pools that use FILE device classes, did you compare the maximum number of client sessions to the number of volumes that are defined?	Always maintain enough usable volumes in the storage pools to allow for the expected peak number of client sessions that run at one time. The volumes might be scratch volumes, empty volumes, or partly filled volumes.	For storage pools that use FILE device classes, only one session or process can write to a volume at the same time.
For storage pools that use FILE device classes, have you set the MOUNTLIMIT parameter of the device class to a value that is high enough to account for the number of volumes that might be mounted in parallel?	For storage pools that use data deduplication, the MOUNTLIMIT parameter is typically in the range of 500 - 1000. Set the value for MOUNTLIMIT to the maximum number of mount points that are needed for all active sessions. Consider parameters that affect the maximum number of mount points that are needed: <ul style="list-style-type: none">• The MAXSESSIONS server option, which is the maximum number of IBM Storage Protect sessions that can run concurrently.• The MAXNUMMP parameter, which sets the maximum number of mount points that each client node can use. For example, if the maximum number of client node backup sessions is typically 100 and each of the nodes has MAXNUMMP =2, multiply 100 nodes by the 2 mount points for each node to get the value of 200 for the MOUNTLIMIT parameter.	Use the REGISTER NODE or UPDATE NODE server command to set the MAXNUMMP parameter for client nodes.

Question	Tasks, characteristics, options, or settings	More information
For storage pools that use DISK device classes, have you determined how many storage pool volumes to put on each file system?	<p>How you configure the storage for a storage pool that uses a DISK device class depends on whether you are using RAID for the disk system.</p> <p>If you are not using RAID, then configure one file system per physical disk, and define one storage pool volume for each file system.</p> <p>If you are using RAID 5 with $n + 1$ volumes, configure the storage in one of the following ways:</p> <ul style="list-style-type: none"> • Configure n file systems on the LUN and define one storage pool volume per file system. • Configure one file system and n storage pool volumes for the LUN. 	For an example layout that follows this guideline, see Figure 29 on page 186 .
Did you create your storage pools to distribute I/O across multiple file systems?	<p>Ensure that each file system is on a different LUN on the disk system.</p> <p>Typically, having 10 - 30 file systems is a good goal, but ensure that the file systems are no smaller than approximately 250 GB.</p>	<p>For details, see the following topics:</p> <ul style="list-style-type: none"> • Tuning disk storage for the server • Tuning and configuring storage pools and volumes
Have you scheduled audit operations to identify corrupted files in storage pools?	To schedule audit operations, use the DEFINE SCHEDULE command to run AUDIT VOLUME FIX=NO commands.	

Related tasks

[Tuning and configuring storage pools and volumes](#)

Logical storage pools and storage volumes are the principal components in the IBM Storage Protect model of data storage. By manipulating the properties of these objects, you can optimize the use of storage devices.

[Choosing the correct type of storage technology for IBM Storage Protect](#)

Storage devices have different capacity and performance characteristics. These characteristics affect which devices are better for use with IBM Storage Protect.

Checklist for IBM Storage Protect server configuration

Evaluate key configuration settings and scheduling that can affect performance for the IBM Storage Protect server.

Question	Tasks, characteristics, options, or settings	More information
Have you set up server schedules so that critical operations do not interfere with each other?	<p>Schedule operations that might otherwise be automatically started:</p> <ul style="list-style-type: none">• Disable automatic expiration by setting the EXPINTERVAL server option to 0.• Configure storage pools so that storage-pool migration, reclamation, and duplicate-identification operations are not automatically started.• Schedule each type of server data maintenance task with controlled start times and durations so that they do not overlap with each other. <p>Schedule data replication to avoid or minimize overlap with client backups.</p>	<p>“Tuning the schedule for daily operations” on page 149</p>

Question	Tasks, characteristics, options, or settings	More information
<p>Are you running enough processes to handle data operations in your environment?</p>	<p>Verify that the number of processes for an operation is enough to complete the workload. For example, if performance for reclamation seems slow, tune the number of parallel processes that are allocated for that operation.</p> <p>Use the following commands and parameters to control processes for different operations:</p> <ul style="list-style-type: none"> • Storage pool backup processes: MAXPROCESS parameter on the BACKUP STGPPOOL command • Storage pool parallel processes for the storage rule: MAXPROCESS parameter on the DEFINE STGRULE and UPDATE STGRULE commands • Duplicate identification processes: NUMPROCESS parameter on the IDENTIFY DUPLICATES command. • Migration activity: MIGPROCESS parameter on the DEFINE STGPPOOL command • Parallel expiration activity: RESOURCES parameter on the EXPIRE INVENTORY command • Reclamation processes: RECLAIMPROCESS parameter on the DEFINE STGPPOOL command <p>Continue increasing parallel processes to the point where a resource on the server becomes saturated.</p>	<p>For more information, see the following topics:</p> <ul style="list-style-type: none"> • “Improving the speed of database backups” on page 130 • “Multiple session backup and restore” on page 218

Question	Tasks, characteristics, options, or settings	More information
Are client backup schedules configured so that backups are spread across the available time?	<p>Schedule client backups in a backup window that is isolated from all data maintenance processes, such as storage-pool migration, reclamation, and duplicate identification processing.</p> <p>If possible, schedule client backups so that they are not all started at one time. You might not have to stagger schedules if sufficient server resources are available to process all client backups.</p> <p>Also, if you are using client-side data deduplication and there is commonality in the data that is being backed up, you might not have to stagger schedules.</p>	<p>“Avoiding contention for server resources during client operations” on page 159</p>
Are server option values updated from the defaults for optimum performance?	<p>Set the EXPINTERVAL server option to 0 and schedule inventory expiration processing.</p> <p>If the data is being predominantly stored in the container storage pools, remove the REORGDURATION from dsmserv.opt and restart SP instance to make sure that the reorganization of database tables related to server deduplication can be performed without limit.</p> <p>Set the MAXSESSIONS server option to a value no higher than 1000, which is the maximum that was tested in IBM labs. Setting the value higher than what is required for the maximum number of sessions that you expect can unnecessarily consume memory on the server.</p>	<p>“Limits for the server database size and peak client sessions” on page 69</p>

Question	Tasks, characteristics, options, or settings	More information
<p>Have you set the environment variable <i>MALLOCOPTIONS</i> in the .profile file of the AIX instance user?</p>	<p>In the .profile of the instance user, make sure that the environment variables that follows are set:</p> <ul style="list-style-type: none"> For extra small, small and medium blueprint systems set environment variable <i>MALLOCOPTIONS</i> to: <pre>MALLOCOPTIONS=multiheap:16</pre> <ul style="list-style-type: none"> For larger blueprint systems set environment variable <i>MALLOCOPTIONS</i> to: <pre>MALLOCOPTIONS=multiheap:20</pre>	
<p>Have you set up a schedule for your database backups?</p> <p>Have you configured backup operations appropriately for the size of your database?</p>	<p>When you set up a schedule for database backup, you have more control over when server resources are engaged. Schedule database backups to run after both the client backup, and if used, the storage pool backup.</p> <p>Perform only full database backups, not incremental backups.</p> <p>For databases over 500 GB, use multistreaming for database backups to improve performance.</p> <p>Make the archive log directory for the database large enough so that you do not run out of space between database backups and so that only one database backup is required every 24 hours. Under normal conditions, do not back up the database at unscheduled times.</p>	<p>For more information, see the following topics:</p> <ul style="list-style-type: none"> “Tuning the schedule for daily operations” on page 149 “Improving the speed of database backups” on page 130
<p>Have you formatted disk storage pool volumes sequentially if they are placed on the same file system?</p>	<p>Formatting volumes sequentially helps avoid disk fragmentation and improves sequential read and write performance.</p> <p>To format multiple sequential disk pool volumes, use the DEFINE VOLUME command and specify a value for the NUMBEROFVOLUMES parameter. Each volume is allocated sequentially to avoid fragmentation.</p>	<p>“Checklist for storage pools on DISK or FILE” on page 27</p>

Checklist for data deduplication

Data deduplication requires more processing resources on the server or client. Use the checklist to verify that hardware and your IBM Storage Protect configuration have characteristics that are key to good performance.

Question	Tasks, characteristics, options, or settings	More information
Are you using fast disk storage for the IBM Storage Protect database as measured in terms of input/output operations per second (IOPS)?	<p>Use a high-performance disk for the IBM Storage Protect database. At a minimum, use 10,000 rpm drives for smaller databases that are 200 GB or less. For databases over 500 GB, use 15,000 rpm drives or solid-state drives.</p> <p>Ensure that the IBM Storage Protect database has a minimum capability of 3000 IOPS. For each TB of data that is backed up daily (before data deduplication and compression), include an extra 1000 IOPS to this minimum.</p> <p>For example:</p> <ul style="list-style-type: none">• For daily ingest of data of 1 TB, the server needs 4000 IOPS.• For daily ingest of data of 10 TB, the server needs 13000 IOPS.• For daily ingest of data of 10 TB - 30 TB, the server needs 13000 IOPS - 33000 IOPS.• For daily ingest of data of 30 TB - 100 TB, the server needs 33000 IOPS - 100300 IOPS. <div>$3000 \text{ IOPS minimum} + 30000 (30 \text{ TB} \times 1000 \text{ IOPS}) = 33000 \text{ IOPS}$</div>	<ul style="list-style-type: none">• “Checklist for server database disks” on page 14• For more information about IOPS, go to the IBM Storage Protect Blueprint and locate the Blueprint for your operating system.

Question	Tasks, characteristics, options, or settings	More information
Do you have enough memory for the size of your database?	<p>Use a minimum of 64 GB of system memory for IBM Storage Protect servers that are deduplicating data. If the retained capacity of backup data grows, the memory requirement might need to be higher.</p> <p>Monitor memory usage regularly to determine whether more memory is required.</p> <p>Use more system memory to improve caching of database pages. The following memory size guidelines are based on the daily amount of new data that you back up:</p> <ul style="list-style-type: none"> • For daily ingest of data less than 10 TB, you need 64 GB of memory. • For daily ingest of data of 10 TB - 30 TB, you need to 192 GB of memory. • For daily ingest of data unto 100 TB, you need 384 GB of memory. 	Memory requirements

Question	Tasks, characteristics, options, or settings	More information
<p>Have you properly sized the storage capacity for the database active log and archive log?</p>	<p>The suggested starting size for the active log is 16 GB.</p> <p>Configure the server to have a maximum active log size that is appropriate for the size of your deployment. For example, for a small to medium sized deployment, you can configure a maximum active log size of 128 GB by setting the ACTIVELOGSIZE server option to a value of 131072. For more information about active log size, go to the IBM Storage Protect Blueprints and locate the Blueprint for your operating system.</p> <p>ACTIVELOGSIZE scaling is covered in the Summary of configured elements table.</p> <p>The suggested starting size for the archive log is 48 GB. The size of the archive log is limited by the size of the file system on which it is located, and not by a server option. Make the archive log at least as large as the active log.</p> <p>Use a directory for the database archive logs with an initial free capacity of at least 500 GB. Specify the directory by using the ARCHLOGDIRECTORY server option.</p> <p>Define space for the archive failover log by using the ARCHFAILOVERLOGDIRECTORY server option.</p>	
<p>Are the IBM Storage Protect database and logs on separate disk volumes (LUNs)?</p> <p>Is the disk that is used for the database configured according to best practices for a transactional database?</p>	<p>The database must not share disk volumes with IBM Storage Protect database logs or storage pools, or with any other application or file system.</p>	<p>Server database and recovery log configuration and tuning</p>

Question	Tasks, characteristics, options, or settings	More information
Are you using a minimum of eight (2.2 GHz or equivalent) processor cores for each IBM Storage Protect server that you plan to use with data deduplication?	If you are planning to use client-side data deduplication, verify that client systems have adequate resources available during a backup operation to complete data deduplication processing. Use a processor that is at least the minimum equivalent of one 2.2 GHz processor core per backup process with client-side data deduplication.	Container Storage Pools Best Practices
Have you properly sized disk space for storage pools?	<p>For a rough estimate, plan for 100 GB of database storage for every 10 TB of data that is to be protected in deduplicated storage pools. <i>Protected data</i> is the amount of data before deduplication, including all versions of objects stored.</p> <p>As a best practice, define a new container storage pool exclusively for data deduplication. Data deduplication occurs at the storage-pool level, and all data within a storage pool, except encrypted data, is deduplicated.</p>	Checklist for container storage pools
Have you estimated storage pool capacity to configure enough space for the size of your environment?	<p>You can estimate capacity requirements for a deduplicated storage pool by using the following technique:</p> <ol style="list-style-type: none"> 1. Estimate the base size of the source data. 2. Estimate the daily backup size by using an estimated change and growth rate. 3. Determine retention requirements. 4. Estimate the total amount of source data by factoring in the base size, daily backup size, and retention requirements. 5. Apply the deduplication ratio factor. 6. Apply the compression ratio factor. 7. Round up the estimate to consider transient storage pool usage. 	Container Storage Pools Best Practices

Question	Tasks, characteristics, options, or settings	More information
Have you distributed disk I/O over many disk devices and controllers?	<p>Use arrays that consist of as many disks as possible, which is sometimes referred to as wide striping.</p> <p>When I/O bandwidth is available and the files are large, for example 1 MB, the process of finding duplicates can occupy the resources of an entire processor during a session or process. When files are smaller, other bottlenecks can occur.</p> <p>Use the following guidelines to specify file systems for the deduplicated storage pool device class so that I/O is distributed across as many LUNs and physical devices as possible:</p> <ul style="list-style-type: none"> • For daily ingest of data less than 10 TB, you need 8 or more file systems. • For daily ingest of data of 10 - 30 TB, you need 12 or more file systems. • For daily ingest of data up to 100 TB, you need 32 or more file systems. 	<p>Checklist for storage pools on DISK or FILE.</p> <p>For more information, refer to IBM Storage Protect Blueprint and locate the Blueprint for your operating system.</p>

Question	Tasks, characteristics, options, or settings	More information
<p>Have you scheduled data deduplication processing based on your backup strategy?</p>	<p>If you are not creating a secondary copy of backup data or if you are using data replication for the second copy, client backup and duplicate identification can be overlapped. This overlap can reduce the total elapsed time for these operations, but might increase the time that is required for client backup.</p> <p>If you are using storage pool backup, do not overlap client backup and duplicate identification. The best practice sequence of operations is client backup, storage pool backup, and then duplicate identification.</p> <p>For data that is not stored with client-side data deduplication, schedule storage-pool backup operations to complete before you start data deduplication processing. Set up your schedule this way to avoid reconstructing objects that are deduplicated to make a non-deduplicated copy to a different storage pool.</p> <p>Consider doubling the time that you allow for backups when you use client-side data deduplication in an environment that is not limited by the network.</p> <p>Ensure that you schedule data deduplication before you schedule compression.</p>	<p>Scheduling data deduplication and node replication processes.</p>

Question	Tasks, characteristics, options, or settings	More information
Are the processes for identifying duplicates able to handle all new data that is backed up each day?	<p>If the process completes, or goes into an idle state before the next scheduled operation begins, then all new data is being processed.</p> <p>The duplicate identification (IDENTIFY) processes can increase the workload on the processor and system memory.</p> <p>If you use a container storage pool for data deduplication, duplicate identification processing is not required.</p> <p>If you update an existing storage pool, you can specify 0 - 20 duplicate identification processes to start automatically. If you do not specify any duplicate-identification processes, you must start and stop processes manually.</p>	
Is reclamation able to run to a sufficiently low threshold?	<p>If a low threshold cannot be reached, consider the following actions:</p> <ul style="list-style-type: none"> • Increase the number of processes that are used for reclamation. • Upgrade to faster hardware. 	
Do you have enough storage to manage the Db2 lock list?	<p>If you deduplicate data that includes large files or large numbers of files concurrently, the process can result in insufficient storage space. When the lock list storage is insufficient, backup failures, data management process failures, or server outages can occur.</p> <p>File sizes greater than 500 GB that are processed by data deduplication are most likely to deplete storage space. However, if many backup operations use client-side data deduplication, this problem can also occur with smaller-sized files.</p>	For more information about tuning the Db2 LOCKLIST parameter, see Tuning server-side data deduplication .

Question	Tasks, characteristics, options, or settings	More information
Is deduplication cleanup processing able to clean out the dereferenced extents to free disk space before the start of the next backup cycle?	<p>Run the SHOW DEDUPDELETE command. The output shows that all threads are idle when the workload is complete.</p> <p>If cleanup processing cannot complete, consider the following actions:</p> <ul style="list-style-type: none"> • Increase the number of processes that are used for duplicate identification. • Upgrade to faster hardware. • Determine whether the IBM Storage Protect server is ingesting more data than it can process with data deduplication and consider deploying an extra IBM Storage Protect server. 	
Is sufficient bandwidth available to transfer data to an IBM Storage Protect server?	Use client-side data deduplication and compression to reduce the bandwidth that is required to transfer data to an IBM Storage Protect server.	For more information, see the enablededupcache client option.
Are you deduplicating data from an Oracle or SAP database?	<p>If you are deduplicating data from an Oracle or SAP database, and the average extent size is less than 100 KB, consider specifying an extent size that is larger than the default value of 50 KB.</p> <p>Tip: To determine the average extent size, generate data deduplication statistics by using the GENERATE DEDUPSTATS command and view the statistics by using the QUERY DEDUPSTATS command. Based on the output, calculate the average extent size by using the following formula:</p> <p>Total Protected Data/ (Compressed Extent Count + Uncompressed Extent Count)</p>	For more information about changing the extent size, see the MINIMUMEXTENTSIZE parameter section for the REGISTER NODE and UPDATE NODE commands.

Related tasks

[Evaluating data deduplication results](#)

You can evaluate the effectiveness of IBM Storage Protect data deduplication by examining various queries or reports. Actual data reduction results can show whether the expected storage savings are achieved. You can also evaluate other key operational factors, such as database utilization, to ensure that they are consistent with expectations.

[Tuning server-side data deduplication](#)

Tune settings and configuration for different operations to ensure that the performance of server-side data deduplication is efficient.

Tuning client-side data deduplication

The performance of client-side data deduplication can be affected by processor requirements and deduplication configuration.

Checklist for implementing data replication

A successful implementation of data replication relies on sufficient, dedicated hardware resources. Increased amounts of memory and processor cores are required. The database and its logs must be appropriately sized to ensure that transactions can complete. A dedicated network, with enough bandwidth to handle the amount of data you intend to replicate, is required.

Use the checklist to verify that hardware and your IBM Storage Protect configuration have characteristics that are key to good performance.

Question	Tasks, characteristics, options, or settings	More information
Are you using a high-performance disk for the IBM Storage Protect database?	<p>Ensure that the disks that are configured for the IBM Storage Protect database have a minimum capability of 3,000 I/O operations per second (IOPS). For each TB of data that is backed up daily (before data deduplication), add 1,000 IOPS to this minimum.</p> <p>For example:</p> <ul style="list-style-type: none">• For daily ingest of data of 1 TB, the server needs 4000 IOPS.• For daily ingest of data of 10 TB, the server needs 13000 IOPS.• For daily ingest of data of 10 TB - 30 TB, the server needs 13000 IOPS - 33000 IOPS.• For daily ingest of data of 30 TB - 100 TB, the server needs 33000 IOPS - 100300 IOPS. <div>$3000 \text{ IOPS minimum} + 30000 (30 \text{ TB} \times 1000 \text{ IOPS}) = 33000 \text{ IOPS}$</div>	“Checklist for server database disks” on page 14
Are you using enough processor cores and memory for replication operations and, optionally, data deduplication?	<p>If you are replicating data without deduplication, use a minimum of 6 processor cores and 64 GB of RAM for the source and each target replication servers.</p> <p>For any server that is configured for data replication and data deduplication, use a minimum of 10 processor cores and 128 GB of RAM.</p>	

Question	Tasks, characteristics, options, or settings	More information
Have you properly sized your disk space for the database, logs, and storage pools?	<p>To determine whether your database can handle the additional space requirements, you must first estimate how much more database space data replication uses.</p> <p>For the active log, use a minimum size of 64 GB for data replication. Use the maximum allowed size for the active log, which is 128 GB if you are also using data deduplication.</p> <p>Make the archive log space at least as large as the space defined for the active log. Also, specify a directory for the archive failover login in case it is needed.</p>	Determining server database requirements for node replication (version 7.1.1)
Is your network capable of handling the additional traffic for the amount of data that you intend to replicate between source and target replication servers?	<p>For data replication, network bandwidth must be greater than the planned maximum throughput.</p> <p>You can estimate network bandwidth that is based on the amount of data that you are replicating.</p>	Estimating network bandwidth for node replication (version 7.1.1)
If your IBM Storage Protect server replicates nodes or protects storage pools to a remote server, did you determine whether Aspera® Fast Adaptive Secure Protocol (FASP®) technology can improve data throughput?	<p>Restrictions:</p> <ul style="list-style-type: none"> • Use Aspera FASP technology when your wide area network (WAN) shows signs of high packet loss, data transfer delays that are caused by network impairment, or both. If WAN performance meets your business needs, do not enable Aspera FASP technology. • To enable Aspera FASP technology for replication operations, the data must be stored in a directory-container storage pool. • Aspera FASP technology is available only on Linux x86_64 operating systems. • Before you enable Aspera FASP technology, you must obtain the appropriate licenses. Both evaluation and full licenses are available. 	See Determining whether Aspera FASP technology can optimize data transfer in your system environment.

Question	Tasks, characteristics, options, or settings	More information
Are you using replication storage rules to replicate data?	If you implemented replication by using the REPLICATE NODE command, consider a transition to replication storage rules and subrules, as it can improve the performance of replication operations.	Defining a replication storage rule
Are you using data deduplication with data replication?	Consider using data deduplication if you are not already doing so. By using data deduplication with replication operation, you reduce the bandwidth that is required for replication operations. Data deduplication reduces the amount of data that is sent to the target replication server during replication operation.	Measuring effects of data deduplication on node replication processing (version 7.1.1)
Have you scheduled data replication at the optimum time in the daily schedule?	<p>Before replication: Schedule data deduplication and client backup processing before replication processing. (Should client backup occur before data deduplication? Or does it matter?)</p> <p>After replication: Schedule compression.</p>	<p>For more information, see the following topics:</p> <ul style="list-style-type: none"> • “Scheduling data deduplication and replication processes” on page 152 • “Compatibility and resource usage for server processes” on page 154
Have you optimized the number of sessions that are used for sending data to the target replication server?	<p>You can improve replication performance by using the MAXSESSIONS parameter on replication storage rules or the REPLICATE NODE command to specify data sessions.</p> <p>The number of sessions that are used for replication depends on the amount of data that you are replicating.</p>	Managing the number of replication sessions (version 7.1.1)
Do you have enough mount points to avoid stalled replication servers and other server processes?	<p>Determine the number of logical and physical drives that can be dedicated to the replication process. For example, if a library has 10 tape drives and four of the drives are used for another task, there are six available drives for replication operation.</p> <p>Specify the number of mount points you require and ensure that there are drives available to complete replication operation.</p>	Typically tape is not used for data replication except for the initial replication.

Question	Tasks, characteristics, options, or settings	More information
Does the replication operation completely replicate all newly ingested data before the beginning of the next backup cycle?	<p>If the replication processes cannot finish before the start of the next backup cycle, consider the following actions:</p> <ul style="list-style-type: none"> • Ensure that there are sufficient mount points and drives available for replication processes to complete. • Increase the number of data sessions that are used for replication operation. • Upgrade to faster hardware and more bandwidth for the source and target replication servers. 	
If you are using data deduplication with replication operation, do the processes for identifying duplicates complete before the start of replication operation so that data deduplication is used to its full advantage?	<p>If the process completes, or goes into an idle state before replication operation begins, then all new data is being processed.</p>	

For more information about data replication, see [Defining a replication storage rule](#).

Chapter 3. Configuring clients for optimal performance

You can configure the client system to ensure that the client is set up for good performance.

Procedure

Use the information in the following table to help you configure the client for optimal performance.

Action	More information
Ensure that the client system meets the minimum hardware and software requirements.	For information about client requirements, see Client environment requirements .
Ensure that you are using the appropriate method to back up the data in your environment.	See Selecting the optimal client backup method .
If client options were changed from the default values, note them for further analysis. Some problems can be resolved by setting the client option values to the default values.	For information about client options, see Processing options .
Look for solutions for common client performance problems.	For information about solving common client performance problems, see Common client performance problems .
Fine-tune the client by adjusting the values of client options that affect performance.	See Chapter 13, “ Tuning client performance ,” on page 189

Chapter 4. Monitoring and maintaining the environment for performance

By monitoring server and client operations constantly, you can find problems early and more easily identify causes. Keep records of monitoring reports for up to a year to help you identify trends and plan for growth.

Procedure

- Use the servermon component, which is automatically installed and configured as part of the server installation, to collect data at regular intervals.
- Use monitoring tools to verify that client and server operations are completed successfully and within a time that meets your business requirements.

For monitoring tips, see the [Daily monitoring checklist](#).

- a) Verify that server maintenance processes such as database backup and storage pool backup are completed successfully. Investigate failed or incomplete processes.
 - b) Verify that client backups are successful and complete. Investigate failed or incomplete client backups, especially for clients that are most critical to your overall operations.
- If you are using data deduplication, data replication, or both, verify that the processes that are related to these features are finishing.

For example:

- For server-side data deduplication, verify that the duplicate identification processes can handle all of the new data that the server stores each day. If the processes finish or go to an idle state before subsequent operations such as reclamation begin, you know that the processes can handle the new data.
 - For data replication, verify that replication processes finish replicating all newly stored data before the start of client backups on the next day.
 - If you are using both data deduplication and data replication, verify that the duplicate identification processes can finish before the start of data replication. This sequence ensures that you are taking full advantage of data deduplication.
- Keep your IBM Storage Protect software up to date. Go to the [IBM Support](#) website and search for fix packs that might be useful to apply to your server, clients, or both.
 - Keep other software and hardware products in your environment up to date. For software and hardware products in your environment other than IBM Storage Protect, review service levels and firmware periodically and update them as needed. Completing such a review every six months is a typical goal. However, review and apply security fixes monthly or as needed (for example, on an emergency basis).

Related concepts

[Checklist for data deduplication](#)

Data deduplication requires more processing resources on the server or client. Use the checklist to verify that hardware and your IBM Storage Protect configuration have characteristics that are key to good performance.

[Checklist for implementing data replication](#)

A successful implementation of data replication relies on sufficient, dedicated hardware resources. Increased amounts of memory and processor cores are required. The database and its logs must be appropriately sized to ensure that transactions can complete. A dedicated network, with enough bandwidth to handle the amount of data you intend to replicate, is required.

Monitoring performance with operating system tools

Monitor your IBM Storage Protect solution so that you know when you must investigate performance changes. Operating systems have different tools that are available for monitoring performance. Simulating workloads to test performance is another useful task to learn.

Procedure

- To monitor system processors and storage for usage and the effects of IBM Storage Protect activities, you can use the following commands and tools.

AIX systems

For information about a command, search for the command in the product information for the version of AIX that you are using.

Command or tool	Purpose	More information
iostat command	Statistics about input/output for the entire system and for devices that are attached to the system	
lparstat command	Reports about logical partition (LPAR) configuration and statistics	
nmon command	System monitoring reports	
nstress package of tools	Stress testing of the system	
perfpnr script	A data collection script, typically used before you report a problem to IBM Software Support	Search for information about the script in the product information for the version of AIX that you are using.
sar command	System activity monitoring	
vmstat command	Virtual memory statistics	
Iometer, an open source tool	Measurement and characterization of the input/output capabilities of a system	For information, see www.iometer.org .
Netperf benchmark	Tools to help you measure bandwidth and latency of networks	For information, see the Netperf Homepage .
iperf tool	Measurement of the network performance and evaluation of the different TCPWINDOWSIZE values	For information, see the iperf homepage .

Linux systems

For information about the commands, see the operating system documentation.

Command or tool	Purpose
iostat command	Statistics about input/output for the entire system and for devices that are attached to the system
nmon command	System monitoring reports
sar command	System activity monitoring
Iometer, an open source tool	Measurement and characterization of the input/output capabilities of a system For information, see www.iometer.org .

Command or tool	Purpose
Netperf benchmark	Tools to help you measure bandwidth and latency of networks. For more information, see the Netperf Homepage

Windows systems

Command or tool	Purpose
Windows Performance Monitor (perfmon command)	Performance monitoring of the system and attached devices For information, see the operating system documentation.
Iometer, an open source tool	Measurement and characterization of the input/output capabilities of a system For information, see www.iometer.org .
Netperf benchmark	Tools to help you measure bandwidth and latency of networks. For more information, see the Netperf Homepage

- To help you understand the performance of the IBM Storage Protect environment, consider the following tools. These tools can be useful to determine performance under ideal conditions. However, they simulate only some of the operations that occur during IBM Storage Protect activities.

FTP

FTP is available on almost any system. You can use FTP to estimate the throughput that IBM Storage Protect might get on a backup or restore operation. The results of the test are only an approximation.

To use FTP to evaluate performance, create or use an existing file and use FTP to transfer it. Use a single file larger than 200 MB for the test. If the operation involves smaller files, then the results from a test with FTP might not be accurate.

You might have to manually time how long the transfer takes to calculate throughput. Include these operations in the time estimates:

- Read from disk
- Send over network
- Write to disk

An alternative to FTP is SCP. However, SCP encrypts the data so it might not perform as well as FTP.

dd

The command is available on systems such as AIX and Linux to initiate disk reads or writes.

Related tasks

Analyzing data flow with the dd command

You can use the **dd** command as a quick test to estimate the best possible results for data flow to disks. The command is available for operating systems such as AIX or Linux.

Chapter 5. Sending and restoring data with the object agent service

You can configure the IBM Storage Protect object agent service to provide an Amazon Simple Storage Service (S3) API endpoint that can send data to and restore data from the server.

When you send data to or restore data from a server, the implementation depends on the Amazon S3 client application or the Amazon S3 software development kit (SDK) that is used.

Sending data to the server with the object agent service

When you send data to the IBM Storage Protect server with the object agent service, performance generally improves in relation to the size of the input/output operation.

Guidelines for optimal performance

The following circumstances help to improve performance:

- Client objects are stored by using a single HTTP PUT operation.
- A multipart upload process is used as part of the S3 protocol. A client object is separated into multiple parts that are transferred by using individual HTTP PUT operations, which can occur serially or in parallel.

Table 3 on page 53 lists the guidelines for optimal performance when you send data to the IBM Storage Protect server by using the object agent service.

Table 3. Guidelines for sending data to the server with the object agent service	
Network type	Guidelines for optimal performance
Low-latency networks with low or no data loss, such as a local area network (LAN)	<p>Client objects of less than 100 MB</p> <p>For client objects that are less than 100 MB, send the client object to the server with a single HTTP PUT operation.</p> <p>Client objects of 100 MB or greater</p> <p>For client objects that are 100 MB or greater, use the S3 protocol multipart upload service to send the client object to the server.</p>
Higher-latency networks, such as a wide area network (WAN), for example, the internet	<p>Client objects of 50 MB</p> <p>For client objects that are less than 50 MB, store the client object with a single HTTP PUT operation.</p> <p>Client objects of 50 MB or greater</p> <p>For client objects that are greater than 50 MB, use the S3 protocol multipart upload service to send the client object to the server.</p> <p>For networks where data loss is a concern, use the minimum part size of 5 MB to avoid retransmitting large amounts of data. Retransmission of parts might be necessary for the client application if data loss occurs.</p>

Guidelines for multipart upload

When you use the S3 protocol multipart upload service to send the client object to the server, consider the following service characteristics:

- The minimum part size for multipart upload with the S3 protocol is 5 MB.
- Throughput performance generally increases with larger part sizes up to approximately 50 MB.

Table 4 on page 54 lists guidelines and examples for sending data to the server with the S3 protocol multipart upload service.

Table 4. Guidelines for sending data to the server with the S3 protocol multipart upload service		
Upload characteristics	Guidelines for optimal performance	Examples for multipart upload
<ul style="list-style-type: none">• Network data loss is not a concern.• Multipart upload is done serially.• A single client application thread sends the parts of a client object sequentially.	Use the largest part size less than or equal to 50 MB or 50% of the client object size, whichever is smaller.	Client object size of 100 MB Send client objects in 50 MB parts. Client object size of 50 MB Send client objects in 25 MB parts.
<ul style="list-style-type: none">• Network data loss is not a concern.• Multipart upload is done in parallel.• Multiple threads send the different parts of a front-end object in parallel.	Use smaller part sizes so that more threads can send data at the same time.	For example, when you use the S3 protocol transfer manager service, a thread pool of many threads can be used to transfer object parts in parallel. For a thread pool of 10 threads and an object size of 100 MB, a part size of 10 MB can provide for optimal parallelism across threads. For parallel multipart upload, use a part size greater than or equal to the minimum of 5 MB. In this way, you can more effectively distribute the client object across all of the threads.

Restoring data to the server with the object agent service

When you restore data to the IBM Storage Protect server with the object agent service, performance generally improves in relation to the size of the input/output operation.

Guidelines for optimal performance

Table 5 on page 54 lists the guidelines for optimal performance when you restore data to the IBM Storage Protect server by using the object agent service.

Table 5. Guidelines for restoring data to the server with the object agent service	
Object type	Guidelines for optimal performance
Most objects for which the client application requires the entire object content.	Use a single HTTP GET request to restore the object.

<i>Table 5. Guidelines for restoring data to the server with the object agent service (continued)</i>	
Object type	Guidelines for optimal performance
The object is large, such as several gigabytes or terabytes in size, and network data loss is a concern.	To avoid data retransmissions over the network, use the range read function with a series of sequential HTTP GET requests.
Objects for which only a portion, or range, of the object is needed by the client application.	<p>Use the largest possible HTTP GET requests with the range read function for best performance.</p> <p>Example scenario</p> <p>A 10 MB portion of a 100 MB object is needed by the client application.</p> <p>Example action</p> <p>Use a 10 MB HTTP GET operation with a Range header to read this data.</p>

Part 3. Solving performance problems

When you notice degradation in performance of an IBM Storage Protect solution, start by verifying that conditions outside of the server and client are not the cause. Learn about problem symptoms and causes, and how to use tools to identify them.

Chapter 6. Performance tuning and bottleneck identification

When you tune an IBM Storage Protect solution, you must look at all of its components and their configuration. Performance degradation in key operations, including client backups, storage pool migration, and expiration, can be the result of insufficient resources, poor configuration, or both.

The following variables must be examined:

- Server hardware and its configuration
 - Processors, memory, cache, and storage backplane
 - Internal and external storage resources, including disk systems for the server database, recovery logs, and storage pools

Any hardware component that is in the data path might be the bottleneck. For illustrations of the data paths and the possible components, see [“Potential bottlenecks in the data flow for IBM Storage Protect operations” on page 61.](#)

- The network that is used for communications and data transfers among the components
- Client hardware and its configuration, and the characteristics of the client data that is being protected

The best way to begin performance tuning is to provide sufficient resources for and optimally configure the server and clients. For example, for a server provide enough system memory, disk systems that are properly sized and configured to handle the workload, database and logs that are properly separated, and correct operating system settings. For backup-archive clients, key items include enough memory, adequate network bandwidth, and careful choice and configuration of backup methods.

To find bottlenecks and identify ways to improve performance, you can use both built-in tools for systems and storage devices, and IBM Storage Protect tools.

This performance information provides guidelines for the best performance. Also included are procedures and information about analysis tools to identify performance problems.

Standard subscription and support services from IBM do not include extensive performance analysis and tuning. The comprehensive analysis of a performance problem is a billable service that is offered to IBM Storage Protect customers. For more information, see the IBM Software Support Handbook at <http://www.ibm.com/support/customer/sas/f/handbook/home.html>.

Performance tuning guidelines and expectations

Performance tuning is not a one-time task but an ongoing effort. Because the system environment and client workloads change, you must continually monitor and periodically tune the solution.

Because the performance of an IBM Storage Protect solution can be affected by so many factors, make changes in a controlled way. Evaluate the effects of the changes that you introduce by collecting appropriate measurements before and after changes.

For example, the following approach can be effective:

1. Track performance of the solution over time by establishing an initial baseline of measurements of operational performance. Periodically collect the same measurements and compare the results.
2. Implement a method to track all changes that are made to the IBM Storage Protect solution.
 - Use strict change control to help you understand the performance impact of any change.
 - Limit the changes that you make at one time so that you can more easily determine what makes a difference.
3. After a change and before you make further changes, observe system operations and performance over enough time to evaluate the full effect of the changes.

Observe the system over a time span that is based on typical cycles of operations. For example, if you always have a once-a-week peak in client backup operations, be sure that peak time is included in your observations.

4. Evaluate the results before you make further changes.

Most performance tuning yields limited improvements. Carefully consider how much time is reasonable to spend on improving system performance. Service Level Agreements are an excellent way to set performance goals.

Running a system near its limits can have negative consequences. In such a case, 10 percent more workload might degrade response times dramatically more than an expected 10 percent. In this situation, you must determine which component or process is the bottleneck and eliminate it.

After a system is properly tuned, improving system performance can be achieved only by reducing the workload or adding appropriate resources. You might have to revise your goals and expectations. For significant performance improvements, you must find the bottleneck and then consider one or more of the following actions:

- Use faster processors
- Add processors
- Add system memory
- Use faster communication links

Consider a dedicated backup network, if client backup operations use a LAN that is shared with many other types of operations.

- Add disk storage
- Create a new server on a different system

Symptoms and causes of performance problems

When the performance of an IBM Storage Protect environment is less than expected, there can be one or more causes. Identifying the bottleneck in your environment can explain the performance degradation.

The following symptoms can indicate poor IBM Storage Protect performance:

- Processes or client backups take more time to run than normal
- Slow response times occur for commands that are issued
- Slow response times occur, and the system or process might seem hung
- Unexpected changes occur in response times or resource usage
- Throughput on the system is not as expected
- Processor usage is higher than normal for a certain process
- Network problems occur that are related to load, firewall, or routers

Performance problems can occur when changes are made in the environment. For example, changes to any of the following items can affect performance:

- Hardware configuration: Adding, removing, or changing configurations such as how disks are connected
- Operating system: Installing or updating a file set, installing fix packs, and changing parameters
- Applications: Installing new versions and fixes, configuring or changing data placement, or installing or upgrading device drivers or firmware
- Network: Any changes to the network, packet loss, or intermittent connectivity
- Disk units that are aging, or are damaged
- Options that are used to tune the operating system or an application
- Scheduling of processes or backups during times of high usage
- Unexpected increase in usage of a shared resource like the network or disks

You can collect data on the IBM Storage Protect server, client, or both at the same time to help diagnose where the problem is occurring in the environment and what the problem is.

Potential bottlenecks in the data flow for IBM Storage Protect operations

In operations such as client backup and storage pool migration, data moves through many physical components that can affect the speed of the operations. Understanding the characteristics of these components can help when you are working to improve performance.

Data flow for client backup operations over a LAN

Figure 1 on page 61 shows data flow in a typical configuration for client backup operations over a local area network (LAN). For a client backup operation, the data flow starts at the client disk (item 1 in the graphic and table) and ends at one of the devices for the server storage pools (item 10 or 12).

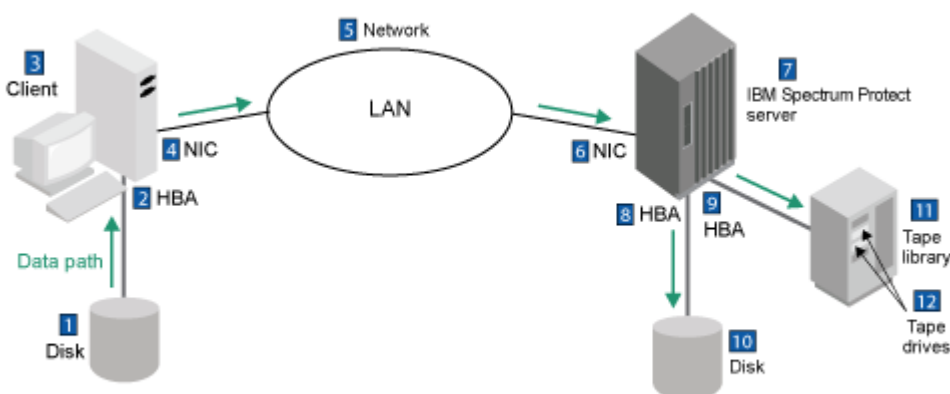


Figure 1. Data flow for client backup operations over a LAN

The data for backup operations flows through many hardware components, any of which are potential bottlenecks. The following table describes characteristics of the hardware that can affect performance.

Item in Figure 1 on page 61	Item	Key characteristics that can affect performance	Details
1	Disk for the client system	Disk type and rotational speed	
2	Host bus adapter (HBA) that connects the disk to the client system	HBA type and its capabilities	

Item in Figure 1 on page 61	Item	Key characteristics that can affect performance	Details
3	Client system	The speed of the I/O bus, the number of processors, the speed of the processors, and the amount and speed of RAM	<p>Use of data compression, data deduplication, and encryption, including Secure Sockets Layer (SSL) protocol, can affect processor performance on the client system. If processor usage is too high on the system, consider adding more processors, or turning off the options that enable compression, data deduplication, or encryption. Then, see whether performance improves.</p> <p>For information about tuning for client memory limits, see “Reduce client memory use” on page 206.</p> <p>Software such as firewalls and antivirus programs might affect the efficiency of client operations. For example, during a restore operation, an antivirus program might scan the contents of each restored object, checking for virus signatures. If you suspect that a firewall or antivirus program is slowing client operations, consider temporarily turning the firewall or antivirus program off to see whether performance improves. For tips to minimize the impact of firewall or antivirus programs on other applications, see the documentation for those programs.</p>
4	Network interface card (NIC) that connects the client system to the LAN	NIC type and its capabilities	<p>A fast network interface card (NIC) improves network throughput. If you cannot use the latest NIC in your configuration, consider adjusting the client TCPWINDOWSIZE option to improve network throughput, particularly on client systems that are geographically distant from the server. Adjust the TCPWINDOWSIZE options in small increments; a window size that is larger than the buffer space on the network interface adapter might actually degrade throughput. For more network considerations, see Chapter 14, “Tuning network performance,” on page 237.</p>

Item in Figure 1 on page 61	Item	Key characteristics that can affect performance	Details
5	Network	The many components on a network, and the effective speed of data transfer over the network, which is limited by its slowest component	
6	NIC that connects the server to the LAN	NIC type and its capabilities	
7	Server system	The speed of the I/O bus, the number of processors, the speed of the processors, and the amount and speed of RAM	
8	HBA that connects the server to the disk	HBA type and its capabilities	See “Tuning HBA capacity” on page 172.
9	HBA that connects the server to the tape library	HBA type and its capabilities	
10	Disk for the server storage pool	Disk type and rotational speed	
11	Tape library for the server storage pool	Number of drives and mount point availability for the operation	
12	Tape drive for the server storage pool	Tape type and sustainable speed	

Data flow for client backup operations over SAN

Figure 2 on page 64 shows data flow in a typical configuration for client backup operations over a storage area network (SAN). Metadata for the backup operation flows over the LAN. For a client backup operation, the data flow starts at the client disk (item 1 in the graphic and table) and ends at one of the devices for the server storage pools (item 11 or 13).

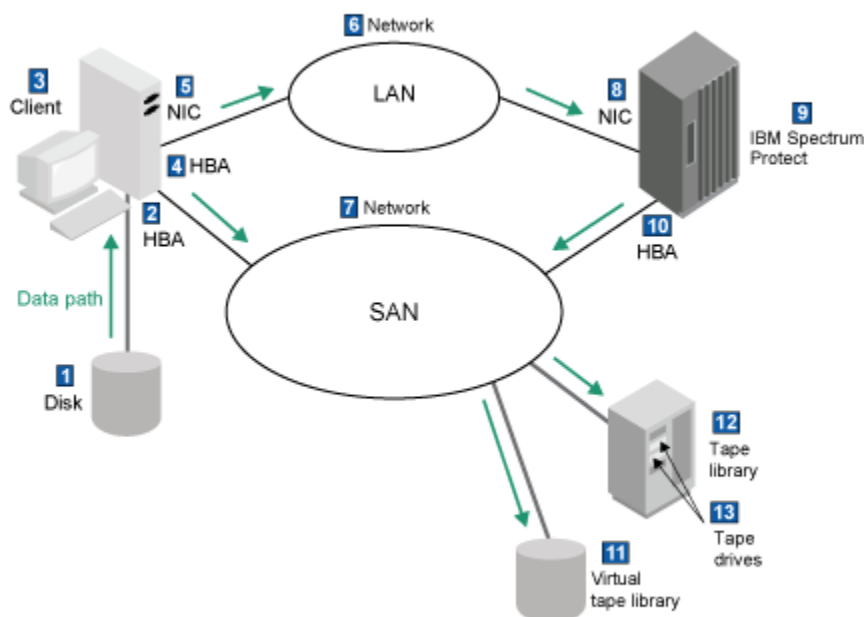


Figure 2. Data flow for client backup operations over the SAN

The data for backup operations flows through many hardware components, any of which are potential bottlenecks. The following table describes characteristics of the hardware that can affect performance.

Item in Figure 2 on page 64	Item	Key characteristics that can affect performance	Details
1	Disk for the client system	Disk type and rotational speed	
2	Host bus adapter (HBA) that connects the disk to the client system	HBA type and its capabilities	

Item in Figure 2 on page 64	Item	Key characteristics that can affect performance	Details
3	Client system	The speed of the I/O bus, the number of processors, the speed of the processors, and the amount and speed of RAM	<p>Use of data compression, data deduplication, and encryption, including Secure Sockets Layer (SSL) protocol, can affect processor performance on the client system. If processor usage is too high on the system, consider adding more processors, or turning off the options that enable compression, data deduplication, or encryption. Then, see whether performance improves.</p> <p>For information about tuning for client memory limits, see “Reduce client memory use” on page 206.</p> <p>Software such as firewalls and antivirus programs might affect the efficiency of client operations. For example, during a restore operation, an antivirus program might scan the contents of each restored object, checking for virus signatures. If you suspect that a firewall or antivirus program is slowing client operations, consider temporarily turning the firewall or antivirus program off to see whether performance improves. For tips to minimize the impact of firewall or antivirus programs on other applications, see the documentation for those programs.</p>
4	HBA that connects the client system to the SAN	HBA type and its capabilities	

Item in Figure 2 on page 64	Item	Key characteristics that can affect performance	Details
5	Network interface card (NIC) that connects the client system to the LAN	NIC type and its capabilities	A fast network interface card (NIC) improves network throughput. If you cannot use the latest NIC in your configuration, consider adjusting the client TCPWINDOWSIZE option to improve network throughput, particularly on client systems that are geographically distant from the server. Adjust the TCPWINDOWSIZE options in small increments; a window size that is larger than the buffer space on the network interface adapter might actually degrade throughput. For more network considerations, see Chapter 14, “ Tuning network performance ,” on page 237.
6	Network: LAN	The many components on a network, and the effective speed of data transfer over the network, which is limited by its slowest component	
7	Network: SAN	The many components on a network, and the effective speed of data transfer over the network, which is limited by its slowest component	
8	NIC that connects the server to the LAN	NIC type and its capabilities	
9	Server system	The speed of the I/O bus, the number of processors, the speed of the processors, and the amount and speed of RAM	
10	HBA that connects the server to the SAN	HBA type and its capabilities	See “ Tuning HBA capacity ” on page 172.
11	Virtual tape library (VTL) for the server storage pool	VTL model characteristics that affect the performance of the operation	
12	Tape library for the server storage pool	Number of drives and mount point availability for the operation	
13	Tape drive for the server storage pool	Tape type and sustainable speed	

Server storage data flow

Figure 3 on page 67 shows data flow in the storage backplane in a server system. The data flow might be for an operation such as migration of storage pool data from disk storage pools to other storage pools. For a migration operation, the data flow starts at the source storage pool (item 1 in the graphic and table) and ends at the device for the target storage pool (item 6 or 8).

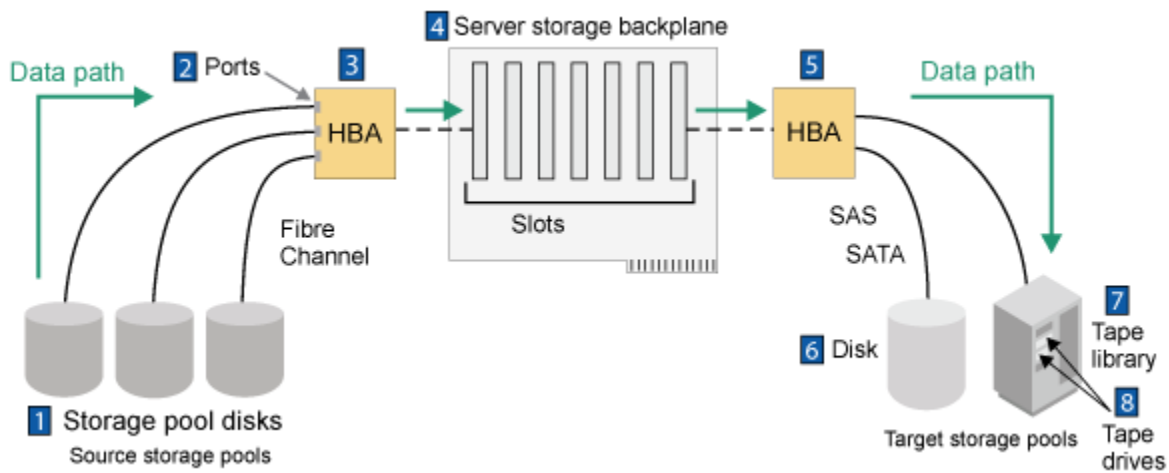


Figure 3. Data flow through the server storage backplane

The following table describes characteristics of the hardware that can affect performance of the operation.

Item in Figure 3 on page 67	Item	Key characteristics that can affect performance
1	Disks for source storage pools	Disk type and rotational speed
2	Ports	Multiple connection points to devices
3	HBA	These devices can have multiple ports. The total amount of data that is being transferred by disks at one time cannot exceed the overall throughput of the HBA.
4	Server storage backplane	The total of the speeds of all cards that are attached to the backplane cannot exceed the speed of the bus.
5	HBA	These devices can have multiple ports. The total amount of data that is being transferred by disks at one time cannot exceed the overall throughput of the HBA.
6	Disks for target storage pools	Disk type and rotational speed
7	Tape library for target storage pools	Number of drives and mount point availability for the operation
8	Tape drives for target storage pools	Tape type and sustainable speed

Related concepts

[Reduce client data flow with compression](#)

The backup-archive client can compress data before it sends the data to the server. Enabling compression on the client reduces the amount of data that is sent over the network and the space that is needed to store it on the server and storage pools. Two client options determine when and if the client compresses data: **compression** and **compressalways**.

Related tasks

Analyzing data flow with the dd command

You can use the **dd** command as a quick test to estimate the best possible results for data flow to disks.

The command is available for operating systems such as AIX or Linux.

Workloads for the server

The ability of a server to handle workload is directly related to the server's resources, including system processors, memory, and I/O bandwidth. The ability of a server to efficiently process daily tasks determines how large a server can be.

Because any system has finite resources, the maximum workload on the server is limited by the recovery objectives. For example, if you decrease the frequency of server database backups to lessen the workload, you increase the time between recovery points for the system. Less frequent server database backups might cause the system to miss the recovery point objectives (RPO).

Ensure that a server can complete essential tasks in a 24-hour cycle:

- Complete the client workload.

The client workload is the amount of data that is backed up or archived during the backup window. The backup window is typically a span of time at night or in the early morning. The ability of the server to store this data within the backup window can be limited by various factors:

- Server storage capacity
- I/O throughput to the storage devices
- Network bandwidth
- Other system attributes such as available memory or processors for the server
- Characteristics of the client systems that are being backed up, including the following characteristics:
 - Processor speeds and memory for the client systems
 - Disk speeds on client systems
 - The total amount of data from all clients
 - The total number of clients that request services from the server at one time

- Complete the essential server maintenance operations.

Daily completion of the following operations keeps the server environment in good working condition and helps you prepare for disaster recovery of the server. These operations are key to effective data maintenance and management:

- Expiration
- Database backup
- Reclamation

Additional daily operations are required depending on the configuration of the solution and the features that are used:

- Storage pool migration
- Storage pool backup
- Duplicate identification processes
- Data replication processes

For examples of how a solution can be configured to handle workloads, see the sample architecture documents in the IBM Storage Protect wiki on Service Management Connect at [Sample Architectures](#).

Limits for the server database size and peak client sessions

IBM tests the IBM Storage Protect server to a specific database size and peak number of concurrent client sessions. However, you must consider the tested values along with other operational factors in your specific environment. Experiences that are reported by other users are also helpful.

Database size

Testing shows that databases with utilization as much as 4 TB are possible.

The practical limit for the database size depends on the performance characteristics of the server system and the time that is required to back up or restore the database. For many users, a 1 - 2 TB server database allows completion of database backup and database restore operations in a time that fits their maintenance window.

Consider deploying another server if the following conditions occur:

- Performance degrades to an unacceptable level as the database grows
- The time that is required to complete server maintenance such as database backup exceeds the total window of time for server maintenance

When you add a server, either balance the existing workload across the servers, or assign any new workload to the new server.

Peak number of client sessions

Testing shows that the IBM Storage Protect server can handle up to 1000 concurrent client sessions. If this value is exceeded, depending on memory or other system limitations, the server performance might degrade or operations might become unresponsive.

The actual number of concurrent sessions where performance problems arise depends on the resources that are available to the server. What the sessions are doing also influences the practical limit on sessions. For example, sessions that move data have a larger effect on the amount of I/O to the target storage pool compared to incremental backup sessions that mostly send queries without sending many files. Also, sessions that perform client-side deduplication drive more I/O to the server database than other sessions.

To reduce the peak-session workload, it might be appropriate to deploy another server or to adjust client scheduling.

Set the **MAXSESSIONS** server option no higher than the tested limit of 1000. Setting the maximum number of sessions higher than necessary uses extra RAM on the server system, but might not have a significant impact.

Related reference

[Running concurrent client sessions](#)

Running two or more client program instances at the same time on the same system might provide better overall throughput than a single client instance, depending on the available resources.

Samples of data protection solutions

Samples of data protection solutions that use IBM Storage Protect for selected scenarios are available on the Service Management Connect wiki. The samples describe specific hardware and software configurations and give performance measurements that were obtained in IBM test labs.

Related concepts

[Resolving common client performance problems](#)

The table contains common client problems and actions that can help to improve client performance.

Related tasks

[Selecting the optimal client backup method](#)

You can use several techniques with the backup-archive client to help you ensure the best performance during various types of backup processing.

Chapter 7. Taking the first steps for solving performance problems

Start by looking for problems that are indicated by system errors or obvious deficiencies outside of IBM Storage Protect. After you are satisfied that no such problems exist, continue by reviewing information about identifying performance problems in the IBM Storage Protect environment. Throughout the process, monitor for errors and changes in performance.

Procedure

1. Review the best practices that are described in [Part 2, “Configuration best practices,”](#) on page 7. Make changes as needed, based on the information.
2. Verify that there are no problems or errors outside of the server software. Problems in the server hardware, the operating system, the network, and attached storage devices can severely affect operations. Fix all errors outside of the server software before you diagnose server performance issues.
 - a) Review error logs for the operating system to find errors that can affect the server.
For example, for AIX systems use the **errpt** command to view errors. For Linux systems, look in the `/var/log` path.
 - b) Verify that attached storage devices such as disk systems are operational and have no errors.
 - c) Verify that storage and local area networks do not have frequent port errors.
3. Review the server activity log and client error logs.
4. Review the log for the server database, the `db2diag.log` file. To find the file, see [Locating Db2 diagnostic log files](#).

What to do next

Implement modifications and fix any problems as described in the preceding steps, and then determine whether you must further analyze performance. Use techniques that are described in [Chapter 8, “Identifying performance bottlenecks,”](#) on page 73 to analyze your system for performance bottlenecks.

Chapter 8. Identifying performance bottlenecks

When you suspect that you have a performance problem or want to improve performance, you can use the provided flowcharts to help identify sources of performance issues. The charts suggest the use of tools and scripts to help you measure performance on your system.

Procedure

- For client problems, or if you are not sure what performance issue you have, start at [“Diagnosing backup and restore performance”](#) on page 73.
- For server problems, start at [“Identifying server performance problems”](#) on page 75.

Diagnosing backup and restore performance

Use the table to diagnose issues with backup and restore performance. The table includes links to diagnostic tasks.

Table 6. Explanation of the decisions and tasks for backup and restore performance		
Step	Question	Diagnostic tasks
1	Do operating system tools show a disk bottleneck on the client or server? For more information, see “Identifying disk bottlenecks for IBM Storage Protect servers” on page 79.	Yes Fix any disk issues. If applicable, check the load on the system for applications that are not related to IBM Storage Protect. See Chapter 12, “Tuning disk storage for the server,” on page 177. No Collect client and server instrumentation. For more information, see “Client instrumentation report” on page 99. Go to question 2.
2	Does client instrumentation show high file I/O or directory process time?	Yes Go to question 3. No Go to question 4.
3	Do you have many files or a complex directory structure?	Yes Use file system or operating system tools to help improve backup and restore operations. See “File space tuning” on page 228. No Fix any disk issues. If applicable, check the load on the system for applications that are not related to IBM Storage Protect. See Chapter 12, “Tuning disk storage for the server,” on page 177.

Table 6. Explanation of the decisions and tasks for backup and restore performance (continued)

Step	Question	Diagnostic tasks
4	Does server instrumentation show slow disk or tape I/O rates?	<p>Yes</p> <p>Fix any disk issues. If applicable, check the load on the system for applications that are not related to IBM Storage Protect.</p> <p>See Chapter 12, “Tuning disk storage for the server,” on page 177.</p> <p>No</p> <p>Find the session thread for this node in server instrumentation. The thread might include the node name or you can see if the session thread ID is listed in the Activity Log to determine which node is affected.</p> <p>Go to question 5.</p>
5	Does the client show high Data Verb time and does the server show high network time for that node’s session thread?	<p>Yes</p> <p>Investigate network speeds and issues and fix any problems.</p> <p>See Chapter 14, “Tuning network performance,” on page 237.</p> <p>No</p> <p>Go to question 6.</p>
6	Does server instrumentation show that most time is spent in database operations for that session thread?	<p>Yes</p> <p>Investigate database issues, including disk and buffer pool settings. Fix any problems.</p> <p>See “Identifying server performance problems” on page 75.</p> <p>No</p> <p>Go to question 7.</p>
7	<p>Are you deduplicating data from an Oracle or SAP database, and the average extent size is less than 100 KB?</p> <p>Tip: To determine the average extent size, generate data deduplication statistics by using the GENERATE DEDUPSTATS command and view the statistics by using the QUERY DEDUPSTATS command. Based on the output, calculate the average extent size by using the following formula:</p> <p>Total Protected Data/ (Compressed Extent Count + Uncompressed Extent Count)</p>	<p>Yes</p> <p>Consider specifying an extent size that is larger than the default value of 50 KB. For more information about changing the extent size, see the MINIMUMEXTENTSIZE parameter section of the REGISTER NODE and UPDATE NODE command documentation.</p> <p>No</p> <p>IBM Support can help in diagnosing performance problems by requesting traces and other information from the environment. The comprehensive analysis of a performance problem is a billable service that is offered to IBM Storage Protect customers. For more information, see the IBM Software Support Handbook at http://www.ibm.com/support/customercare/sas/f/handbook/home.html.</p> <p>See also Chapter 9, “Collecting and analyzing data for performance problems,” on page 83.</p>

Related concepts

[Potential bottlenecks in the data flow for IBM Storage Protect operations](#)

In operations such as client backup and storage pool migration, data moves through many physical components that can affect the speed of the operations. Understanding the characteristics of these components can help when you are working to improve performance.

Identifying server performance problems

Use the flowchart to diagnose issues with server operations. The table after the flowchart provides more information and links to diagnostic tasks and tools.

Tip: Before you review the flowchart, ensure that you answer all questions and fix any issues that are described in the [“Checklist for the server hardware and the operating system” on page 10](#) and the [“Checklist for IBM Storage Protect server configuration” on page 31](#).

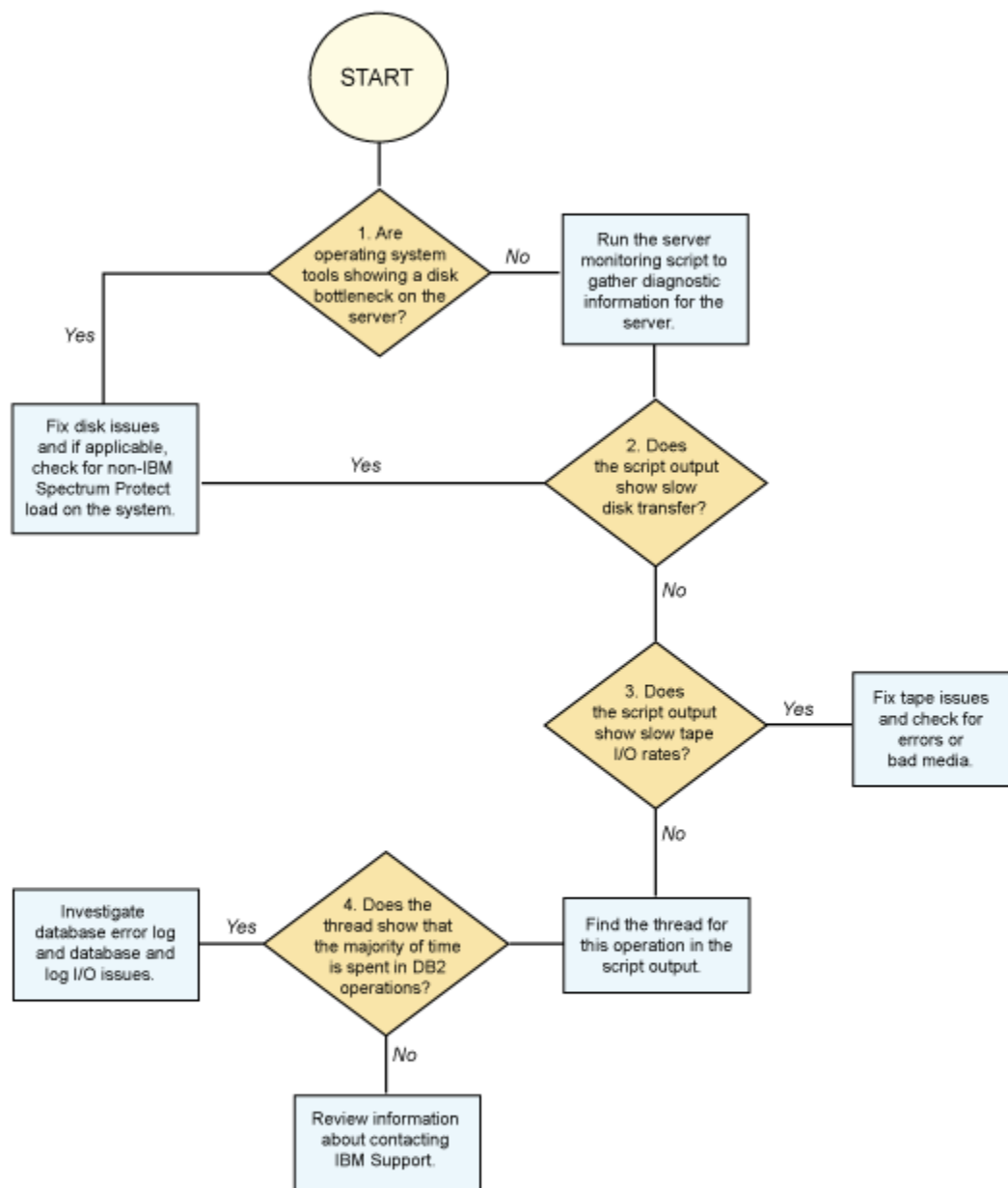


Figure 4. Flowchart to solve issues with server performance

Table 7. Flowchart questions and tasks to help you diagnose and fix server performance problems

Step	Question	Diagnostic tasks
1	Do operating system tools show a disk bottleneck on the server? For more information, see “Identifying disk bottlenecks for IBM Storage Protect servers” on page 79.	<p>Yes Fix any disk issues. If applicable, check the load on the system for applications that are not related to IBM Storage Protect. See Chapter 12, “Tuning disk storage for the server,” on page 177.</p> <p>No Check the servermon component to view diagnostic information for the server. Go to question 2.</p>
2	Does the server monitoring script output show slow disk transfer?	<p>Yes Fix any disk issues. If applicable, check the load on the system for applications that are not related to IBM Storage Protect. See “Analyzing the basic performance of disk systems” on page 82.</p> <p>No Go to question 3.</p>
3	Does the server monitoring script output show slow tape I/O rates?	<p>Yes Fix tape issues and check for errors or bad media, if applicable. See “Tuning tape drive performance” on page 170.</p> <p>No Find the thread for the operation in the script output. Go to question 4.</p>
4	Does the thread show that most of time is spent in Db2 operations?	<p>Yes Investigate database error log and database and log I/O issues. See “Server database and recovery log configuration and tuning” on page 127.</p> <p>No IBM support personnel can help in diagnosing performance problems by requesting certain traces and other information from the environment. The comprehensive analysis of a performance problem is a billable service that is offered to IBM Storage Protect customers. For information about IBM Support and gathering data for problems, see the Software Support Handbook at http://www.ibm.com/support/customer/care/sas/f/handbook/home.html See also Chapter 9, “Collecting and analyzing data for performance problems,” on page 83.</p>

Related concepts

[Potential bottlenecks in the data flow for IBM Storage Protect operations](#)

In operations such as client backup and storage pool migration, data moves through many physical components that can affect the speed of the operations. Understanding the characteristics of these components can help when you are working to improve performance.

Related tasks

Reporting performance problems

Before you report a problem, you can collect information in advance to facilitate the problem investigation.

Evaluating data deduplication results

You can evaluate the effectiveness of IBM Storage Protect data deduplication by examining various queries or reports. Actual data reduction results can show whether the expected storage savings are achieved. You can also evaluate other key operational factors, such as database utilization, to ensure that they are consistent with expectations.

Before you begin

Consider the following factors when you are evaluating data deduplication results:

- When you are using data deduplication, you might not see immediate space savings on the server.
- Because data deduplication includes multiple backup operations across multiple clients, processing will become more effective over time.

Therefore, it is important to collect results at regular intervals to record valid results.

Procedure

- Use the following commands and tools to help you evaluate data deduplication effectiveness:

Action	Explanation
Use the QUERY STGPOOL server command to quickly check deduplication results.	<p>The Duplicate Data Not Stored field shows the actual reduction of data, in megabytes or gigabytes, and the percentage of reduction of the storage pool. For example, issue the following command:</p> <pre>query stgpool format=detailed</pre> <p>If the query is run before reclamation of the storage pool, the Duplicate Data Not Stored value is not accurate because it does not reflect the most recent data reduction. If reclamation did not yet take place, issue the following command to show the amount of data to be removed:</p> <pre>show deduppending backkuppool-file</pre> <p>Where backkuppool-file is the name of the deduplicated storage pool.</p>
Use the QUERY OCCUPANCY server command.	This command shows the logical amount of storage per file space when a file space is backed up to a deduplicated storage pool.

Action	Explanation
Examine the IBM Storage Protect client backup reports to see the data reduction for a backup operation that is run with client-side data deduplication and compression.	<p>The backup reports are available upon the completion of backup operations.</p> <p>Over time, if the backup reports repeatedly show little to no data reduction after many backups, consider redirecting the client node to a non-deduplication storage pool if one is available. This way, the client is not wasting time by processing data that are not good candidates for data deduplication.</p>
Run the deduplication report script to show information about the effectiveness of data deduplication.	<p>The report provides details of deduplication-related utilization of the IBM Storage Protect database. You can also use it to gather diagnostic information when the deduplication results are not consistent with your expectations.</p> <p>To obtain the script and usage instructions for the script, see technote 476911.</p>

What to do next

For more information, see [Container Storage Pools Best Practices](#).

Related concepts

[Checklist for data deduplication](#)

Data deduplication requires more processing resources on the server or client. Use the checklist to verify that hardware and your IBM Storage Protect configuration have characteristics that are key to good performance.

Related tasks

[Tuning server-side data deduplication](#)

Tune settings and configuration for different operations to ensure that the performance of server-side data deduplication is efficient.

[Tuning client-side data deduplication](#)

The performance of client-side data deduplication can be affected by processor requirements and deduplication configuration.

Identifying disk bottlenecks for IBM Storage Protect servers

Tools can help you identify where disk storage that is used for IBM Storage Protect servers might have bottlenecks.

Before you begin

Before you begin this task, review information about optimal disk storage configuration for the server database, recovery logs, and storage pools.

Procedure

To identify disk bottlenecks, you can use one or both of the following methods:

- Use the servermon component, that is automatically installed and configured as part of the server installation, to collect data at regular intervals.
- Use analysis tools that are provided by third parties. Such tools can be effective to analyze storage systems for basic performance characteristics before they are used for IBM Storage Protect storage.

See [“Analyzing disk performance by using system tools”](#) on page 80.

Related concepts

[Potential bottlenecks in the data flow for IBM Storage Protect operations](#)

In operations such as client backup and storage pool migration, data moves through many physical components that can affect the speed of the operations. Understanding the characteristics of these components can help when you are working to improve performance.

Related reference

[Checklist for server database disks](#)

Use the checklist to verify that the system where the server is installed meets requirements for hardware and software configuration.

[Checklist for server recovery log disks](#)

The recovery log for the server consists of the active log, the archive log, and optional logs for mirroring and failover. Use the checklist to verify that disk systems that are being used for the logs have the characteristics and configuration that are key to good performance.

[Checklist for storage pools on DISK or FILE](#)

Use the checklist to review how your disk storage pools are set up. This checklist includes tips for storage pools that use DISK or FILE device classes.

Analyzing disk performance by using system tools

You can use system tools to monitor I/O for disk storage to help you identify disk bottlenecks. For example, use tools such as **nmon** for AIX and Linux operating systems, and the Performance Monitor for Windows systems.

About this task

Other tools can be used but are not documented here. For example, for operating systems such as AIX and Linux, you can use the **sar** command to collect information about system activity.

Related tasks

[Analyzing data flow with the dd command](#)

You can use the **dd** command as a quick test to estimate the best possible results for data flow to disks. The command is available for operating systems such as AIX or Linux.

Analyzing disk performance by using the nmon command

You can use the **nmon** command on AIX and Linux operating systems. Use the command to display local system statistics in interactive mode and to record system statistics in recording mode.

Procedure

1. Run the command as root.

You can run the command from any directory, but you must be logged in as the root user. The command starts a background process to capture the specified number of snapshots of system statistics at the specified intervals, and writes the output file.

Important: Do not use the **kill** command to end the process because that causes the file to be incomplete and unusable for analysis.

To run the command, use the following parameters:

-f

Specifies that the output is written to a file. The file is created in the directory where you run the command, with the name *hostname_YYMMDD_HHMM.nmon*.

-s nnnn

Specifies the seconds between snapshots of statistics.

-c nnn

Specifies the number of snapshots.

Typically to create a report for performance analysis, you specify 15 minutes between snapshots (900 sec) over 24 hours (96 snapshots). For example, issue the following command:

```
nmon -f -s 900 -c 96
```

To collect a snapshot of the system every hour for seven days, issue the following command:

```
/home/hm12857/netmon/nmon -f -s 3600 -c 168
```

To collect a snapshot of the system every 30 minutes for five days, issue the following command:

```
/home/hm12857/netmon/nmon -f -s 1800 -c 240
```

2. Analyze the data by using the **nmon** Analyzer spreadsheet tool. Focus on the disk-busy statistics (Disk %Busy). Look for disks that are consistently over 80% busy (weighted average). The weighted average is shown in red in the chart on the diskbusy tab.

Analyzing disk performance by using Windows Performance Monitor (perfmon)

Examine disk statistics by using performance counters that are available in the Performance Monitor.

Procedure

1. Start the Performance Monitor.
From a command prompt, enter: **perfmon**.
2. Create a Data Collector Set to collect data about disks.
Select the following performance counters from the **Physical Disk** category:
 - **Avg. Disk Sec./Transfer**
 - **Avg. Disk Queue Length**
 - **Avg Disk Bytes/Transfer**
 - **Disk Bytes/sec**
 - **Split IO/sec**
3. Run the **perfmon** tool while you are experiencing the performance problem. Compare the results with the guidance in the following table.

Performance counter	Guidance
Physical Disk: Avg. Disk Sec./Transfer	Values less than 25 ms are good.
Physical Disk: Avg Disk Queue Length	A value that is 2 or 3 times the number of disks in the array is optimal.
Physical Disk: Avg Disk Bytes/Transfer	The goal is that the stripe size for the array is at least the average of this counter.
Physical Disk: Disk Bytes/sec	Optimal result is that the sum of values for all disks that are attached to a single controller is less than 70% of the theoretical throughput.
Physical Disk: Split IO/sec	A nonzero value for this counter indicates possible disk fragmentation.

Analyzing the basic performance of disk systems

To verify that a storage system can meet the workload requirements for IBM Storage Protect operations, run stress tests. You might also want to analyze disk performance when changes are made to the host or the network backbone.

About this task

Various tools are available for analysis and stress tests for disk characteristics such as I/O operations per second (IOPS).

Procedure

- For AIX, you can use the **ndisk64** command.
Search for the latest **nstress** package at [AIX performance analysis and tuning](#).
- You can use third-party tools such as Iometer, which is available for Windows and other operating systems.
For information about the Iometer tool, see <http://www.iometer.org>.
- For operating systems such as AIX and Linux, you can use the **dd** command for simple tests of capabilities.

Analyzing data flow with the dd command

You can use the **dd** command as a quick test to estimate the best possible results for data flow to disks. The command is available for operating systems such as AIX or Linux.

About this task

The **dd** command can be useful if you do not have or do not want to install more powerful tools. To estimate peak performance under ideal conditions, use the **dd** command to time how long a write to a device takes. Then, time how long a read from the device takes.

Procedure

1. To run a write test, issue the following command.

```
time dd if=/dev/zero of=/device_path/filename bs=262144 count=40960
```

where *device_path* is the name of the file system that you want to test, and *filename* is the name of a file.

Important: The *filename* file must not exist in the file system. If the file does exist, the command overwrites it with zeros.

The output of this command gives you the time that is required to write a 10 GB file in 256 KB blocks.

2. To run a read test of the file that was written, issue the following command.

```
time dd if=/device_path/filename of=/dev/null bs=262144 count=40960
```

When you evaluate the command results, keep in mind that if you just ran the write test that the data might still be in the disk cache. The time reported by the command for the read operation is therefore less than what you can expect for typical IBM Storage Protect server operations. For typical IBM Storage Protect server operations, data is not likely to be in the cache and is read from the disk itself.

Chapter 9. Collecting and analyzing data for performance problems

Capturing specific performance metrics, as the problem occurs in your environment, is essential to help IBM support personnel with the analysis.

Most performance problems appear as unacceptable response times or resource usage. Performance problems can slowly develop over time, as a result of dwindling resources, or suddenly as the result of a hardware or software change in the environment.

As part of the standard product support package, IBM will help to determine whether a performance problem is a result of a product defect. Gathering key performance metrics, from the customer environment will be a key part of this activity. The comprehensive analysis of a performance problem is a billable service that is offered to IBM Storage Protect customers. For more information, see the Software Support Handbook at <http://www.ibm.com/support/customer/sas/f/handbook/home.html>.

Measuring the baseline

Performance problems are often reported immediately following some change to system hardware or software. Unless there is a pre-change baseline measurement with which to compare post-change performance, it can be difficult to qualify the problem.

About this task

Any change in the environment, including software, hardware, or network changes, can affect the performance of operations in your environment.

As a best practice, monitor the environment before and after each change. The alternative is to take the baseline measurements at regular intervals, for example, once a month, and save the output. When a problem is found, you can use the previous measurements for comparison. It is a good idea to collect a series of measurements, which can help you diagnose a possible performance problem.

To maximize performance diagnosis, collect data for various periods of the day, week, or month when performance is likely to be an issue. For example, you might have workload peaks during these times:

- In the middle of the mornings for online users
- During a late-night batch run
- During the end-of-month processing
- During major data loads

Collect data for each peak in workload because a performance problem might cause problems only during one of these periods and not during other times.

Restriction: Using any tool to collect the baseline data can impact the performance of the system that is being measured.

Procedure

To collect baseline data, you can use the following tools:

- On any operating system, you can use the IBM Storage Protect servermon component. The servermon component is automatically installed and configured as part of the server installation to collect data at regular intervals.
- On AIX and Linux operating systems, you can use the nmon utility.
- On Windows operating systems, you can use the perfmon utility to gather a set of performance counters.

Periodically record your baseline measurements so that you can use the data for comparisons after an unexpected performance degradation. If you collect baseline data before a performance problem is detected, IBM Support can use the data to help you resolve performance problems.

Describing performance problems

Support personnel often receive information that is insufficient to accurately determine the nature of a performance problem. You must be able to describe the problem in as much detail as possible.

About this task

Always obtain as much detail as possible before you collect or analyze data, by asking the following questions about the performance problem:

- Can the problem be demonstrated by running a specific command or reconstructing a sequence of events? What is the least complex example of the problem.
- Is the slow performance intermittent? Does it get slow at certain times, but then returns to normal? Does it occur at certain times of the day or in relation to some specific activity?
- Is everything slow or only some things? What aspect is slow? For example, time to run a command, or elapsed time to complete a process, or time to paint the screen?
- When did the problem start occurring? Was the situation the same since the system was first installed or went into production? Did anything change on the system before the problem occurred (such as adding more users or migrating more data to the system)?
- If the issue is client and server, can the problem be demonstrated locally on the server (network versus server issue)?
- If network related, how are the network segments configured (including bandwidth such as 100 Mb/sec or 10 Mb/sec)? Are there any routers between the client and server?
- What vendor applications are running on the system, and are those applications involved in the performance issue?
- What is the impact of the performance problem on the users?

Reporting performance problems

Before you report a problem, you can collect information in advance to facilitate the problem investigation.

About this task

When you report a performance problem, it is not enough just to gather data and analyze it. Without knowing the nature of the performance problem, you might waste time and resources when you analyze data that might have nothing to do with the problem that is being reported.

Your local support personnel can use this information to help solve the performance problem with you.

For information about IBM Support and gathering data for problems, see the Software Support Handbook at <http://www.ibm.com/support/customercare/sas/f/handbook/home.html>

Procedure

To help get your problem resolved more quickly, complete the following tasks:

1. Gather information about the performance issue to help you prepare a problem description:
 - For backup-archive client performance issues, run the client instrumentation. See [“Collecting instrumentation data with the client”](#) on page 99.
 - For server performance issues, use the servermon component that is automatically installed and configured as part of the server installation to collect data at regular intervals.

- Gather detailed information about LUN layout, cache size and setting information, disk system information, file systems type, RAID type, and other setup details. Because many performance issues are I/O related, this information is important.
 - Collect a list of the hardware information such as host bus adapter type, processor type, and amount of RAM you have on the client and server.
 - Gather network and SAN zoning information.
2. Provide a statement of a simple specific instance of the problem. Separate the symptoms and facts from the theories, ideas, and your own conclusions. Problem Management Records that report the system is slow statements can require extensive investigation to determine what is meant by slow, how it is measured, and what is acceptable performance.
 3. Gather information about everything that changed on the system in the weeks before the problem. Missing something that changed can block a possible investigation path and might delay finding a resolution. If all the facts are available, IBM Support can eliminate the unrelated facts.

Tip: Be sure that you collect the information from the correct system. In large sites, it is easy to accidentally collect the data on the wrong system, which makes it difficult to investigate the problem.

4. Provide the following information:
 - A problem description that can be used to search the problem-history database to see whether a similar problem was reported.
 - Describe the aspect of your analysis that led you to conclude that the problem is caused by a defect in the operating system.
 - Describe the hardware and software configuration in which the problem is occurring:
 - Is the problem confined to a single system, or does it affect multiple systems?
 - What are the models, memory sizes, and number and size of disks on the affected systems?
 - What kinds of LAN and other communications media are connected to the systems?
 - Does the overall configuration include other operating systems?
 - Describe the characteristics of the program or workload that is experiencing the problem.
 - Does an analysis with operating system tools indicate that it is processor-limited or I/O-limited?
 - What is the workload that is run on the affected systems?
 - Describe the performance objectives that are not being met.
 - Is the primary objective console or terminal response time, throughput, or real-time responsiveness?
 - Were the objectives derived from measurements on another system? If so, what was its configuration?
5. If this report is the first report of the problem, you will receive a PMR number for use in identifying any additional data that you supply and for future reference. Include all of the following items when the supporting information and the performance data are gathered:
 - A means of reproducing the problem:
 - If possible, include a program or shell script that demonstrates the problem.
 - At a minimum, a detailed description of the conditions under which the problem occurs is needed.
 - The application that is experiencing the problem:
 - If the application is, or depends on, any software product, identify the exact version and release of that product.
 - If the source code of a user-written application cannot be released, document the exact set of compiler parameters that are used to create the executable program.

Collecting instrumentation data for the IBM Storage Protect server, client, and API

IBM Storage Protect instrumentation can collect data to help isolate performance bottlenecks in the IBM Storage Protect client, server, or the network.

IBM Storage Protect instrumentation is available for the IBM Storage Protect server, client, and API. It is intended to be used for performance tuning and problem determination. You can use the instrumentation as an alternative to collecting data from the traditional IBM Storage Protect **trace** command.

The following steps are a basic approach to working on performance bottlenecks:

1. Use IBM Storage Protect instrumentation data to determine which backup component (client, server, or network) accrues the most time during the IBM Storage Protect process.
2. After you isolate the component that takes the most time, try to determine whether the operation is restricted by a hardware or software resource.
3. Change how that resource is used, or augment it. For example, upgrade processors, or increase memory, disks, or tape drives.
4. Repeat this process as needed to reduce the bottleneck to an acceptable level.

Benefits of IBM Storage Protect instrumentation

There are benefits that are associated with using IBM Storage Protect instrumentation function over using the traditional IBM Storage Protect **trace** command.

The advantages of using the IBM Storage Protect instrumentation function are as follows:

- IBM Storage Protect traces potentially produce large trace files, frequently causing out-of-space conditions on file systems, and can cause significant performance degradation. With the IBM Storage Protect instrumentation function, no huge trace files are created, causing minimal performance impact.
- IBM Storage Protect instrumentation generates concise reports that collate and summarize significant performance information. The report files are typically small in size, often less than 1 MB, and are designed to have minimal impact on performance. The data is stored in memory until the instrumentation session ends.

How processes are tracked

Instrumentation tracks operations that can affect performance.

For example, the following operations are tracked:

- Disk I/O
- Network I/O
- Tape I/O

Each IBM Storage Protect process can have multiple threads. All threads can operate on different processors. The IBM Storage Protect server can have hundreds of active threads at a time. You can use the **show threads** command to see a snapshot of the active threads.

For example, a backup operation uses at least two threads. A **SessionThread** thread receives data from the client and sends it to an **SsAuxSinkThread** thread. When you back up data to a sequential device, the **AgentThread** thread moves the data from the **SsAuxSinkThread** thread and writes the data to tape. When you back up data to random disk on IBM AIX, Linux, and UNIX systems, a **DiskServerThread** writes the data to the device. When you back up data to disk on Microsoft Windows systems, the data moves directly to the random disk from the **SsAuxSinkThread** thread.

IBM Storage Protect instrumentation tracks processes in the following manner:

- Operations are tracked on a thread-by-thread basis
- Most sessions and processes use more than one thread

- Results are stored in memory until instrumentation is ended

Server instrumentation for performance analysis

You can use server instrumentation to track operations, such as backup and restore, and to help identify where performance problems originate.

The servermon component that is automatically installed and configured as part of the server installation collects data at regular intervals.

Related tasks

Starting and stopping server instrumentation

You can start server instrumentation from an administrative command line or from an administrative client. After you stop server instrumentation, you can use the results to determine where performance problems are occurring.

Server instrumentation categories

IBM Storage Protect server instrumentation can report on the elapsed times for the process categories that are documented in the table. Server instrumentation tracks all input and output on a thread-by-thread basis for the categories.

Table 8 on page 87 lists the server instrumentation categories that are tracked and the activity that is timed.

<i>Table 8. Server instrumentation categories</i>	
Category	Activity
Acquire Latch	The amount of time to acquire a database page from disk or a buffer pool.
Acquire XLatch	The amount of time to acquire a database page for update (from disk or a buffer pool).
Cloud Attach	The amount of time to attach a thread to the internal Oracle Java™ virtual machine (JVM) that makes Java Native Interface (JNI) method calls.
Cloud Close	The amount of time to close a connection to an IBM Cloud Object Storage endpoint.
Cloud Connect	The amount of time to connect to an IBM Cloud Object Storage endpoint.
Cloud Delete	The amount of time to run an HTTP DELETE operation that deletes cloud object data.
Cloud Detach	The amount of time to detach a thread from the internal JVM that makes JNI method calls.
Cloud Stats	The amount of time to retrieve statistics that are related to the internal JVM and HTTP request methods that are used on the server. The methods can include PUT, GET, POST, and DELETE.

Table 8. Server instrumentation categories (continued)

Category	Activity
Cloud Stream	The amount of time to stream data buffers to object storage in the following situations: <ul style="list-style-type: none"> • When you use a storage tiering rule to move data from a directory-container storage pool to a cloud-container storage pool. • When you use a cloud reclamation rule to reclaim space in a cloud-container storage pool.
Cloud Upload	The amount of time to upload a file from the cloud accelerator cache to object storage in a cloud-container storage pool.
Compression	The amount of time to compress data by using the LZ4 compression algorithm.
CRC Processing	The amount of time to compute or compare cyclic redundancy check (CRC) values in storage pools.
Data Copy	The amount of time to copy data to buffers in memory.
Db2 Commit	The amount of time to commit a Db2 transaction.
Db2 Connect	The amount of time to connect to a Db2 process.
Db2 CR Exec	The amount of time to run an SQL statement that counts rows.
Db2 CR Prep	The amount of time to prepare an SQL statement that counts rows.
Db2 Delet Exec	The amount of time for Db2 to run an SQL statement that deletes a row.
Db2 Delet Prep	The amount of time for Db2 to parse an SQL statement that deletes a row.
Db2 Fetch	The amount of time to prepare an SQL statement that retrieves one row from Db2.
Db2 Fetch Exec	The amount of time for Db2 to run an SQL statement that returns one row.
Db2 Fetch Prep	The amount of time for Db2 to prepare an SQL statement that returns one row.
Db2 Inser Exec	The amount of time for Db2 to run an SQL statement that inserts a row.
Db2 Inser Prep	The amount of time for Db2 to parse an SQL statement that inserts a row.
Db2 MFetch	The amount of time to prepare an SQL statement that retrieves many rows from Db2.
Db2 MFtch Exec	The amount of time for Db2 to run an SQL statement that returns many rows.
Db2 MFtch Prep	The amount of time for Db2 to prepare an SQL statement that returns many rows.

Table 8. Server instrumentation categories (continued)

Category	Activity
Db2 Reg Exec	The amount of time for Db2 to run complex SQL statements.
Db2 Reg Fetch	The amount of time for Db2 to retrieve rows for a complex SQL statement.
Db2 Reg Prep	The amount of time for Db2 to prepare complex SQL statements.
Db2 Rollback	The amount of time for Db2 to roll back a transaction instead of committing the transaction.
Db2 Updat Exec	The amount of time for Db2 to run an SQL statement that updates a row.
Db2 Updat Prep	The amount of time for Db2 to parse an SQL statement that updates a row.
Decrypt Data	The amount of time to decrypt data buffers during system-to-system communication.
Decrypt Final	The amount of time to decrypt the final piece of data during system-to-system communication.
Decrypt Init	The amount of time to initialize a system-to-system communication cipher for decryption. For example, the amount of time to initialize the AES-GCM mode of operation for decryption.
Disk Commit	The amount of time to run the FSYNC command or other system call to ensure that write operations are completed to disk.
Disk Read	The amount of time to read data from disk.
Disk Write	The amount of time to write data to disk.
	You can combine this amount with the Disk Commit amount to get the total write time.
DNS Lookup	The amount of time to run a system DNS lookup during TCP/IP session communication.
Encrypt Data	The amount of time to encrypt data buffers for system-to-system communication in a secure session. For example, sessions that use the AES-GCM mode of operation for encryption.
Encrypt Final	The amount of time to encrypt the final piece of data during system-to-system communication.
Encrypt Init	The amount of time to initialize a system-to-system communication cipher for encryption. For example, the amount of time to initialize the AES-GCM mode of operation for encryption.
Fingerprint	The amount of time used to find extent boundaries for data deduplication.
IBM C®rypto for C Digest	The amount of time to run an algorithm for data deduplication extents.

Table 8. Server instrumentation categories (continued)

Category	Activity
MD5	The amount of time to compute the MD5 digest for a data buffer.
Namedpipe Recv	The amount of time to receive data on a named pipe.
Namedpipe Send	The amount of time to send data on a named pipe.
Network Recv	The amount of time to receive data on a network from a client.
Network Send	The amount of time to send data on a network to a client.
Queue Wait Get	The amount of time for an empty queue to become available for the next item in the queue. For example, the amount of time for a producer to put a work item in a worker thread queue. Long wait times might indicate that the producer is not creating work fast enough or that too many workers are involved in processing.
Queue Wait Put	The amount of time to put data into a first-in first-out queue with a fixed capacity that is full. A long wait time might indicate that worker threads are processing work slower than the producers can generate work. Adding worker threads might improve processing speed for the queue and reduce the wait time.
SHA1	The amount of time to compute the SHA1 digest for a data buffer.
Shmem Copy	The amount of time to copy data to and from a shared memory segment.
Shmem Read	The amount of time to read data from a shared memory buffer.
Shmem Write	The amount of time to write data to a shared memory buffer.
Sleep	The amount of time to delay execution of a thread.
SSL Receive	The amount of time to receive data over a socket that is encrypted by using the SSL or TLS protocol.
SSL Send	The amount of time to send data over a socket that is encrypted by using the SSL or TLS protocol.
Tape Commit	The amount of time to synchronize tape to ensure that data is written from device buffers to media.
Tape Data Copy	The amount of time to copy data to tape buffers in memory.
Tape Locate	The amount of time to locate a tape block for read/write operations.

Table 8. Server instrumentation categories (continued)

Category	Activity
Tape Misc	The amount of time to process tape that is not tracked in another tape category (operations such as open or rewind).
Tape Read	The amount of time to read data from tape.
Tape Write	The amount of time to write data to tape.
TCP Read	The amount of time to read data from server.
Thread Wait	The amount of time to wait for another thread.
Tm Lock Wait	The amount of time to acquire the transaction manager lock.
Uncompress	The amount of time to uncompress data.
Unknown	The amount of time for an activity that is not tracked by another category.
xxHash	The amount of time to compute an xxHash digest for a data buffer.

Server threads in instrumentation output

The server program divides its operations into threads. In instrumentation output, the names of the threads identify the operations.

Only some of the threads in the instrumentation output are useful for diagnosing performance problems. The most important threads are for reclamation of storage pool volumes, migration of data from random-access storage pools, and backup of storage pools.

Reclamation of storage pool volumes

The main thread for a reclamation operation for a storage pool volume is called `AfRclmVolumeThread`. The main thread starts one or two child threads. Each child thread controls a thread that is called `AgentThread`. Data movement operations start with an `AgentThread` that reads an object from a volume that is being reclaimed. See [Figure 5 on page 92](#).

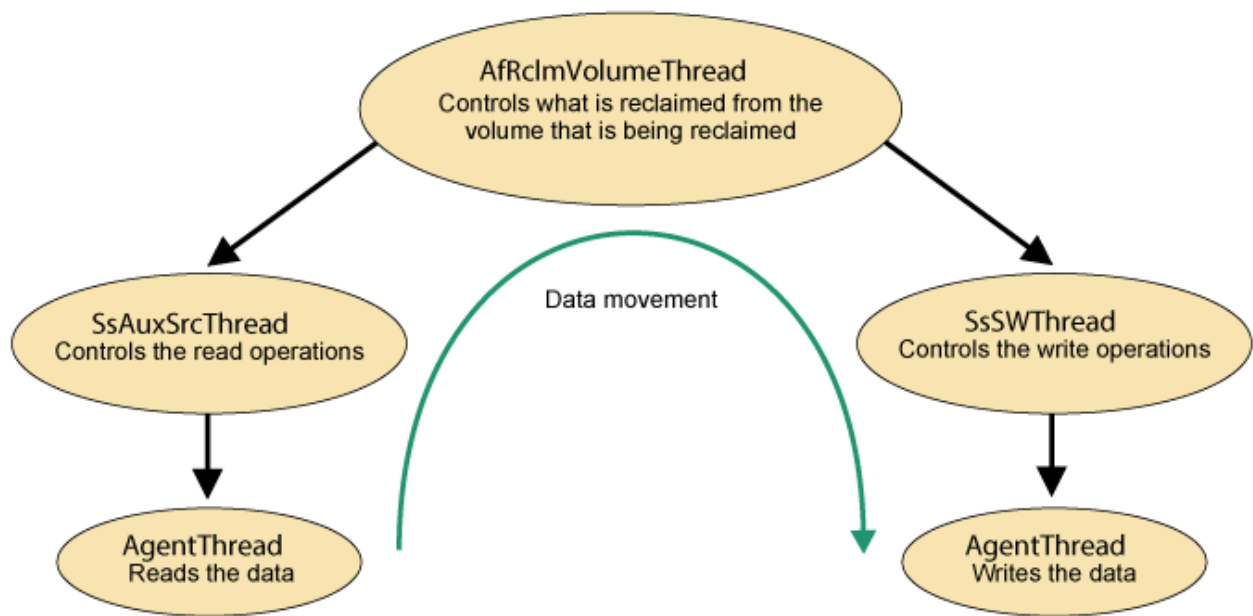


Figure 5. Threads for reclamation of storage pool volumes

A typical data movement operation starts with an AgentThread that reads an object from a volume that is being reclaimed. That data is processed through the SsAuxSrcThread, AfRclmVolumeThread, and SsSWThread threads. Data movement ends when the data is written on the target volume by the AgentThread thread that writes the data.

Migration of data from random-access storage pools

The main thread for a migration operation for a random-access storage pool is DfMigrationThread. The child threads to complete the migration operation differ by operating system.

AIX and Linux

The main thread, DfMigrationThread, does the work of selecting the data for migration and the volumes that are read from and written to. The thread starts two child threads: SsAuxSrcThread, which controls the read operations, and SsSWThread, which controls the write operations. See [Figure 6 on page 93](#).

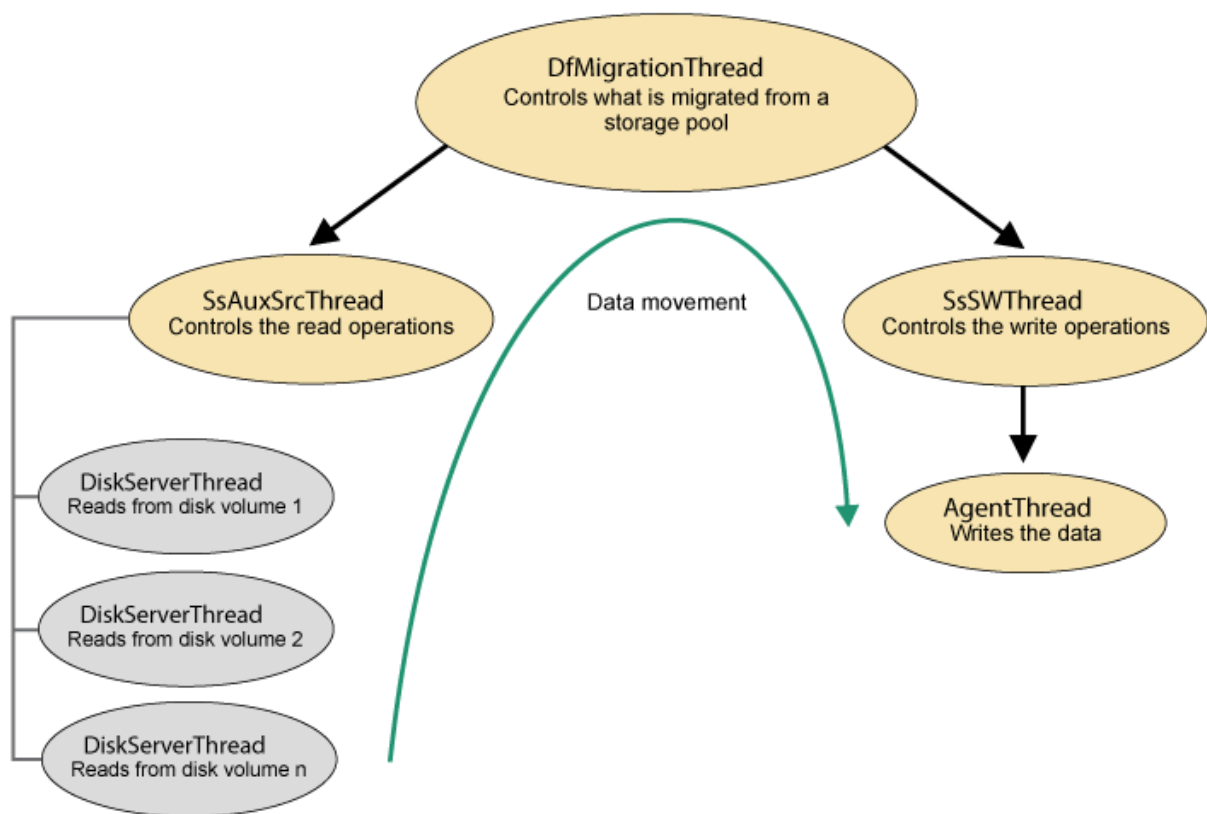


Figure 6. Threads for storage pool migration on AIX and Linux systems

To read the data, the SsAuxSrcThread thread uses a DiskServerThread thread for each volume that must be read. The SsAuxThread thread uses multiple DiskServerThread threads if the data that is being migrated is on more than one volume.

DiskServerThread threads are independent of the SsAuxSrcThread thread. For each volume in a random-access storage pool, a DiskServerThread runs constantly to read and write from that particular volume. For example, if the storage pool has 10 disk volumes, then 10 DiskServerThread threads are always running. Because the SsAuxThread is not a parent for the DiskServerThread threads, you cannot use the ID of the SsAuxThread thread to find a DiskServerThread that is being used.

To write the data, the SsSWThread thread controls a child thread that is called AgentThread, which writes the data to the target volume.

Data movement starts with the DiskServerThread that reads the data from the volume that has the data to be migrated. That data is processed through the SsAuxSrcThread, DfMigrationThread, and SsSWThread threads. Data movement ends when the data is written on the target volume by the AgentThread thread that writes the data.

Windows

The main thread, DfMigrationThread, does the work of selecting the data for migration and the volumes that are read from and written to. The thread starts two child threads: SsAuxSrcThread, which controls the read operations, and SsSWThread, which controls the write operations. The SsAuxSrcThread thread reads data directly from the disks, without using other threads. For writing the data, the SsSWThread thread controls a separate child thread that is called AgentThread, which writes the data to the target volume.

See [Figure 7](#) on page 94.

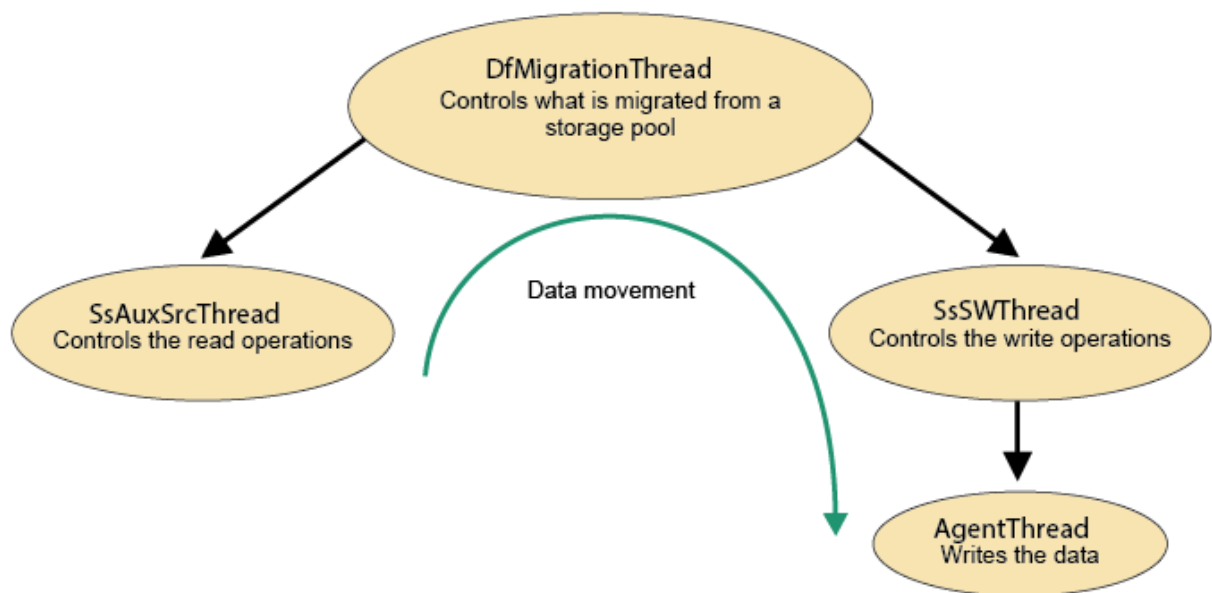


Figure 7. Threads for storage pool migration on Windows systems

Data movement starts with the SsAuxSrcThread that reads the data from the volume that has the data to be migrated. That data is processed through the DfMigrationThread and SsSWThread threads. Data movement ends when the data is written on the target volume by the AgentThread thread that writes the data.

Backups for random-access storage pools

The main thread for a backup operation for a random-access storage pool is DfBackupPoolThread. The threads for reading from the random-access storage pool differ by operating system.

AIX and Linux

The main thread, DfBackupPoolThread, controls the work for the backup operation, including selection of volumes and reading and writing the data. The thread starts two child threads: SsAuxSrcThread, which controls the read operations, and SsSWThread, which controls the write operations. See [Figure 8 on page 95](#).

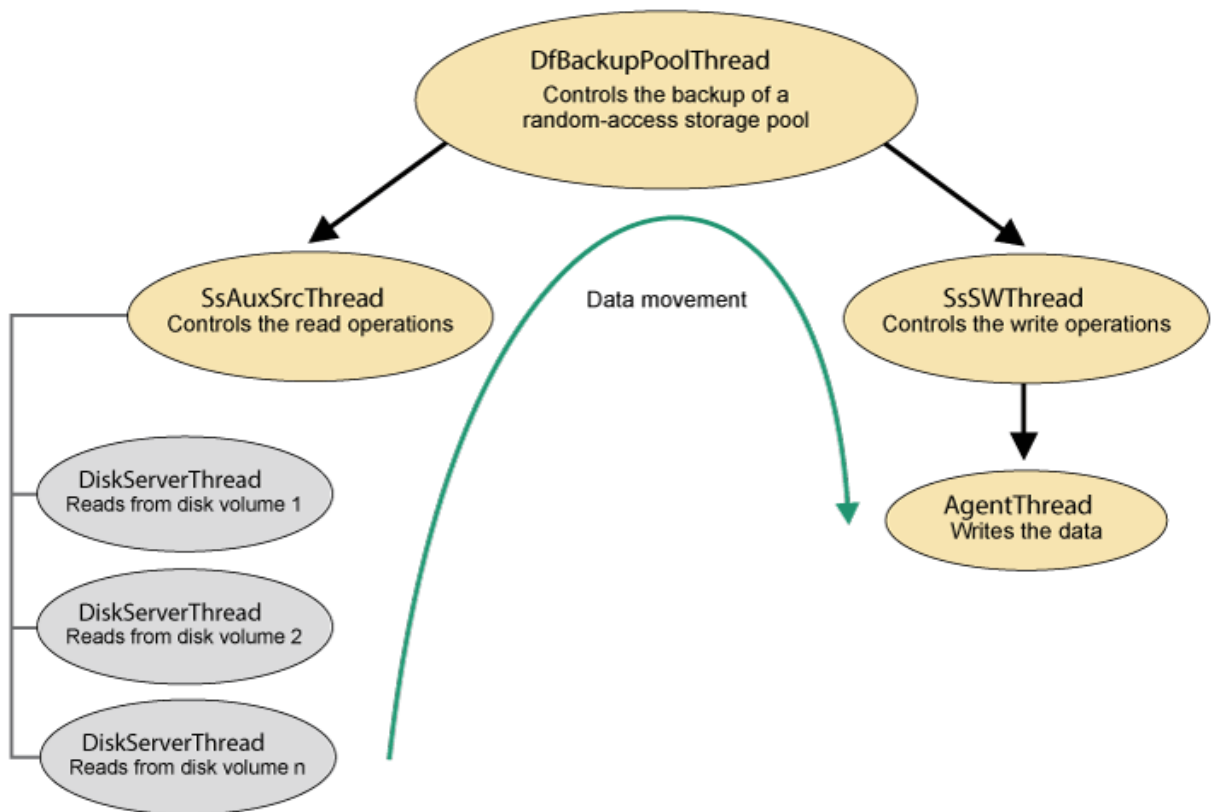


Figure 8. Threads for backup of random-access storage pools on AIX and Linux systems

To read the data, the SsAuxSrcThread thread uses a DiskServerThread thread for each volume that must be read. The SsAuxThread thread uses multiple DiskServerThread threads if the data that is being migrated is on more than one volume.

DiskServerThread threads are independent of the SsAuxSrcThread thread. For each volume in a random-access storage pool, a DiskServerThread runs constantly to read and write from that particular volume. For example, if the storage pool has 10 disk volumes, then 10 DiskServerThread threads are always running. Because the SsAuxThread is not a parent for the DiskServerThread threads, you cannot use the ID of the SsAuxThread thread to find a DiskServerThread that is being used.

To write the data, the SsSWThread thread controls a child thread that is called AgentThread, which writes the data to the target volume.

Data movement starts with the DiskServerThread that reads the data from the volume that has the data to be backed up. That data is processed through the SsAuxSrcThread, DfBackupPoolThread, and SsSWThread threads. Data movement ends when the data is written on the target volume by the AgentThread thread that writes the data.

Windows

The main thread, DfBackupPoolThread, controls the work for the backup operation, including selection of volumes and reading and writing the data. The thread starts two child threads: SsAuxSrcThread, which controls the reading of data, and SsSWThread, which controls the writing of data. The SsAuxSrcThread thread reads the data directly from the disks, without using other threads. For writing the data, the SsSWThread thread controls a separate child thread that is called AgentThread, which writes the data to the target volume. See [Figure 9 on page 96](#).

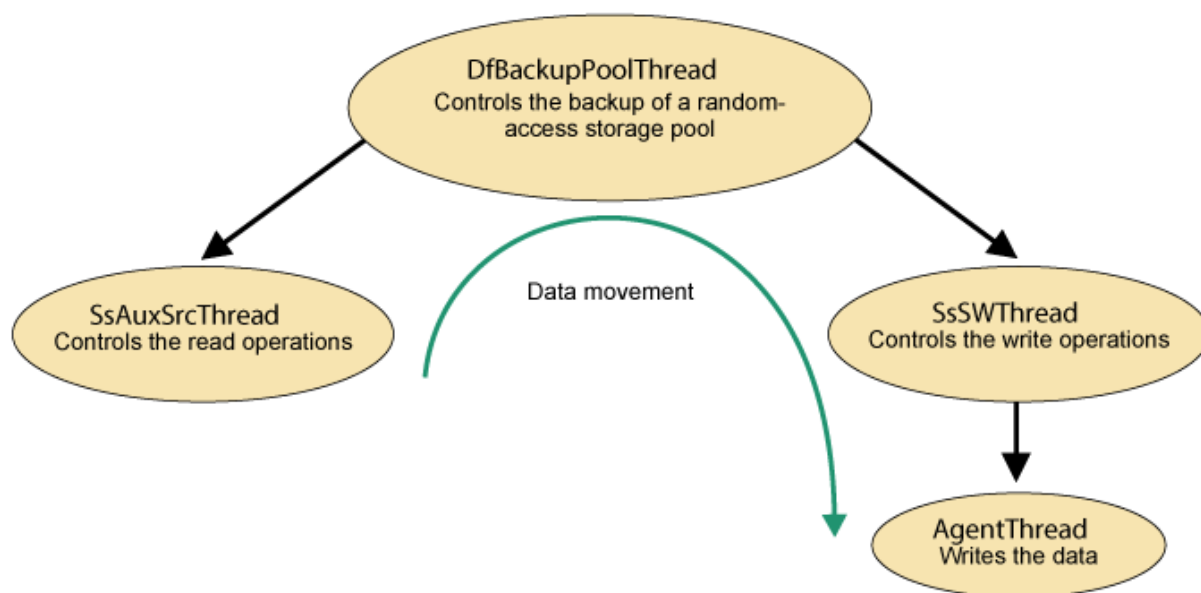


Figure 9. Threads for backup of random-access storage pools on Windows systems

Data movement starts with the SsAuxSrcThread that reads the data from the volume that has the data to be backed up. That data is processed through the DfBackupPoolThread and SsSWThread threads. Data movement ends when the data is written on the target volume by the AgentThread thread that writes the data.

Backups for sequential-access storage pools

The main thread for a backup operation for a sequential-access storage pool is AfBackupPoolThread. This thread controls the work for the backup operation, including selection of volumes and reading and writing the data. The main thread starts two child threads: SsAuxSrcThread, which controls the read operations, and SsSWThread, which controls the write operations. Each of these child threads controls a separate child thread that is called AgentThread, which either reads or writes the data. See [Figure 10 on page 96](#).

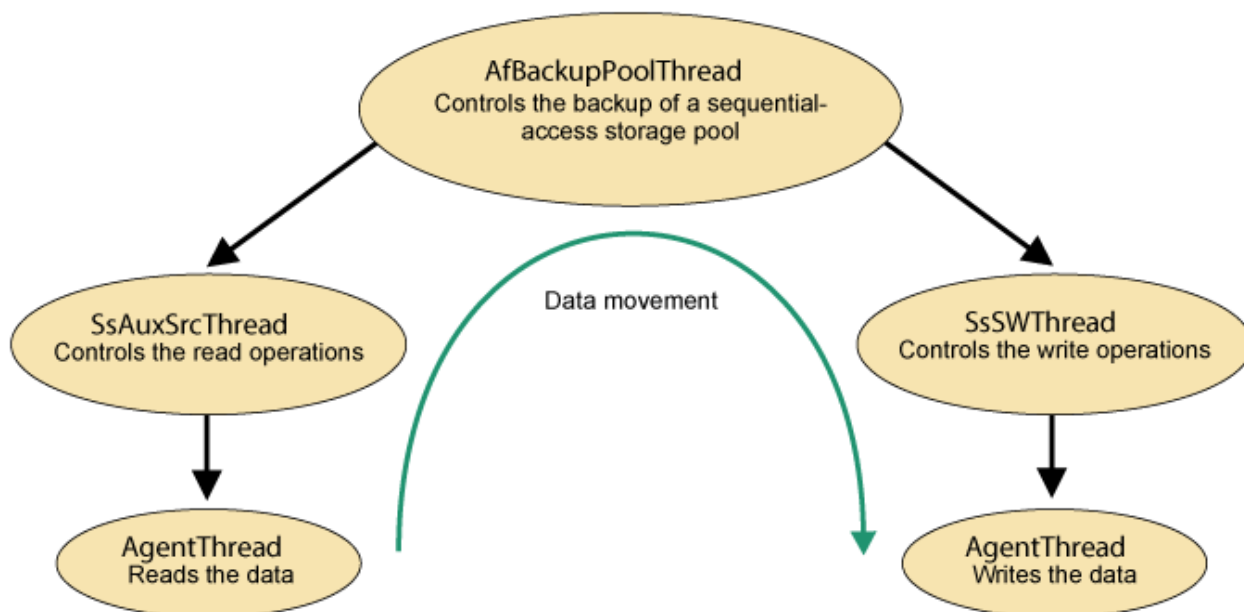


Figure 10. Threads for backup of sequential-access storage pools

Data movement starts with the AgentThread that reads the data from the volume that is being backed up. That data is processed through the SsAuxSrcThread, AfBackupPoolThread, and SsSWThread threads.

Data movement ends when the data is written on the target volume by the AgentThread thread that writes the data.

Copying active data for storage pool volumes

The main thread for a copy operation for a storage pool volume is called DfCopyActiveDataThread. The main thread starts one or two child threads. Each child thread controls a thread that is called AgentThread. See [Figure 11 on page 97](#).

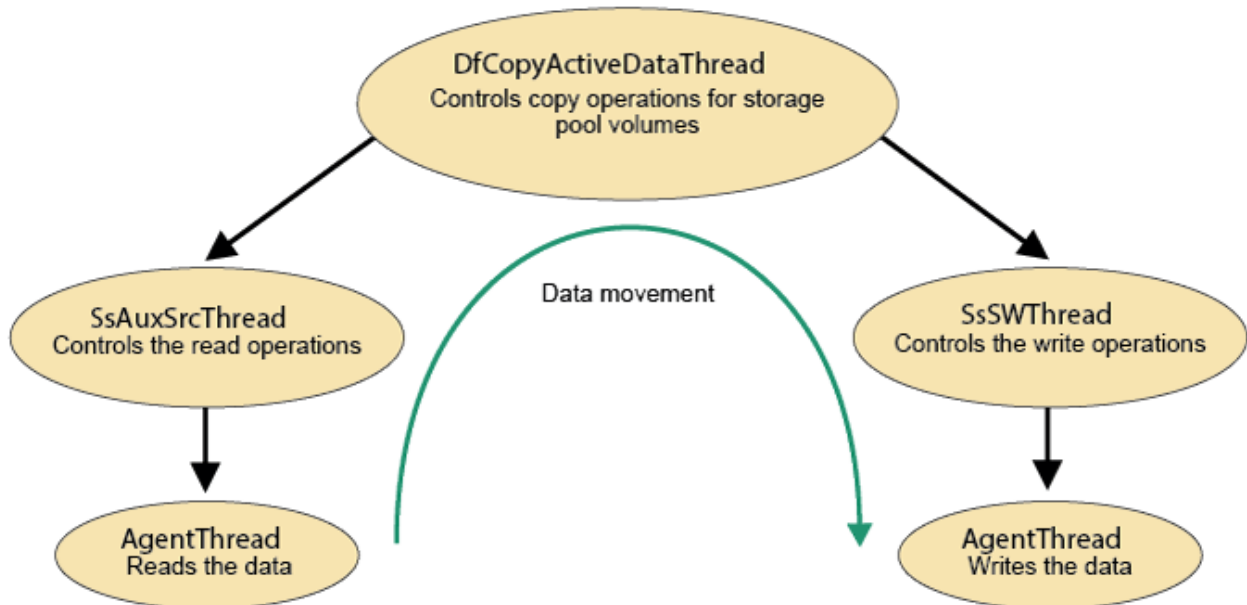


Figure 11. Threads for copying active data for storage pools

A typical data movement operation starts with an AgentThread that reads an object from a volume that is being copied. That data is processed through the SsAuxSrcThread, DfCopyActiveDataThread, and SsSWThread threads. Data movement ends when the data is written on the target volume by the AgentThread thread that writes the data.

Replicating data from a source replication server

The main threads for replicating data from a source replication server to a target replication server are NrReplicateFilespace. This thread determines the data that must be replicated, which is a database heavy task and it is expected that database activities are dominant. NrReplicateBatch threads then send the data to the target replication server over the network. To read the data, the NrReplicateBatch threads start a child thread, SsAuxSrcThread, which controls the read operations. The NrReplicateBatch thread sends the data that is identified by the NrReplicateFilespace threads to the target replication server. See [Figure 12 on page 98](#).

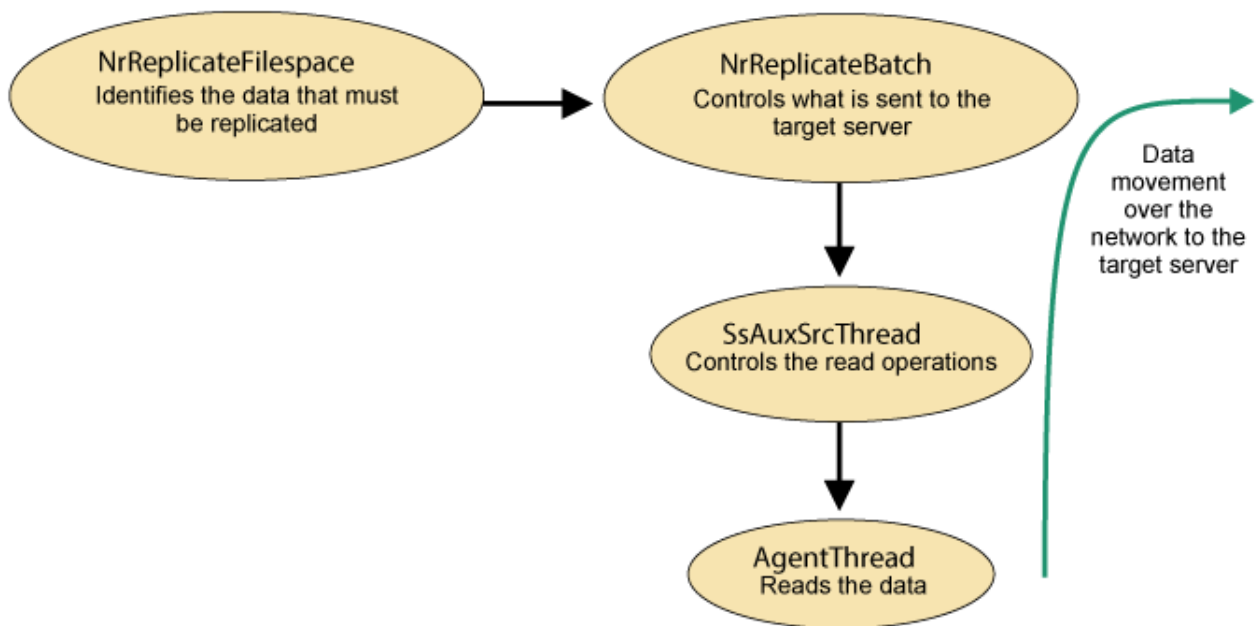


Figure 12. Threads for replicating data from a source replication server

Time spent in "Thread Wait" by the **NrReplicateBatch** threads can be spent in one of the following ways:

- Waiting for the **NrReplicateFilespace** thread to provide lists of files to replicate
- Waiting for the **SsAuxSrcThread** to read the source data from storage

The **NrReplicateBatch** threads control the network messages that are sent to the target replication server and the database.

Expire inventory

The main thread for expire inventory is **ExpirationProcessThread**. Expire inventory does not move data and is a database intensive operation. It is expected that database operations are dominant in this thread.

There might be several of these threads active, depending on the **RESOURCE** option that is used. See [Figure 13 on page 99](#).

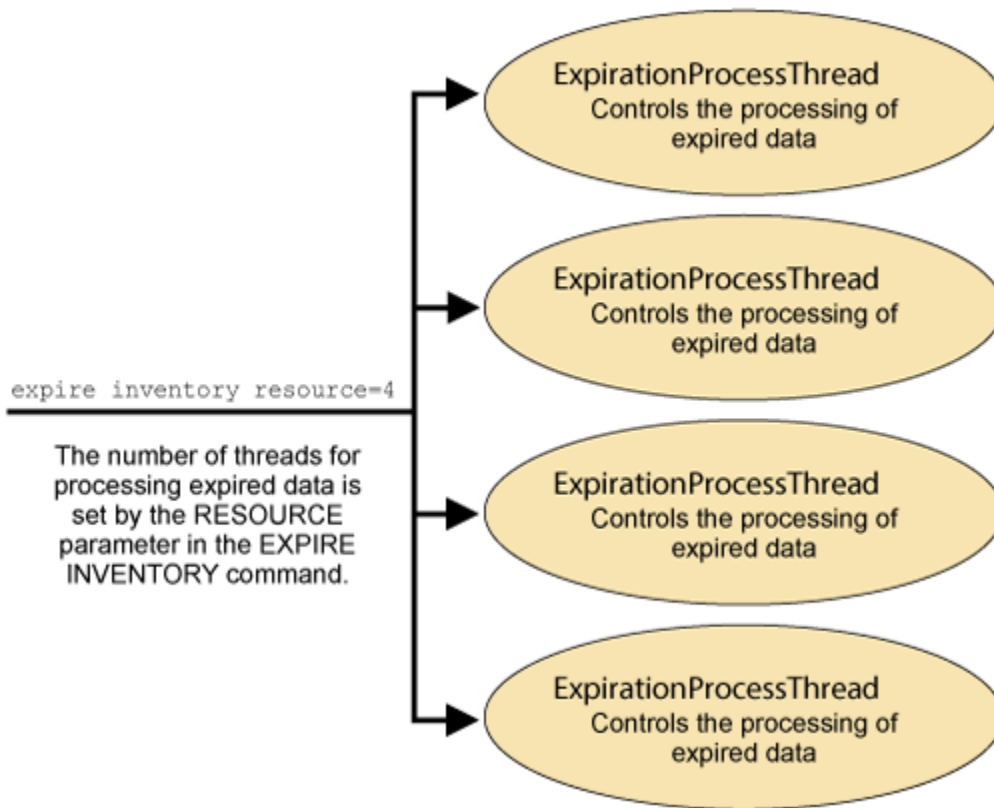


Figure 13. Threads for expire inventory

Client instrumentation report

Use client instrumentation to collect performance data on the IBM Storage Protect backup-archive client.

Collecting instrumentation data with the client

IBM Storage Protect client instrumentation identifies the elapsed time that is spent performing specific activities. By default, instrumentation data is automatically collected by the backup-archive client during backup or restore processing.

About this task

To disable or later enable instrumentation, use the `enableinstrumentation` option.

With this option enabled, you do not have to wait for a customer service representative to direct you to collect performance data when a problem occurs. Instead, the data is collected whenever you run a backup or restore operation. This feature can be helpful because you do not have to re-create the problem just to collect performance data. The information is already collected by the client.

The default setting of this option is `yes`, which means that instrumentation data is collected even if you do not specify this option. Typically, collecting instrumentation data by default has no measurable impact on regular performance.

By default, the output is appended to the instrumentation log file (`dsminstr.log`) in the directory that is specified by the `DSM_LOG` environment variable. If you did not set the `DSM_LOG` environment variable, the instrumentation log file is stored in the current directory (the directory where you started the **dsmc** command).

You can optionally change the name and location of the instrumentation log file by using the `instrlogname` option. You can also control the size of the log file by specifying the `instrlogmax` option.

Instrumentation data is not collected for the backup-archive client GUI or web client GUI.

The `enableinstrumentation` option replaces the `-TESTFLAG=instrument:detail`, `-TESTFLAG=instrument:API`, and `-TESTFLAG=instrumentation:detail/API` options that are used in previous versions of the client.

Procedure

To collect client instrumentation data, use one of the following methods:

- By default, client instrumentation data is collected during backup or restore processing so you do not need to update the client options file.

However, if you need to turn off the collection of instrumentation data, set the `enableinstrumentation no` option in the client options file (`dsm.opt` on Windows or `dsm.sys` on UNIX and Linux clients).

To later enable instrumentation, set `enableinstrumentation yes` or remove the option from the client options file.

- If the `enableinstrumentation no` option is set in the client options file, you can start client instrumentation when you run a backup or restore operation by including the `-enableinstrumentation=yes` option in a command.

For example, on Windows clients, start a selective backup and client instrumentation by issuing the following command:

```
dsmc sel c:\mydir\* -subdir=yes -enableinstrumentation=yes
```

For example, on UNIX and Linux clients, start a selective backup and client instrumentation by issuing the following command:

```
dsmc sel /home/mydir/* -subdir=yes -enableinstrumentation=yes
```

Similarly, you can turn off client instrumentation when you run a backup or restore operation by including the `-enableinstrumentation=no` option in a command.

Example

The following example shows the type of data that is collected:

```
PROCESS[4428] Starting TSM Instrumentation Report: Mon Apr 18 10:58:05 2016
=====>PROCESS[4428] NEW COMMENCE REPORT<=====
-----
PROCESS[4428] TSM Client final instrumentation statistics: Mon Apr 18 10:58:05 2016
Instrumentation class: Client detail
Completion status: Success
-----
-----
No instrumented activity reported for thread 4420
-----
-----
Detailed Instrumentation statistics for
Thread: 5076 Elapsed time = 510.979 sec
Section                Actual(sec)    Average(msec)    Frequency used
-----
Compute                0.218          0.0              27535
BeginTxn Verb          0.000          0.0              32
Transaction            0.374          11.7             32
File I/O               2.668          0.1             20702
Compression            32.105         1.2             27520
Data Verb              445.225        64.3            6927
Confirm Verb           0.000          0.0              1
```

EndTxn Verb	0.000	0.0	32
TCP Read	29.422	198.8	148
Thread Wait	0.905	904.8	1
Other	0.062	0.0	0

Detailed Instrumentation statistics for

Thread: 5532 Elapsed time = 438.018 sec

Section	Actual(sec)	Average(msec)	Frequency used
Process Dirs	0.140	9.4	15
Solve Tree	0.000	0.0	1
Sleep	0.062	62.4	1
TCP Read	0.546	39.0	14
Thread Wait	437.206	950.4	460
Other	0.062	0.0	0

Detailed Instrumentation statistics for

Thread: 5620 Elapsed time = 512.383 sec

Section	Actual(sec)	Average(msec)	Frequency used
Sleep	0.125	62.4	2
TCP Read	0.796	44.2	18
Thread Wait	510.495	1012.9	504
Other	0.967	0.0	0

No instrumented class associated with thread 6108

Current command:

sel c:\fileLoad* -subdir=yes

IBM Tivoli Storage Manager

Command Line Backup-Archive Client Interface

Client Version 7, Release 1, Level 6.18 20160418A

Client date/time: 04/18/2016 10:58:05

Options settings:

```

    BACKUPREGISTRY: YES
    CHANGINGRETRIES: 4
    COLLOCATEBYFILESPEC: NO
    COMMMETHOD: TCP/IP
    COMPRESSALWAYS: YES
    COMPRESSION: YES
    DEDUPCACHEPATH: c:\Program Files\Tivoli\tsm\baclient
    DEDUPCACHESIZE: 256
    DEDUPLICATION: NO
    DISKBUFFSIZE: 32
    ENABLEDEDUPCACHE: YES
    ENABLELANFREE: NO
    ENCRYPTIONTYPE: AES128
    FOLLOWSYMBOLIC: CLC
    IMAGEGAPSIZE: 32
    LANFREECOMMMETHOD: NAMED PIPE
    MAKESPARSEFILE: YES
    MAXCMDRETRIES: 2
    MEMORYEFFICIENTBACKUP: NO
    NODENAME: OEMTEST10
    PASSWORDACCESS: TRUE
    PRESERVELASTACCESSDATE: NO
    PROCESSORUTILIZATION: 0
    REPLACE: TRUE
    RESOURCEUTILIZATION: 2
    SKIPMIGRATED: NO
    SKIPNTPERMISSIONS: NO
    SKIPNTSECURITYCRC: NO
    SNAPSHOTCACHESIZE: 1
    SUBDIR: TRUE

    TAPEPROMPT: NO
    TCPBUFFSIZE: 32 KB
    TCPNODELAY: YES
    TCPSENBUFFSIZE: 0 KB
    TCPWINDOWSIZE: 63 KB
    TXNBYTELIMIT: 25600K
    VERBOSE: VERBOSE

```

```

-----
Session established with server ARC1: AIX
Server Version 7, Release 1, Level 4.100
Server date/time: 04/18/2016 08:54:40 Last access: 04/18/2016 08:37:01

```

```

Total number of objects inspected:      79
Total number of objects backed up:      79
Total number of objects updated:        0
Total number of journal objects:        0
Total number of objects rebound:        0
Total number of objects deleted:        0
Total number of objects expired:        0
Total number of objects failed:         0
Total number of objects encrypted:      0
Total number of bytes transferred:      212.71 MB
LanFree data bytes:                     0 B
Data transfer time:                     445.11 sec
Network data transfer rate:              489.35 KB/sec
Aggregate data transfer rate:            426.23 KB/sec
Total number of bytes pre-compress:     671,102,861
Total number of bytes post-compress:    222,963,689
Total number of objects grew:           0
Total number of retries:                0
Objects compressed by:                  67%
Total data reduction ratio:              66.77%
Elapsed processing time:                 00:08:31
Average file size:                      8.10 MB

```

```

PROCESS[4428] Ended TSM Instrumentation Report: Mon Apr 18 11:06:38 2016
-----

```

Related information

[Enableinstrumentation](#)

[Instrlogname](#)

[Instrlogmax](#)

Client instrumentation categories

IBM Storage Protect client instrumentation reports the elapsed times for many process categories.

The following table lists the client instrumentation categories that are tracked and the activity that is timed.

Table 9. Client instrumentation categories

Category	Activity
Query Server Dirs	Receiving the server inventory directories for incremental backup
Query Server Files	Receiving the server inventory files for incremental backup
Process Dirs	Scanning for files to back up
Cache Examine	Scanning the local disk cache database for files to expire
Solve Tree	Determining directory structure
Compute	Computing throughput and compression ratio
BeginTxn Verb	Building transactions
Transaction	File open, close, and other miscellaneous operations
File I/O	File read and write

Table 9. Client instrumentation categories (continued)

Category	Activity
Compression	Compressing and uncompressing data
Encryption	Encrypting and decrypting data
CRC	Computing and comparing CRC values
Data Verb	Sending and receiving data to and from the server (points to the network or IBM Storage Protect server)
Confirm Verb	Response time during backup for server confirm verb
TCP Read	The amount of time to read data from server.
EndTxn Verb	Server transaction commit and tape synchronization (points to the IBM Storage Protect server)
Other	Everything else that is not tracked already

Cloud instrumentation processes

IBM Storage Protect reports the time that it takes for certain processes that run in a cloud environment to complete.

The following table lists the cloud instrumentation processes that are tracked and the activity that is timed.

Table 10. Cloud instrumentation processes

Process	Time that is spent on this activity
INST_CLOUD_CONNECT	Connecting to the cloud.
INST_CLOUD_CONT	Creating, deleting, or managing cloud containers.
INST_CLOUD_DELETE	Deleting objects from cloud containers.
INST_CLOUD_ATCH	Attaching to the IBM Storage Protect server Java virtual machine (JVM).
INST_CLOUD_DTCH	Detaching from the IBM Storage Protect server JVM.
INST_CLOUD_STATS	Collecting and reporting cloud statistics for the Operations Center.
INST_CLOUD_READ	Read operations from the specified cloud provider.
INST_CLOUD_WRITE	Write operations to the specified cloud provider.

Virtual machine instrumentation categories

IBM Storage Protect virtual machine (VM) instrumentation reports the elapsed times for many process categories.

The following table lists the virtual machine instrumentation categories that are tracked and the activity that is timed.

Table 11. Virtual machine instrumentation categories

Category	Activity
VM Snapshot	Time that is spent generating and removing a VM guest snapshot by using the VMware Infrastructure Software Development Kit (VI SDK). Some of the work is completed asynchronously, for example, snapshot delete.
VM Send Data	Time that is spent sending data to the IBM Storage Protect server. Data processing includes client-side data deduplication and the Network Send phase.
VM Get Data	Time that is spent retrieving data from the IBM Storage Protect server. This category includes the following activities: <ul style="list-style-type: none"> Retrieving control files from the IBM Storage Protect server during incremental backup. Buffering data that is received during VM guest restore, and is then written out by using VM I/O writes to VMware.
VM Query	Time that is spent querying the IBM Storage Protect server to determine the following: <ul style="list-style-type: none"> Whether data deduplication and compression, or both, are enabled. Whether file space queries for nodes on virtual machines are used.
VM Query VE	Time that is spent querying whether VMware VM guest disks are changed. This category uses the VMware VI SDK to identify a set of changed blocks.
VM Assign	Time that is spent assigning file groups to the IBM Storage Protect server, by using the <code>dsmGroupHandler</code> function.
VM VCM Lock	Time that is spent locking a semaphore during Volume Control Manager (VCMLIB) API calls. The time is spent on the following functions: <ul style="list-style-type: none"> Locking for volume control data reads or updates Manipulating data by block control Retrieving megablocks in a volume
VM Transaction	Time that is spent processing transactions with the IBM Storage Protect server.
VM I/O	Time that is spent reading and writing data to and from the VIX Disk Library for Virtual Disk Development Kit (VDDK) disks for a VM. Performance can vary depending on whether thin or thick-provisioned disks are used and if the disks are lazy zeroed.
VM Control File I/O	Time that is spent reading and writing VM guest control (CTL) files during VM guest backup and restore operations.
Thread Wait	Time that is spent opening and closing VDDK disks on IBM Storage Protect for Virtual Environments clients. <p>Opening and closing of VM disks is serialized for an IBM Storage Protect client instance. The time that it takes to open and close the VM disks includes mounting the disk on the client or proxy system.</p>

API instrumentation report

Use API instrumentation to collect performance data on applications that use the IBM Storage Protect API.

API instrumentation identifies the time that elapsed during application activities. It is used for applications and products that use the API. The following products use the API:

- IBM Storage Protect Snapshot
- IBM Storage Protect for Mail
- IBM Storage Protect for Databases
- IBM Storage Protect for Virtual Environments
- IBM Storage Protect for Enterprise Resource Planning

Collecting instrumentation data with the API

API instrumentation identifies the time that elapsed during application activities. It is used for applications and products that use the IBM Storage Protect API. By default, instrumentation data is automatically collected by the API during backup or restore processing.

About this task

To disable or later enable instrumentation, use the `enableinstrumentation` option.

With this option enabled, you do not have to wait for a customer service representative to direct you to collect performance data when a problem occurs. Instead, the data is collected whenever you run a backup or restore operation. This feature can be helpful because you do not have to re-create the problem just to collect performance data. The information is already collected by the API.

The default setting of this option is `yes`, which means that instrumentation data is collected even if you do not specify this option. Typically, collecting instrumentation data by default has no measurable impact on regular performance.

By default, the output is appended to the instrumentation log file (`dsminstr.log`) in the directory that is specified by the `DSM_LOG` environment variable (or the `DSMI_LOG` environment variable for API-dependent products such as IBM Storage Protect for Databases: Data Protection for Microsoft SQL Server and IBM Storage Protect for Mail: Data Protection for Microsoft Exchange Server). If you did not set the `DSM_LOG` environment variable, the instrumentation log file is stored in the current directory (the directory where you started the **dsmc** command).

You can optionally change the name and location of the instrumentation log file by using the `instrlogname` option. You can also control the size of the log file by specifying the `instrlogmax` option.

The `enableinstrumentation` option replaces the `-TESTFLAG=instrument:API` option that is used in previous versions of the API.

Procedure

To collect API instrumentation data, use one of the following methods:

- By default, API instrumentation data is automatically collected during backup or restore processing so you do not need to update the client options file.

However, if you need to turn off the collection of instrumentation data, set the `enableinstrumentation no` option in the client options file (`dsm.opt` on Windows or `dsm.sys` on UNIX and Linux clients).

To later enable instrumentation, set `enableinstrumentation yes` or remove the option from the client options file.

- To turn API instrumentation on the command-line interface, append the following option to the end of a command:

```
-enableinstrumentation=yes
```

To turn off API instrumentation on the command-line interface, append the following option to the end of a command:

```
-enableinstrumentation=no
```

Results

The categories for tracking API activities are different from the client instrumentation categories.

Example

The following example shows the type of data that is collected:

```
PROCESS[4120] Starting TSM Instrumentation Report: Mon Apr 18 10:43:13 2016
=====>PROCESS[4120] NEW COMMENCE REPORT<=====
-----
PROCESS[4120] TSM Client final instrumentation statistics: Mon Apr 18 10:43:13 2016
Instrumentation class: API
Completion status: Success
-----
-----

Detailed Instrumentation statistics for
Thread: 5472 Elapsed time =      3.354 sec

Section                Actual(sec)    Average(msec)    Frequency used
-----
Waiting on App          3.354          838.5            4
API Send Data           0.000           0.0             3
Other                   0.000           0.0             0
-----

Detailed Instrumentation statistics for
Thread: 4208 Elapsed time =      9.703 sec

Section                Actual(sec)    Average(msec)    Frequency used
-----
Waiting on App          4.009          167.1            24
API Send Data           4.914          614.3             8
API Query                0.062           31.2             2
API End Txn             0.499          166.4             3
API Misc                 0.218           72.8             3
Other                   0.000           0.0             0
-----

Detailed Instrumentation statistics for
Thread: 2268 Elapsed time =     10.109 sec

Section                Actual(sec)    Average(msec)    Frequency used
-----
Waiting on App          9.532         1361.7             7
API Query                0.312           52.0             6
API End Txn             0.187          187.2             1
API Misc                 0.078           78.0             1
Other                   0.000           0.0             0
-----

Detailed Instrumentation statistics for
Thread: 4276 Elapsed time =     18.502 sec

Section                Actual(sec)    Average(msec)    Frequency used
-----
```

Waiting on App	16.193	476.3	34
API Query	0.842	49.6	17
API Misc	1.466	209.5	7
Other	0.000	0.0	0

PROCESS[4120] Ended TSM Instrumentation Report: Mon Apr 18 10:43:32 2016

Related information

[Enableinstrumentation](#)

[Instrlogname](#)

[Instrlogmax](#)

API instrumentation categories

IBM Storage Protect API client instrumentation reports the elapsed times for many process categories.

The following table lists the API client instrumentation categories that are tracked and the activity that is timed.

Table 12. API instrumentation categories

Category	Activity
Waiting on App	Time that the IBM Storage Protect API is waiting on the application that is sending IBM Storage Protect data. For example, the time spent waiting for a database application to send IBM Storage Protect data. If this value is high, focus your performance analysis on the application that is sending the data as well as the disk performance.
API Send Data	Time that is spent sending data to the IBM Storage Protect server. If the value is high, there might be a network problem or a storage pool performance problem on the IBM Storage Protect server.
API Query	Time that is spent querying the IBM Storage Protect server for information.
API Get Data	Time spent retrieving data from the IBM Storage Protect server. A high value might represent network problems between the server and client or server storage pool performance problems. For example, slow disk speeds or tape mount times.
API End Txn	Time that is spent in committing the current transaction to the IBM Storage Protect server. If the value is high, consider changing the settings that might use larger client transactions, or examine the server active log write performance.
API Dedup fingerprint	Time that is spent calculating the segmentation sizes of the incoming data. This category is a CPU intensive operation.
API ICC Digest (dedup)	Time that is spent computing the hash for the deduplication segments. This category is a CPU intensive operation.

Table 12. API instrumentation categories (continued)

Category	Activity
API Query Dedup Cache	Time that is spent querying the deduplication cache on local disk for deduplication segments.
API Query Server Dedup	Time that is spent querying the IBM Storage Protect server for data deduplication segments. If this value is high, examine database performance on the IBM Storage Protect server.
API Misc	Other minor activities of the IBM Storage Protect API client.

Scenarios for analyzing instrumentation data

Scenarios can help to illustrate how to use and interpret instrumentation reports.

Scenario: Improving performance of client backups

The scenario illustrates how to interpret client instrumentation reports to resolve a problem with backup operations.

Problem

Dave, an IBM Storage Protect administrator, is experiencing a performance degradation of backup operations on a file server. He is backing up the file server from an AIX backup-archive client to an AIX server, and the performance has degraded from an average throughput of 32 MB per second to 15 MB per second in the last few days.

Goal

Dave expects to restore his previous throughput level during backups.

Data collection

Dave collects the following data:

- Client instrumentation data
- Server monitoring data, which is generated by the servermon component

Analysis and determination of the bottleneck

During the backup operation, the data flows from the client, over the network, to the IBM Storage Protect server. The data is transferred from the disk system through a host bus adapter (HBA) that connects the disk to the IBM Storage Protect client. The client backs up the data to the server through a local area network (LAN) connection. A Network Interface Card (NIC) connects the client to the LAN and a separate NIC connects the LAN to the server. The data is backed up from the server to disk and to a tape library through separate HBA devices.

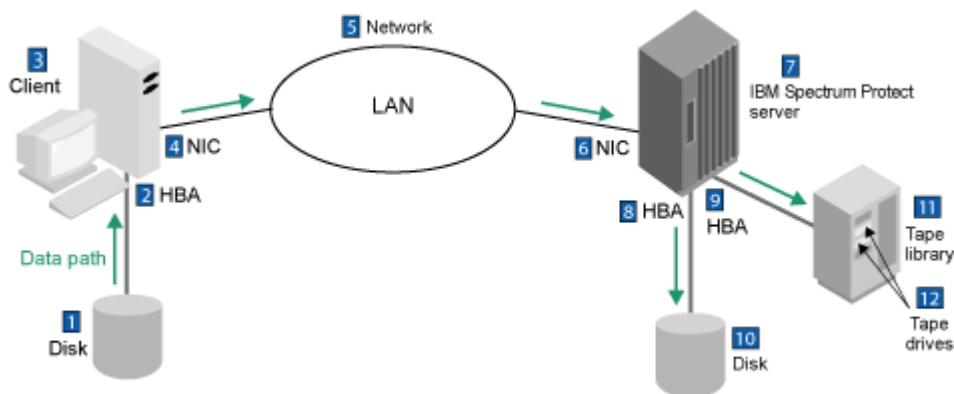


Figure 14. Data flow for client backup operations over a local area network

Following the flow of data, Dave looks at the client data first.

The most recent backup operation finished with the following statistics:

```

Total number of objects inspected:      1
Total number of objects backed up:     1
...
Total number of bytes transferred:    11.80 GB
LanFree data bytes:                   11.80 GB
Server-Free data bytes:               0 B
Data transfer time:                   216.01 sec
Network data transfer rate:           57,294.91 KB/sec
Aggregate data transfer rate:         16,542.69 KB/sec
Elapsed processing time:              00:12:28
Average file size:                   11.66 GB

```

To determine where the slowdown occurs, Dave runs a test backup of the file server with the `testflag=instrument:detail` option. By setting this option, Dave collects client instrumentation data about the backup operation. He reviews the instrumentation output. The client instrumentation report shows that, during the backup, a large amount of time was spent on File I/O operations. The elapsed time of the backup was 746.666 seconds, and 524.380 seconds were spent in File I/O. By looking at the client instrumentation categories in [Table 9 on page 102](#), Dave knows that the File I/O category represents the elapsed time for file read and write operations.

Section	Actual (sec)	Average(msec)	Frequency used
Thread: 2571 Elapsed time	746.666 sec		
Process Dirs	0.000	0.0	0
Solve Tree	0.000	0.0	0
Compute	0.234	0.0	48345
BeginTxn Verb	0.000	0.1	2
Transaction	0.715	357.5	2
File I/O	524.380	10.8	48346
Compression	0.000	0.0	0
Encryption	0.000	0.0	0
CRC	128.042	2.6	48398
Delta	0.000	0.0	0
Data Verb	87.912	1.8	48345
Confirm Verb	0.136	8.5	16
EndTxn Verb	2.234	1117.0	2
Other	4.513	0.0	0

The rate of data transfer is calculated by dividing the file size by the time that is recorded in the File I/O field in the instrumentation output:

```
(11.66 GB x 1024 MB/GB) / 524.380s = 22.8 MB/s
```

The backup reads the data only from the client disk, so Dave suspects a problem with the disk on the client system.

Problem resolution

Upon further investigation, Dave discovers that the AIX file system was recently mounted with the `cio` option, which enabled concurrent I/O in the file system. He concludes that mounting the file system with the AIX concurrent I/O option caused a degradation in backup performance. Concurrent I/O prevents file system read-ahead operations.

Dave modified the system settings to ensure that the AIX file system is not mounted with the `cio` option. As a result, the performance of backup operations is restored to its previous level.

Scenario: Improving performance of migration operations

This scenario illustrates how to interpret server monitoring data to improve migration operations.

Problem

Kate is an IBM Storage Protect administrator, and she noticed that the disk-to-tape storage pool migrations run slowly on her Windows server.

Goal

Kate expects the write time to be a value close to the capabilities of the drive.

Data collection

Kate collects the server monitoring data, which is automatically collected by the `servermon` component.

Analysis and determination of the bottleneck

During the migration operation, data is read from the disk and written to tape. The data is transferred from the disk through a host bus adapter (HBA) to the IBM Storage Protect server. The data is transferred from the server through a separate HBA to the tape system.

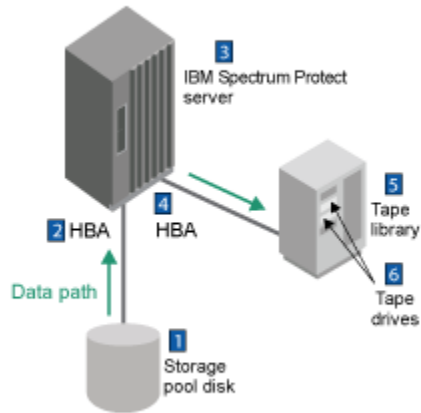


Figure 15. Data flow for migration operations

Following the flow of data, Kate looks at the disk read activity first.

As Kate investigates, she notices that all tape operations are slow, including client backups, reclamation, and database backups to tape. Kate suspects a problem with the tape system, which might be causing slow write times. The tape system is an LTO library with SCSI-attached tape drives.

To determine where the slowdown occurs, Kate runs a disk-to-tape storage pool migration test and collects server instrumentation data, by using the data that is automatically collected in the servermon component. She locates the files that were collected by the servermon component, and finds the files that are time-stamped during a time when migration was slow. She looks for the files that are in the form: YYYYMMDD-HHMM-YYYYMMDD-HHMM-instr.txt, where the time stamps represent the start and end times when the output was gathered.

Kate reviews the file, and finds Thread 61, which processed a total of 966912 KB. Then, she finds the corresponding thread that has a Tape Write category in it, which processed about the same amount of data. She found Thread 34, which processed 968192 KB.

Thread 61 DfMigrationThread (Win Thread ID 4436) 17:39:076-->17:47:38							
Operation	Count	Tottime	Avgtime	Min-time	Max-time	Inst Tput	Total KB
Disk Read	3777	22.680	0.006	0.000	0.031	42632.8	966912
Thread Wait	3778	487.450	0.129	0.016	0.313		
Unknown		0.061					
Total		510.191				1895.2	966912
Thread 34 AgentThread (Win Thread ID 5340) 17:39:07.816-->17:47:38.007							
Operation	Count	Tottime	Avgtime	Min-time	Max-time	Inst Tput	Total KB
Tape Write	30257	508.816	0.017	0.000	0.141	1902.8	968192
Tape Data Copy	31661	0.863	0.000	0.000	0.016		
Thread Wait	3777	0.220	0.000	0.000	0.016		
Unknown		0.292					
Total		510.191				1897.7	968192

Kate used the server instrumentation categories in [“Server instrumentation categories”](#) on page 87 to understand the meaning of the Disk Read, Tape Write, and Thread Wait categories. The server instrumentation output shows the following characteristics:

- The Thread 61 data shows a high value for the Thread Wait field (487.450 seconds), and low value for the Disk Read field (22.680 seconds). This result indicates that Thread 61 is not a concern because

the instantaneous throughput is acceptable and the thread was in a waiting phase. The disk is not a problem.

Tip: The instantaneous throughput, `Inst Tput`, is calculated by dividing the `Total KB` value by `Tottime` value.

- The `Thread 34` data shows that there is a low value for the `Thread Wait` field (0.220 seconds), and a high value for the `Tape Write` field (508.816 seconds). The output also shows that the `Inst Tput` value for the tape write is very slow (1902.8 KB per second). She concluded that the problem is related to the tape system, as shown by the high `Tape Write` value.

Problem resolution

Kate investigates the following probable sources of the tape system problem:

- The tape attachment path
- The device driver level of the tape drive
- The SCSI adapter driver level
- The SCSI adapter settings

After the investigation, Kate upgrades the SCSI adapter device driver. The disk-to-tape storage pool migration improves to 75% of the native capacity. Client backups to tape are also much faster.

Tip: All server speeds depend on the hardware that is used. The values in this scenario might not apply to your system. Refer to the performance characteristics of your tape drive or disk system to determine an acceptable level of performance.

Scenario: Improving performance of database backups

This scenario illustrates how to interpret API client instrumentation data and server monitoring data to improve database backup operations.

Problem

Dave, a system administrator, notices a performance degradation in server database backups after an IBM Storage Protect upgrade.

Goal

Dave expects the time to complete a database backup to be the same as it was before the upgrade.

Data collection

Dave collects the following data:

- API client instrumentation data
- Server monitoring data, which is generated by the `servermon` component

Dave runs the API client instrumentation for the entire duration of the database backup.

Analysis and determination of the bottleneck

During database backup operations, the database is read and then backed up to the target storage pool. The data is backed up from the disk where the IBM Storage Protect database is located through a host bus adapter (HBA) to the IBM Storage Protect server. The data is transferred from the server through a separate HBA to the tape system.

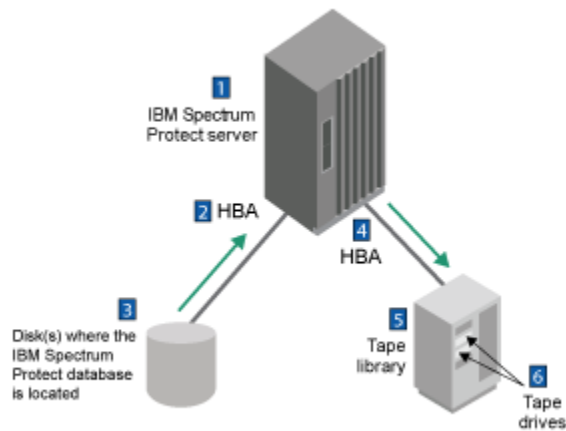


Figure 16. Data flow for database backup operations

Dave starts by reviewing data from the API client instrumentation. He finds that Thread 1 shows a backup period of almost 5 hours (17954.687 seconds). About 99% of the time was spent in the API send data category. By reviewing the API client categories in “API instrumentation categories” on page 107, Dave determines that most of the time was spent either sending data to the IBM Storage Protect server through the network, or writing the data to the backup device.

Thread: 1 Elapsed time = 17954.687 sec (1)

Section	Actual(sec)	Average(msec)	Frequency used
Waiting on App	168.155	2.8	59246
API Send Data	17786.518	300.3	59236 (2)
API Query	0.001	0.1	4
API End Txn	0.004	4.1	1
API Misc	0.009	3.0	3
Other	0.000	0.0	0

Dave reviews the server instrumentation data. He determines the time when the database backup was started by examining the IBM Storage Protect activity log. He locates the files that were collected by the servermon component, and finds the files that are time-stamped during the database backup. The file names have the following structure: YYYYMMDD-HHMM-show.txt. By examining a file that was time-stamped during the database backup, Dave finds the thread number for the database backup.

Tip: For backups of server databases, the associated node name is always \$\$_TSMDBMGR_\$\$_

Dave searches the data for a node that is named \$\$_TSMDBMGR_\$\$_, and finds the information for the database backup session:

```

Session 22486:  Type=Node,  Id=$$_TSMDBMGR_$$_
Platform=DB2/AIX64, NodeId=1, Owner=tsminst
SessType=4, Index=0, TermReason=0
threadId=24431
ProxyByAgent False
RecvWaitTime=0.000 (samples=0)
Backup Objects ( bytes ) Inserted: 0 ( 0.0 )
Backup Objects ( bytes ) Restored: 0 ( 0.0 )
Archive Objects ( bytes ) Inserted: 0 ( 0.0 )
Archive Objects ( bytes ) Retrieved: 0 ( 0.0 )
Last Verb ( ConfirmResp ), Last Verb State ( Sent )
Global id reports 0 mount points in use
Write MP count 0, read MP count 0 keepUsedMP = No.
  
```

The information shows that Thread 24431 is associated with the database backup. Dave opens the server instrumentation report from the same time frame. The file names for server instrumentation reports have

the format YYYYMMDD-HHMM-YYYYMMDD-HHMM-instr.txt, where the time stamps represent the start and end times when the output was gathered. He searches for Thread 24431, and finds the following output:

```
Thread 24431 00:21:34.695-->00:43:20.577
Operation      Count  Tottime  Avgtime  Mintime  Maxtime  InstTput  Total KB
-----
Network Recv   660678 1190.148    0.002    0.000    64.847  15556.7  18514797(3)
Network Send     21    0.000    0.000    0.000    0.000      0
Thread Wait    72323  112.404    0.002    0.000    33.003
Unknown              3.328
-----
Total              1305.881                14178.0  18514797
```

Dave finds that most of the time was spent in the Network Recv phase. By reviewing the server instrumentation categories in [“Server instrumentation categories” on page 87](#), he determines that most of the time was spent receiving data from the network.

Dave determines that the network is causing performance degradation. The client and server reports show lengthy times for sending and receiving data over the network

Problem resolution

Dave identifies network settings that were incorrectly set as a result of the upgrade. Dave corrects the settings, and the performance of the database backup time achieves the same level as before the upgrade.

Scenario: Improving performance of restore operations for database applications

This scenario illustrates how to interpret API client instrumentation data and server monitoring data to improve database application restore operations.

Problem

Kate, an IBM Storage Protect administrator, notices a performance degradation in restore operations for an SAP application database.

Goal

Kate expects to improve the performance of database restore operations.

Data collection

Kate collects the following data:

- API client instrumentation data
- Server monitoring data, which is generated by the servermon component

Kate runs the API client instrumentation and the servermon component for the entire duration of the restore operation. She runs a restore operation and collects the instrumentation data during a time of little or no activity on the IBM Storage Protect server. Kate uses IBM Storage Protect for Enterprise Resource Planning to back up the database for an SAP application.

Analysis and determination of the bottleneck

During the restore operation, data is read from the tape drives, sent over the network to the IBM Storage Protect client, and is written to the client disk. The data is transferred from the tape system through a host bus adapter (HBA) that connects the tape library to the server. The server transfers the data to the client system through a local area network (LAN) connection. A Network Interface Card (NIC) connects

the server to the LAN and a separate NIC connects the LAN to the client. The data is restored to the disk from the client through a separate HBA device.

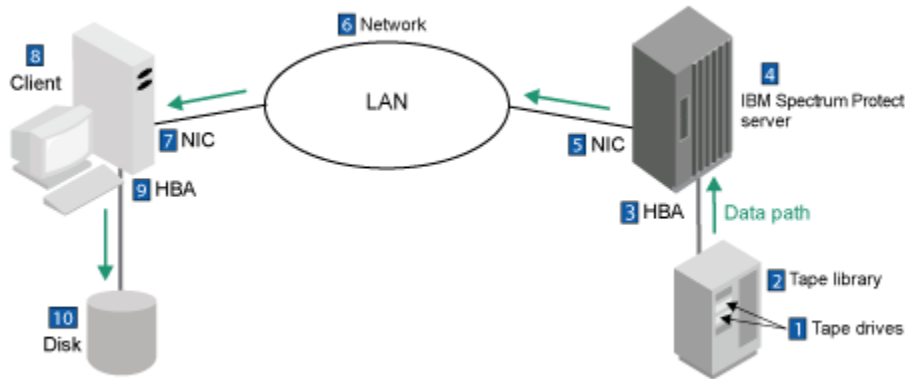


Figure 17. Data flow for restore operations for database applications

Following the flow of data, Kate starts to review the server instrumentation data. First, Kate notes the time when the restore was started by reviewing the IBM Storage Protect activity log. Then, she locates the files that were collected by the servermon component, and finds the files that are time-stamped around the time when the restore operation was slow. The instrumentation file names have the following structure: YYYYMMDD-HHMM-YYYYMMDD-HHMM-instr.txt.

Kate examines the instrumentation output to find a thread that was reading from tape. She finds Thread 131787. Data for the thread shows that 9.100 seconds was spent in the Tape Read category and has a fast instantaneous throughput (InstTput) of 269584.5 KB per sec. She reviews the server instrumentation categories in [“Server instrumentation categories”](#) on page 87 and finds that there is no problem reading from tape. The instantaneous throughput and high amount of thread wait time indicate that tape read is not the bottleneck. Kate sees that 99% of the time was spent in Thread Wait (1199.192 seconds):

Thread 131787 AgentThread parent=131782 13:16:25.938-->13:36:34.274							
Operation	Count	Tottime	Avgtime	Mintime	Maxtime	InstTput	Total KB
Tape Read	9583	9.100	0.001	0.000	0.354	269584.5	2453248
Thread Wait	9585	1199.192	0.125	0.000	267.561		
Unknown		0.042					
Total		1208.335				2030.3	2453248

Following the data flow, she looks for a thread with a Network Send category and a thread that processed approximately the same amount of data (2453248 KB). She finds Thread 131781, which processed a total of 2452368 KB. Thread 131781 is a psSessionThread thread, which shows the timing of the data as it was sent from the server, through the network to the client. Kate sees that the InstTput for the thread was 2052.8 KB per second, which was slower than expected. The Thread Wait phase took 0.786 seconds, which was less than 1% of the total time. The data indicates that most of the time was spent in the Network Send phase. The findings suggest performance degradation of the network on the client side.

```

Thread 131781 psSessionThread parent=299 13:16:25.938-->13:36:34.274
Operation      Count  Tottime  Avgtime  Mintime  Maxtime  InstTput  Total KB
-----
Data Copy      1      0.000    0.000    0.000    0.000
Network Recv   5      12.778    2.556    0.001    12.719      0.0      0
Network Send  19170  1194.666    0.062    0.000    267.561    2052.8    2452368
DB2 Fetch Prep  1      0.003    0.003    0.003    0.003
DB2 Fetch Exec  8      0.003    0.000    0.000    0.002
DB2 MFtch Exec  3      0.008    0.003    0.000    0.004
DB2 Fetch      8      0.000    0.000    0.000    0.000
DB2 MFtch      5      0.000    0.000    0.000    0.000
DB2 Commit     2      0.000    0.000    0.000    0.000
DB2 Reg Prep   2      0.003    0.002    0.000    0.003
DB2 Reg Exec   2      0.000    0.000    0.000    0.000
DB2 Reg Fetch  2      0.000    0.000    0.000    0.000
Thread Wait    9584    0.786    0.000    0.000    0.351
Unknown
-----
Total          1208.335                                2029.5    2452368

```

Next, Kate looks at the API client instrumentation report.

Detailed Instrumentation statistics for

Thread: 2911 Elapsed time = 1669.061 sec

Section	Actual(sec)	Average(msec)	Frequency used
Waiting on App	1435.153	94.8	15131
API Get Data	233.909	15.5	15131
Other	0.000	0.0	0

Detailed Instrumentation statistics for

Thread: 2902 Elapsed time = 1913.868 sec

Section	Actual(sec)	Average(msec)	Frequency used
Waiting on App	1681.437	110.4	15224
API Get Data	232.432	15.3	15224
Other	0.000	0.0	0

Detailed Instrumentation statistics for

Thread: 2893 Elapsed time = 2093.726 sec

Section	Actual(sec)	Average(msec)	Frequency used
Waiting on App	1926.681	127.1	15153
API Get Data	167.045	11.0	15153
Other	0.000	0.0	0

By reviewing the data from the API client instrumentation, Kate sees that the longest amount of time was spent on the `Waiting on App` category. This category shows the time that was spent by the application to process the restore data. After Kate reviews the API report, she determines that most of the restore time is spent waiting for the application database. As a next step, Kate could verify the performance of the disk system.

Problem resolution

Kate contacts the storage administrator and the SAP administrator to help resolve the problem. After further investigation, the storage administrator determines that the database layout is the source of the problem, and corrects it.

Part 4. Tuning the components

After you verify that you are following the guidelines for optimal configuration, learn about options and parameters that can help you to tune performance.

Chapter 10. Tuning Operations Center performance

The performance of the IBM Storage Protect Operations Center depends on the system resources of the hub and spoke servers and the processor and memory on the computer where the Operations Center is installed. Performance also depends on the number of client nodes and virtual machine file spaces that are being monitored by all servers.

Before you begin

Review the system requirements for the Operations Center before you install it.

About this task

Use the Operations Center System Requirements Calculator at [technote 225373](#) to estimate the system requirements for running the Operations Center and the hub and spoke servers that are monitored by the Operations Center.

Resource usage on the Operations Center computer

The number of administrators who are logged in to the Operations Center affects resource usage on the computer where the Operations Center runs.

Resource usage for each logged-in administrator

The number of administrators who are logged in at the same time, and the number of tasks that each is working on, affects the resource usage by the Operations Center. For example, the following situation can be typical:

- Running the Operations Center uses about 3% of a processor core (based on an Intel X7550 2.00 GHz or equivalent processor).
- Each administrator who is logged in to the Operations Center uses 2% - 3% of a processor core on the computer where the Operations Center runs. This usage level assumes that each administrator completes about 40 tasks per hour.
- The Operations Center might have 8 - 10 administrators who are logged in and completing tasks at the same time. A task might be one of the following activities:
 - Opening and viewing a page, window, or wizard in the interface, for example, a storage pool **Properties** page, a server **Quick Look** window, an **Alert Details** window, or the **Add Client** wizard.
 - Configuring settings on the **Settings** page
 - Issuing a command from the Operations Center command line

A typical user might complete 20 - 120 of these tasks per hour.

Memory for the JVM heap size

The Operations Center requires at least 128 MB of memory for the maximum Oracle Java virtual machine (JVM) heap size for up to eight concurrently logged-in administrators. This memory requirement assumes that each administrator completes about 40 tasks per hour within the user interface. The JVM uses a maximum of 1 GB, or 25% of the physical memory of the system (with systems that have more than 192 MB of memory). The heap size starts at a smaller amount but can increase up to the maximum.

If the system that is running the Operations Center has more than 1 GB of memory, no adjustments are needed. If you must adjust the memory for the JVM heap size, use the `-Xmx` option in the `jvm.options` file for the web server JVM. The `jvm.options` file is in the following directory, where `installation_dir` represents the directory where the Operations Center is installed:

- **AIX** | **Linux** `installation_dir/ui/Liberty/usr/servers/guiServer/`
- **Windows** `installation_dir\ui\Liberty\usr\servers\guiServer\`

If the `jvm.options` file is not in the directory, and you must adjust the memory, create the file in the directory.

Effect of the network on performance

The network that connects the hub server and the system where the Operations Center is installed can affect the performance of the Operations Center.

To achieve better system performance, use one of the following configurations:

- The Operations Center is on the same system as the hub server.
- The Operations Center is on a system that is located physically near the hub server system.

Also, consider facilitating upgrades and maintenance by designating a physical or virtual system that is separate from the production environment as the hub server.

Network latency

Network latency is the time interval between the following operations:

- The initiation of a send operation from a source system
- The completion of the matching receive operation by the target system

Latency between the Operations Center web server and web browsers

For the best responsiveness when logged in to the Operations Center, use a network connection with a round-trip latency that is no greater than 5 ms. This latency can typically be achieved when the systems are on the same local area network (LAN). Higher latencies might be acceptable but can cause degraded responsiveness. For example, the responsiveness across a wide area network (WAN) might not be acceptable to users of the Operations Center.

Latency between the Operations Center web server and the hub server

For the best responsiveness, use a network connection with a round-trip latency that is no greater than 10 ms. Lowest latency is achieved when both of these servers are installed on the same system or on the same LAN.

Latency between the hub server and spoke servers

Round-trip latency can be high, for example, 100 ms, or low, for example, 10 ms. However, with high latency, Operations Center pages that show details about an individual client, policy domain, server, or storage pool might take a longer time to be displayed. Also, if connection timeout issues occur, you might be required to adjust the **ADMINCOMMTIMEOUT** value on the hub and spoke servers. For information about what to do if connection timeout issues occur, see the following known issue on the IBM Support Portal at [technote 497591](#).

Estimating the network latency

You can use a **ping** command to estimate the network latency. To use a **ping** command, complete the following steps:

1. From the source system, ping the target system by issuing the following commands, where *remoteHost* represents the address for the target system:
 - **AIX** | **Linux** `ping -c 20 remoteHost`
 - **Windows** `ping -n 20 remoteHost`
2. Calculate the average of the intervals for all successful responses from the remote host. This calculation is an estimate of the round-trip latency.

Effect of status monitoring on performance

When you configure a server as a hub or spoke server, status monitoring is automatically enabled. Status monitoring requires extra resources on each server on which it is enabled.

Tip: In this topic, the term *client* represents both client nodes and virtual machine file spaces. Also, the resource estimates are approximate.

Server resource requirements for status monitoring

The resources that are required depend primarily on the number of clients that are managed by the hub and spoke servers.

Also, the hub server requires fewer resources if the spoke servers are running Tivoli® Storage Manager 7.1 or later, or IBM Storage Protect 7.1.3 or later, than it does if the spoke servers are running version 6.3.4 or a later modification of version 6.3.

Table 13 on page 121 summarizes the resource requirements for a server on which status monitoring is enabled.

Table 13. Resource requirements for a server on which status monitoring is enabled			
Resource requirement	Resource usage for the base level of up to 1000 clients	Resource usage for every 1000 clients over the base level	Example: Resource usage for a spoke server with 2000 clients
Processor usage The value is based on lab measurements that used the Intel X7550 2.00 GHz core.	1.1 processor cores	0.1 processor cores	1.2 processor cores
More space in the server database	2 GB if the server is at version 7.1 1 GB if the server is at version 7.1.1 or later	2 GB if the server is at version 7.1 1 GB if the server is at version 7.1.1 or later	4 GB if the server is at version 7.1 2 GB if the server is at version 7.1.1 or later
More space for the server archive log The value assumes that a full database backup is completed every 24 hours.	10 GB	10 GB	20 GB
Spoke server at version 6.3.4 or a later modification of version 6.3: Data transfer to the hub server over the network	30 - 60 MB per hour	30 - 60 MB per hour	60 - 120 MB per hour
Spoke server at version 7.1 or later: Data transfer to the hub server over the network	5 - 10 MB per hour	5 - 10 MB per hour	10 - 20 MB per hour

Consider adding a buffer of 25% - 50% to the database and log requirements for a server that has a heavy workload. For example:

- A server that is scheduled to back up hundreds of client nodes or virtual-machine file spaces daily
- A server that has many I/O operations per second (IOPS) due to operations such as data deduplication

Extra resource requirements for a hub server

The hub server must have adequate resources for the number of clients that it manages directly, according to the estimates in [Table 13 on page 121](#). For managing spoke servers, the hub server must also have the extra resources that are described in [Table 14 on page 122](#). Use the table to monitor items that are defined on the hub server and the spoke servers that are connected to it.

<i>Table 14. Extra resource requirements for a hub server</i>		
Resource requirement	For managing spoke servers at version 7.1 or later	For managing spoke servers at version 6.3.4 or a later modification of version 6.3
Processor usage The value is based on lab measurements that used the Intel X7550 2.00 GHz core.	Negligible	More processor resources, equal to 0.1 processor cores for every 1000 clients on all monitored servers (all clients on all spoke servers at version 6.3.4 or a later modification of version 6.3).
More space in the server database	Negligible	If the hub server is at version 7.1: More disk space for the database, equal to 2 GB for every 1000 clients on all monitored spoke servers at version 6.3.4 or a later modification of version 6.3. If the hub server is at version 7.1.1 or later: More disk space for the database, equal to 1 GB for every 1000 clients across all monitored spoke servers at version 6.3.4 or a later modification of version 6.3.
More space for the server archive log The value assumes that a full database backup is completed every 24 hours.	More disk space for the archive log, equal to 600 MB for every 1000 clients on all monitored spoke servers at version 7.1 or later.	More disk space for the archive log, equal to 10 GB for every 1000 clients on all monitored spoke servers at version 6.3.4 or a later modification of version 6.3.
IOPS capacity for the server database on the hub server	More I/O capability for the database volumes, to support 50 IOPS for every 1000 clients on spoke servers at version 7.1 or later. The estimate is based on an average I/O size of 8 KB.	More I/O capability for the database volumes, to support 200 IOPS for every 1000 clients on spoke servers at version 6.3.4 or a later modification of version 6.3. The estimate is based on an average I/O size of 8 KB.

For a hub server that manages spoke servers, you get optimal performance if the server database is on disks that can process 8 KB operations at the rate of at least 1000 IOPS. To get this IOPS capacity, use a single enterprise-level solid-state drive (SSD). If SSD is not an option, you might want to use a SAN-attached array of 15000-rpm serial-attached SCSI (SAS) hard disk drives, each capable of handling hundreds of 8 KB IOPS. The choice depends on the overall workload of the hub server.

Example of resource requirements for a hub server

Table 15 on page 123 shows a resource estimate for a hub server with spoke servers that are at version 6.3.4 or a later modification of version 6.3. Table 16 on page 124 shows a resource estimate for a hub server with spoke servers that are at version 7.1 or later. In both examples, the hub server has 1000 clients, and each of the five spoke servers has 2000 clients.

Table 15. Example of resource requirements for a hub server with spoke servers that are at version 6.3.4 or a later modification of version 6.3

Resource requirement	Resource usage for 1000 clients that are managed on the hub server	Resource usage on the hub server for the 10,000 clients that are managed on five spoke servers that are at version 6.3.4 or a later modification of version 6.3 (2000 clients on each)	Total estimated resource usage
Processor usage The value is based on lab measurements that used the Intel X7550 2.00 GHz core.	1.1 processor cores	1 processor core The estimate is based on 0.1 processor core for every 1000 clients on the spoke servers.	2.1 processor cores
More space in the server database	2 GB if the hub server is at version 7.1 1 GB if the hub server is at version 7.1.1 or later	20 GB if the hub server is at version 7.1 10 GB if the hub server is at version 7.1.1 or later	22 GB if the hub server is at version 7.1 11 GB if the hub server is at version 7.1.1 or later
More space for the server archive log The value assumes that a full database backup is completed every 24 hours.	10 GB	100 GB	110 GB
Spoke server: Data transfer to the hub server over the network	Not applicable	300 - 600 MB per hour The estimate is based on 30 - 60 MB per hour for every 1000 clients on the spoke servers.	300 - 600 MB per hour
IOPS capacity for the server database on the hub server	200 IOPS	2000 IOPS The estimate is based on 200 IOPS for every 1000 clients on the spoke servers.	2200 IOPS

Table 16. Example of resource requirements for a hub server with spoke servers that are at version 7.1 or later

Resource requirement	Resource usage for 1000 clients that are managed on the hub server	Resource usage on the hub server for the 10,000 clients that are managed on five spoke servers at version 7.1 or later (2000 clients on each)	Total estimated resource usage
Processor usage The value is based on lab measurements that used the Intel X7550 2.00 GHz core.	1.1 processor cores	Negligible	1.1 processor cores
More space in the server database	2 GB if the hub server is at version 7.1 1 GB if the hub server is at version 7.1.1 or later	Negligible	2 GB if the hub server is at version 7.1 1 GB if the hub server is at version 7.1.1 or later
More space for the server archive log The value assumes that a full database backup is completed every 24 hours.	10 GB	6 GB The estimate is based on 600 MB for every 1000 clients on the spoke servers.	16 GB
Spoke server: Data transfer to the hub server over the network	Not applicable	50 - 100 MB per hour The estimate is based on 5 - 10 MB per hour for every 1000 clients on the spoke servers.	50 - 100 MB per hour
IOPS capacity for the server database on the hub server	200 IOPS	500 IOPS The estimate is based on 50 IOPS for every 1000 clients on the spoke servers.	700 IOPS Consider establishing a baseline capacity of 1000 IOPS for the hub server database if the hub server manages any spoke servers.

Effect of the status refresh interval on performance

In the Operations Center, the status refresh interval is the number of minutes between status collection refreshes. Changing the default value of this interval can affect the performance of the Operations Center and the servers that are using status monitoring.

You can set this interval on the **Settings** page of the Operations Center or by issuing the **SET STATUSREFRESHINTERVAL** command on each hub or spoke server. Use the same interval on the hub and spoke servers. Using different intervals can reduce the accuracy of the information that is shown in the Operations Center.

Effect of decreasing the interval

If you decrease the interval to get more frequent refreshes of collected data, more data is processed and maintained, which uses more space in the server database. More frequent refreshes can also mean higher processor usage. The disks where the server database is located might also require higher input/output operations per second (IOPS).

If you decrease the interval by half, the server database and the archive log space that is required for status monitoring is doubled.

Also, do not decrease the interval to less than 5 minutes. An interval of less than 5 minutes can cause the following issues:

- Operations Center data that is supposed to be refreshed after the defined interval takes a longer time to be refreshed.
- Operations Center data that is supposed to be refreshed almost immediately when a related change occurs in the storage environment also takes a longer time to be refreshed.

Effect of increasing the interval

If you increase the interval to get less frequent refreshes of collected data, resource requirements are reduced. However, the data that is shown in some Operations Center views might not be current.

If you double the interval, the server database and the archive log space that is required for status monitoring is reduced by half.

If a hub server is connected to spoke servers over a network with high latency, consider increasing the interval to reduce the data that is sent over the network for status monitoring.

Chapter 11. Tuning server performance

Many factors must be considered when you are tuning the configuration of your IBM Storage Protect server for optimum performance. Review this information to evaluate settings for your operating system, key IBM Storage Protect operations, scheduling for the server and client workloads, and configuration for functions that require more server resources.

Server database and recovery log configuration and tuning

How you configure and size the database and recovery logs is essential to IBM Storage Protect performance.

Before you begin

The server records changes that are made to the database in the recovery log. The recovery log is used to maintain the database in a transactionally consistent state, and to maintain consistency across server startup operations. The recovery log is composed of an active log, an archive log, and optional logs, including the active log mirror and archive failover log. The following figure shows the IBM Storage Protect server, database, and recovery log in relation to each other.

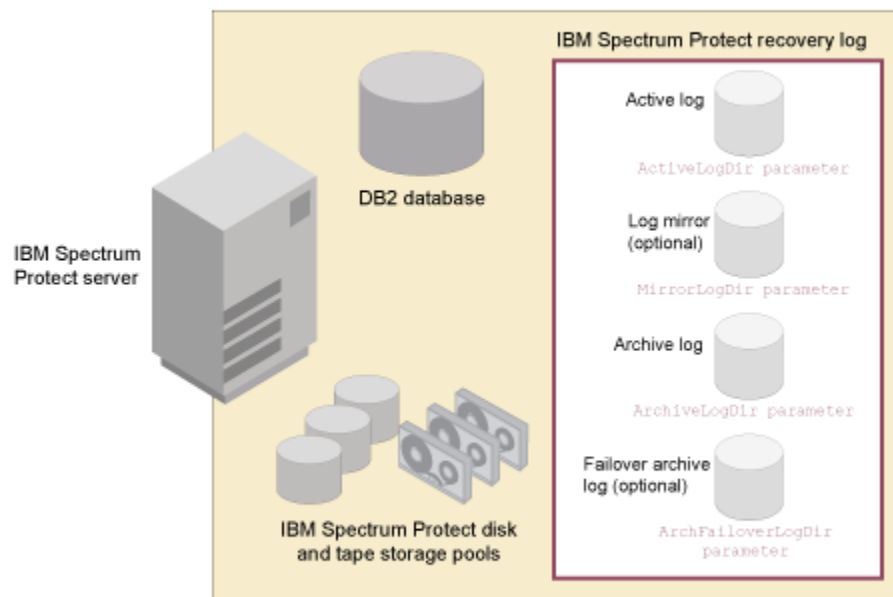


Figure 18. Components of the IBM Storage Protect database and recovery log

Database configuration and tuning

The IBM Storage Protect database contains information that is needed for server operations and information about client data. Ensure that you configure adequately for your space requirements and plan ahead for future database growth.

Before you begin

Tip: Complete the [“Checklist for server database disks” on page 14](#) to verify that your server database hardware is configured for good performance.

Configuring the server database directories for optimal performance

The *database containers* are the directories that Db2 uses to store the server database. Overall database size and the amount of data that is handled on a daily basis are the most important factors for how you configure the database directories.

Plan for growth with more directories at the start

When you are setting up your database space, ensure that you configure for future growth by including extra directories at the outset. If you are starting out with a moderate amount of data but plan to manage more, make sure that you plan for the larger amount of data to avoid reconfiguration work later.

If you must add directories after your initial configuration, make sure that you create directories that are equally sized. You can use the **EXTEND DBSPACE** command to add new directories for the database to use.

Plan ahead when you want to add space to the database. After new directories are added, data is redistributed and space is reclaimed for the system. This process can take considerable server resources. For more information, see the **EXTEND DBSPACE** command.

Use multiple database directories

How you spread the database directories across available disk storage has a strong effect on performance. Follow these guidelines for the database directories that you use:

- Use at least four directories initially for the database, spread across four LUNs or physical disks. For large IBM Storage Protect servers, use eight directories or more. You can use up to 128 directories for the server database.

For 2 TB servers for which data deduplication is planned, use eight directories or more for the server database. With data deduplication, the load on the database becomes greater because there are frequent queries to the database to determine what deduplicated extents are on the server.
- Make all directories that are used for the database the same size to ensure parallelism.
- Place each database directory in a different file system. This placement improves performance because Db2 stripes the database data across the various directories.
- Place the directories on disks that have the same capacity and I/O characteristics. For example, do not mix 10000 rpm and 15000 rpm drives for the database directories.
- For most disk systems, performance is best if one database directory is on one LUN, which has one logical volume.

The following image illustrates how to distribute database directories for data deduplication, by using eight disks.

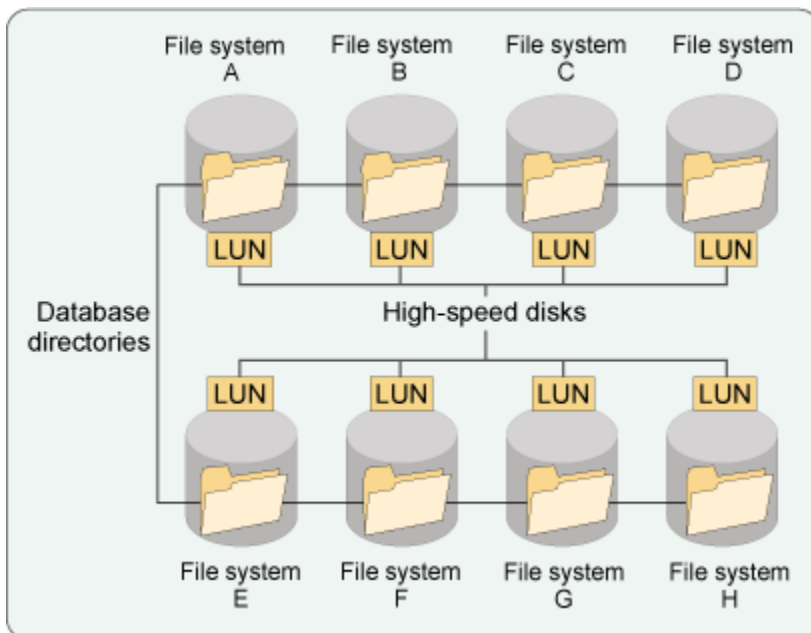


Figure 19. Distribution of IBM Storage Protect database directories

Reorganizing database tables and indexes

Server performance can degrade when database tables and table indexes are not reorganized.

About this task

Over time, database information becomes fragmented, causing unexpected database growth and increased space requirements for the active and archive logs. When you reorganize tables and table indexes, it reduces fragmentation and reclaims space.

If you are not running data deduplication on the IBM Storage Protect server, you might have to reorganize only the tables. However, you must monitor the database growth to avoid server degradation. If you are running data deduplication, reorganize both tables and indexes for best results.

You can reorganize tables and indexes online when the IBM Storage Protect server is running, or offline when the server is halted. Depending on your server workloads, you might have to disable both table and index reorganization to maintain server stability, and to complete daily server activities. If you experience unacceptable database growth or server performance degradation when reorganization is disabled, reorganize the tables offline.

IBM Storage Protect reorganizes tables and indexes by default. Allow server-initiated reorganization of both tables and indexes to run. If automatic reorganization is affecting server performance, you can manually schedule reorganization.

If you are manually reorganizing tables and indexes, make sure to determine whether a specific table must be reorganized, refer to the **DB2 REORGCHK** documentation. Running the **DB2 REORGCHK** command with the **CURRENT STATISTICS** parameter rather than the **UPDATE STATISTICS** parameter (the default) is recommended because of the potential for impacting server performance. Make sure that you are logged in with the DB2 instance user ID before you run the **REORGCHK** command. The tables and indexes should be manually reorganized with the help of IBM software support.

If IBM Storage Protect container storage pools are being used predominantly in the environment, then set the server option **REORGDURATION** to a value of 0 so that automatic reorganization of tables and indexes related to data deduplication can be performed without limitation. This can be checked with **QUERY OPTION REORGDURATION**.

Improving the speed of database backups

You can use multiple, concurrent data streams for database backups, which reduce the time that is required for a backup or restore operation to complete.

About this task

The amount of time that you can save depends on the size of the database. In general, the best performance benefits occur when you update your configuration from one to two streams and when the database is more than 100 GB.

A potential disadvantage when you use multistream processing is that more volumes and drives are required than for single-stream processing. For example, if the backup of an 850 GB database requires a single Linear Tape-Open (LTO) volume, switching to four data streams requires four drives. Volumes might be partially filled, especially if you use high-capacity volumes and device compression. If a backup of the database uses only 30% of a tape volume after compression, then four-stream processing results in even larger amounts of wasted space.

In some cases, particularly on AIX, problems with the TCP loopback interface can cause database backups to be unacceptably slow. Try configuring shared memory for database backups to see whether the speed of the database backup improves.

Guidelines for reducing database size by pruning the SDRS table

Over time, data extents accumulate in the *SDRS* table in the server database. You can delete expired and unreferenced data entries from the SDRS table and then reclaim unused space to reduce database size and optimize server storage space. You can also delete data entries with their associated *SHA1* digests.

When is it appropriate to delete data entries from the SDRS table to reclaim unused space?

Note: Because the **PROTECT STGPOOL** command is the only operation that uses the SDRS table, the following guidelines apply only if you use the **PROTECT STGPOOL** command. Also, if you plan to transition from using the **PROTECT STGPOOL** command to using replication storage rules, you can follow these instructions to delete data entries from the SDRS table.

- After you transition from using the **PROTECT STGPOOL** command to replication storage rules, you can gradually delete data entries from the SDRS table. The SDRS table is used during **PROTECT STGPOOL** and **REPAIR STGPOOL** operations, but it is not used if the data is replicated by using replication storage rules. After you successfully replicate all data by using replication storage rules, you can gradually delete data entries from the SDRS table.
- To review the volume of data in the SDRS table and decide whether a cleanup is warranted, you can issue the following IBM Db2 command:

```
db2 "select count from tsmdb1.sd_replicated_chunks for read only with ur"
```

Option 1: Schedule daily cleanup of the SDRS table to gradually delete deduplicated and non-deduplicated data entries

You can gradually delete expired and unreferenced data entries from the SDRS table by scheduling the **PROTECT STGPOOL** command to run with the **PURGEDATA=DELETED** parameter value. You can either define a new schedule for the **PROTECT STGPOOL** command with the **PURGEDATA=DELETED** parameter value or update an existing **PROTECT STGPOOL** command schedule with the **PURGEDATA=DELETED** parameter value. For more information about the **PROTECT STGPOOL** command, see *PROTECT STGPOOL (Protect data that belongs to a storage pool)* in IBM Documentation. You can schedule the **PROTECT STGPOOL** command to automatically run once daily. For more information about the **DEFINE SCHEDULE** command, see *DEFINE SCHEDULE (Define a schedule for an administrative command)* in IBM Documentation.

Option 2: Delete deduplicated and non-deduplicated data entries from the SDRC table all at once

To delete all the deduplicate and non-deduplicated data entries from the SDRC table in a single session, complete the following steps:

1. Run an **AUDIT CONTAINER** operation for the container storage pools on the source replication server. By running the audit container operation, any inconsistencies between database information and directory-container storage pools in the server will be identified and tagged. Follow the steps to run an **AUDIT CONTAINER** operation on the server:

- a. Issue the following Db2 command on the source replication server to generate a Db2 macro:

```
db2 -x "SELECT 'AUDIT CONTAINER ' || cntname || ' a=scanall W=Y' from  
tsmdb1.sd_containers where type=2"
```

- b. After the Db2 macro is generated, run the macro on the source replication server to scan for any inconsistencies between database information and directory-container storage pools.
2. Run an **AUDIT CONTAINER** operation for the container storage pools on the target replication server.

Tip: By running an audit operation, on the source and target replication servers, each non-deduplicated extent is populated with an SHA1 digest, which is necessary for repairing any damaged extents without the **PROTECT STGPOOL** command (without using SDRC table entries).

3. Issue the following Db2 command on the source and target replication servers to verify whether the SHA1 digest is calculated for all non-deduplicated data extents:

```
db2 "select count from tsmdb1.sd_non_dedup_locations where chunk_digest_type=0 for read only  
with ur"
```

Verify that the command output produces a value of "0". This value indicates that there are no non-deduplicated data extents without an SHA1 digest.



Attention: If the command output does not display "0", do not proceed to the next step. Instead, repeat [step 1](#) and [step 2](#) to ensure that all non-deduplicated extents are populated with an SHA1 digest.

4. Replicate all data from the source replication server to the target replication server by using replication storage rules. Ensure that you replicate all data from the source container storage pool to the target container storage pool.
5. Issue the **PROTECT STGPOOL** command with the **PURGEDATA=ALL** parameter value on the source replication server. When the command is processed, all data entries are deleted from the SDRC table.
6. Verify whether data entries remain in the SDRC table by issuing the following Db2 command:

```
db2 "select count from tsmdb1.sd_replicated_chunks for read only with ur"
```

If the output includes a value of "0", no data entries remain in the SDRC table. If the command output does not display "0", repeat [step 5](#) to delete remaining data entries from the SDRC table.

Configuring and tuning the recovery log

The management of the recovery log is important to normal operations for the server.

Before you begin

Tip: Complete the [checklist for server recovery log disks](#) so that the disk systems that are being used for the logs are optimally configured. Then, review the following information about how to configure the recovery log.

Procedure

- For the active log and active log mirror, follow these guidelines:

- Ensure that the directory that holds the active log is as large as, or larger than, the size of the active log. A directory that is larger than the active log can accommodate failovers if they occur.

Creating an active log mirror provides higher reliability, but comes at a cost. When the log is mirrored, it can affect performance because of the doubled I/O activity that is required to maintain the mirror. The additional space that the log mirror requires is another factor to consider. If you are using separate physical disks for the active log, use the **MIRRORLOGDIR** parameter in the **DSMSERV FORMAT** command to find the mirror log directory. After you install the server, change the mirror log directory location by changing the value of the **MIRRORLOGDIR** option in the server options file and restarting the server.

- Use disk system read-ahead processing to archive active logs more quickly.
- If there are active log performance issues, you can set the LOGBUFSZ server option in the dsmserve.opt file to 256 and restart the server.
- For the archive log and failover archive log, follow these practices:
 - You can create an archive failover log to store archive log files when the archive log directory is full. For best results, maintain an archive failover log.
 - Do not share archive log space with other applications, including other IBM Storage Protect servers. Other applications can write to the archive log directory and use the space that is required by the archive log. Ensure that each server has a separate storage location that is owned and managed by that specific server.
 - Set aside space in the failover archive log for possible emergency use.
 - Compress the archive log to reduce the need to run full database backups to clear the archive log. By compressing archived logs, you can reduce the amount of disk space that is required for the database environment. In high performance environments, such as large blueprint configurations, failure to compress archive logs can cause bottlenecks and performance degradation.

Related reference

[Optimization of disk system read-ahead processes](#)

Most advanced disk systems can automatically optimize the performance of read operations when the disk system can detect sequential reads. When the disk system detects sequential reads, it can have the data for the next read in cache, or at least have the read in progress.

Recovery log sizing

Ensure that the size of the active and archive logs meets the requirements of the server workload.

The active log is limited to a maximum size of 512 GB. The archive log is limited in size only by the size of the file system where it is located. It is better to create the active and archive log directories too large as opposed to too small. Consider the minimum values in the following table. Using these values or higher values can prevent log space problems for a server.

Storage pool deduplication enabled?	Active log directory: Minimum size	Archive log directory: Minimum size
No	16 GB	48 GB
Yes	Use the maximum size of 128 GB	128 GB

For more information about sizing the recovery log and examples that are based on different server loads, see [Upgrading the server](#).

Related reference

[Checklist for server recovery log disks](#)

The recovery log for the server consists of the active log, the archive log, and optional logs for mirroring and failover. Use the checklist to verify that disk systems that are being used for the logs have the characteristics and configuration that are key to good performance.

Tuning and configuring storage pools and volumes

Logical storage pools and storage volumes are the principal components in the IBM Storage Protect model of data storage. By manipulating the properties of these objects, you can optimize the use of storage devices.

Compressing data to save storage space

You can use server-side data compression to increase the amount of available space in a storage pool.

About this task

Review the following table to compare client-side and server-side compression:

Type of compression	Advantages	Disadvantages
Client-side	<ul style="list-style-type: none">• Reduced load on the network• Reduces data that is stored in the storage pool	<ul style="list-style-type: none">• Higher processor usage by the client• Potentially longer elapsed time for client operations such as backup
Server-side	<ul style="list-style-type: none">• Uses inline compression to compress data as it is written to a container storage pool• Reduces the amount of space that is required to store the data• Does not affect on client operations such as backup	<ul style="list-style-type: none">• No load reduction on the network• Higher processor usage by the server

Procedure

Inline compression is enabled by default. To disable compression, issue the **UPDATE STGPPOOL** command and specify the **COMPRESSION=NO** parameter.

Related concepts

[Reduce client data flow with compression](#)

The backup-archive client can compress data before it sends the data to the server. Enabling compression on the client reduces the amount of data that is sent over the network and the space that is needed to store it on the server and storage pools. Two client options determine when and if the client compresses data: **compression** and **compressalways**.

Optimizing data organization for restore and retrieval operations

How you organize data affects how efficiently and quickly IBM Storage Protect can access it and perform retrieval operations.

Grouping data by using collocation in server storage pools

Use collocation to improve IBM Storage Protect performance and maintain optimal data organization.

Before you begin

Tip: The following information does not apply to container storage pools.

When you use collocation, the performance of restore operations for large amounts of data can be significantly improved because fewer mounts are required to locate the necessary files. Generation of backup sets and export operations are also faster. In addition, collocation decreases the chance for media contention with other clients. While performance is improved by using collocation, enabling it increases both the amount of server time that is needed to collocate files for storing and the number of volumes that are required for data storage.

You can enable collocation by node, group, or file space. Collocation by group is the default. Each option provides different benefits and considerations for performance.

Table 17. Collocation trade-offs			
Type	Volume usage	Volume mounts	Restore time
No collocation	Low volume usage	Few number of mounts for migration and reclamation	Longest restore time
Collocated by node	High volume usage	High number of mounts for migration and reclamation	Good restore time, but not optimized for multi-session restore
Collocated by group	Low volume usage	Few mounts for migration and reclamation	Good restore time
Collocated by file space	High volume usage	High number of mounts for migration and reclamation	Good restore time, but not optimized for multi-session restore

About this task

Consider the following information when you are determining which type of collocation you want to use:

- Collocation by group provides the best balance of restore performance versus tape volume efficiency and it is the best practice choice for most situations. Collocation by group results in a reduction in unused tape capacity, which allows more collocated data on individual tapes. If collocation is needed to improve restore performance, use collocation by group. Manage the number of nodes in the groups so that backup data for the entire group is spread over a manageable number of volumes.
- For primary storage pools on tape, use collocation by group:
 - To get the full benefit of collocation by group, you must define the collocation groups and their nodes.
 - Nodes that are not grouped are collocated by node.
- For nodes with two or more large file spaces that might get close to filling a tape volume, use collocation by file space.
- Use an active data pool to collocate active data.
- Group nodes that have a low chance of being restored at the same time to avoid volume contention.
- Group nodes that are backed up to disk at the same time.

To enable collocation, use the **COLLOCATE** parameter on the **DEFINE STGPOOL** command when you are defining a primary sequential-access, copy, or active-data storage pool. You can use the **UPDATE STGPOOL** command to enable collocation for an existing storage pool.

Maintaining active-data storage pools

Setting up active-data storage pools can be useful for fast client restore operations. By keeping only active data in a storage pool, you can reduce the number of onsite or offsite storage volumes that you use, or reduce bandwidth when you copy or restore files that are vaulted electronically in a remote location.

Before you begin

One of the main benefits of using active-data storage pools is that it requires less disk space because the active-data pool contains only active file versions. Inactive data can be moved to tape.

The benefits of an active-data pool depend on the specific device type that is associated with the pool. For example, active-data pools that are associated with a FILE device class are ideal for fast client restore operations for the following reasons:

- FILE volumes do not have to be physically mounted
- The server does not position past inactive files that do not have to be restored

In addition, client sessions that are restoring from FILE volumes in an active-data pool can access the volumes concurrently, which also improves restore performance.

Active-data pools that use removable media, such as tape or optical, offer similar benefits. Although tapes must be mounted, the server does not have to position past inactive files. However, the primary benefit of using removable media in active-data pools is the reduction of the number of volumes that are used for onsite and offsite storage. If you vault data electronically to a remote location, you can save bandwidth by copying and restoring only active data. Use an active-data pool that is associated with a SERVER device class to copy and restore data when you want to save bandwidth.

For optimal efficiency during point-in-time restores and to avoid switching between active-data pools and primary or copy storage pools, the server retrieves both active and inactive versions from the same storage pool and volumes.

Improving file restore performance with caching in disk storage pools

You can use caching for IBM Storage Protect random-access (DISK) storage pools to improve restore performance of files.

Restriction: This procedure is applicable only, if you use IBM Storage Protect random-access (DISK) storage pools to improve restore performance of the files. The use of IBM Storage Protect random-access (DISK) storage pools is not recommended for modern servers.

When you enable caching for a storage pool, the migration process leaves a cached copy of a file in the storage pool after migrating files to the next storage pool. If your disk storage pool is large enough to hold the backed-up data from one day, then caching is a good method. When you restore files that were recently backed up to a disk storage pool, the benefit of caching is evident. The time that is required to retrieve files is reduced.

Trade-offs for the use of caching include the following performance impacts:

- Disk caching can affect backup throughput. When cached files must be deleted to create space for file backups, the deletion operations require database updates.
- If you use caching with larger storage pools, they might become fragmented and response time can increase.

Use the **DEFINE STGPOOL** or **UPDATE STGPOOL** command with the **CACHE=YES** parameter to enable caching. If you enable disk storage pool caching, and then suspect that it is affecting performance, try turning off caching. To turn off caching, issue the **UPDATE STGPOOL** command with **CACHE=NO** to see whether there is a positive effect.

Improving data restore performance with caching in cloud-container storage pools

You might be able to improve data restore performance for cloud-container storage pools by enabling the **CLOUDREADCACHE** parameter and configuring the related CLOUDREADCACHEMAXUSAGE and CLOUDREADCACHERETENTIONTIME server options.

Before you begin

For an introduction to the use of caches for cloud-container storage pools, see [Cloud-container storage pool data cache](#).

For guidance about specifying settings, review the preferred cloud cache settings for light restore workloads, moderate restore workloads, and disaster recovery events in [Use case configurations for caching cloud-container storage pool data](#).

Restriction: The **CLOUDREADCACHE** parameter is available only for non-Swift cloud types.

To use a cloud cache, you must first define the local disk storage pool directories where you want to store read cache data. Use the **DEFINE STGPOOLDIRECTORY** command to define local disk storage pool directories for your cloud-container storage pool.

Procedure

1. When you define or update a cloud-container storage pool by using the **DEFINE STGPOOL** or **UPDATE STGPOOL** command, enable the **CLOUDREADCACHE** parameter by specifying one of the following values:

Value	Description
ON	The system downloads and stages read cache data in the cloud cache for cloud-container storage pool objects with significant read activity. As client restore activity is processed, the IBM Storage Protect server analyzes container data read requests for the cloud-container storage pool to determine whether to download data from object storage and stage it to the cloud cache. For data that the system decides to stage, client restore operations retrieve data from the cloud cache. For data that was not staged, client restore operations retrieve data from object storage.
ONPREFERINGEST	<p>The ONPREFERINGEST value enables similar functionality as the ON value. However, with the ONPREFERINGEST value, if the system detects a low-space condition when data is ingested, read cache data is removed from the cloud cache so that the space can instead be used for ingested data.</p> <p>A low-space condition can occur when data is ingested during a cloud-container restore operation on the same cloud cache. Ingest activity might include direct backups, archives, and node replication target activity.</p>

The default value is OFF.

2. Specify how much space will be used for read cache data by setting a value for the CLOUDREADCACHEMAXUSAGE option. You can specify a value in the range 0 - 100. The default value is 95.

For example, if you have a cloud cache with 1000 GB of usable space and you specify the value 90 (that is, 90%), the system can use up to 900 GB of space in the cloud cache for read cache data (1000 GB x 90% = 900 GB). In the following example, a value of 50 is specified to ensure that the system uses no more than 50% of the cloud cache for read cache data.

```
cloudreadcachemaxusage 50
```

Important: When a storage pool directory file system has 50 GB or less of usable space, the IBM Storage Protect server does not add more read cache data to the cloud cache. The server maintains 50

GB of space to use for ingested data. If none of the directory file systems in a storage pool directory have more than 50 GB of usable space, no further read cache data is added to the cloud cache for the entire storage pool.

3. Use the **CLOUDREADCACHERETENTIONTIME** server option to specify how many seconds read cache data is retained in the cloud cache after the last read operation ends. You can specify a value in the range 1 - 2147483647 (the maximum 32-bit integer). The default value is 600 (10 minutes). For example, if you keep the default value of 1200 and the system does not read from the cloud cache for 1200 seconds or more, the read cache data in the cloud cache is deleted.

```
cloudreadcacheretentiontime 1200
```

Consider the following guidance about specifying a small, medium, or large value for the **CLOUDREADCACHERETENTIONTIME** server option:

Value size	Use case
Small	Specify a small value, such as 120 (2 minutes) if one or more of the following conditions are met: <ul style="list-style-type: none"> • The workload that is being restored is highly unique (that is, it has a low deduplication ratio in the storage pool). • The cloud cache is smaller than the total amount of read cache data that is being restored.
Medium	Specify a medium value, such as 1800 (30 minutes) or 3600 (60 minutes), when all of the following conditions are met: <ul style="list-style-type: none"> • The workload that is being restored is not highly unique (that is, it has a high ratio of data deduplication in the storage pool). • The file system space is large enough to copy the total amount of read cache data that is being restored to the cloud cache while also handling data ingestion for the cloud-container storage pool.
Large	Specify a large value, such as 86400 (1 day) or 432000 (5 days), to either restore a client workload over a period of many hours or days or to restore the same client data repeatedly. <p>Tip: If the restore workload is considerably larger than the capacity of the cloud cache, consider reducing the CLOUDREADCACHERETENTIONTIME parameter value so that older read cache data is removed faster from the cloud cache. By removing old data, you provide space for new data. When the cloud cache runs out of space, the IBM Storage Protect server does not add more read cache data to the cloud cache. The cloud cache can run out of space if the CLOUDREADCACHEMAXUSAGE parameter value is reached for all storage pools or if 50 GB or less of usable space is in the storage pool.</p>

Related concepts

[Cloud-container storage pool data cache](#)

IBM Storage Protect 8.1.12 introduced the **CLOUDREADCACHE** parameter, which can be used while you are defining or updating a cloud-container storage pool to help improve performance when you restore a large amount of data from cloud-container storage pools. When this parameter is enabled, the system can download a copy of cloud-container object data that is being restored, and then stage a copy of that data locally in the cloud cache. As a result, the cloud cache can store both ingested data and read cache data.

Cloud-container storage pool data cache

IBM Storage Protect 8.1.12 introduced the **CLOUDREADCACHE** parameter, which can be used while you are defining or updating a cloud-container storage pool to help improve performance when you restore a large amount of data from cloud-container storage pools. When this parameter is enabled, the system can

download a copy of cloud-container object data that is being restored, and then stage a copy of that data locally in the cloud cache. As a result, the cloud cache can store both ingested data and read cache data.

With cloud-container storage pools, unique data extents are stored in a deduplicated, and optionally compressed and encrypted, format in cloud-container objects that are within object storage. Each client file that is backed up to the IBM Storage Protect server is composed of a sequence of deduplicated extents that are stored in these cloud-container objects within object storage. Depending on how client data deduplicates with other client data, individual client files might share data extents in these cloud containers with other files or data that is stored in the cloud-container storage pool.

When you enable the **CLOUDREADCACHE** parameter for a cloud-container storage pool, the IBM Storage Protect server analyzes container data read requests for the storage pool. Read requests are initiated by client restore and retrieve activity. If significant read activity occurs within a short time for a cloud-container object, the server downloads the requested data from object storage, and then stages it as a copy to the cloud cache. The cloud cache is in local disk storage pool directories. The system uses the read cache data that was downloaded to the cloud cache to meet further client read requests for that cloud-container object. By using a cloud cache, you can improve the performance of restore operations in some circumstances because reading data from local disk storage pool directories can be much faster than sending small read requests to object storage over the Ethernet network.

Enabling the **CLOUDREADCACHE** parameter might be useful in the following circumstances:

- The IBM Storage Protect server is connected to the object storage system on a network with high bandwidth. For example, the network bandwidth is 5 Gbps, 10 Gbps, or 40 Gbps.
- The IBM Storage Protect server is connected to the object storage system on a network with high latency. For example, the object storage system might be an off-premises public cloud with 10 milliseconds of round-trip latency.
- You must restore more than 10 GB of data.
- You have an extent workload that is predominantly composed of small files, is stored with a low ratio of data deduplication, or both. For example, you have one of the following types of workload:
 - IBM Storage Protect backup-archive client, file server type workloads
 - Oracle database workloads
- You would like to restore data multiple times to different clients or for different purposes. For example, you might want to recover a virtual machine or database multiple times when you are setting up test environments.

Enabling the **CLOUDREADCACHE** parameter might not be useful in the following circumstances:

- The IBM Storage Protect server is connected to the object storage system on a network with low bandwidth. For example, the network bandwidth is 1 Gbps or 100 Mbps bandwidth.
- You must restore only a small amount of data (less than 10 GB).
- You have a large deduplicated extent workload and the IBM Storage Protect server is connected to the object storage system on a network with low latency. For example, the object storage system might be on the same local area network (LAN) as the cloud-container storage pool, and the round-trip latency might be less than 10 milliseconds.

Read cache data invalidation:

Read cache data in the cloud cache is not preserved when the IBM Storage Protect server is restarted. Instead, read cache data is invalidated on a server restart and is no longer used. The cloud cache is initialized, and it prepares storage pool directories.

Read cache data in the cloud cache is invalidated in all of the following situations:

- The **CLOUDREADCACHE** parameter value is changed from ON to OFF.
- The **CLOUDREADCACHE** parameter value is specified as either ON or ONPREFERINGEST and the IBM Storage Protect server is restarted.
- The **CLOUDREADCACHE** parameter value is changed from OFF to either ON or ONPREFERINGEST.

Related tasks

Improving data restore performance with caching in cloud-container storage pools

You might be able to improve data restore performance for cloud-container storage pools by enabling the **CLOUDREADCACHE** parameter and configuring the related CLOUDREADCACHEMAXUSAGE and CLOUDREADCACHERETENTIONTIME server options.

Use case configurations for caching cloud-container storage pool data

For guidance about configuring cloud-container storage pool caching for your deployment, review the use cases for light restore workloads, moderate restore workloads, and disaster recovery events.

Cloud cache settings for light restore workloads

The following table lists the preferred cloud cache settings to use for light restore workloads. With a light restore workload, data is ingested daily to a cloud-container storage pool with occasional restore activity that is limited in data volume and duration. The total quantity of restored data is much smaller than the total cloud cache capacity.

Table 18. Cloud cache configuration for light restore workloads		
Ethernet bandwidth to object storage	Other considerations	Preferred settings
High bandwidth (for example, 5 Gbps or 10 Gbps)	None	CLOUDREADCACHE = ON CLOUDREADCACHEMAXUSAGE = 95 CLOUDREADCACHERETENTIONTIME = 600
Low bandwidth (for example, 100 Mbps or 1 Gbps)	None	CLOUDREADCACHE = OFF

Cloud cache settings for moderate restore workloads

The following table lists the preferred cloud cache settings to use for moderate restore workloads. With a moderate restore workload, data is ingested daily to a cloud-container storage pool with restore activity that might use 50% or more of the cloud cache. The total quantity of restored data is 50% or greater than the total cloud cache capacity.

Table 19. Cloud cache configuration for moderate restore workloads		
Ethernet bandwidth to object storage	Other considerations	Preferred settings
High bandwidth (for example, 5 Gbps or 10 Gbps)	Allow the cloud cache to use as much space as required to store read cache data, but prioritize data ingest performance.	CLOUDREADCACHE = ONPREFERINGEST CLOUDREADCACHEMAXUSAGE = 95 CLOUDREADCACHERETENTIONTIME = 600
High bandwidth (for example, 5 Gbps or 10 Gbps)	Allow the cloud cache to use as much space as required to store read cache data and prioritize restore performance.	CLOUDREADCACHE = ON CLOUDREADCACHEMAXUSAGE = 95 CLOUDREADCACHERETENTIONTIME = 600

Table 19. Cloud cache configuration for moderate restore workloads (continued)

Ethernet bandwidth to object storage	Other considerations	Preferred settings
High bandwidth (for example, 5 Gbps or 10 Gbps)	Restrict the amount of cloud cache space that can be used to store read cache data. For example, specifying a value of 50 ensures that no more than 50% of the cloud cache is used to store read cache data.	CLOUDREADCACHE = ON CLOUDREADCACHEMAXUSAGE = 50 (or other appropriate percentage value) CLOUDREADCACHERETENTIONTIME = 600
Low bandwidth (for example, 100 Mbps or 1 Gbps)	None	CLOUDREADCACHE = OFF

Tip: If you are planning to recover a large amount of data that might not fit in a cloud cache, consider temporarily increasing the capacity of the cloud cache by defining extra storage pool directories for the cloud-container storage pool. For example, if you must restore 100 TB of client data from object storage, but your cloud cache is configured for only 10 TB, you can temporarily add storage so that you have enough space to hold the read cache data on disk.

Cloud cache settings for disaster recovery events

The following table lists the preferred cloud cache settings to use for disaster recovery events. During disaster recovery, a large volume of data must be restored from cloud-container storage pools in object storage. The total quantity of restored data is much larger than the current cloud cache capacity.

Table 20. Cloud cache configuration for disaster recovery events

Bandwidth and requirements for repeated restore operations	Other considerations	Preferred cloud cache settings
High bandwidth (for example, 5 Gbps or 10 Gbps), a requirement to stage data to disk for repeated restore operations, or both	The client restore workload is larger than the capacity of the cloud cache.	CLOUDREADCACHE = ON CLOUDREADCACHEMAXUSAGE = 95 CLOUDREADCACHERETENTIONTIME = Specify a value that is less than the default. For example, specify 60 seconds or 120 seconds.
High bandwidth (for example, 5 Gbps or 10 Gbps), a requirement to stage data to disk for repeated restore operations, or both	Cloud cache capacity was increased to accommodate the restore workload and you want to keep the restored data staged to disk for repeated restore operations.	CLOUDREADCACHE = ON CLOUDREADCACHEMAXUSAGE = 95 CLOUDREADCACHERETENTIONTIME = Specify a value that is greater than the default. Specify how many seconds you want the data to stay in the cache. For example, specify multiple hours or days.

Table 20. Cloud cache configuration for disaster recovery events (continued)

Bandwidth and requirements for repeated restore operations	Other considerations	Preferred cloud cache settings
Low bandwidth (for example, 100 Mbps or 1 Gbps) and no requirement to stage data to disk for repeated restore operations	None	CLOUDREADCACHE = OFF

Using file system cache for storage pools

By default, the IBM Storage Protect server reads and writes storage pool data with non-buffered I/O, which bypasses the file system cache. Using file system cache can be beneficial in certain situations but it is not a best practice.

About this task

If you switch to file system cache, change the default with care. The use of file system cache can decrease overall server throughput and increase server processor usage. Use file system cache for storage pool I/O only if it results in a significant performance improvement.

You might see benefits from using file system cache in the following environments:

- A disk storage system that you are using for storage pools has a relatively small cache memory size.
- The disk storage system does not provide read-ahead capability. You must improve storage pool migration throughput from the disk volumes by using the operating system read-ahead capability.
- Data that is stored on the IBM Storage Protect server is expected to be immediately retrieved through some other process.

Restriction: Enabling the server to use file system cache is not a preferred practice. If you contact IBM Software Support for a performance issue, you must disclose that you are using this option. You might be asked to disable this option before IBM can resolve the issue.

Procedure

To use file system cache for storage pool I/O, complete the following steps:

1. Specify one of the following options in the `dsmserv.opt` file:

- Container storage pools: `DIOENABLED NO`
- Other types of storage pools: `DIRECTIO NO`

Tips: For better performance, make the following adjustments:

- Disable direct I/O for IBM Storage Scale
- Enable IBM Elastic Storage Server

2. Restart the server.

3. Observe operations for the effects of the change and determine whether to keep the `DIRECTIO NO` entry in the options file or remove it.

Tip: To display the current value for the **DIRECTIO** option, specify the option name when you issue the **QUERY OPTION** command.

File system fragmentation for storage pools on disk

Fragmentation problems can occur when you are writing to scratch volumes in multiple storage pools simultaneously, in one file system. Because unfragmented volumes generally perform better than fragmented volumes, use the **DEFINE VOLUME** command to preallocate volumes for sequential disk storage pools (**DEVTYPE** is **FILE**).

In some operating system environments, preallocating several **FILE**-device class volumes or random-disk pool volumes in parallel in the same file system can also result in fragmentation. If you are running a Windows system or using JFS2 file systems on AIX or ext4 file systems on Linux, fragmentation is not a problem with preallocated volumes. IBM Storage Protect uses operating system tools to allocate files without fragmentation, even when created in parallel.

If you encounter slow performance, you can use operating system tools to help determine how fragmented a storage pool volume file might be. The tools report the number of fragments per file. To achieve high read performance from a volume, the number of fragments per megabyte should not exceed 0.25. For example, if a 50 GB volume has more than 12,800 fragments, then read performance for the volume might be poor because of fragmentation.

Use the following tools to determine fragmentation level:

- On AIX, issue the **fileplace** command.
- On Linux, issue the **filefrag** command.
- On Windows, use the **contig** utility. You can find details and usage information about the **contig** utility on the [Microsoft TechNet site](#).

To avoid fragmentation, use preallocated volumes and use more file systems per storage pool. By using a greater number of file systems, you can usually avoid fragmentation because writing to different volumes is more likely to occur in different file systems.

Optimal number and size of volumes for storage pools that use disk

When you calculate the size and number of volumes for storage pools that are defined with a **FILE** or **DISK** device class, consider the type of storage, how much storage is available, and other variables.

The optimal size for storage pool volumes depends on these elements:

- The type of storage pool, whether random access (**DISK**) or sequential access (**FILE**)
- How much total storage is available
- How many concurrent writes are expected (client sessions and server processes)
- How many storage pool volumes can be managed
- The storage pool collocation settings
- Whether you use preallocated or scratch volumes
- The average object size that is stored in the storage pool (having the volume size larger than the average object)

If you do not have the information to estimate a size for **FILE** device class volumes, start with volumes that are 50 GB.

The **FILE** device class allows only one session or process to write to a volume at a time. Therefore, you must have at least as many volumes as you have sessions and processes that you expect to write to the pool at the same time.

If you are using collocation, a minimum number of volumes is required. For example, if you are running collocation by node, use at least one volume per node. When you use collocation with preallocated volumes, consider the potential unused storage within those volumes. Collocation with preallocated volumes requires more space. If you use collocation by group, you can reduce the amount of unused storage by designing groups to minimize this waste.

You can use preallocated volumes to avoid file system fragmentation, but you must estimate how much space is required for the storage pool and allocate enough volumes to handle that load. This estimate

assumes that the amount of data you are backing up does not fluctuate. For example, processing that is done at the end of the month might use more storage pool space than regular daily processing. If you use scratch volumes, fluctuation is not a factor because IBM Storage Protect allocates what is needed for each day's processing. If you are using preallocated volumes at a recovery site, some additional time is required for the recovery process because volumes must be preallocated and formatted.

Example: Choosing volume size for a sequential file storage pool

In this example, you must determine a volume size for a sequential-access FILE storage pool with 100 TB of available space. Collocation is not used, therefore you do not have to consider the empty space that is required in preallocated volumes. The maximum number of simultaneous backup sessions that are used during the backup window is 250. To prepare for growth of the system and the occurrences where 250 sessions are exceeded, double that amount to 500. Using these values, the maximum volume size that is needed is 204 GB.

Related tasks

[Grouping data by using collocation in server storage pools](#)

Use collocation to improve IBM Storage Protect performance and maintain optimal data organization.

Configuring and tuning the server

How you configure and tune different IBM Storage Protect server settings and operations can affect performance of the system.

Memory requirements

You must have sufficient memory available for server operations such as replication and data deduplication. Memory requirements are based on the projected maximum database size. Extra memory improves database query performance by increasing the amount of memory that is available for the database buffer pools.

Use the following table to determine the minimum amount of memory you must allocate on the server to complete operations, based on the database size. The size that is required depends on the functions that you use.

<i>Table 21. Minimum size requirements based on database size</i>			
Database size	Standard server operations (no data deduplication or replication required)	Either data deduplication or replication required	Both data deduplication and replication required
500 GB	16 GB	24 GB	32 GB
1000 GB	24 GB	40 GB	56 GB
1500 GB	32 GB	56 GB	80 GB
2000 GB	40 GB	72 GB	104 GB
2500 GB	48 GB	88 GB	128 GB
3000 GB	56 GB	104 GB	152 GB
3500 GB	64 GB	120 GB	176 GB
4000 GB	72 GB	136 GB	200 GB

Sizing a cloud cache to optimize backup operations

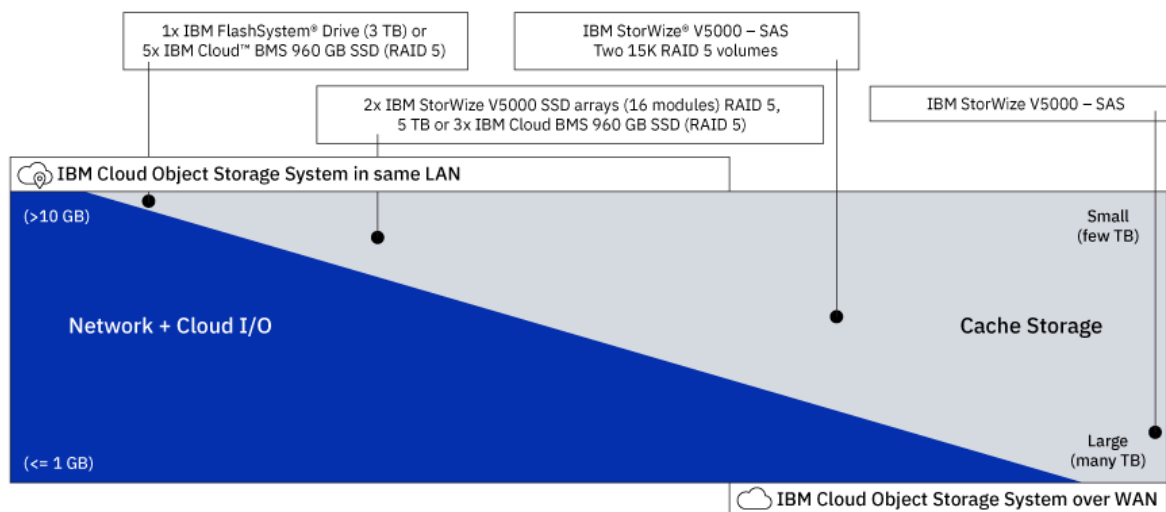
Before you back up data to a cloud-container storage pool, size the cloud cache. By correctly sizing the cache, you can help to improve the throughput of backup operations and reduce the risk of backup failures.

Before you begin

Tip: If you plan to use cloud tiering to move data from one or more directory-container storage pools to a cloud-container storage pool, you do not have to size the cloud cache. Size the cloud cache only for cloud-container storage pools that are the target of backup operations.

Review the following figure to learn about disk technologies that you can use to size the cloud cache. The optimum size of the cloud cache depends on the speed of the network connection and the throughput capability of the object storage system. The figure shows a range of 1 GB to 10 GB Ethernet network connection. For systems with slower network connections and less throughput capability to cloud object storage, you can use slower and larger disk technologies for the cloud cache. To optimize throughput, use fast disk technology that can write and read data at your network speed simultaneously with a 256 KiB input/output (I/O) size. The following example storage systems are used:

- Fast network connection (more than 10 GB): 1 IBM FlashSystem® Drive (3 TB) or 5 IBM Cloud® BMS 960 GB Solid State Disk (SSD) arrays, for example, RAID 5 arrays
- Fast network connection (8 - 10 GB): 2 IBM StorWize V5000 SSD arrays, for example, 5 TB RAID 5 arrays or 3 IBM Cloud BMS 960 GB SSDs, for example, RAID 5 arrays
- Medium network connection (4 - 8 GB): 1 IBM StorWize V5000 serial-attached SCSI (SAS) and 2 15,000 RAID 5 volumes
- Slow network connection (1 - 4 GB): 1 IBM StorWize V5000 SAS



About this task

For cloud-container storage pools, the data is held only temporarily on disk. After the data is transferred to the cloud, the data is deleted from the cloud cache.

Client data is ingested to the cloud cache at the same time that other ingested data is transferred to the cloud. To prevent the failure of backup operations, size the cloud cache in terms of I/O operations per second (IOPS) and capacity.

Procedure

To size the cloud cache, use the following guidelines:

- Size the cloud cache twice as large as the largest number of data backups that are ingested at the same time. If twice as large is not possible, size the cloud cache at least as large as the largest number of data backups that occur at once.
- To optimize data ingestion, choose fast disk technology, such as SSD, for the cloud cache that can both write and read data at your network speed simultaneously. If you size the cloud cache too large and on slow disk technology, only one copy of the data backup exists. If the cloud cache fails, you lose the data backup. Slow disk technology can cause the cloud cache to be a bottleneck to the system's ingest capability.
- Use RAID 5, RAID 6, or other disk protection for the cloud cache to avoid data loss.
- Use benchmarking tools to help you to size the cloud cache. The cloud benchmarking tools and example benchmarking tests are available on the wiki page at [Cloud Blueprints](#). Benchmark the IOPS capability of the cloud cache and the throughput capability of the network and object storage. For cloud cache benchmarking, the cloud benchmarking tools use an I/O size of 256 KB. The 256 KB size supports simultaneous read-and-write operations with many threads.

What to do next

1. Configure a storage pool directory for each cloud-container storage pool. The storage pool directory will define the size of the disk cache. For instructions, see [Optimizing performance for cloud object storage](#).
2. Use a single file system, with striping across multiple volumes. Use a stripe size of 16 KB for the cloud cache and ensure that striping is enabled across all physical volumes. Use a single storage pool directory and file system for the cloud cache to avoid disk hot spots with the overlapped I/O of client data ingestion and transfer to cloud.

Sizing a cold-data-cache storage pool

Before you back up object client data to a cold-data-cache storage pool, size the cold-data-cache storage pool. The *cold-data-cache storage pool* acts as the initial disk-based storage location for object client data that is copied to IBM Storage Protect for archiving to tape storage. By correctly sizing the cold-data-cache storage pool, you can help to improve the throughput of archive operations, reduce the risk of archive failures, and ensure that enough storage capacity is available for data ingestion and restore operations.

Before you begin

An object client must be an IBM Storage Protect Plus server. Review the following information:

- For instructions about setting up IBM Storage Protect Plus as an object client to the IBM Storage Protect server, see [Copying data from IBM Storage Protect Plus to IBM Storage Protect](#).
- For instructions about copying data from IBM Storage Protect Plus to tape storage, see [Configuring operations for copying data to tape](#).
- For instructions about restoring archived data from tape to IBM Storage Protect Plus, see [Restoring data from tape to IBM Storage Protect Plus](#).

Tip: In previous releases, the process of copying data from IBM Storage Protect Plus to secondary backup storage was known as *offloading data*. Beginning with IBM Storage Protect 8.1.9, the process is known as *copying data*.

About this task

Data that is copied from IBM Storage Protect Plus is stored temporarily on disk in file volumes that are specified for the cold-data-cache storage pool. Then, data is migrated to the next storage pool that is defined on the **DEFINE STGPOOL** command for the cold-data-cache storage pool. After the data is migrated to a tape storage pool, the data is deleted from the cold-data-cache storage pool.

Tip: The tape storage pool is defined as a next storage pool by specifying the **NEXTSTGPOOL** parameter on the **DEFINE STGPOOL** command for the cold-data-cache storage pool.

Similarly, during a restore operation, the object data is restored temporarily to the cold-data-cache pool before the data can be read by an object client. When IBM Storage Protect Plus issues a request to restore the object data from tape storage, the IBM Storage Protect server copies the data from the tape storage pool to the cold-data-cache storage pool temporarily. Then, IBM Storage Protect Plus can restore the data. Requested data is stored on the cold-data-cache storage pool for a specified number of days before deletion.

Consider the following guidelines for running migration processes on cold-data-cache storage pools:

- Data becomes eligible for migration from the cold-data-cache storage pool as file volumes become full or are closed.
- Processes to ingest new data and migrate eligible data to next storage pools can occur in parallel. As the data is migrated, it is deleted from the cold-data-cache storage pool. You can configure the number of parallel processes by specifying the **MIGPROCESS** parameter on the **DEFINE STGPOOL** command for the cold-data-cache storage pool. The number of parallel processes might be limited by the number of drives that are available for migration on the tape storage pool.
- Migration performance can be limited by the throughput capability of the tape storage pool drives. For example, throughput rates of 300-400 MBs per second are common with LTO-8 tape drives and volumes during migration.

To accommodate both recently copied data and data copies that are staged for restore operations back to the object client, adequate space must be provisioned for the cold-data-cache storage pool. The IBM Storage Protect server reads and writes to the cold-data-cache storage pool predominantly in 256 KB blocks.

Procedure

To size and tune the cold-data-cache storage pool, follow the guidelines:

- Use the **tsmdiskperf.pl** Perl script as a benchmarking tool to size the cold-data-cache storage pool.
 - Benchmark the directory paths to be used for the cold-data-cache storage pool with an overlapped, sequential read-and-write workload with a 256 KB block size.
 - To run the script, issue the following command:

```
perl tsmdiskperf.pl workload=stgpool fslist=directory_list
```

where *directory_list* is a comma-separated list of directory paths.

- Ensure that the data ingestion rate that is obtainable for these directory locations satisfies the speed requirements for data-ingestion operations in your environment.

For benchmarking tools and sample benchmarking tests, see the [IBM Storage Protect Blueprints](#). The benchmarking tool **tsmdiskperf.pl** is available in the *Blueprint configuration scripts* package.

- Ensure that the cold-data-cache storage pool is large enough to hold the daily volume of data from a copy operation. In this way, if an issue with the next tape storage pool prevents or slows migration, sufficient space is available to contain the daily workload and avoid failures.
- Where possible, optimize disk system performance by configuring the disk system for random read/write operations rather than sequential read/write operations.
- Use RAID 5, RAID 6, or other disk protection for the cold-data-cache directory file system disks to avoid data loss.
- On the **DEFINE STGPOOL** or **UPDATE STGPOOL** commands for the cold-data-cache storage pool, set the **MIGPROCESS** parameter value to match the number of tape drives from the next tape storage pool that can be used for migration activities. To optimize migration performance and ensure that the cold-data-cache storage pool releases space as quickly as possible, set the **MIGPROCESS** parameter with as high a value as possible. You can enter a value in the range 1 - 999.

Tip: When you specify the **MIGPROCESS** parameter, consider other uses of the tape storage pool that might compete for resources. For example, you might use the tape storage pool to back up the IBM Storage Protect database.

- For optimal throughput for the object client node that is running the backup and restore operations to the cold-data-cache storage pool, set the **MAXNUMMP** parameter on the **REGISTER NODE** or **UPDATE NODE** commands to a value of at least 100.

Tip: This parameter limits how many mount points a node can use on the server. The IBM Storage Protect object agent can distribute backup and restore data movement across as many as 100 sessions for a single client node.

- On the **DEFINE STGPOOL** or **UPDATE STGPOOL** commands for the tape storage pool, set the **COLLOCATE** parameter to match your requirements. By default, group level collocation is used for sequential-access storage pools. If no collocation groups exist on the server, collocation by node is used by default. Each migration process from the cold-data-cache storage pool attempts to use a drive on the next tape storage pool, if available. When collocation is used, the IBM Storage Protect server attempts to store group, node, or file space data together on as few tape volumes as possible.

Tip: During an operation to restore data from tape storage, the IBM Storage Protect server might attempt to use multiple tape volume mounts, depending on the number of tape volumes in use. By default, the IBM Storage Protect server attempts to use up to four processes to restore data from tape volumes. The number of volumes limits the number of processes.

- To release space and allow for the ingestion of recently copied data to preempt data restore operations, specify the **REMOVERESTORED COPYBEFORELIFETIMEEND=YES** setting on the **DEFINE STGPOOL** or **UPDATE STGPOOL** commands of the cold-data-cache storage pool. When this parameter is set to YES, IBM Storage Protect removes certain restored data copies (that are eligible for early deletion according to defined conditions) to create space for new data copy operations.
- By default, the **MAXSCRATCH** parameter on the **DEFINE STGPOOL** command is set to 5000 for a cold-data-cache storage pool. This parameter controls the maximum number of scratch file volumes that can be created in the storage pool during data ingestion and restore operations. By default, the device class that is created when you define the cold-data-cache storage pool has a volume size of 10 GB for an overall default capacity of 50,000 GB.

If a larger capacity is needed, use the **UPDATE STGPOOL** command to increase the **MAXSCRATCH** parameter value for the cold-data-cache storage pool. The maximum value for this parameter is 9999. If more capacity is needed, you can also increase the cold-data-cache storage pool's device class volume size by issuing the **UPDATE DEVCLASS** command.

Example architecture of data flows for copy and restore operations

The following image shows an example of a typical data flow to copy data from IBM Storage Protect Plus to the cold-data-cache storage pool on an IBM Storage Protect server so that the server can move the data to tape storage.

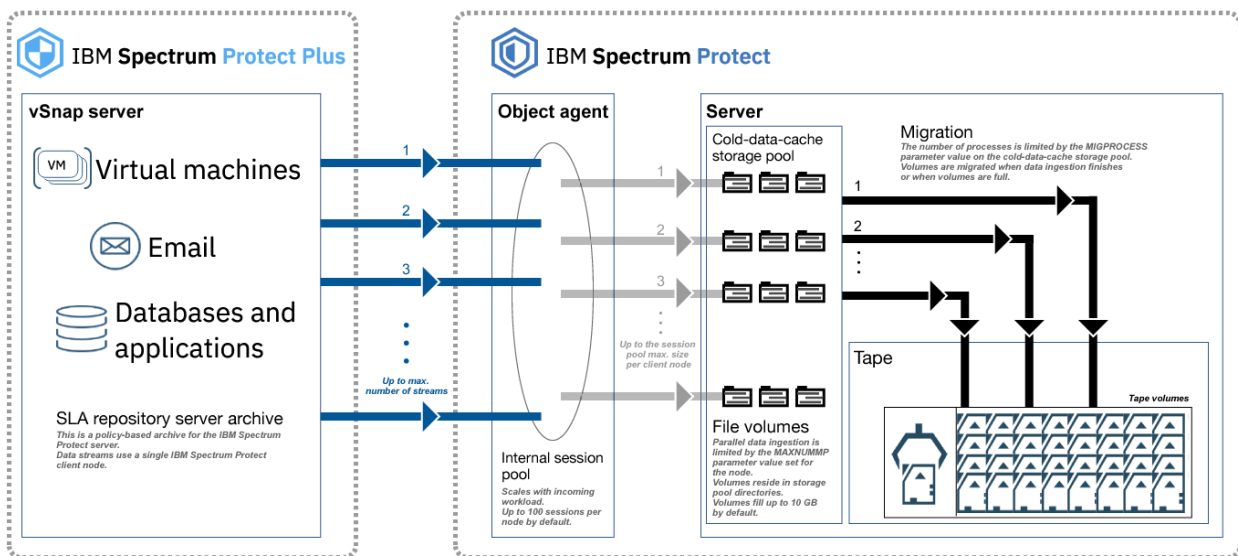


Figure 20. Data flow for copying data

Tip: For detailed instructions, see [Configuring operations for copying data to tape](#).

The following image shows an example of a typical data flow to restore data from tape storage to the IBM Storage Protect Plus object client by using cold-data-cache storage pools on the IBM Storage Protect server.

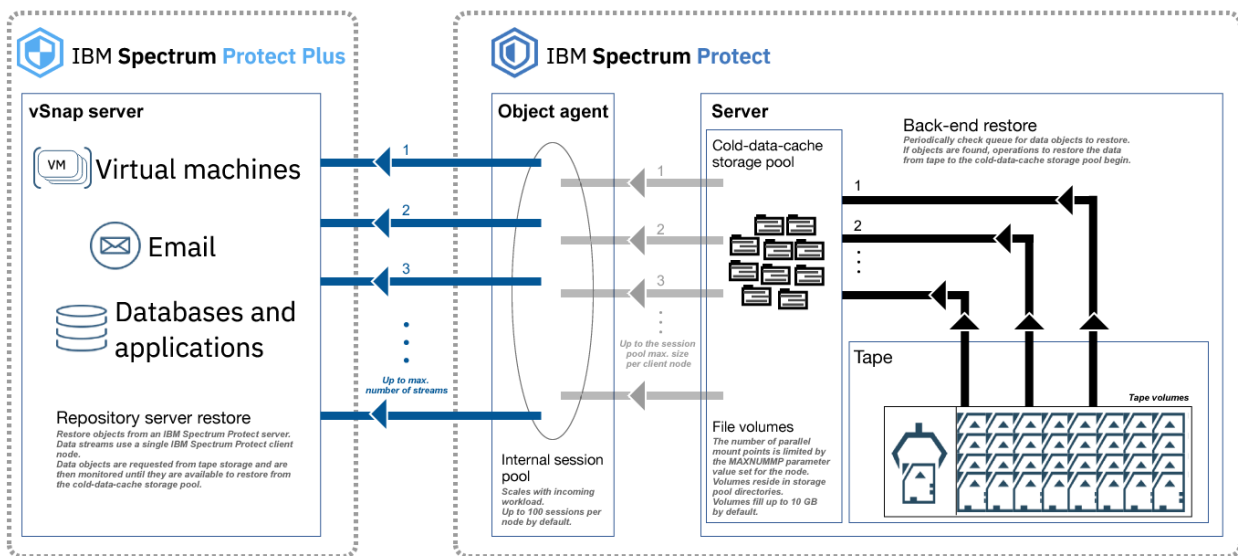


Figure 21. Data flow for restoring data

Tip: For detailed instructions, see [Restoring data from tape to IBM Storage Protect Plus](#).

What to do next

- Monitor used space within the cold-data-cache storage pool. If the storage pool frequently runs out of space, the performance of disk-read and tape-write operations might be insufficient to handle the target data ingestion workload.

Tuning the schedule for daily operations

Typically, backup operations must be completed every day for all clients. Certain server maintenance processes must also run every day. Ensuring that resources for these critical operations are available when needed requires planning and tuning.

About this task

During the client workload phase, server resources support client operations. Client workloads are predominantly client backup and archive activities. Typically, these operations are completed during the nightly schedule window. During the server workload phase, server resources are dedicated to managing the recently received data from the client workload and performing the following activities, which are necessary to manage, protect, and maintain the server:

- Protecting client data by backing up storage pools
- Allocating data in the storage hierarchy
- Keeping the database, storage hierarchy, and server operations working efficiently
- Preparing for the next schedule cycle

Carefully schedule client and server workloads to achieve the best performance for your environment. If the client and server operations overlap or are not allotted enough time and resources to be processed, the environment might be impacted in the following ways:

- Unprotected data, incomplete copies of data, or both
- Less processing power and memory available to support an operation
- Performance degradation
- Insufficient space for data storage
- Issues with data placement
- Failed operations

For optimum performance, split the tasks of backing up and archiving client data, and performing server data maintenance into separate time windows. Most operations on the server have an optimal order, and in some cases must be completed without overlap to avoid resource contention problems.

What to do next

In addition to tasks that are completed in all IBM Storage Protect environments, you might have to plan for optional processes.

Daily operations for directory-container storage pools

Schedule daily operations for the server depending on the type of storage pool that you use. You can complete specific tasks with directory-container storage pools.

About this task

The following image illustrates how IBM Storage Protect tasks fit into the daily schedule.

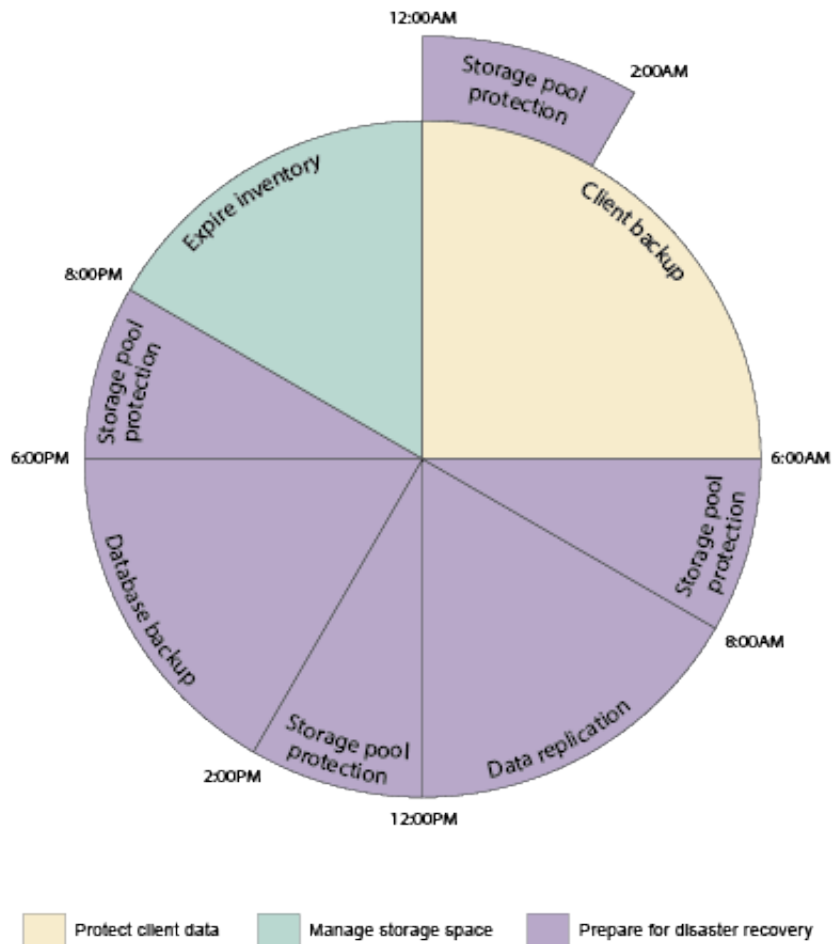


Figure 22. Daily schedule of operations for directory-container storage pools

You can schedule daily activities for IBM Storage Protect by using the Operations Center. The Operations Center creates the storage pool protection schedules when you use the wizards to configure replication or add a directory-container storage pool. You can also use the Operations Center to schedule client backups.

To manually create a schedule for daily operations, use the **DEFINE SCHEDULE** command.

Procedure

1. Perform an incremental backup of all clients on the network by using the **incremental** client command or use another supported method for client backup operations.
2. Create a DR copy of the IBM Storage Protect database by using the **BACKUP DB** command.
3. Protect data in directory-container storage pools. You can choose either of the following methods to protect data:
 - Use the **PROTECT STGPOOL** command. This method reduces data replication time.
 - Use replication storage rules. If you use replication storage rules to replicate data, you don't have to issue the **PROTECT STGPOOL** command separately. Replication storage rules combine the functionality of the **REPLICATE NODE** and **PROTECT STGPOOL** commands into one replication operation, which is designed to improve the overall performance of the replication operation.
4. Perform replication to create a secondary copy of the client data on another IBM Storage Protect server. You can choose either of the following methods to replicate data:
 - Use the **REPLICATE NODE** command if you are replicating to a single target server. Perform node replication to create a secondary copy of the client data on another IBM Storage Protect server by using the **REPLICATE NODE** command.

- Define replication storage rules if you have one or more target replication servers. Perform replication to replicate data from source replication server to one or more target replication servers by defining replication storage rules.
5. Remove objects that exceed their allowed retention period by using the **EXPIRE INVENTORY** command.

Daily operations for storage pools on FILE and DISK devices

Schedule daily operations for the server depending on the type of storage pools you use. You can complete specific tasks with storage pools on FILE and DISK devices.

About this task

The following image illustrates how IBM Storage Protect operations fit into the daily schedule.

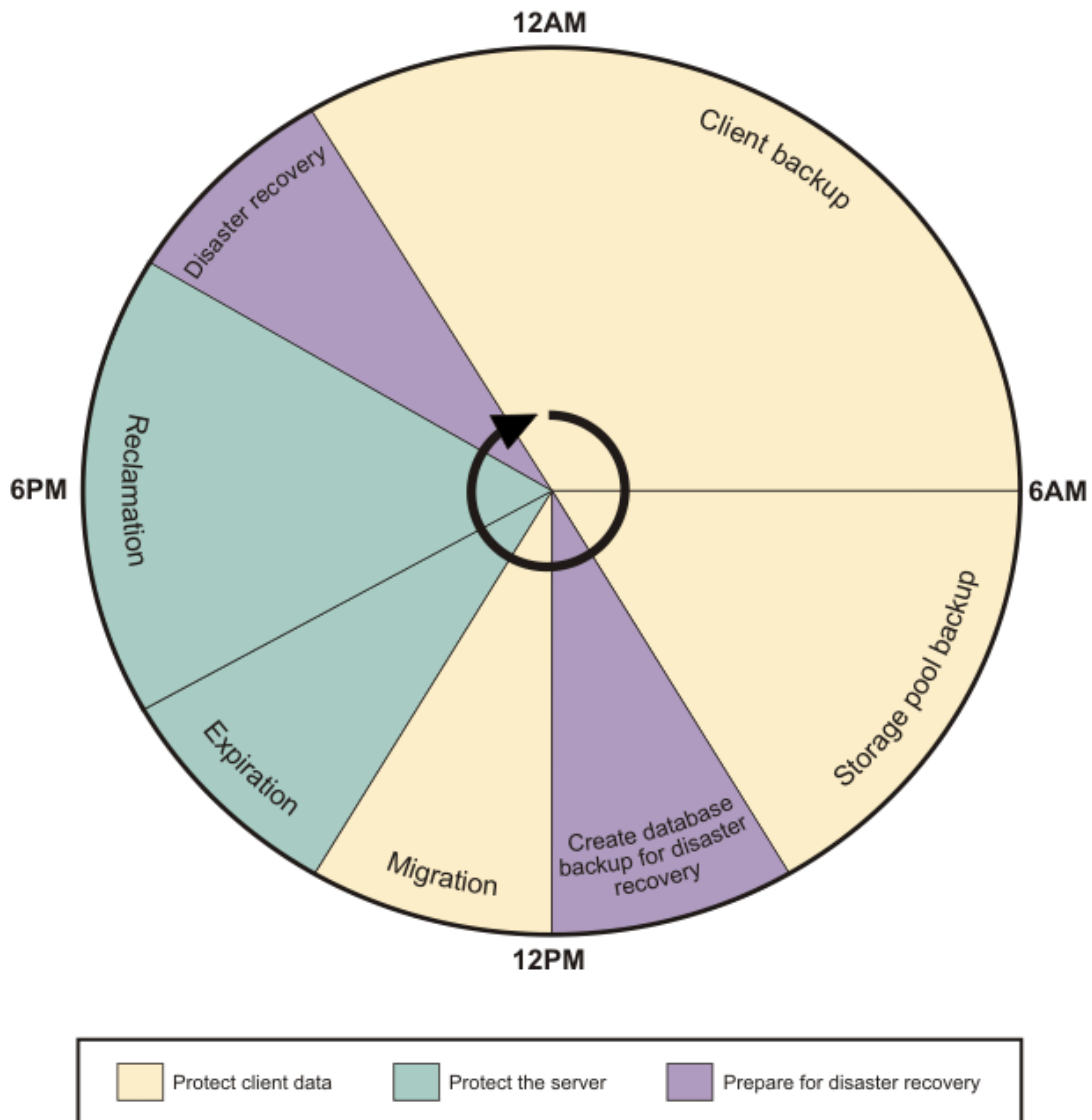


Figure 23. Daily schedule of server operations for storage pools on FILE or DISK devices

To follow best practices, schedule daily operations for IBM Storage Protect. Sample commands are provided to implement each step. All listed commands are server commands unless otherwise noted.

Procedure

1. Perform an incremental backup of all clients on the network by using the **incremental** client command or use another supported method for client backup operations.
2. Create a secondary disaster recovery (DR) copy of your client data by using the **BACKUP STGPOOL** command.
If you are copying active data, complete this operation during the storage pool backup window.
3. Create a DR copy of the IBM Storage Protect database by using the **BACKUP DB** command. In addition, use the **BACKUP VOLHISTORY** and **BACKUP DEVCONFIG** commands to create DR copies of the volume history and device configuration files.
4. Migrate data from disk storage pools to tape storage pools with the **MIGRATE STGPOOL** command.
5. Remove objects that exceed their allowed retention period by using the **EXPIRE INVENTORY** command.
6. Reclaim unused space from storage pool volumes that are released through processes like data deduplication and inventory expiration by using the **RECLAIM STGPOOL** command.
7. Complete disaster recovery preparations.
For example, if you are using the IBM Storage Protect disaster recovery manager function (DRM), issue the following commands:
 - **DELETE VOLHISTORY** to remove older versions of database backups, which are no longer required.
 - **MOVE DRMEDIA** to track database backup and copy storage pool volumes that are to be moved offsite and to identify the expired or empty volumes that are to be moved onsite.
 - **PREPARE** to create a recovery plan file.

Scheduling data deduplication and replication processes

Data deduplication and data replication are optional functions that can be used with IBM Storage Protect. They provide added benefits but also require additional resources and consideration for the daily schedule.

About this task

Depending on your environment, using data deduplication and data replication can change the tasks that are required for the daily schedule. If you are replicating data to one or more replication servers to create the backup copy of your data, then storage pool backups are not needed. Likewise, you do not need to migrate your data to tape storage pools for the creation of offsite backup media.

The following image illustrates how to schedule data deduplication and data replication processes to achieve the best performance. Tasks that overlap in the image can be run at the same time.

Restriction: The amount of duplicate identification processes that can be overlapped is based on the processor capability of the IBM Storage Protect server and the I/O capability of the storage pool disk.

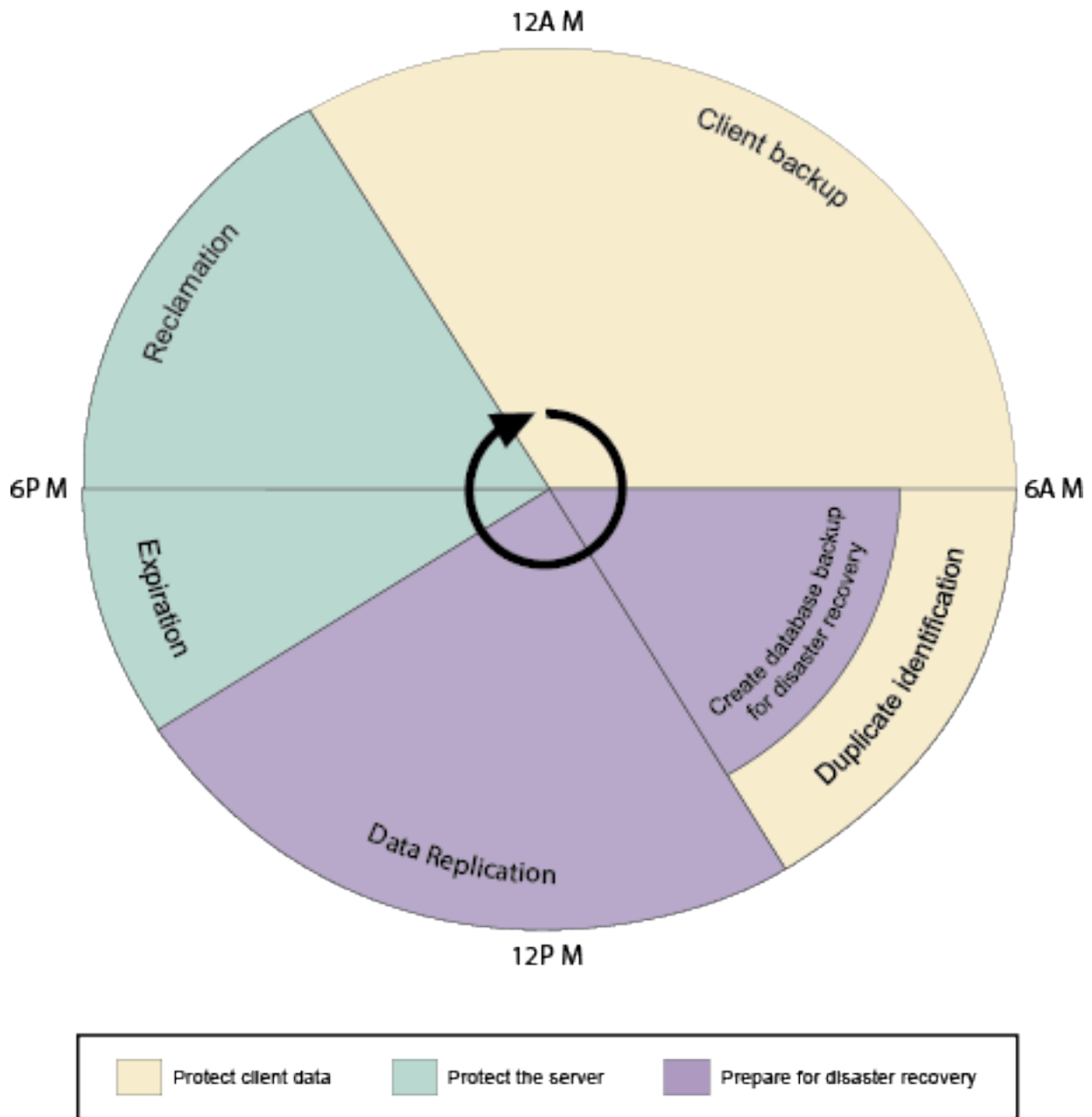


Figure 24. Daily schedule when data deduplication and data replication are used

The following steps include commands to implement the schedule that is shown in the image. For this example, tape is not used in the environment.

Procedure

1. Perform an incremental backup of all clients on the network to a deduplicated file storage pool by using the **incremental** client command or use another supported method for client backup.
2. You can run the following tasks in parallel:
 - a) Perform server-side duplicate identification by running the **IDENTIFY DUPLICATES** command. If you are not using client-side data deduplication, this step processes data that was not already deduplicated on your clients.
 - b) Create a disaster recovery (DR) copy of the IBM Storage Protect database by running the **BACKUP DB** command. In addition, run the **BACKUP VOLHISTORY** and **BACKUP DEVCONFIG** commands to create DR copies of the volume history and device configuration files.

3. Perform replication to create a secondary copy of the client data on another IBM Storage Protect server. You can either use the **REPLICATE NODE** command or define replication storage rules to replicate client data to target replication servers. Consider using replication storage rules for replicating data as it offers new features which can improve overall replication operations.
By performing data replication after duplicate identification processing, you can take advantage of data reduction during replication.
4. Remove objects that exceed their allowed retention by using the **EXPIRE INVENTORY** command.
5. Reclaim unused space from storage pool volumes that are released through data deduplication and inventory expiration by using the **RECLAIM STGPOOL** command.

Related concepts

Checklist for data deduplication

Data deduplication requires more processing resources on the server or client. Use the checklist to verify that hardware and your IBM Storage Protect configuration have characteristics that are key to good performance.

Checklist for implementing data replication

A successful implementation of data replication relies on sufficient, dedicated hardware resources. Increased amounts of memory and processor cores are required. The database and its logs must be appropriately sized to ensure that transactions can complete. A dedicated network, with enough bandwidth to handle the amount of data you intend to replicate, is required.

Compatibility and resource usage for server processes

Review this information about resource requirements and compatibility issues to help plan your daily schedule and run processes in the optimal order.

The table contains the following information for server tasks:

Process

Lists the process or operation that is performed by the IBM Storage Protect server.

Requirements and recommendations

Lists any requirements that must be met before a process can be performed. Best practice information is also covered where it is applicable.

Compatibility issues

Lists any compatibility issues that might arise when you are running processes together.

Prerequisite tasks

Lists tasks that must be completed before the process is performed.

Resource implications

Lists resources that are required to run the process and provides guidance on how much usage can be expected:

Low

Resource usage is low. Running the process does not affect other operations.

Moderate

Resource usage is moderate. Running the process might affect other operations.

High

Resource usage is high. Dedicate the resource to running the process until it completes.

Tip: Mount points and volumes are used for most server processes. Because the use of these resources is highly variable, depending on environment configuration, the table does not include a usage designation.

For operations that use file mount points with a device class of type FILE, set the mount limit parameter of the device class high enough to accommodate all simultaneous mounts. For example, the number of parallel backup sessions for a database backup are typically not more than 5, but for client backup, the mount point requirement can be in the range of 500 - 1000.

For operations that use physical tape mounts, the mount points are limited by the number of actual tape drives. When you are backing up storage pools to tape, plan to use parallel storage pool backup processes that do not exceed the number of available tape drives, and possibly leave some drives unused to remain available for client restores.

<i>Table 22. Server process requirements</i>				
Process	Requirements and recommendations	Compatibility issues	Prerequisite tasks	Resource implications
Backing up the database	None	None	Backing up storage pools	<ul style="list-style-type: none"> – Mount points and volumes – Processor (low) – Memory (low) – Disk or tape I/O (moderate) – Database I/O (high)
Backing up or archiving client data	<p>Requirement: Define and configure client nodes within the IBM Storage Protect server.</p> <p>Recommendation: Back up storage pools immediately after the main client backup or archive operation is finished to ensure that a complete copy is created for the primary storage pool.</p>	<p>Expiring inventory</p> <p>Running inventory expiration while you are backing up clients can cause resource contention problems. If expiration is processing a node that is being backed up, performance degradation is usually the result.</p> <p>Backing up storage pools</p> <p>Wait for client backups to finish before you start a storage pool backup. Otherwise, the storage pool backup copy does not include the entire client backup.</p>	None	<ul style="list-style-type: none"> – Mount points and volumes – Locking (moderate) – Processor (moderate) – Memory (moderate) – Disk or tape I/O (moderate) – Database I/O (moderate) – Network (moderate to high)
Backing up storage pool	<p>Requirement: Store new data in the primary storage pool.</p>	None	Backing up client data	<ul style="list-style-type: none"> – Mount points and volumes – Locking (low) – Processor (moderate) – Memory (moderate) – Disk or tape I/O (moderate) – Database I/O (moderate)

Table 22. Server process requirements (continued)

Process	Requirements and recommendations	Compatibility issues	Prerequisite tasks	Resource implications
Copying active data	Requirement: Store new active data in the primary storage pool.	None	Backing up client data	<ul style="list-style-type: none"> – Mount points and volumes – Locking (low) – Processor (moderate) – Memory (moderate) – Disk or tape I/O (moderate) – Database I/O (moderate)
Expiring inventory	Requirement: Deactivated data must exist on the server. Recommendation: Run inventory expiration within its own processing window as much as possible. In addition, run inventory expiration before the reclamation process to ensure that the process reclaims as much space as possible, considering policy definitions.	Backing up client data Expiring inventory while you are backing up clients can cause resource contention problems. If expiration is processing a node that is being backed up, performance degradation is usually the result.	None	<ul style="list-style-type: none"> – Locking (high) – Processor (high) – Memory (moderate) – Database I/O (high)
Generating backup sets	Requirement: Store data in at least one primary storage pool.	None	None	<ul style="list-style-type: none"> – Mount points and volumes – Locking (low) – Processor (moderate) – Memory (moderate) – Disk or tape I/O (moderate) – Database I/O (moderate)

Table 22. Server process requirements (continued)

Process	Requirements and recommendations	Compatibility issues	Prerequisite tasks	Resource implications
Identifying duplicates	<p>Requirement:</p> <p>Store new data that is not deduplicated from client-side deduplication in a primary storage pool that is enabled for server-side deduplication.</p> <p>Recommendation:</p> <p>Run duplicate identification before reclamation (as much as possible).</p>	None	<p>Potential prerequisite:</p> <p>If you are backing up storage pools, the process might not run at optimal speed against objects that are already identified. In heavy deduplication environments, it can be beneficial to back up storage pools before you run duplicate identification.</p>	<ul style="list-style-type: none"> – Mount points and volumes – Locking (moderate) – Processor (moderate) – Memory (moderate) – Disk or tape I/O (moderate) – Database I/O (moderate)
Migrating storage pools	<p>Requirement:</p> <p>Store data in at least one primary storage pool.</p>	None	<p>Potential prerequisite:</p> <p>If data deduplication is being used in the storage pool that is being migrated, and the target storage pool is deduplicated, run duplicate identification before you move or migrate that data.</p>	<ul style="list-style-type: none"> – Mount points and volumes – Locking (high) – Processor (moderate) – Memory (moderate) – Disk or tape I/O (moderate) – Database I/O (moderate)
Moving data	<p>Requirement:</p> <p>Store data in at least one primary storage pool.</p>	None	<p>Potential prerequisite:</p> <p>If data deduplication is being used in the storage pool that is being migrated, and the target storage pool is deduplicated, run duplicate identification before you move or migrate that data.</p>	<ul style="list-style-type: none"> – Mount points and volumes – Locking (high) – Processor (moderate) – Memory (moderate) – Disk or tape I/O (moderate) – Database I/O (moderate)

Table 22. Server process requirements (continued)

Process	Requirements and recommendations	Compatibility issues	Prerequisite tasks	Resource implications
Moving data by node	Requirement: Store data in at least one primary storage pool.	None	Potential prerequisite: If data deduplication is being used in the storage pool that is being migrated, and the target storage pool is deduplicated, run duplicate identification before you move or migrate that data.	<ul style="list-style-type: none"> – Mount points and volumes – Locking (high) – Processor (moderate) – Memory (moderate) – Disk or tape I/O (moderate) – Database I/O (moderate)
Reclaiming volumes in an onsite storage pool	Requirement: Store data on storage pool volumes that are expired. In addition, put data on storage pool volumes that are identified as duplicated (through the identify duplicates process).	None	<p>Expire inventory before you reclaim volumes in an onsite storage pool.</p> <p>Potential prerequisite: If deduplication is being used for the storage pool that is being reclaimed, complete duplicate identification and a storage pool backup before deduplicating data.</p>	<ul style="list-style-type: none"> – Mount points and volumes – Locking (high) – Processor (moderate) – Memory (moderate) – Disk or tape I/O (moderate) – Database I/O (moderate)
Reclaiming volumes in an offsite storage pool	Requirement: Store data on storage pool volumes that are expired. In addition, data is on storage pool volumes that are identified as duplicated (through the identify duplicates process). The data must be in a copy storage pool that is flagged as offsite.	None	<p>Expire inventory before you reclaim volumes in an offsite storage pool.</p> <p>Potential prerequisite: If deduplication is being used for the storage pool that is being reclaimed, complete duplicate identification and a storage pool backup before deduplicating data.</p>	<ul style="list-style-type: none"> – Mount points and volumes – Locking (high) – Processor (moderate) – Memory (moderate) – Disk or tape I/O (moderate) – Database I/O (moderate)

Table 22. Server process requirements (continued)				
Process	Requirements and recommendations	Compatibility issues	Prerequisite tasks	Resource implications
Replicating nodes	<p>Requirement:</p> <p>Store data in at least the primary storage pools and define and prepare a target server for replication.</p> <p>Recommendation:</p> <p>If you are using data deduplication for the replication process, run identify duplicates to completion in the primary storage pools before you run replication. This recommendation can be ignored if you are using client-side data deduplication for your entire environment.</p>	None	<p>Back up client data before you replicate nodes</p> <p>Potential prerequisite:</p> <p>If the replication process relies on data that is being deduplicated, run duplicate identification against all data that is being replicated.</p>	<ul style="list-style-type: none"> – Mount points and volumes – Locking (moderate) – Processor (moderate) – Memory (moderate) – Disk or tape I/O (moderate) – Database I/O (moderate) – Network (moderate to high)

Avoiding contention for server resources during client operations

IBM Storage Protect requires many resources to complete client backup, archive, or hierarchical storage management migration operations. If multiple sessions are in contention for server resources, system performance can be affected.

Server resources that are used during client operations include database logs, server locks, drives, mount points, or volumes. For example, a client backup session might need a mount point, tape drive, and tape volume on which to store data. After these resources are allocated to the backup session, another client session to restore data that is on the tape volume might start. The restore session is delayed until the backup session unmounts the volume and releases it.

Resource contention has a direct effect on performance and the ability to complete an operation in a timely manner. The problem of resource contention is more critical with long-running client sessions or server processes. Because new database entries are stored in the recovery log until they are committed to the database, long-running sessions or processes can insert many of these entries under a single transaction, pinning the recovery log. A pinned recovery log prevents all transactions from applying database changes and causes server processes to run slowly.

You can schedule client sessions and server processes at different times to avoid resource contention and delays. When you are setting up schedules, you might need to prevent some server processes from starting automatically. For example, disable expiration, migration, reclamation, and duplicate-identification so that they can be scheduled later. Use administrator-command schedules to run these operations daily.

Related reference

[Compatibility and resource usage for server processes](#)

Review this information about resource requirements and compatibility issues to help plan your daily schedule and run processes in the optimal order.

Disabling automatic processes and setting schedules

Disable automatic processes like inventory expiration, migration, reclamation, and identification of duplicate data and set up schedules so that you can control when these operations are completed during the daily schedule.

About this task

Review the following sections for examples of disabling automatic operations and setting up schedules. Some general setup steps are omitted and the example uses a particular order for the server processes. However, you can reorder the processes to better fit your environment.

Procedure

1. Disable automatic inventory expiration by setting the **EXPINTERVAL** server option to zero.

```
setopt expinterval 0
```

2. Disable automatic migration and reclamation processes by using the **DEFINE STGPOOL** command to set the **HIGHMIG** and **RECLAIM** parameters to a value of **100**.

You might have to increase the number of allowed processes for migration and reclamation in order for them to complete in a reasonable amount of time. The actual number of processes depends on available tape drives. If you have already defined storage pools, you can change values for the **MIGPROCESS** and **RECLAIMPROCESS** parameters by using the **UPDATE STGPOOL** command.

```
def devc LARGEFILE devt=file mountlimit=500 maxcap=20480m dir=/tsmfile
def stg FILEPOOL LARGEFILE maxscratch=200 reclaim=100 hi=100 lo=0 migpr=4
reclaimpr=20 next=tapepool
```

3. If you have storage pools that are defined with data deduplication enabled, disable the duplicate identification processes:

```
def stg FILEPOOL LARGEFILE maxscratch=200 reclaim=100 hi=100 lo=0 dedup=yes
identifypr=0 migpr=4 reclaimpr=4
```

Example: Setting a schedule for the client backup window

This example initiates an incremental backup of all associated nodes in the **STANDARD** domains.

The schedule starts daily at 8:00 PM by using server-prompted scheduling mode. The long-running schedules continue past the duration, so a shorter duration can be used to force those schedules to start close to the beginning of the start window.

```
def schedule standard nightly_backups description="Nightly backups of nodes in
domain standard" starttime=20:00 duration=5 durunits=hours period=1 perunits=days
```

Example: Setting up the server maintenance schedule

Schedule server maintenance operations to run outside of the client backup window, with as little overlap as possible.

You can control the timing of schedules for maintenance tasks by setting the start time in combination with the duration time for each operation. Here is an example of how you might time each process:

08:00 - end

Storage pool backup.

11:00 to 13:00

Identification of duplicates.

13:00 - 15:00

Inventory expiration.

14:00 - 16:00

Storage pool migration.

16:00 - 18:00

Reclamation processing.

18:00 - end

Database backup, including volume history and device configuration backup.

20:00 - end

Client backup.

After you determine a timeline, use the **DEFINE SCHEDULE** command to create schedules for each process. You can include scripts in each schedule so that commands are processed automatically. Use the **DEFINE SCRIPT** command to create a script and the **UPDATE SCRIPT** command to add lines.

The following scripts are examples of how to define each server task:

Storage pool backup

```
/*-----*/
/* Storage Pool Backup */
/*-----*/
def script STGBACKUP "/* Run stg pool backups */"
upd script STGBACKUP "backup stg archivepool cypool maxproc=4
wait=yes" line=005
upd script STGBACKUP "backup stg backuppool cypool maxproc=4
wait=yes" line=010
upd script STGBACKUP "backup stg filepool cypool maxproc=4 wait=yes"
line=020
upd script STGBACKUP "backup stg filepool2 cypool maxproc=4 wait=yes"
line=025
upd script STGBACKUP "backup stg tapepool cypool maxproc=3 wait=yes"
line=030
def sched STGBACKUP type=admin cmd="run STGBACKUP" active=yes desc="Run
all stg pool backups." \
startdate=today starttime=08:00:00 dur=45 duru=minutes per=1 peru=day
commit
```

Identification of duplicates

```
/*-----*/
/* Deduplication */
/*-----*/
def script DEDUP "/* Run identify duplicate processes. */"
upd script DEDUP "identify duplicates FILEPOOL numpr=4 duration=120" \
line=010
upd script DEDUP "identify duplicates FILEPOOL2 numpr=2 duration=120" \
line=015
def sched DEDUP type=admin cmd="run DEDUP" active=yes desc="Run identify
duplicates." \
startdate=today starttime=11:00:00 dur=45 duru=minutes per=1 peru=day
commit
```

Inventory expiration

```
/*-----*/
/* Expiration */
/*-----*/
def script EXPIRE "/* Run expiration processes. */"
upd script EXPIRE "expire inventory wait=yes duration=120" line=010
def sched EXPIRATION type=admin cmd="run expire" active=yes desc="Run
expiration." \
startdate=today starttime=13:00:00 dur=45 duru=minutes per=1 peru=day
commit
```

Storage pool migration

```
/*-----*/
/* Storage Pool Migration */
/*-----*/
def script MIGRATE "/* Run stg pool migration */"
upd script MIGRATE "migrate stg archivepool duration=30 wait=yes" line=005
```

```

upd script MIGRATE "migrate stg backuppool duration=30 wait=yes" line=010
upd script MIGRATE "migrate stg filepool2 duration=60 wait=yes" line=015
def sched MIGRATE type=admin cmd="run MIGRATE" active=yes desc="Migrate
data to tape pools ." \
    startdate=today starttime=14:00 dur=45 duru=minutes per=1 peru=day
commit

```

Reclamation processing

```

/*-----*/
/* Storage Pool Reclamation */
/*-----*/

def script RECLAIM "/* Run stg pool reclamation */"
upd script RECLAIM "reclaim stg filepool threshold=40 duration=120
wait=yes" line=005
upd script RECLAIM "reclaim stg filepool2 threshold=40 duration=120
wait=yes" line=008
upd script RECLAIM "reclaim stg tapepool threshold=60 duration=60 wait=yes"
line=010
def sched RECLAIM type=admin cmd="run RECLAIM" active=yes desc="Reclaim
space from FILEPOOL and TAPEPOOL." \
    startdate=today starttime=16:00 dur=45 duru=minutes per=1 peru=day
commit

```

Database backup, including volume history and device configuration backup

```

/*-----*/
/* Database Backup */
/*-----*/

def script DBBACKUP "/* Run DB backups */"
upd script DBBACKUP "backup db devc=ts3310devc type=full wait=yes" line=005
upd script DBBACKUP "backup volhist" line=010
upd script DBBACKUP "backup devconf" line=015
def sched DBBACKUP type=admin cmd="run DBBACKUP" active=yes desc="Run
database backup." \
    startdate=today starttime=18:00:00 dur=45 duru=minutes per=1 peru=day
commit

```

Tuning database backups to cloud object storage

You can back up a database to, and restore a database from, cloud object storage for disaster recovery purposes.

When you use a CLOUD device class for IBM Storage Protect database backup operations, the following files are copied to object storage:

- Database volumes
- Device configuration file
- Volume history file
- Master encryption key for the server

Large items, such as database volumes, are copied to object storage by using multipart upload. By specifying multiple, concurrent data streams, you can reduce the time that is required to back up the database. The number of data streams that are used for the database backup operation is the same as the number of data streams that are required for any subsequent database restore. The number of data streams affects throughput. Each database backup operation uses the following separate data resources:

- A session connection from Db2 to the IBM Storage Protect server
- A server thread that sends data from the server to object storage

When you back up the database to cloud object storage, consider the following questions:

Are you using object storage endpoints?

To optimize performance, use multiple object storage endpoints, such as IBM Cloud Object Storage Accessers, instead of a load balancer.

How many IBM Cloud Object Storage Accessers are you using?

Use the following number of IBM Cloud Object Storage Accessers for small, medium, and large Blueprint systems:

- Small system: 1 IBM Cloud Object Storage Accesser
- Medium system: 2 IBM Cloud Object Storage Accessers
- Large system: 3 - 4 IBM Cloud Object Storage Accessers

Tip: The IBM Cloud Object Storage Accessers can be used for other storage requirements in addition to the IBM Storage Protect server.

Are the disks configured for optimal performance?

The following items can impact the performance of database backup and restore operations:

- Database disks
- Object storage system
- Network to the object storage system

Use benchmarking tools to determine the throughput capability of the network, object storage, and database disks. For more information, see [“Optimizing database backup operations to cloud object storage” on page 164.](#)

Is network bandwidth greater than the planned maximum throughput for backup operations?

For system operations, such as backups, network bandwidth must be greater than the planned maximum throughput. The system must complete operations on schedule to meet service level commitments.

When the TCP/IP network link to object storage shows signs of packet loss, the performance for database backup and restore operations requires improvement. Packet loss of 2% or more, dropped packets, or both, can cause a large degradation in throughput for database backup or restore operations with object storage.

When you use an off-premises IBM Cloud Object Storage system with larger IBM Storage Protect servers, use a dedicated network link to the object storage resource. A 1 Gb network might be enough for a small Blueprint system. A 10 Gb network is necessary for medium and large Blueprint systems with larger databases when regular database backup and time-intensive restore operations must be completed in an acceptable amount of time. For example, a 1 Gb network link might process only 100 MB per second of throughput. A 1 TB database backup operation might take 3 or more hours to complete.

Are you using Db2 database encryption?

Db2 database encryption is used by default for database backup operations to cloud device classes. You can specify encryption or compression for a database backup operation to cloud, but not both. If you specify compression, for a database backup to cloud, encryption is disabled.

Compression impacts backup performance and limits front-end throughput to approximately 0.5 TB per hour, or less. Use compression with database restore operations to improve performance. For smaller IBM Storage Protect servers with smaller databases, use compression when the following conditions are met:

- The network link to object storage is 1 Gb or less.
- Database encryption is not necessary.
- Compression savings are required.

How many streams are you using for the database backup?

Depending on IBM Storage Protect server size, use the following stream quantities for database backup operations for small, medium, and large Blueprint systems:

- Small system: 10 streams
- Medium system: 25 streams
- Large system: 50 streams

Adjust the number of data streams until you achieve optimal throughput.

Backing up a database to cloud object storage uses approximately 20 MB per data stream on the IBM Storage Protect server. For example, a 50-stream database backup operation consumes approximately 1000 MB of memory on the server.

Restriction: If the process of backing up the database to cloud object storage takes longer than expected, check the Operations Center or activity log to determine whether the database backup operation failed and whether the operation was retried. When a database backup operation is retried, the operation uses a single backup stream, which might be insufficient for your throughput requirements. To optimize throughput, cancel the database backup operation and try to manually back up the database with a larger stream count.

Is throughput sufficient for database backup operations based on your server maintenance schedule?

Throughput requirements for database backup operations depend on the IBM Storage Protect server maintenance schedule. A typical full database backup window is 2 hours daily. For example, an 8 TB database must back up at least 4 TB per hour to comply with the backup window. Four TB per hour is approximately the limit of a single 10 Gb Ethernet link. The database disks must manage approximately 1200 MB per second of larger (256 – 512 KB) input/output operations per second (IOPS). More throughput is required if concurrent operations occur on the IBM Storage Protect server in addition to database backup operation. Longer scheduled windows can be used to accommodate slower throughput.

Optimizing database backup operations to cloud object storage

You can back up an IBM Storage Protect database to cloud object storage. By having a database backup, you can simplify the recovery process in case of a disaster and help to ensure high availability of your system. Ensure that you take steps to optimize the backup operation.

Procedure

1. Determine the maximum read throughput of the database disk. Complete the steps in [“Calculating throughput for IBM Storage Protect database disks”](#) on page 164.
2. Determine the maximum achievable throughput for data transfer to the object storage system. Complete the steps in [“Calculating throughput for object storage”](#) on page 165.
3. Configure the database backup operation with the optimal number of data streams to ensure that the database backup operation can be completed within the specified time window.
4. Schedule a full database backup to cloud daily. The starting number of streams that you use depends on the size of the IBM Storage Protect server.

What to do next

Monitor the throughput of daily database backup operations to cloud object storage. Incrementally use more or fewer streams until you achieve the optimal daily throughput or when the disk or object storage maximum throughput is reached. Increase or decrease the streams incrementally, for example, by 5 streams, and record data points for multiple days at a particular setting for precise estimates.

Calculating throughput for IBM Storage Protect database disks

You can use benchmarking tools to measure throughput capability of the IBM Storage Protect database disks during full database backup operations.

Procedure

1. Use the **tsmdiskperf.pl** Perl script as a benchmarking tool to determine the input/output operations per second (IOPS) capability of the database disks:
 - a. Benchmark the directory paths that are used for the read-only workload of database sequential disks with a 256 KB block size.

b. To run the script, issue the following command:

```
perl tsmdiskperf.pl workload=stgpool mode=readonly fslist=directory_list
```

where *directory_list* is a comma-separated list of database directories.

Tip: Specify the **workload=stgpool** parameter to ensure that sequential reads occur during database backup operations.

c. Ensure that the read-only data ingestion rate that is obtainable for these database paths meets the speed requirements for full database backups to finish within the scheduled timeframe.

For benchmarking tools and sample benchmarking tests, see the [IBM Storage Protect Blueprint](#). The benchmarking tool **tsmdiskperf.pl** is available in the *Blueprint configuration scripts* package.

2. Run the benchmarking tool again and include more database disk subdirectories until throughput stabilizes or decreases.
3. Use the highest throughput value as the maximum to estimate database disk throughput during database backup operations to cloud object storage.
4. If the database disk throughput estimate is too low to achieve adequate throughput for database backup operations, reconfigure the database disks. For example, you can get greater random and sequential throughput by provisioning more physical disk volumes for database disk volume groups.

Calculating throughput for object storage

You can calculate the throughput capability of an object storage system and network for optimal performance.

Before you begin

Use a memory-mapped file system location to provide the source data for benchmarking. If a memory-mapped file system is not possible in your system environment, use subdirectories on the database disks. The preferred method to eliminate bottlenecks on the source disk is to use memory-mapped file system locations, such as the temporary file system (tmpfs) on Linux operating systems.

Procedure

1. Use both of the following benchmarking tools to measure throughput capability:

- `tsmobjperf.pl` Perl script
- `SPObjBench.jar` Java application

To obtain the benchmarking tools, see the [Cloud Blueprints](#). For instructions about using the benchmarking tools, see the *Cloud Cache and Object Storage Benchmarking.pdf* guide, which is included with the Cloud Blueprints.

2. Populate a set of 10 1-GB files in a memory-mapped file system. For example, on a Linux system with at least 11 GB of free RAM, issue the following commands:

```
mkdir /mnt/ramdisk
mount -t tmpfs -o size=11g tmpfs /mnt/ramdisk
for I in `seq 10`; do dd if=/dev/urandom of=/mnt/ramdisk/file.$I bs=1048576 count=1024;
done
```

3. To run a set of automated tests that scales from 1 - 100 data threads, issue the following command:

```
perl tsmobjperf.pl type=type endpoints=endpoint user="user"
pass="pass" bucket=bucket min=1 max=100 step=10
fslist=comma_delimited_source_files_list
```

where:

type

Specifies the the Simple Storage Service (S3) protocol such as IBM Cloud Object Storage, Amazon Simple Storage Service (Amazon S3), or other approved S3 protocol object storage systems. Use Azure for Microsoft Azure Blob storage.

endpoints

Specifies a comma-delimited list of one or more IP addresses or URLs for the object storage endpoints. Use the same endpoints as planned for database backup operations to cloud object storage. For Microsoft Azure based systems, specify the URL of a user's blob storage account.

user

For S3, *user* specifies a public key ID. For Azure, *user* specifies the blob storage account name. Enclose the parameter value in double quotation marks.

pass

For S3, *pass* specifies the secret key for a user who has valid S3 credentials to create buckets and PUT and GET objects in the region that is specified in the endpoint URL. For Azure, the *pass* value must be a shared access signature (SAS) token with sufficient read/write access rights to the blob storage account over an HTTPS or HTTP connection. Enclose the parameter value in double quotation marks.

bucket

Identifies an S3 bucket, vault name, or an Azure container name for which a credentialed user has PUT and GET access. The bucket name must exist in the object storage system.

min and max

The *min* and *max* values specify the minimum and maximum thread counts that are tested.

step

Specifies the increase in thread count from test to test.

fslist

Specifies a comma-delimited list of source files that are used for multipart upload. Use the source files that you created earlier.

Tips:

- Each thread count test uploads 10 1-GB objects per thread. The tool does not remove objects that are created during the test. After the test, you must manually remove objects that are created on the object storage system.
- When you estimate the object storage throughput for the system, use the highest throughput value that was achieved by the automated tests.
- If the throughput estimate is too low to achieve adequate throughput for database backup operations, use more object storage endpoints. Investigate object storage system-related bottlenecks or reconfigure the network to the object storage system. For example, consider a dedicated link for off-premises clouds, use larger capacity Ethernet ports (10 Gb instead of 1 Gb), or use additional bonded ports to achieve greater throughput.

Tuning the performance of data replication

After you replicate data, you can measure the effectiveness of your configuration and tune the speed of the replication processes.

About this task

You can use commands that are specific to data replication to tune performance.

Measuring the effectiveness of a replication configuration

A replication configuration is optimal if the number of replicated files that are stored on a target replication server is equal to the number of files that are stored on the source replication server. Use the **QUERY REPLNODE** command to display the number of files that are stored on source and target replication servers.

Increasing the speed of a data replication process

You can set the **REPLBATCHSIZE** server option together with the **REPLSIZEETHRESH** server option to increase the speed of processing between two replicating servers. These options specify how many files to include in a batch transaction and define a threshold for the size of the batch, in megabytes.

The default value for each option, which is 4096, is the best practice setting. If you want to improve the performance of a data replication process, try tuning the **REPLBATCHSIZE** and the **REPLSIZETHRESH** server options. Change the default values only after you monitor the data replication performance over several sessions. When the default values of the options are increased, the server requires more space in the active log. You might need to allocate space for the active log that is two or more times larger than an active log size that uses the default size of 4096. In addition, the server might require a longer initialization time when you start.

Use a trial and error method to increase the server options. You can increase the server options in any order. Start by increasing one of the options incrementally by 10%. If replication performance does not improve, revert the setting to its original value. Increase the other option incrementally by 10%. Ensure that you monitor active log usage during the first few replication operations to ensure sufficient active log space is available. Larger transactions run longer and use more active log space that causes other server processes to run slowly. If server processes run slowly, reduce the options until the replication and other server processes can complete.

Tuning server-side data deduplication

Tune settings and configuration for different operations to ensure that the performance of server-side data deduplication is efficient.

Procedure

Tip: The following steps do not apply to container storage pools.

1. Control processor resources by setting the number of duplicate identification processes that you want to use.

Do not exceed the number of processor cores available on your IBM Storage Protect server when you set the **NUMPROCESS** value. Define a duration limit for the **IDENTIFY DUPLICATES** command, otherwise, processes that are running after the command is issued run indefinitely.

2. Determine the threshold for reclamation of a deduplicated storage pool.

A deduplicated storage pool is typically reclaimed to a threshold that is less than the default of 60 to allow more of the identified duplicate extents to be removed. Experiment with the setting of this value to find a threshold that can be completed within the available time.

3. Determine how many reclamation processes to run.

Tip: A reclamation setting of more than 25 and less than 40 is sufficient.

4. Schedule data deduplication processing that is based on how you create a second copy of your data.

If you are backing up your storage pool, do not overlap client backup and duplicate identification. Complete the storage pool backup before the identify process. If the storage pool backup is not complete, the copy process takes longer because it requires the deduplicated data to be reassembled before the backup.

You can overlap duplicate identification and client backup operations in the following scenarios:

- You are not backing up your storage pool.
- You are using data replication to create a secondary copy of your data.

Running these operations together can reduce the time that is required to finish processing, but might increase the time for client backup.

5. To prevent deadlocks in the IBM Storage Protect server, you might need to modify the Db2 **LOCKLIST** parameter before you deduplicate a large amount of data.

When the amount of concurrent data movement activity is high, deadlocks can occur in the server. If the amount of concurrent data that is moved exceeds 500 GB at a time, adjust the Db2 **LOCKLIST** parameter as follows:

Table 23. Tuning Db2 LOCKLIST parameter values	
Amount of data	LOCKLIST parameter value
500 GB	122000
1 TB	244000
5 TB	1220000

Related concepts

Checklist for data deduplication

Data deduplication requires more processing resources on the server or client. Use the checklist to verify that hardware and your IBM Storage Protect configuration have characteristics that are key to good performance.

Related tasks

Scheduling data deduplication and replication processes

Data deduplication and data replication are optional functions that can be used with IBM Storage Protect. They provide added benefits but also require additional resources and consideration for the daily schedule.

Evaluating data deduplication results

You can evaluate the effectiveness of IBM Storage Protect data deduplication by examining various queries or reports. Actual data reduction results can show whether the expected storage savings are achieved. You can also evaluate other key operational factors, such as database utilization, to ensure that they are consistent with expectations.

Tuning client-side data deduplication

The performance of client-side data deduplication can be affected by processor requirements and deduplication configuration.

Restoration of deduplicated data

Restore operations that require data to be reassembled from a sequential-access disk (FILE) storage pool that is set up for data deduplication have different performance characteristics than restore operations from a FILE storage pool that is not set up for deduplication.

In a FILE storage pool that is not set up for data deduplication, files are typically restored in a sequential process. In a FILE storage pool that is set up for data deduplication, however, data is distributed throughout the storage pool. As a result, the input/output (I/O) is more random, which can lead to slower restore times. In addition, more server processor resources are consumed when data is restored from a deduplicated storage pool. This occurs because the data is checked to ensure that it has been reassembled properly by using MD5 algorithms.

Although restore operations of small files from a deduplicated storage pool might be relatively slow, these operations are still typically faster than restore operations of small files from tape because of the added tape mount and locate time.

Improving read performance for deduplicated storage pools

To obtain the different extents that make up a file from a deduplicated storage pool, client restore operations and certain server processes might require opening and closing FILE volumes multiple times. The frequency with which FILE volumes are opened and closed during a session can severely affect performance.

About this task

Opening and closing volumes multiple times can affect the following server processes that read data from a deduplicated storage pool:

- Volume reclamation
- **MOVE DATA** or **MOVE NODEDATA**

- **EXPORT**
- **AUDIT VOLUME**
- Storage-pool restore operation
- Volume restore operation
- Data migration

To reduce the number of times a volume is opened and closed, IBM Storage Protect allows multiple input FILE volumes in a deduplicated storage pool to remain open at the same time during a session. To specify the number of open FILE volumes in deduplicated storage pools that can remain open, use the **NUMOPENVOLSALLOWED** server option. Set this option in the server options file or by using the **SETOPT** command.

During a client-restore operation, volumes can remain open for as long as the client session is active. During a no-query restore operation, the volumes remain open until the no-query restore completes. Then, all volumes are closed and released. However, for a standard restore operation started in interactive mode, the volumes might remain open at the end of the restore operation. The volumes are closed and released when the next classic restore operation is requested.

Procedure

This option can significantly increase the number of volumes and mount points in use at any one time. To optimize performance, complete the following tasks:

- To set **NUMOPENVOLSALLOWED**:
 - a. Select a beginning value. The default is 10. A small increase to this option can be beneficial but that might not be true in all environments.
 - b. Monitor client sessions and server processes.
 - c. Note the highest number of volumes open for a single session or process. If the highest number of open volumes is equal to the value specified by **NUMOPENVOLSALLOWED**, increase the setting of **NUMOPENVOLSALLOWED**.
- To prevent sessions or processes from having to wait for a mount point:
 - a. Increase the value of the **MOUNTLIMIT** parameter in the device-class definition.
 - b. Set the value of the **MOUNTLIMIT** parameter high enough to allow all client sessions and server processes that are using deduplicated storage pools to open the number of volumes that are specified by the **NUMOPENVOLSALLOWED** option.
 - c. Check the following results:
 - For client sessions, check the destination in the copy group definition to determine how many nodes are storing data in the deduplicated storage pool.
 - For server processes, check the number of processes that are allowed for each process for the storage pool.
- For any node that is backing up or archiving data into a deduplicated storage pool, set the value of the **MAXNUMMP** parameter in the client-node definition to a value at least as high as the **NUMOPENVOLSALLOWED** option. Increase this value if you notice that the node is failing client operations because the **MAXNUMMP** value is being exceeded.

Results

Each session within a client operation or server process can have as many open FILE volumes as specified by this option. A session is initiated by a client operation or by a server process. Multiple sessions can be started within each.

Tuning server operations for client backups

When possible, limit the number of versions of any backup file to the minimum required.

About this task

File backup performance is degraded when there are many versions of an object. Use the **DEFINE COPYGROUP** command and modify the VEREXISTS parameter to control the number of versions, or use the **UPDATE COPYGROUP** command. The default number of backup versions is 2.

If the retention requirements in your environment differ among client systems, use different copy groups rather than taking the lowest common denominator. For example, if your accounting systems require records to be kept for seven years, but other systems need data to be kept for only two years, do not specify seven for all of them. Instead, create two separate copy groups. Not only are backups potentially faster, but you also use less storage because you are not keeping data that you do not need.

Similarly, you can set up a separate copy group for system state backups to avoid keeping unnecessary operating system files. For example, if you want to retain system state data for a week and all other data for a year, create a separate copy group for the system state data.

Tuning operations for automatic backup-archive client deployment

You can take actions that might improve backup-archive client deployment performance.

Procedure

- Deploy the client when the clients, server, and network have minimal activity. Do not deploy during client backup operations.
- To avoid the cost of retrieving a wrong package, use separate schedules for each client architecture (for example, x86, x64, ia64).
- If multiple clients are being upgraded concurrently, store the deployment packages in a random-access DISK or in a sequential-access FILE storage pool. Both types of storage pools support read access to the same object from multiple sessions concurrently.

If the storage pool uses tape, the server serializes read access to the storage pool volume that contains the upgrade package data. Serial access is also used for a storage pool volume in a virtual tape library (VTL), even if the data is stored on physical disks.

- Provide sufficient cache memory on the disk system that is used by the random-access DISK or sequential-access FILE storage pools that contain the deployment packages. Deployment packages are read from the storage pool volumes during retrieval by using direct input/output (I/O). This means that the server operating system does not cache the data in memory, and each I/O must come from the disk system. When the server is deploying many clients at the same time, the disk system sees a high read-hit ratio on these data blocks, which leads to faster throughput.
- Balance multiple clients across the network interfaces on the server. Such balancing might be done to optimize backup performance as well.

Tuning tape drive performance

There are a few basic procedures for maintaining the performance of your tape drives.

Configuring enough tape drives

You must configure enough tape drives for operations that occur at the same time in your environment:

- The maximum number of IBM Storage Protect client sessions that are backing up directly to tape at any time during the peak backup window.
- Extra tape drives for other functions that run during the backup window. For example, storage pool migration, storage pool backup, and reclamation.

Cleaning tape drives

Cleaning the tape drive according to the manufacturer's specifications is important to ensure maximum tape drive performance. Failure to clean tape drives can cause read and write errors, drive failures, and poor performance.

Enabling tape compression

In most cases, the preferred method of enabling compression at the tape drive improves IBM Storage Protect throughput.

You can use the **FORMAT** parameter of the **DEFINE DEVCLASS** command to specify the appropriate recording format to be used when you write data to sequential access media. The default is **DRIVE**, which specifies that IBM Storage Protect selects the highest format that can be supported by the sequential access drive on which a volume is mounted. This setting usually allows the tape control unit to perform compression.

Tip: Avoid specifying the **DRIVE** value when a mixture of devices is used in the same library. For example, if you have drives that support recording formats superior to other drives in a library, do not specify the **FORMAT** parameter with the **DRIVE** value.

If you do not use compression at the client and your data is compressible, you should achieve higher system throughput if you use compression at the tape control unit unless your network is slow.

Tape drive transfer rate

Many factors affect the sustained transfer rate of IBM Storage Protect when you use tape drives. The sustained transfer rate takes into account the net effect of all these factors.

The following factors affect the sustained transfer rate:

- Native transfer rate
- Compression ratio
- File size
- Server attachment
- Server attachment host bus adapter (HBA) type
- Disk transfer rate
- Network bandwidth
- Server utilization
- Start/stop performance
- Application control file activity
- IBM Storage Protect transaction size
- Bus bandwidth
- Quality of the media
- Use of collocation for restore operations
- Distribution of data on the tape volume

Tape drive streaming rate performance

Streaming rate is the rate at which a tape drive can read and write, not including any start and stop operations. Most uses of tape include some start and stop operations, which slow down the sustained rate at which the drive operates.

When you are writing to a tape drive, the drive normally returns control to the application when the data is in the tape drive buffer but before the data is written to tape. This mode of operation provides all tape drives a significant performance improvement. However, the drive's buffer is volatile. For the application to ensure that the data is written to tape, the application must flush the buffer. Flushing the buffer

causes the tape drive to back hitch (start/stop). The IBM Storage Protect parameters **TXNBYTELIMIT** and **TXNGROUPMAX** control how frequently IBM Storage Protect issues this buffer flush command.

When you are writing to a tape drive, you must consider network bandwidth. For example, 1-gigabit Ethernet can sustain 60 - 70 MB per second. Therefore, you cannot back up to a tape drive faster than that rate.

Using high performance tape drives

When you use high performance tape drives with IBM Storage Protect, it is important to use the appropriate server and client options to enhance performance.

Consider these settings for the best performance:

Server options

```
TXNGROUPMAX 4096
MOVESIZETHRESH 32768
MOVEBATCHSIZE 1000
```

Client options

```
TXNBYTELIMIT 10G
```

If on average, IBM Storage Protect clients have files smaller than 100 KB, back these clients up to a disk storage pool for later migration to tape. This allows more efficient data movement to tape.

Tuning HBA capacity

The server system must have enough host bus adapters (HBAs) to handle data operations that IBM Storage Protect runs simultaneously.

About this task

Ensure that you have enough HBA bandwidth to handle peak loads on the system. When you are planning for peak loads, consider all operations that can occur simultaneously in your environment.

For example, if you are backing up to a disk pool, you need enough network bandwidth for the client backup operation. You also need a similar amount of outgoing bandwidth to disk through fibre, SAS, or other HBA for that backup data. More bandwidth is required if storage pool migration tends to run during the backup window. In addition to the bandwidth needed for the backup operation, you need bandwidth to read the data from disk and to write to tape. If you consider only the bandwidth that is needed for the backup operation, your available bandwidth is limited when storage pool migration starts.

Related concepts

[Potential bottlenecks in the data flow for IBM Storage Protect operations](#)

In operations such as client backup and storage pool migration, data moves through many physical components that can affect the speed of the operations. Understanding the characteristics of these components can help when you are working to improve performance.

Related reference

[Compatibility and resource usage for server processes](#)

Review this information about resource requirements and compatibility issues to help plan your daily schedule and run processes in the optimal order.

Tuning tasks for operating systems and other applications

Review this information for guidance on improving operating system performance for the server and for impacts that are associated with applications other than IBM Storage Protect.

Tuning AIX systems for IBM Storage Protect server performance

There are a number of actions that can improve performance for an IBM Storage Protect server that is running in an AIX environment.

About this task

You can use an IBM Storage Protect server in System p logical partitions (LPARs).

Procedure

- Use the `rbw` mount option to release memory from the file system cache. For more information about the release-behind sequential read and write (**rbw**) option, see the [AIX product information](#).

AIX systems can cache much file system data, which can take away memory that is needed for IBM Storage Protect server and Db2 processes. To avoid paging with the AIX server, use the `rbw` mount option for the JFS2 file system. Less memory is used for the file system cache and more is available for IBM Storage Protect.

Concurrent I/O (CIO) options are not necessary to access or mount database or log file systems. IBM Storage Protect completes mounts automatically. Also, CIO disables the read-ahead feature of JFS2 file systems, decreasing the read performance of the database during backup. Do not use the file system mount options, CIO and Direct I/O (DIO), for file systems that contain the IBM Storage Protect database, logs, or storage pool volumes. These options can cause performance degradation of many server operations.

IBM Storage Protect and Db2 can still use DIO where it is beneficial to do so, but IBM Storage Protect does not require the mount options to selectively take advantage of these techniques.
- Use the Portable Operating System Interface (POSIX) time zone specification for best system performance.
- IBM Storage Protect supports sharing Fibre Channel ports through NPIV in a pSeries LPAR. While you can share these ports, be sure that the port has adequate bandwidth for all LPARs that use that port. Sharing resources with other logical partitions can affect the performance of the IBM Storage Protect server. When there are other logical partitions on the system, you can dedicate resources to the IBM Storage Protect server partition.
- IBM Storage Protect can use 10-gigabit Ethernet ports that are shared through the VIO server. However, when you use a shared VIO port, it does not always provide full 10-gigabit throughput. If full 10-gigabit Ethernet bandwidth is needed, you can use the logical Host Ethernet Adapter (LHEA) method of sharing.

Tuning AIX virtual memory

The AIX virtual address space is managed by the Virtual Memory Manager (VMM). Monitor the paging statistics to identify problems with virtual memory usage.

Procedure

To monitor paging statistics and identify potential issues, complete the following steps:

1. Run the **vmstat** command. Review the paging statistics in the `pi` and `po` columns. An occasional nonzero value is not a concern because paging is the main principle of virtual memory. If the values are constantly nonzero, there might be a memory bottleneck.

2. If paging occurs constantly, verify whether the problem is caused by heavy use of the file system cache. Review the output from the following command:

```
vmstat -I 5
```

3. If the values for `pi` and `po` columns are high, and the values are similar to the values for `fi` and `fo` columns, consider using the `irbw` mount option on all active JFS2 file systems to reduce or eliminate the paging problem.
4. If the paging problem persists after you use the `irbw` mount option for JFS2 file systems, run the AIX **vmo** command to tune the virtual memory system. For more information about using the **vmo** command, see the documentation for the AIX operating system.

Related tasks

[Monitoring performance with operating system tools](#)

Monitor your IBM Storage Protect solution so that you know when you must investigate performance changes. Operating systems have different tools that are available for monitoring performance. Simulating workloads to test performance is another useful task to learn.

[Configuring AIX systems for disk performance](#)

Use JFS2 file systems for IBM Storage Protect in most cases. Examine the queue depth for disks that are in use.

Tuning Linux systems for IBM Storage Protect server performance

There are a number of actions that can improve performance for an IBM Storage Protect server that is running in a Linux environment.

Before you begin

Review system requirements for the installation of an IBM Storage Protect server to ensure that you have the required specifications for your operating system. For more information, see [technote 84861](#).

Procedure

- Most enterprise distributions come with many features, however most of the time only a small subset of these features are used. Disable features that are not used.
- Use the **vm.pagecache_limit_mb** and **vm.swappiness** options to release memory from the file system cache.

Important: The **vm.pagecache_limit_mb** kernel parameter is deprecated for the following versions of Linux operating systems:

- Red Hat® Enterprise Linux 8.1 or later RHEL 8 levels
- Red Hat Enterprise Linux 7.6 or later RHEL 7 levels
- SUSE Linux Enterprise Server 15, SP1 or later SLES 15 levels

For the earlier versions of the RHEL and SUSE Linux systems, set the **vm.pagecache_limit_mb** parameter value to 0.

Linux systems can cache much file system data, which can take away memory that is needed for IBM Storage Protect server and Db2 processes. As a root user on Linux, you can limit the amount of memory that is allowed for caching file data by setting the **vm.pagecache_limit_mb** kernel parameter to 1024. Also, set the **vm.swappiness** kernel parameter to 5. For example:

```
linuxbox:/ # sysctl vm.pagecache_limit_mb          # to display current value
vm.pagecache_limit_mb = 0                          # (0 means no limit)
linuxbox:/ # sysctl -w vm.pagecache_limit_mb=0     # to change at runtime
vm.pagecache_limit_mb = 1024
linuxbox:/ # sysctl vm.swappiness
vm.swappiness = 60
linuxbox:/ # sysctl -w vm.swappiness=0
vm.swappiness = 5
```

To apply these changes for all restarts of the operating system, edit the `/etc/sysctl.conf` file and add `vm.pagecache_limit_mb=1024` and `vm.swappiness=5`.

Tuning Linux on System z systems for IBM Storage Protect server performance

You can use a number of methods to improve the performance of an IBM Storage Protect server on a Linux on System z system.

Procedure

The following steps can help to improve the performance of an IBM Storage Protect server on a Linux on System z system:

- Upgrade to SUSE Linux Enterprise Server 11 Service Pack 1 (SLES11 SP1) for better disk and network performance.
- Use Fibre Channel SCSI disks instead of fiber-connected DASD (direct access storage device), if possible, for greater throughput.
- Dedicate all IBM Storage Protect database, log, and storage disks to the Linux guest if the system is running under z/VM®.
- Use logical volume manager (LVM) for disk storage-pool logical volumes. Using LVM striping improves throughput on operations such as storage pool backup and migration.
- Use the ext4 file system for improved performance for defining storage pool volumes.
- For the IBM Storage Protect database and logs, use either the ext3 or ext4 file system.

As a best practice, use the following file system that is appropriate for your operating system and level:

- For Red Hat Enterprise Linux x86_64, use the ext3 or ext4 file systems. Use the ext4 file system only if Red Hat Enterprise Linux 6.4 or later is installed.
- For SUSE Linux Enterprise Server and for Red Hat Enterprise Linux ppc64, use the ext3 file system.
- Use OSA-Express3 network adapters instead of OSA-Express2 for greater throughput and reduced processor use. See the networking suggestions for OSA-Express3 at [IBM z Systems® - Networking features](#).
- For external network connections to a Linux guest under z/VM, attach the OSA adapter directly to the Linux guest.

Tuning Windows systems for IBM Storage Protect server performance

You can take a number of actions to improve performance for an IBM Storage Protect server that is running in a Windows environment.

Procedure

The following actions can help improve performance:

- Disable NTFS file compression on disk volumes. Because of the potential for performance degradation, do not use NTFS file compression on disk volumes that are used by the IBM Storage Protect server.
- Use the shared memory communication method when you are using a local client. For optimal backup and restore performance when you are using a local client on a Windows system, use the shared memory communication method. The method is used by including the **COMMETHOD** option set to **SHAREDMEM** in the server options file and in the client options file.
- Use the VMXNET 3 network adapter type when the IBM Storage Protect server is in a VMware guest environment. Provide all disks that are used for the server database, log file, and storage as mapped raw LUNs rather than using virtual disks in a VMware datastore.
- Additional actions can affect IBM Storage Protect client and server performance.

- Windows 8 Defender can significantly degrade IBM Storage Protect backup and restore throughput, especially for smaller files. To improve backup and restore performance on Windows 8, which might increase the security risk to the system, use one of the following methods:
 - Disable Windows 8 Defender. Click **Start > Administrative Tools > Computer Management > Services and Applications > Services**. Locate Windows 8 Defender in the list of services. Right-click **Windows Defender** and then select **Properties**. Change the Startup Type attribute to **Disabled**.
 - Without disabling Windows 8 Defender Service, exclude a specific drive that has backup or restore errors. Use this method when there are multiple logical drives on the system. Excluding a drive is less of a security risk than disabling Windows 8 Defender Service.
- Antivirus software can negatively affect backup performance.
- Disable or do not install unused services.
- Disable or do not install unused network protocols.
- Give preference to background application performance.
- Avoid screen savers.
- Ensure that the paging file is not fragmented.
- Ensure that any device driver is current, especially for new hardware.

Secure Sockets Layer (SSL) effects on server performance

Secure Sockets Layer (SSL) provides secure communications between the IBM Storage Protect client and server, but it can affect system performance.

If SSL is needed, use it only for sessions where it is necessary and add processor resources on the IBM Storage Protect server system to handle the increased requirements. Or, try other options, such as networking devices like routers and switches, that provide SSL function instead.

LDAP directory server usage: effects on performance

If you are using an LDAP server to authenticate administrator and node passwords, there might be some performance impacts.

More processor resources are used when you are authenticating with a Lightweight Directory Access Protocol (LDAP) server instead of using local authentication. Tests in IBM labs show that LDAP has about a 5% impact.

If you are using Secure Sockets Layer (SSL) sessions in combination with LDAP server authentication, the additional performance impact for sessions that transfer small amounts of data is negligible. For sessions that transfer large amounts of data, you can expect a significant performance impact because SSL must encrypt all of the data.

Chapter 12. Tuning disk storage for the server

Disk storage systems have different operating characteristics that can be configured and tuned for improved performance for an IBM Storage Protect server.

About this task

Review the information about how to configure your disk storage systems, and the operating system for your IBM Storage Protect server.

Principles for tuning disk systems for IBM Storage Protect

Many aspects of disk storage can be considered for optimizing operations. For most systems, the separation of server database, logs, and storage pools is key to an IBM Storage Protect configuration that has good performance.

The following principles are key to achieving better disk storage performance:

- Select and configure disk storage for both performance and capacity. Adequate capacity is not the only factor to consider.
- For most disk systems, separate the primary IBM Storage Protect server components from each other. Ensure that the server database, active log, archive logs, and storage pools are each in separate locations.
- Monitor the systems. Workloads on systems often increase, and such increases might trigger the need for more storage or for configuration changes. Enforce strict change control to help in troubleshooting any performance degradation.
- Limit mirroring to one type of mirroring. For example, if the operating system is set up to do mirroring, do not configure the IBM Storage Protect server to do mirroring of the active log (**MIRRORLOGDIRECTORY** server option).
- Ensure that server maintenance operations are running, such as expiration and database table and index reorganization. See [“Tuning the schedule for daily operations” on page 149](#).

Understand the entire picture of disk storage and how it relates to operations in your environment. You must examine not only on the configuration of the disks, but the entire configuration as it relates to performance. All of the following items have an impact:

Disk storage system choice and configuration

- Disk type and speed. Larger disks might not be better if the rotational speed is not also higher.
- Disk layout.
- Type of RAID.
- File system type and mount options.

Server system hardware and its configuration

- Speed and number of processors and amount of memory.
- Whether multiple instances of IBM Storage Protect run on the same system and use the same disk storage systems.
- Host bus adapter (HBA) speed.
- Whether HBAs are dedicated to disk operations. An HBA that is shared by disk and tape can have performance problems.
- Whether disks are shared with other systems or applications.

Related reference

[Checklist for server database disks](#)

Use the checklist to verify that the system where the server is installed meets requirements for hardware and software configuration.

Checklist for server recovery log disks

The recovery log for the server consists of the active log, the archive log, and optional logs for mirroring and failover. Use the checklist to verify that disk systems that are being used for the logs have the characteristics and configuration that are key to good performance.

Disk system types

The choice of disk system affects the configuration options that are available. How the disk system is configured affects the resulting performance of the IBM Storage Protect server. For example, disk systems vary by how the individual disk units can be organized to create RAID arrays.

The following list shows the variety of system types that can be used for IBM Storage Protect disk storage:

Solid-state drive (SSD) systems

Solid-state drive technology (sometimes called *flash memory*) provides the highest level of performance, with much higher rates of I/O operations per second (IOPS) than other storage systems. An SSD works at speeds much closer to that of memory than disk drives. SSD has no delay from a spinning platter or from an arm that must move to the correct position.

If you use SSD for IBM Storage Protect storage, ensure that you are using enterprise-level quality SSD.

IBM DS8000® series

The DS8000 series are high-performance disk systems that accept several different RAID types, including RAID 5 and RAID 10. The size of the arrays, in terms of the quantity of disk units, is fixed. Therefore, a DS8000 series unit has a fixed number of RAID arrays (ranks).

IBM DS5000 series

The DS5000 series of midrange disk systems can manage a wide variety of disk configurations. You can create RAID arrays with as few as two and as many as several dozen disks. You can have smaller RAID arrays for the IBM Storage Protect database and larger arrays for the IBM Storage Protect disk storage pool.

IBM Storwize V7000

The Storwize V7000 system is a midrange system that virtualizes RAID storage. A system consists of a set of drive enclosures. You configure the drives into arrays, and create volumes from those arrays. You can configure the system with multiple device types. With multiple device types, you can use fast disk or SSD for the server database and lower-cost, higher-capacity disk for storage pools.

IBM Storwize® V3700 is an entry-level system that has similar characteristics as the Storwize V7000.

Optimization of disk system read-ahead processes

Most advanced disk systems can automatically optimize the performance of read operations when the disk system can detect sequential reads. When the disk system detects sequential reads, it can have the data for the next read in cache, or at least have the read in progress.

Disk systems detect sequential reads on a LUN-by-LUN basis. However, a sequential read might not be detected if more than one read is in progress for the same LUN. Disk systems do not perceive file systems or files within a LUN and distinguish only the blocks that are being accessed. When two sequential reads are in progress on one LUN, the blocks that are being accessed no longer respond as sequential blocks. The blocks appear to come from different places, and the read-ahead optimizations are typically stopped.

Choosing the correct type of storage technology for IBM Storage Protect

Storage devices have different capacity and performance characteristics. These characteristics affect which devices are better for use with IBM Storage Protect.

Procedure

- Review the following table to help you to choose the correct type of storage technology for the storage resources that the server requires.

Table 24. Storage technology types for IBM Storage Protect storage requirements				
Storage technology type	Database	Active log	Archive log and archive failover log	Storage pools
Solid-state disk (SSD)	Place the database on SSD in the following circumstances: <ul style="list-style-type: none">– You are using IBM Storage Protect data deduplication.– You are backing up more than 8 TB of new data daily.	If you place the IBM Storage Protect database on an SSD, as a best practice, place the active log on an SSD. If space is not available, use high-performance disk instead.	Save SSDs for use with the database and active log. The archive log and archive failover logs can be placed on slower storage technology types.	Save SSDs for use with the database and active log. Storage pools can be placed on slower storage technology types.
High-performance disk with the following characteristics: <ul style="list-style-type: none">– 15k rpm disk– Fibre Channel or serial-attached SCSI (SAS) interface	Use high-performance disks in the following circumstances: <ul style="list-style-type: none">– The server does not do data deduplication.– The server does not do data replication. Isolate the server database from its logs and storage pools, and from data for other applications.	Use high-performance disks in the following circumstances: <ul style="list-style-type: none">– The server does not do data deduplication.– The server does not do data replication. For performance and availability, isolate the active log from the server database, archive logs, and storage pools.	You can use high-performance disks for the archive log and archive failover logs. For availability, isolate these logs from the database and active log.	Use high-performance disks for storage pools in the following circumstances: <ul style="list-style-type: none">– Data is frequently read.– Data is frequently written. For performance and availability, isolate storage pool data from the server database and logs, and from data for other applications.

Table 24. Storage technology types for IBM Storage Protect storage requirements (continued)

Storage technology type	Database	Active log	Archive log and archive failover log	Storage pools
Medium-performance or high-performance disk with the following characteristics: <ul style="list-style-type: none"> – 10k rpm disk – Fibre Channel or SAS interface 	If the disk system has a mix of disk technologies, use the faster disks for the database and active log. Isolate the server database from its logs and storage pools, and from data for other applications.	If the disk system has a mix of disk technologies, use the faster disks for the database and active log. For performance and availability, isolate the active log from the server database, archive logs, and storage pools.	You can use medium-performance or high-performance disk for the archive log and archive failover logs. For availability, isolate these logs from the database and active log.	Use medium-performance or high-performance disk for storage pools in the following circumstances: <ul style="list-style-type: none"> – Data is frequently read. – Data is frequently written. For performance and availability, isolate storage pool data from the server database and logs, and from data for other applications.
SATA, network-attached storage	Do not use this storage for the database. Do not place the database on XIV® storage systems.	Do not use this storage for the active log.	Use of this slower storage technology is acceptable because these logs are written once and infrequently read.	Use this slower storage technology in the following circumstances: <ul style="list-style-type: none"> – Data is infrequently written, for example written once. – Data is infrequently read.
Tape and virtual tape				Use for long-term retention or if data is infrequently used.

Tuning System Storage DS8000 series storage systems

IBM System Storage DS8000 series storage systems are designed to be fast and manage heavy I/O.

About this task

A DS8000 series storage system can be configured in RAID 5 and RAID 10 arrays. The quantity of disk units in an array is fixed. Stripe or segment sizes are set automatically. You cannot tune cache for a DS8000 series storage system; for example, you cannot disable cache for the LUNs that are used for the server database. The inability to tune cache is usually not a problem because these systems typically have large amounts of cache.

Procedure

For best performance, follow these guidelines:

- Place the IBM Storage Protect database, active log, archive log, and disk storage pools on separate extent pools.

This arrangement uses more storage, but achieves better performance.

- Distribute the server components over as many ranks as possible. Ranks are on different device adapter pairs.
- Minimize other applications that are sharing the ranks that IBM Storage Protect uses.
- Use as many adapters as possible to access LUNs.
- If you are using a DS8000 series storage system for multiple IBM Storage Protect servers, put all the server databases on one set of ranks. Put all the storage pools for the servers on another set of ranks.
- Test and monitor the results for the disk system configuration in your environment.

Tuning System Storage DS5000 series and other IBM midrange storage systems

IBM System Storage DS5000 series and other IBM midrange storage systems offer great flexibility as to how they can be configured for use with IBM Storage Protect.

About this task

The systems have these characteristics:

- Many types of RAID can be used.
- The number of disks per LUN is flexible.
- Segment or stripe sizes and cache can be set for each LUN.
- Different models have different disk types (Fibre Channel or SATA).
- Different amounts of system cache are available, though typically less than in systems such as System Storage DS8000 series.

Procedure

- For best performance with IBM midrange disk systems, separate the IBM Storage Protect database, recovery log, and storage pools so that they are on different physical spindles. Samples show how to configure these types of disk systems:
 - [“Sample layouts for a server database on DS5000 series disks” on page 182](#)
 - [“Sample layouts for server recovery logs on DS5000 series disks” on page 184](#)
 - [“Sample layout of server storage pools on DS5000 series disks” on page 185](#)

Disk I/O characteristics for IBM Storage Protect operations

Generally, IBM Storage Protect storage pools are written to and read from by using I/O sizes of 256 KB.

With page sizes of 8 KB and 32 KB, the I/O that is used by the IBM Storage Protect database varies. The database manager might, at times, prefetch larger amounts.

While IBM Storage Protect might request the 8 KB and 32 KB I/O sizes, the operating system might choose to run I/O differently. Running the I/O differently can result in smaller or larger I/O being sent to the disk system.

IBM Storage Protect tries to use direct I/O in most situations, which avoids the file system cache. The result of avoiding the cache is better processor efficiency and performance. If you use operating system parameters to tune the file system cache, you might notice no effect for this reason.

Sample layouts for a server database on DS5000 series disks

Samples illustrate several ways to follow configuration guidelines for the server database when you use DS5000 series Fibre Channel disks. The samples illustrate advantages and disadvantages of configuration choices.

Remember: Ensure that the server database, recovery logs, and storage pools are on different disks.

Sample 1: Good layout for a small server

By using five disks for the database, you can set up the disks with the following characteristics. See [Figure 25 on page 182](#).

- Configure disks in a 4 + 1 RAID 5 array.
- Set the stripe size to 256 KB.
- Define one directory (also called a container) and one logical volume for the database.
- Set the **DB2_Parallel_IO** environment variable:

```
DB2_Parallel_IO=*:4
```

The IBM Db2 program that is the server's database manager uses this value when it balances load across the disks.

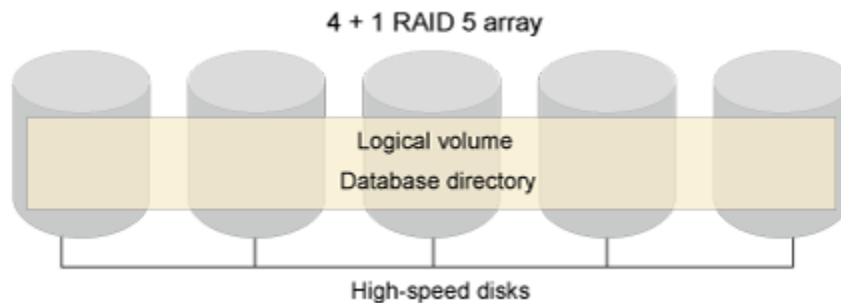


Figure 25. Database layout for a small server

This layout has the following advantages:

- The layout follows the guidelines for optimal stripe size.
- The layout follows the guidelines for having a one-to-one correspondence between logical volumes and containers.

The layout has the following disadvantages:

- Write operations for parity bits can affect performance if the system has a small cache.
- The layout has only one container for the database, which is typically not optimal but might not be a problem for small server workloads.
- Database and database indexes are spread across only five disks.

Sample 2: Better layout that uses RAID 10

By using eight disks for the database, you can set up the disks with the following characteristics. See [Figure 26 on page 183](#).

- Configure disks in a 4 + 4 RAID 10 array.
- Set the stripe size to 256 KB.
- Define one directory (also called a container) and one logical volume for the database.
- Set the **DB2_Parallel_IO** environment variable:

```
DB2_Parallel_IO=*:4
```

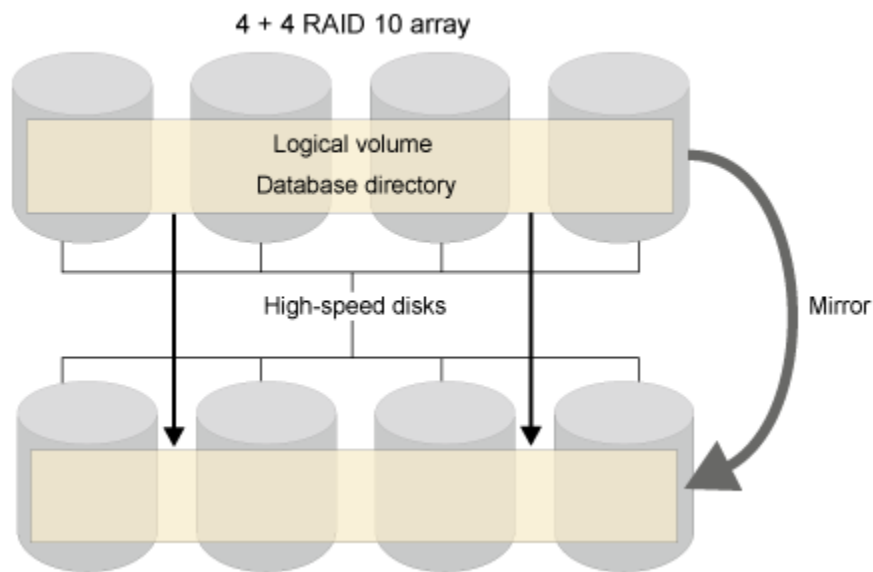


Figure 26. Database layout that uses eight disks in a RAID 10 array

This layout has the following advantages:

- The layout follows the guidelines for optimal stripe size.
- The layout follows the guidelines for having a one-to-one correspondence between logical volumes and containers.
- The system does not have the performance cost of parity write operations.
- RAID 10 is an optimal choice for the server database. With its two sets of disks, this configuration results in faster reads for the database.

The layout has the following disadvantages:

- The layout has only one container for the database, which is typically not optimal but might not be a problem for small server workloads.
- The layout requires twice as many disks as in sample 1 because of the use of RAID 10 instead of RAID 5.

Sample 3: Better layout that uses RAID 10 and more containers

By using 16 disks for the database, you can set up the disks with the following characteristics. See [Figure 27 on page 184](#).

- Configure the disks in two 4 + 4 RAID 10 arrays.
- Set the stripe size to 256 KB.
- Define two directories (also called containers) and two logical volumes for the database.
- Set the **DB2_Parallel_IO** environment variable:

```
DB2_Parallel_IO=*:4
```

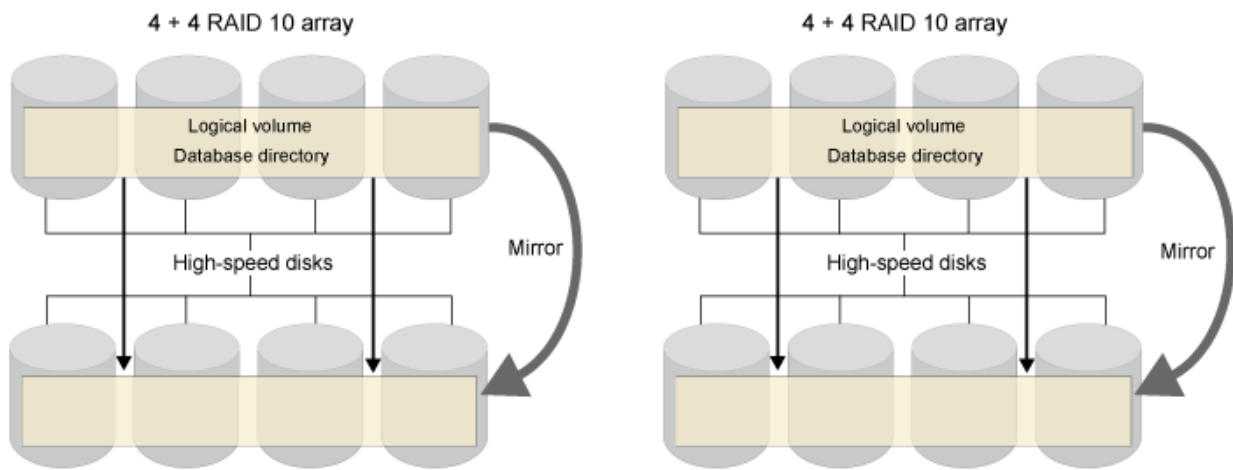


Figure 27. Database layout that uses 16 disks in two RAID 10 arrays

This layout has the following advantages:

- The layout follows the guidelines for optimal stripe size.
- The layout follows the guidelines for having a one-to-one correspondence between logical volumes and containers.
- The system does not have the performance cost of parity write operations.
- RAID 10 is an optimal choice for the server database. With its two sets of disks, this configuration results in faster reads for the database.
- Many more physical spindles means better performance for database read and write operations.
- More database containers means that more data can be prefetched by the Db2 software.

The disadvantage is that this layout requires more disks than the other samples.

Sample layouts for server recovery logs on DS5000 series disks

Samples illustrate several ways to follow configuration guidelines for server recovery logs when you use DS5000 series Fibre Channel disks. The active log is placed on disk with the fastest speed characteristics.

Remember: Ensure that the server database, recovery logs, and storage pools are on different disks.

Sample 1: Good layout without RAID

In this sample, disks that are not configured as a RAID array are used for the recovery logs. This type of arrangement is called *just a bunch of disks* (JBOD). The active log, the archive log, and the failover archive log are placed on separate disks.

The layout has the following advantages:

- Cache read-ahead is used for the disks.
- The separation of the active log, the archive log, and the failover archive log follows guidelines.
- The active log is placed on the fastest disk.

The layout has the following disadvantages:

- This layout has a single point of failure. For example, if the disk for the active log failed, you have no RAID to help you recover.
- All files for the active log are on one disk, which might be slower than if you spread out the files.

Sample 2: Better layout with RAID 1

In this sample, RAID 1 disks are used for the active log and the archive log. This sample has the following features:

- RAID 1 acts as a mirror for the active log. As an alternative, you can use the IBM Storage Protect server option for mirroring the active log, **MIRRORLOGDIRECTORY**.
- RAID 1 is used for the archive log.
- The disk for the archive failover log is not RAID 1 because it is not as critical to server operations as the other logs.

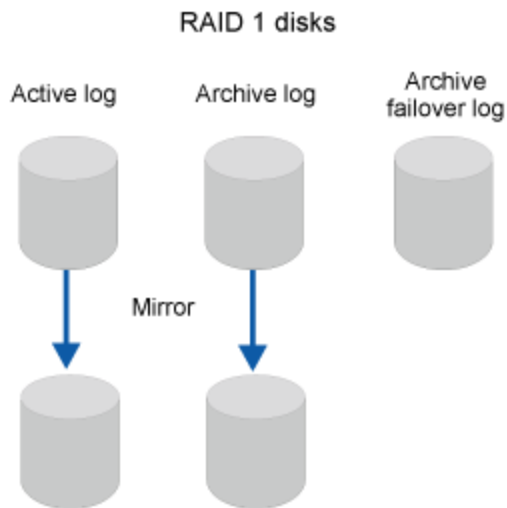


Figure 28. Logs layout with RAID 1

The layout has the following advantages:

- Fast Fibre Channel disks are used.
- Cache read-ahead is used for the disks.
- RAID 1 is used for the active log and archive log locations, which removes the single point of failure for those logs.

Performance might be slower because all of the files for the active log are on one disk.

Sample layout of server storage pools on DS5000 series disks

A sample illustrates how to follow configuration guidelines for storage pools that use the DISK device class and DS5000 series disks.

Remember: Ensure that the server database, recovery logs, and storage pools are on different disks.

Sample layout

In this sample, the Fibre Channel or Serial Advanced Technology Attachment (SATA) disks in a DS5000 series system are configured with these characteristics:

- The disks are configured in a 4 + 1 RAID 5 array. Stripe size is 256 KB.
- Four logical volumes are defined on the disks. In IBM Storage Protect, these volumes are defined as four storage pool volumes for a random-access (DISK) storage pool.

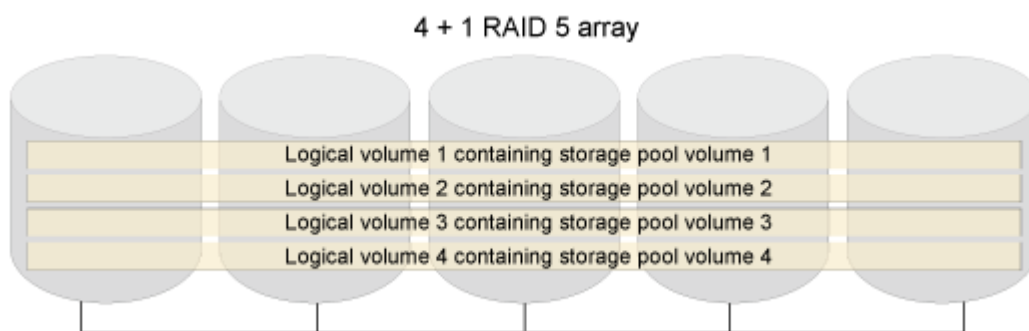


Figure 29. Storage pool layout with 4 + 1 RAID 5

The layout has the following advantages:

- The layout follows the stripe size recommendation (full stripe write).
- The layout follows the guideline that an $n + 1$ RAID array has no more than n volumes.

The disadvantage is that there is a performance cost for the write-parity operations. The write-parity operations might not cause a performance problem if the cache for the disk system is adequate.

Tuning Storwize V7000 and V3700 systems

The IBM Storwize V7000 is an ideal disk system for use with IBM Storage Protect. You can configure the system with multiple device types, so that you can use fast disk or SSD for the server database and lower-cost, higher-capacity disk for storage pools.

About this task

The Storwize V7000 includes the Easy Tier® function, which automatically and nondisruptively moves frequently accessed data from HDD to SSD disks. With this function, you can configure the system to gain most of the benefits of SSD speed for the server database without having the entire database on SSD.

The IBM Storwize V3700 is an entry level system with capabilities that are similar to the Storwize V7000. The Storwize V3700 is also a good disk system to use with IBM Storage Protect.

Example configuration with Storwize V7000 systems

The example illustrates how you can configure IBM Storwize V7000 systems for an IBM Storage Protect server. In this example, the server needs 2 TB for its database and 58 TB for its storage pool on disk.

The following components are used in the configuration:

- One Storwize V7000 Disk Control Enclosure, which can hold 24 2.5-inch disk units
- Three Storwize V7000 Disk Expansion Enclosures, each of which can hold 24 2.5-inch disk units

Together, these enclosures can hold up to 96 disk units. See the following table for the disk specifications and configuration. All disks are 2.5-inch disks.

Server storage requirement	Disk capacity	Disk speed and type	Disk quantity	RAID type	RAID array quantity	Usage
Database	300 GB	15k rpm SAS HDD	12	4 + P RAID 5	2	<p>These disks provide space for the database on two 4 + P RAID 5 arrays, with two spare disks.</p> <p>For optimal performance, use the Easy Tier function with SSD disk units for the database. Use the Easy Tier function for the database only, not for the logs or the storage pool.</p> <p>If you are using data deduplication, data replication, or both, either substitute SSD disk units for the database or add some SSD units and use the Easy Tier function.</p>
Active log and archive log	300 GB	15k rpm SAS HDD	4	RAID 0	2	<p>These disks provide space for the archive log and active log on two RAID 0 arrays. Spare disks are shared with the disks for the database.</p>
Storage pool	900 GB	10k rpm SAS HDD	80	6 + P RAID 5	11	<p>These disks provide space for a 58 TB storage pool. The disks are configured in eleven 6 + P RAID 5 arrays, with three spare disks.</p> <p>If the throughput for the storage pool does not require the higher speed of the 10k rpm drives, you can use 7.2k rpm nearline-SAS HDD drives. Verify that the slower disks can meet throughput requirements.</p> <p>If greater storage pool capacity is required, add more enclosures.</p>

Configuring the operating system for disk performance

The operating system configuration and the types of file systems in use affect the performance of disks. Learn about how to configure these items for the best performance for IBM Storage Protect server.

About this task

Parameters that affect disk performance vary by operating system.

Configuring AIX systems for disk performance

Use JFS2 file systems for IBM Storage Protect in most cases. Examine the queue depth for disks that are in use.

Procedure

- Use JFS2 file systems for the IBM Storage Protect database, recovery logs, and disk storage pools, unless you are using disk for LAN-free operations. If you are using disk for LAN-free operations, then use General Parallel File System (GPFS) for the shared storage pools.

Use the JFS2 **rbw** option on the **mount** command, especially if IBM Storage Protect database backups are stored to a FILE device class.

- The default queue depths for non-IBM disks that are used for IBM Storage Protect are often low by default. If the queue depth is less than 32, see the documentation for the disk system or contact the manufacturer about guidelines for the queue depth. To change the queue depth, see the [AIX product information](#).

Configuring Linux systems for disk performance

When you configure disks for IBM Storage Protect, use the Logical Volume Manager (LVM).

Procedure

- Use Linux Logical Volume Manager (LVM) to create logical volumes on the disk LUNs for all disks that are used for IBM Storage Protect components.

Set the LVM read-ahead to 0 for all logical volumes on disk systems that provide adaptive read-ahead capabilities, for example, enterprise-type disk systems.

If more space is needed, the logical volumes provide an easy way to extend the volumes and file systems. LVM also provides striping, which can be used to improve sequential I/O performance.

- For the IBM Storage Protect database and logs, use either the ext3 or ext4 file system.
As a best practice, use the following file system that is appropriate for your operating system and level:
 - For Red Hat Enterprise Linux x86_64, use the ext3 or ext4 file system. Use the ext4 file system only if Red Hat Enterprise Linux 6.4 or later is installed.
 - For SUSE Linux Enterprise Server and for Red Hat Enterprise Linux ppc64, use the ext3 file system.
- For IBM Storage Protect storage pools, use the ext4 file system.

The ext4 file system has the following advantages for use with storage pools:

- You do not have to write out each block I/O to allocate the storage pool volume, which improves the performance of the **DEFINE VOLUME** command.
- You can avoid file and free space fragmentation, which improves read and write performance.
- When you define new volumes, the IBM Storage Protect server activities that are running are not negatively affected.

Chapter 13. Tuning client performance

You can optimize the performance of IBM Storage Protect clients. Review the methods for backing up data and select the best methods for your environment. Review the information about client options and adjust configuration settings as needed.

Selecting the optimal client backup method

You can use several techniques with the backup-archive client to help you ensure the best performance during various types of backup processing.

About this task

For most situations, incremental backup is the most comprehensive backup method to use. It is the best method for detecting file changes in your local system, and provides the capability to run individual file restores.

However, there are two major factors that can prevent you from completing incremental backups:

- The amount of available memory
- The backup window duration

Incremental backups can be memory intensive because the memory for an incremental backup holds the list of files that are known to the IBM Storage Protect server. Therefore, the memory that is required for an incremental backup is proportional to the number of files in the file system that are being backed up. If the system does not have enough memory, the incremental backups can fail. The time that it takes the client to scan the file system and the amount of changed data can also contribute to backup processing that does not complete within the scheduled backup window. Memory-efficient and journal-based backups do not maintain the whole file list in memory.

Use the following general guidelines to resolve memory and backup window issues.

Guideline	More information
1. Resolve any memory issues first. You must resolve memory issues before you can resolve a backup window issue. For example, you can use the <code>memoryefficientbackup yes</code> or <code>memoryefficient diskcachemethod</code> options to reduce the memory requirements for backing up file systems that contain millions of files.	See “Reduce client memory use” on page 206
2. Resolve any backup window issues. For example, if the number of files that change daily is relatively small, you can use journal-based backup.	See “Tuning journal-based backups” on page 221
3. If you cannot otherwise resolve memory issues, consider using image backups.	See “Image backup techniques” on page 196

Deciding which backup method to use

Many backup techniques are available in the backup-archive client. Begin with progressive incremental backups and move to other types of incremental backups or image backups if necessary.

Procedure

- Use the following table to determine the backup method that you want to use. The table contains common backup scenarios that you might encounter and the suggested backup methods to use.

Scenario	Use this backup method
I want to run the most comprehensive type of file backup on my client system.	<ul style="list-style-type: none"> – “Progressive incremental backup” on page 191 – “Incremental-by-date backup” on page 194
I want to take advantage of the benefits of progressive incremental backups, but I am running into memory issues.	<ul style="list-style-type: none"> – “Memory-efficient backup” on page 192 – “Memory-efficient backup with disk caching” on page 193
I have many small files in my file system with a few changes, but incremental backup processing does not complete within the allotted time.	“Journal-based backup” on page 191
On operating systems such as AIX or Linux, I have large file systems that can be divided into logical partitions. I want to be able to provide a direct path to the files that I want to back up.	“Backup of virtual mount points” on page 193
Scheduled backups do not complete within the allotted time. I have an application that can provide a list of files that changed since the previous backup. I want to improve the speed of the backup process by backing up just this list of changed files.	“File list backup” on page 194
I want to increase the throughput of incremental backup operations and restore operations.	“Multiple session backup” on page 195
I tried to use several types of incremental backups but scheduled backup processing does not complete within the allotted time.	<ul style="list-style-type: none"> – “Image backup” on page 196 – “Image plus incremental backup” on page 198
I want to streamline the backups of the volumes on my NetApp vFiler.	“Snapshot differential backup” on page 199
I want to back up my VMware virtual machines with my IBM Storage Protect for Virtual Environments software.	<ul style="list-style-type: none"> – “Virtual machine incremental-forever-full backup” on page 200 – “Virtual machine incremental-forever-incremental backup” on page 200
I want to back up my Microsoft Hyper-V virtual machine with IBM Storage Protect for Virtual Environments software.	For more information, see Back up virtual machines on a Hyper-V system .
I want to improve the speed of my VMware virtual machine backups by backing up several virtual machines at the same time.	For more information, see Parallel back ups of virtual machines .

File backup techniques

If you are backing up your system on a file-by-file basis, you can use several backup techniques.

Use the following information to determine which file backup technique to use that best meets your needs.

Progressive incremental backup

Progressive incremental backup is the standard method of backup that is used by IBM Storage Protect. Incremental backup processing backs up only those files that changed since the last full or incremental backup, unless the files are excluded from backup.

How it works

The following processes occur during an incremental backup:

- The client queries the IBM Storage Protect server for active backup version metadata.
- The server returns a list of active backup versions for the entire file system.
- The client scans and compares the list with the local file system to determine which files are new or changed since the last backup.
- The client backs up the new and changed files.

When to use

Use incremental backup when the system is not constrained by memory, backup window duration, or other operational issues. Incremental backup is the default backup method.

Advantages

Incremental backup processing has the following advantages:

- This method is the most comprehensive backup method for IBM Storage Protect.
- No redundant backups are made. You back up only what is changed.
- There is less network utilization because unchanged files do not have to be sent over the network.
- This method is a form of single-instance storage because a file is not backed up again if it does not change. Incremental backups are more efficient and save space on the server storage pools.
- Files are easier to restore because you do not have to restore the base backup version first and apply incremental or differential changes.

Disadvantages

Incremental backups processing has the following disadvantages:

- The client system might run out of memory if the number of active backup versions is too large.
- The time that it takes to scan file systems that contain millions of files can exceed the duration of the backup window.

If incremental backup operations do not complete successfully, try other variations of incremental backup:

- [“Journal-based backup” on page 191](#)
- [“Memory-efficient backup” on page 192](#)
- [“Memory-efficient backup with disk caching” on page 193](#)
- [“Backup of virtual mount points” on page 193](#)
- [“Incremental-by-date backup” on page 194](#)
- [“File list backup” on page 194](#)
- [“Multiple session backup” on page 195](#)

Journal-based backup

Journal-based backup is an alternative form of incremental backup that uses a change journal that is maintained by the IBM Storage Protect journal process. On Windows clients, the change journal is maintained by a journal service. On AIX and Linux clients, the change journal is maintained by a journal daemon process.

How it works

The following processes occur during journal-based backup processing:

- Journal-based backup processing uses real-time monitoring of a file system for changed files.

- The names of the changed files are logged to the journal database.
- During backup processing, the client queries the journal for the list of changed files, and then backs up the changed files.

When to use

Use journal-based backup in the following situations:

- The scheduled backups are not completed within the allotted time.
- There are less than 1,000,000 files and a small number of changes between backups (less than 1,000,000).
- There are less than 10,000,000 objects with 10-15% velocity of change. The velocity of change means the rate at which files are changed over a short amount of time (such 1 or 2 seconds).

Advantages

Journal-based backup can often greatly reduce the time that it takes to determine which files changed.

Disadvantages

Journal-based backup processing has the following limitations:

- You must still run incremental backups periodically.
- Journal-based backups are not suitable for file systems where large numbers of files can change over a short time interval, such as changing hundreds or thousands of files in 1 or 2 seconds.
- This method is available only on Windows, AIX, and Linux clients.

Related task: [“Tuning journal-based backups” on page 221](#)

Memory-efficient backup

The performance of incremental backups can degrade if the system is memory-constrained before the backup begins. Run incremental backup with the **memoryefficientbackup** yes option in the client options file. This setting causes the client to process only one directory at a time during incremental backups, which reduces memory consumption but increases backup time.

How it works

The following processes occur during an incremental backup with the memory-efficient setting:

- The client queries the server for the metadata of active backup versions for the first directory to be backed up.
- The server returns a list of active backup versions for the directory.
- The client scans the list and compares it with the local file system, and backs up the new and changed files.
- The client queries the server for the next directory and repeats the process for all directories.

When to use

Use memory-efficient backup when your system has a low amount of memory available for incremental backups.

Advantages

Memory-efficient backup is a comprehensive incremental backup with a smaller backup memory footprint.

Disadvantages

Memory-efficient backup processing has the following disadvantages:

- The backup run time is increased.
- This method does not work for a single directory that contains a large number of files.
- If the system is not memory-constrained, running memory-efficient backup can degrade the backup performance.

Related task: [“Reduce client memory use” on page 206](#)

Memory-efficient backup with disk caching

If your client system is memory-constrained and incremental backups still cannot complete successfully with the **memoryefficientbackup** yes setting, run incremental backups with the **memoryefficientbackup** diskcachemethod option. This setting causes the client to use less memory but requires more disk space on the client system.

How it works

This method is similar to incremental backup processing but the client temporarily stores active backup version metadata on disk instead of memory.

When to use

Use memory-efficient backup with disk caching in the following situations:

- The client is running out of memory with incremental backups and memory-efficient backup is not sufficient.
- Journal-based backup is not available on the operating system.

Advantages

Memory-efficient backup with disk caching is a comprehensive incremental backup operation with a smaller backup memory footprint.

Disadvantages

Memory-efficient backup processing with disk caching has the following disadvantages:

- The backup processing time might be longer because the active backup inventory is on disk instead of in memory.
- Gigabytes of free disk space are required to temporarily cache the active backup inventory.

Related task: [“Reduce client memory use” on page 206](#)

Backup of virtual mount points

You can save processing time when you define a virtual mount point within a file system because it provides a direct path to the files that you want to back up.

How it works

The following processes occur during the backup of virtual mount points:

- Instead of backing up an entire file system to a single file space on the server, you can logically partition a large file system into smaller file systems, and then define mount points for backup processing.
- The file system that are represented by the mount points can be managed as separate file spaces on the server.

When to use

Use virtual mount points to back up large, balanced, AIX, Linux, and Solaris file systems that can be efficiently divided into logical partitions.

Advantages

Backup processing of virtual mount points provides a balanced approach to the backup of large file systems by effectively dividing them into smaller file systems. It is more efficient than defining the file system with the **domain** option, and then specifying the **exclude** option to exclude the files you do not want to back up.

Disadvantages

Backup processing of virtual mount points has the following limitations:

- This method of backup processing is not appropriate for a single directory that contains a large number of files.
- Virtual mount points are static and cannot be changed.

- This method requires monitoring to ensure that new directories are still backed up in one of the virtual mount points, along with other processing that is required to maintain the virtual mount point definitions.
- Command-line restore operations require the use of braces ({ }) to delimit the virtual mount point name in the file specification.
- This method is only available for AIX and Linux operating systems.

Related concept: [“File space tuning” on page 228](#)

Incremental-by-date backup

This backup method backs up new and changed files that have a modification date later than the date of the last incremental backup that is stored at the server, unless the files are excluded from backup.

How it works

The following processes occur during an incremental-by-date backup:

- The client queries the server for the most recent backup of the entire file system.
- The server returns the time stamp of the most recent backup of the entire file system.
- The client scans and compares the list from the server with the local file system and backs up the new and changed files that are based on the time stamp of the most recent backup.

When to use

Use incremental-by-date backup in the following situations:

- The scheduled backups are not completed within the allotted time.
- The changes to the file system are additive or changing, but not deleted.
- You also run weekly (or periodic) full incremental backups.

Advantages

Incremental-by-date backup processing has the following benefits:

- This method reduces the time that it takes to determine which files changed.
- This method removes the processing time on the server that is used to query the database for changed files.
- This method removes the network traffic that is used to communicate the query results.

Disadvantages

Incremental-by-date backup processing has the following disadvantages:

- This method reduces the flexibility of the scope of the backup operation. You must back up the entire file system.
- The files are not backed up if the changes do not affect the date (for example, attribute, mode, ACL, rename, copy, move, and security changes).
- The deleted files are not expired on the server.
- Policy rebinding does not take place.
- The entire file system must be scanned.
- This method cannot be used if the client and server clocks are set to different times or are not in the same time zone.

Related information: [“Incremental-by-date backup” on page 194](#)

File list backup

You can control which files are backed when run a backup with the **filelist** option.

How it works

File list backup can be used in the following manner:

- An application creates a list of files for backup and passes the list to the client.
- The client runs a selective backup of the files that are specified in the list.

When to use

Use file list backup in the following situations:

- The scheduled backups are not completing within the allotted time.
- The list of changed files is known.

Advantages

Selective backup eliminates the query of the server database and the scan of local file system.

Disadvantages

File list backup has the following disadvantages:

- You must find a way to create the file list.
- You must explicitly specify the files. You cannot use wildcard characters or directory recursion in the file list.
- Large file lists have an affect on memory requirements during client restore and retrieve operations.

Related reference: [“File list backup” on page 194](#)

Multiple session backup

The backup-archive client can run concurrent sessions to back up and restore data to help improve performance. During incremental backup processing, the client can process multiple objects in parallel by opening more than one session with the IBM Storage Protect server.

How it works

Multiple sessions are used when you specify multiple file specifications on a **backup**, **restore**, **archive**, or **retrieve** command. For example, you can start a multiple session backup with the following command:

- On the AIX, Linux, Mac OS X, or Solaris client:

```
incr /Volumes/filespace_A /Volumes/filespace_B
```

- On the Windows client:

```
incr c: d:
```

The **resourceutilization** option is used to regulate the level of resources that the IBM Storage Protect server and client can use during processing. The default is to use a maximum of two sessions, one session to query the server and one session to send file data.

When to use

Use multiple backup sessions when you want to increase client performance, and you have sufficient client and server resources and processing capacity. For example, the server and client hardware must have sufficient memory, storage, and processor capacity to support multiple sessions. The network bandwidth must also be sufficient to handle the increased amount of data that flows across the network.

Advantages

Using more than one backup session can often lead to overall improvements in throughput.

Disadvantages

Running multiple backup sessions has the following disadvantages. Some workarounds are included.

- During a multiple session backup operation, files from one file specification might be stored on multiple tapes on the server and interspersed with files from different file specifications. This arrangement can decrease restore performance.

To avoid the performance degradation in restore operations, set the **collocatebyfilespec** option to yes. This setting eliminates the interspersing of files from different file specifications

by limiting the client to one server session for each file specification. Therefore, if the data is stored to tape, the files for each file specification are stored together on one tape, unless another tape is required for more capacity.

- The client might produce multiple accounting records.
- The server might not start enough concurrent sessions. To avoid this situation, the **maxsessions** server parameter must be reviewed and possibly changed.
- A **query node** command might not summarize the client activity.

Related task: [“Improving client performance by using multiple sessions” on page 217](#)

Image backup techniques

If the different variations of progressive incremental backup and file backup operations do not complete successfully, consider running an image backup to reduce the backup window.

Image backup

Image backup processing backs up your file system as a single object. All data in the file system is backed up even if the amount of data that changed since the last backup operation is small.

How it works

During image backup processing, the client sends a logical block image of a file system to the IBM Storage Protect server.

When to use

Use image backup processing in the following situations:

- You cannot resolve system memory problems or progressive incremental backup is otherwise unusable.
- There are too many changes in the file system (greater than 1,000,000 objects) for journal-based backup.
- Most of the file system contains small files (average size less than 1 MB).
- You must have a faster recovery time than what can be achieved with file-level restore.
- For AIX, Linux, and Solaris clients:
 - The file system is at least 60% full.
 - Online image backup is not available, and you can unmount the file system.

Advantages

Image backup processing has the following benefits:

- Backups are faster.
- No scan time is required to determine what changed.
- Overall data movement is faster.
- Restore times are faster.

Disadvantages

You cannot restore individual files directly from the IBM Storage Protect server.

The following variations of image backup are available:

- Offline (static) image backup processing
 - The volumes to be backed up are mounted read-only.
 - This method is available for AIX, Linux x86, Solaris, and Windows operating systems.
 - This method is the most effective backup method for FlashCopy® operations.
- Online (dynamic) image backup processing
 - The volumes to be backed up remain online.

- Fuzzy backup processing occurs when the data is changed during the image backup processing.
- Online image backup by using snapshots
 - The volumes to be backed up remain online.
 - The image backup is made at a single point in time.
 - It is available only for AIX JFS2, Linux x86, and Windows operating systems.

Image plus incremental-by-date image backup

Image backup plus incremental-by-date image backup processing is one of two methods that you can use to run efficient incremental backups of your file system.

How it works

The following processes occur during image plus incremental-by-date image backup processing:

- During a full image backup (for example, when you issue the `dsmc backup image` command), the client sends a logical block image of a file system to the server.
- Subsequent backups are incremental-by-date image backups (for example, when you issue the `dsmc backup image -mode=incremental` command), in which the client queries the server for the last backup of the entire file system.
- The server sends the time stamp of last backup of the entire file system to the client.
- The client scans and compares the time stamp with the local file system, and backs up the new and changed files.

During an image plus incremental-by-date restore operation, the following processes occur:

- The client requests an incremental image restore.
- The server sends the base image to the client.
- The server returns more files that must be applied to the base image to satisfy the recovery point.

When to use

Run image plus incremental-by-date image backup processing in the following situations:

- You need faster backups.
- You must be able to restore files to a specific point in time.

Tip: Periodically run full image backups to maintain a file system image that is close to what existed at the time of the last incremental-by-date image backup. When you periodically run a full image backup, it can also improve restore time.

Advantages

Image plus incremental backup processing has the following advantages:

- Backups are faster.
- No scan time is required to determine what changed.
- Overall data movement is faster.
- Restore times are faster.
- Protection of files that changed after the image backup was created.
- In some cases, recovery time and recovery point objectives are improved.

Disadvantages

Image plus incremental-by-date image backup processing has the following limitations:

- This method reduces the flexibility of the scope of the backup operation. You must back up the entire file system.
- The files are not backed up if the changes do not affect the date (for example, attribute, mode, ACL, rename, copy, move, and security changes).
- The deleted files are not expired on the server.

- Policy rebinding does not take place.
- The entire file system must be scanned.
- This method cannot be used if the client and server clocks are set to different times or are not in the same time zone.
- Deleted files are not reconciled. Deleted files are not expired on the server. Therefore, when you restore an image with the incremental option, files that were deleted after the original image backup are present after the restore.
- More storage space is required on the IBM Storage Protect server.

Image plus incremental backup

Image backup plus file system incremental backup processing is the second method that you can use to run efficient incremental backups of your file system.

How it works

The following processes during image plus incremental backup processing:

- During a full image backup (for example, when you issue the `dsmc backup image` command), the client sends a logical block image of a file system to the server.
- Subsequent backups are progressive incremental backups in which the client queries server for the active backup version metadata.
- The server returns list of active backup versions for the entire file system.
- The client scans and compares the list with the local file system.
- The client backs up the new and changed files.

During an image plus progressive incremental restore operation, the following processes occur:

- The client requests an incremental image restore.
- The server returns the base image.
- The server returns more files that must be applied to the base image to satisfy the recovery point.
- The server optionally returns the list of files that must be deleted from the base image.

When to use

Use image plus incremental backup processing in the following situations:

- You need faster backups.
- You want to restore files to a specific point in time.
- You want to be able to reconcile deleted files.

Tips:

- Run incremental backups of the file system periodically to ensure that the server records additions and deletions accurately.
- Run an image backup periodically to ensure faster restores.

Advantages

Image plus incremental backup processing has the following advantages:

- Backups are faster.
- No scan time is required to determine what changed.
- Overall data movement is faster.
- Restore times are faster.
- Protection of files that changed after the image backup was created.
- In some cases, recovery time and recovery point objectives are improved.

Disadvantages

Image plus incremental backup processing has the following disadvantages:

- More time is required to periodically create image backups.
- More storage space is required on the IBM Storage Protect server.

Snapshot differential backup

If you are backing up NetApp filer or vFiler volumes or N-Series file server volumes, you can use a snapshot differential backup to streamline the incremental backup process.

How it works

The following processes occur during snapshot differential backup processing:

- The first time that you run an incremental backup with the **snaptiff** option, a snapshot is created (the base snapshot) and a traditional incremental backup is run by using this snapshot as the source. The name of the snapshot that is created is recorded in the IBM Storage Protect database.
- The second time an incremental backup is run with the **snaptiff** option, a newer snapshot is either created, or an existing one is used to find the differences between these two snapshots. The second snapshot is called **diffsnapshot**. The client then incrementally backs up the files that are reported as changed by NetApp to the IBM Storage Protect server.

When to use

Use this method to back up NetApp filer or vFiler volumes or N-Series file server volumes on Windows, AIX 64-bit, and Linux x86/86_64 clients.

Advantages

Snapshot differential backup processing can save you time by not having to scan the whole volume for changed files.

Disadvantages

Snapshot differential backup processing has the following limitations:

- On Windows systems, it does not work for any of the NetApp predefined shares, including C\$, because the client cannot determine their mount points programmatically.
- You must periodically take a new base snapshot with the **createnewbase** option to back up any files that might have been skipped.

Virtual machine backup techniques

You can use several backup techniques to back up virtual machines.

Use the following information to determine which virtual machine backup technique to use that best meets your needs and environment.

Windows Hyper-V backup

You can use IBM Storage Protect for Virtual Environments: Data Protection for Microsoft Hyper-V to back up virtual machines that are hosted by a Windows Server operating systems with the Hyper-V role installed.

How it works

For Windows Server 2012, Windows Server 2012 R2, and Windows Server 2016 or later operating systems, Data Protection for Microsoft Hyper-V creates an incremental-forever full or incremental-forever incremental backup of Hyper-V virtual machines in VHDX disk format. The virtual machine is backed up to the IBM Storage Protect server. Microsoft Volume Shadow Copy Service (VSS) is used to take a consistent snapshot of the virtual machine. On Windows Server 2016 or later systems, Resilient® Change Tracking (RCT) is used to create snapshots and provide change block tracking capability for Hyper-V VM disks.

Advantages

Hyper-V backup processing has the following advantages:

- This method can back up data without stopping the virtual machine or stopping any running applications within the virtual machine.
- This method can restore either individual virtual machines or a group of virtual machines that run on a Hyper-V server for disaster recovery purposes.
- This method adds backup and restore functions for guest operating systems without the need for you to install an IBM Storage Protect client on the guest virtual machine.
- This method can be used for both disaster recovery and long-term data backup support.
- You can run individual file restore operations by using the IBM Storage Protect file restore interface in Data Protection for Microsoft Hyper-V 8.1.4 or later.

Disadvantages

Hyper-V backup processing has the following disadvantages:

- Backups are not granular.
- You cannot run individual file restore operations from a full virtual machine backup (in Data Protection for Microsoft Hyper-V 8.1.2 or earlier).

For complete information about creating incremental-forever backups of Hyper-V virtual machines, or for information about using the IBM Storage Protect file restore interface to restore files from a Hyper-V virtual machine backup, see the online product information at <https://www.ibm.com/docs/en/spfve>.

Virtual machine incremental-forever-full backup

Incremental-forever-full virtual machine backup processing backs up all the used blocks on a virtual machine's disks. To run this type of backup, you must have a license for one of the following products:

- IBM Storage Protect for Virtual Environments: Data Protection for VMware.
- IBM Storage Protect for Virtual Environments: Data Protection for Microsoft Hyper-V on Windows Server 2012 or later Windows Server operating systems with the Hyper-V role installed.

How it works

The following processes occur during incremental-forever-full virtual machine backup processing:

- A full virtual machine backup is required only one time.
- Data from incremental backups is combined with data from the full backup to create a synthetic full backup image. This type of full backup is called a synthetic backup because it is created from the data that is stored on the server and not from reading the used blocks on the production disks.
- Each incremental-forever-full virtual machine backup operation reads and copies all of the used blocks, whether the blocks are changed or not since the previous backup.

Advantages

Incremental-forever-full virtual machine backup processing has the following advantages:

- During a restore operation, you can specify options for a point in time and date to recover data. The data is restored from the original full backup and all of the changed blocks that are associated with the data.

Disadvantages

Incremental-forever-full virtual machine backup processing has the following disadvantages:

- If one or more of the progressive incremental backups is corrupted on the server, you might not be able to fully recover a virtual machine. To ensure that you can fully recover a virtual machine, periodically run a full virtual machine backup.
- This method is available only with data movers that run on Linux and Windows clients.

Virtual machine incremental-forever-incremental backup

Incremental-forever-incremental backup processing backs up only the disk blocks that have changed since the last backup. To run this type of backup, you must have a license to use one of the following products:

- IBM Storage Protect for Virtual Environments: Data Protection for VMware.
- IBM Storage Protect for Virtual Environments: Data Protection for Microsoft Hyper-V on Windows Server 2012 or later Windows Server operating systems with the Hyper-V role installed.

How it works

The following processes occur during incremental-forever-incremental backup processing of a virtual machine:

- A full virtual machine backup is required only one time.
- A full virtual machine backup operation copies all of the used disk blocks that are owned by a virtual machine to the IBM Storage Protect server.
- After the initial full backup, all subsequent backup operations of the virtual machine are incremental-forever-incremental backups.
- This method copies only the blocks that changed since the previous backup, regardless of the type of the previous backup.
- The server uses a grouping technology that associates the changed blocks from the most recent backup with data that is already stored on the server from previous backups.
- A new full backup is then effectively created each time changed blocks are copied to the server by an incremental-forever-incremental backup.

Advantages

Incremental-forever-incremental backup processing has the following advantages:

- Periodic full virtual machine backups are no longer necessary.
- This method reduces the amount of data that goes across the network.
- This method reduces data growth because all incremental backups contain only the blocks that changed since the previous backup.
- No comparison with the backup target is required since only changed blocks are identified.
- Impact to the client system is minimized.
- The length of the backup window is reduced.
- Data restore operations are simplified.
- This method optimizes data restore operations.

Disadvantages

Incremental-forever-incremental backup processing has the following disadvantages:

- If one or more of the progressive incremental backups is corrupted on the server, you might not be able to fully recover a virtual machine. To ensure that you can fully recover a virtual machine, periodically run a full virtual machine backup.
- This method is available only with data movers that run on Linux and Windows clients.

Parallel backups of virtual machines

You can improve performance of virtual machine backups by running parallel backups of multiple virtual machines by using a single instance of the data mover. For more information, refer to [Vmmaxrestoresessions](#) and [Vmmaxrestoreparalleldisks](#).

How it works

The following processes occur during parallel backup processing of virtual machines:

- A single IBM Storage Protect data mover node can be used to concurrently back up multiple virtual machines.
- When the backups are initiated, the client establishes parallel sessions to copy the data to the IBM Storage Protect server.

Advantages

Parallel virtual machine backup processing has the following advantages:

- The backup window is reduced.
- You can optimize the backups so that they do not adversely affect the servers that are hosting the virtual machines.

Disadvantages

You must optimize the parallel backups. The number of virtual machines that you can back up in parallel depends on the following factors:

- The processing power of the server that the IBM Storage Protect data mover node runs on.
- The performance of I/O between the client and the IBM Storage Protect server.

Common client performance problems

Typical client performance issues are often related to backup operations that do not complete within the backup window or that send too much data over the network.

Resolving common client performance problems

The table contains common client problems and actions that can help to improve client performance.

Scenario	Solution	More information
During incremental backups, the client gets out of memory errors that cause the operating system to use more virtual memory or use more RAM than the client system can handle. How can I reduce these memory errors and have my backups complete within the backup window?	Update the client system hardware by increasing the system memory. If it is not possible to update the hardware, try running journal-based backups. If more memory is needed, try memory-efficient incremental backups.	For more information, see the following topics: <ul style="list-style-type: none"> • “Tuning journal-based backups” on page 221 • “Reduce client memory use” on page 206
Journal-based backups do not complete within the backup window. What alternatives can I use?	Try one or more of the following actions: <ul style="list-style-type: none"> • Use image backups to back up an entire volume as a snapshot • Examine the design of the file system on AIX, Linux, and Solaris operating systems 	For information about image backups, see “Image backup” on page 196 . For information about tuning client file spaces, see “File space tuning” on page 228 .
How can I reduce the amount of client data that is sent across the network to the IBM Storage Protect server?	Try one or more of the following methods: <ul style="list-style-type: none"> • Use compression during backup operations • Use include-exclude options to exclude files from the backup operation • Use client-side deduplication • Use LAN-free backups 	For more information, see the following topics: <ul style="list-style-type: none"> • “Reduce client data flow with compression” on page 208 • “Reducing client data flow with include and exclude options” on page 213 • “Checklist for data deduplication” on page 35 • “Performance tuning for LAN-free environments” on page 235

Scenario	Solution	More information
Some backup-archive clients back up much of the same data daily. How can I prevent the data that is a duplicate of the data from the previous day from being resent to the server?	Run incremental backups, use client-side data deduplication, or both.	“Checklist for data deduplication” on page 35
We have limited network bandwidth. How can I improve the communication between the client and IBM Storage Protect server?	Fine-tune the network and communication settings.	Chapter 14, “Tuning network performance,” on page 237
What other methods can I use to reduce the time that it takes to back up a client?	Try one of the following actions: <ul style="list-style-type: none"> • Use multiple client sessions for backup operations • Set the resourceutilization option to optimize the number of multiple sessions 	For more information, see the following topics: <ul style="list-style-type: none"> • Chapter 14, “Tuning network performance,” on page 237 • “Running concurrent client sessions” on page 217 • “Multiple session backup and restore” on page 218 • “Optimizing the number of multiple sessions to run ” on page 219

Resolving common performance problems with virtual machine backup operations

The table contains common questions and situations about virtual machine backup operations and solutions that can help to improve the performance.

The information in the following table, unless otherwise stated, applies to virtual machine backup operations on both the standard backup-archive client and IBM Storage Protect for Virtual Environments.

Scenario	Solution	More information
Excluding IBM Storage Protect settings, what settings can improve the efficiency of VMware operations with virtual guests?	<p>The VMware Changed Block Tracking (CBT) feature for virtual machines on the vSphere client is enabled automatically by IBM Storage Protect.</p> <p>This feature can reduce the data that is sent to the IBM Storage Protect server during incremental backups. CBT can reduce the total backup time and the storage pool size on the IBM Storage Protect server.</p> <p>CBT might increase the resource utilization slightly on vSphere hosts.</p>	Review information about CBT. Go to http://www.vmware.com/ and search for information about Changed Block Tracking (CBT) on virtual machines.

Scenario	Solution	More information
What is the best network adapter type to use for a system that is running as a VMware guest?	When you are running an IBM Storage Protect Windows client in a VMware guest environment, use the VMXNET Generation 3 (VMXNET 3) network adapter type. The VMXNET 3 adapter is a virtual network device from VMware that is optimized to provide enhanced hardware and software performance in a virtual environment.	Review information about the VMXNET 3 adapter. Go to http://www.vmware.com/ and search for information about the VMXNET 3 network adapter.
I am running parallel backups of virtual machines. How do I reduce the processor load during parallel backups, and also improve the throughput from the IBM Storage Protect backup-archive client to the IBM Storage Protect server?	Optimize parallel backups by using the following client options: <ul style="list-style-type: none"> • vmmaxparallel can be used with VMware and Microsoft Hyper-V virtual machines • vmmlimitperhost can be used with VMware only • vmmlimitperdatastore can be used with VMware only 	For more information, see the following topics: <ul style="list-style-type: none"> • “Optimizing parallel backups of virtual machines” on page 230 • Vmmaxparallel client option • Vmlimitperhost client option • Vmlimitperdatastore client option
How do I select the best transport mode for virtual backups?	The optimal transport mode to use depends on the composition of the backup environment. Use the vmvstortransport option to specify the preferred order of transport modes to use during backup or restore operations of VMware virtual machines.	For more information, see the following topics: <ul style="list-style-type: none"> • “Selecting a transport mode for VMware backups” on page 232 • Vmvstortransport client option
For IBM Storage Protect for Virtual Environments, how can I fine-tune the scalability of incremental-forever backups of virtual guests?	Virtual machine disk files are stored on the IBM Storage Protect as blocks of data called megablocks. When a change occurs on a disk in an area that is represented by a megablock, an IBM Storage Protect object is created. When a large number of IBM Storage Protect objects exist for the same virtual machine data, excessive demands are placed on the IBM Storage Protect server database. Use the following options to control the number of IBM Storage Protect objects that are created on the server: <ul style="list-style-type: none"> • mbobjrefreshthresh • mbpctrefreshthresh 	For more information, see the following topics: <ul style="list-style-type: none"> • “Adjusting the scalability of virtual machine backup operations” on page 233 • Mbobjrefreshthresh client option • Mbpctrefreshthresh client option • Setting options for an incremental forever backup schedule

Client restart options

The **commrestartduration** and **commrestartinterval** options affect how aggressively the client attempts to reconnect with the server if a network outage or other condition interrupts client and server communications.

Defaults are provided for both options and the defaults work well in most environments. You can experiment with these options to see whether smaller values speed up the reconnection process.

Neither option directly improves client performance, but either option, or both options, can be set to reduce reconnection times so that client data can be processed as soon as possible if the client is disconnected from the server.

commrestartduration client option

The **commrestartduration** option sets the number of seconds that the client waits after a communications failure before it tries to reconnect with a server.

You can set this value to any integer in the range 0 - 9999; the default is 60 seconds.

commrestartinterval client option

The **commrestartinterval** option sets the number of seconds that the client waits between each reconnection attempt.

You can set this value to any integer in the range 0 - 65535; the default is 15 seconds.

Tuning memory

You can tune the client to use less memory during incremental backup operations.

AIX | Linux | Mac OS X Client memory requirements and ulimit settings for incremental backups

The amount of memory that is used for normal incremental backups is proportional to the number of files that are being examined for backup. When you start an incremental backup operation that backs up numerous files, one way to prevent the client from running out of usable memory is by setting the operating system ulimit data value (`ulimit -d`) to **unlimited**.

If your business policies do not support a ulimit value of **unlimited**, you can estimate the system memory that is required for incremental backups by using the following calculations, and then set the ulimit to an appropriate value. Alternatively, you can set the `MEMORYEFFICIENT DISKCACHEMETHOD` client option or use journal-based backups to reduce memory requirements.

Tip: Your file system configuration can affect memory use.

To estimate the memory that is needed for an incremental backup, which is based on the number of objects (files and directories) that exist on the file system, complete the following steps:

1. Multiply the total number of objects by 700 to get an estimated number of bytes in the file system. The 700 is an estimate for the number of bytes in each path. For example, if the number of objects is 500000, then $500000 \times 700 = 350000000$.

Tip: The multiplier that is used in these calculations (700) is an estimate of the amount of memory that is needed per object. If your files and directories have file names that are longer than 80 characters, you might need extra memory.

2. Round up the value by 33% or to the next 100 MB from the previous step. In this example, round the 350 MB value to 500 MB. Convert this value to KB ($400 \times 1024 = 409600$ KB).
3. If you have more than one file system, estimate the memory requirements for each file system, and use the highest of these estimates to set the ulimit data value.

This estimate works when the client `RESOURCEUTILIZATION` option is set to 4 or less. The estimation is part of a *producer session*. A producer session is a producer thread that scans the file system to search for changed, new, or deleted files. A producer session uses memory, and the number of producer sessions is important for calculating random access memory (RAM).

If you use a `RESOURCEUTILIZATION` value of 5, 6, or 7, you can have up to two concurrent producer sessions. `RESOURCEUTILIZATION` values of 8 and 9 can have up to three concurrent producer sessions. If `RESOURCEUTILIZATION 10` is set, you can have up to four concurrent producers. You must base the `ulimit` value on the sum of the number of objects in each of the two, three, or four file systems that have the highest number of objects.

For example, in the following scenario, you have:

- `/fs1` contains 500000 objects
- `/fs2` contains 400000 objects
- `/fs3` contains 50000 objects
- `/fs4` contains 40000 objects

If you specify `RESOURCEUTILIZATION 6`, you can have up to two concurrent producer sessions. Therefore, you must calculate the `ulimit` value for the file systems with the highest number of objects (500000 and 400000):

$(500000 + 400000) * 700 = 630000000 \approx 630 \text{ MB}$. Round up to 700 MB, then convert to KB = 716800. Set the `ulimit` value to 716800.

Tip: If the number of objects on the file system increases, you must readjust the `ulimit` value to accommodate the growth.

If you do not set the `ulimit` value to **unlimited**, or to an estimated value, you can use the following methods to reduce the memory that is needed for incremental backups:

MEMORYEFFICIENTBACKUP DISKCACHEMETHOD

This method uses disk space as if it were system memory. You might be able to use the default `ulimit` value, but you need free disk space that is available to process objects. For more information about estimating the disk space that is needed for this option, see the backup-archive client documentation. If disk space is limited, you can use the **memoryefficientbackup yes** option setting. This option uses less disk space than the `DISKCACHEMETHOD` option, but it does decrease incremental backup performance.

Journal-based backup

This method uses journal-based backups. The journal daemon records changes to an object or its attributes in a journal database. During a journal-based backup, the client obtains a list of files that are eligible for backup from the journal database instead of by scanning the entire file system. Journal-based backups reduce the memory that is needed to process incremental backups.

Reduce client memory use

The **memoryefficientbackup** client option determines how much memory the client uses during incremental backup operations. Restricting the amount of memory that the client can use during incremental backups reduces the efficiency of incremental backup processing. The default setting for the **memoryefficientbackup** option is `no`, which does not limit the memory that the client can use.

During an incremental backup, the client determines which objects are new or changed since the last backup, and which objects must be expired on the server. An object, in this context, is a file or a directory.

By default, the client uses memory to create the list of new, changed, or expired objects to be updated by an incremental backup. Using available memory for this process makes incremental backups more efficient by reducing the time that is needed to prepare the list of objects to include in the backup. On client systems that have either limited memory, or that have applications that are not adversely affected if the client uses up the available memory, you can limit how much memory the client uses during incremental backups.

The following settings are available for the **memoryefficientbackup** option:

memoryefficientbackup no

The client uses an algorithm that does not limit the amount of memory that is used to process an incremental backup. This setting is the default, and it is the most efficient setting for incremental backups in many configurations.

memoryefficientbackup yes

The client uses an algorithm that requires less memory when it is processing incremental backups. This setting can increase the server workload, which, in turn, increases the time that is required to complete incremental backups. This setting can adversely affect incremental backup performance in configurations where there are many clients that are backing up files to the same server, and where each of the client systems has many file system objects.

memoryefficientbackup diskcachemethod

The client uses an algorithm that requires even less memory than `memoryefficientbackup yes`. With this setting, the client keeps the list of objects to back up or expire on disk, so it uses less memory.

In many configurations, the most efficient setting for incremental backups is `memoryefficientbackup no`. However, if memory is limited on the client systems, or if you see IBM Storage Protect error messages that indicate memory errors, consider using a different setting. To determine the setting, review the following guidelines and select the first setting that applies:

- `memoryefficientbackup no`

To determine whether you have enough memory to use `memoryefficientbackup no`, complete the following steps:

1. Determine the number of objects in the client file systems.
2. Round up the number of objects in the file systems, to the next million and divide that number by 1,000,000. Multiply the quotient by 300 MB to determine how to set the **memoryefficientbackup** option.

For example, if the client file systems have 5,202,131 objects, round up this number to 6,000,000. Divide the result by 1 million and assign it to a temporary variable that is called *numfsobjs*. In this example, *numfsobjs*=6 (6,000,000/1,000,000=6). Use the value of *numfsobjs* as described in the following calculations:

32-bit clients

If the value of the *numfsobjs* variable is less than or equal to 5, multiply *numfsobjs* by 300 MB. If the amount of physical memory on the client system is equal to or greater than the product of *numfsobjs* x 300 MB, specify `memoryefficientbackup no` (the default).

64-bit clients

If the amount of physical memory on your client is equal to or greater than the product of *numfsobjs* x 300 MB, specify `memoryefficientbackup no` (the default).

- `memoryefficientbackup diskcachemethod`

If the client has at least the following amount of fast, temporary disk storage available for use by the client process, specify `memoryefficientbackup diskcachemethod`.

- On UNIX and Linux systems, temporary disk space must meet or exceed *numfsobjs* x 300 MB.
- On Windows systems, temporary disk space must meet or exceed *numfsobjs* x 600 MB.
- On Mac OS X systems, temporary disk space must meet or exceed *numfsobjs* x 1200 MB.
- If none of the previous conditions apply, use `memoryefficientbackup yes`.

Alternatives to using the memoryefficientbackup client option

To reduce client memory consumption, you can use the following alternatives instead of the setting `memoryefficientbackup yes`.

- Use the client include and exclude options to back up only what is necessary.

- Use journal-based incremental backup on Windows (NTFS), AIX (JFS2), or Linux (all supported file systems) clients.
- Use the **virtualmountpoint** option to define multiple virtual mount points within a single file system, and back up these mount points, sequentially. Virtual mount points can be used on UNIX and Linux systems, but not on Mac OS X.
- Spread the data across multiple file systems and back up these file systems sequentially.
- Use the image backup function to back up the entire volume. Image backups can take less time and resources than incremental backups, on file systems that have many small files.

Tuning client data throughput

Use client options to improve the throughput of client data to IBM Storage Protect.

Reduce client data flow with compression

The backup-archive client can compress data before it sends the data to the server. Enabling compression on the client reduces the amount of data that is sent over the network and the space that is needed to store it on the server and storage pools. Two client options determine when and if the client compresses data: **compression** and **compressalways**.

In addition to compressing objects, to reduce the amount of data, you can also consider enabling client-side data deduplication. For information about configuring client-side data deduplication, see [“Tuning client-side data deduplication”](#) on page 210.

Related tasks

Compressing data to save storage space

You can use server-side data compression to increase the amount of available space in a storage pool.

compression client option

The **compression** client option specifies whether compression is enabled on the IBM Storage Protect client. For optimal backup and restore performance with many clients, consider enabling client compression.

Compressing the data on the client reduces demand on the network and the IBM Storage Protect server. The reduced amount of data on the server continues to provide performance benefits whenever this data is moved, such as for storage pool migration and storage pool backup. If you use data replication, compressed data remains compressed during the transfer from the source replication server to the target replication server. The data is then stored on the target replication server in the compressed format.

Client compression reduces the performance of each client, and the reduction is more pronounced on the slowest client systems. For optimal backup and restore performance when you have fast clients and a heavily loaded network or server, use client compression. For optimal backup and restore performance when you have a slow client, or a lightly loaded network or server, do not use compression. However, you must consider the trade-off of greater storage requirements on the server when you are not using client compression. The default for the **compression** option is no.

Compression can cause severe performance degradation if attempts to compress a file fail. Compression fails when the compressed file is larger than the original file. The client detects this size difference and stops the compression process, fails the transaction, and resends the entire transaction uncompressed. Compression failure occurs because the file type is not suitable for compression or the file is already compressed. Short of turning off compression, there are two options that you can use to reduce or eliminate compression failures:

- Use the **compressalways yes** option. This default option prevents compression reattempts if the compressed file is larger than the uncompressed file.
- Use the **exclude.compression** option in the client options file. This option disables compression for specific files, for example, all *.gif files or other files that get larger during attempted compression. Excluding these files saves processor cycles because there is no attempt to compress files that cannot

be compressed. Look in the client output (dmsched.log) for files that are causing compression reattempts and exclude those file types.

Use the following values for the **compression** option:

- For a single fast client, a fast network, and a fast server:

```
compression no
```

- For multiple clients, a slow network, or a slow server:

```
compression yes
```

Do not enable the client **compression** option if a client has a built-in file compression feature. For example, if hardware compression is in use for the media where Data Protection for Oracle data is stored, do not enable client compression. Compression, on these types of clients, reduces the amount of data that is backed up to the server.

Restriction: **Windows** Data can be compressed by using NTFS. However, the data must be decompressed before it can be accessed by the IBM Storage Protect server. Therefore, slower backups and higher processor usage might occur if you use NTFS compression.

compressalways client option

The **compressalways** option specifies whether to continue compressing an object if it grows during compression, or resend the object, uncompressed. This option is valid when client compression is enabled by the **compression** option.

The **compressalways** option is used with the **archive**, **incremental**, and **selective** commands. This option can also be defined on the server. If this option is set to yes, the default value, files compression continues even if the file size increases. To stop compression if the file size grows, and resend the file uncompressed, specify **compressalways no**. This option controls compression only if the administrator specifies that the client node determines the selection. To reduce the impact of repeated compression attempts if the compressed file is larger than the original, specify **compressalways yes**.

To prevent unsuccessful compression attempts, you can list files that cannot be compressed on one or more client **exclude.compression** statements. Exclude files that contain graphics; even exclude word-processing files if they contain embedded graphics. Also, exclude audio files, video files, files that are already encrypted, and files that are saved in an archive format, such as .jar files, .zip files, and other compressed file formats.

Using IBM Storage Protect client compression and encryption for the same files is valid. The client first compresses the file data and then encrypts it, so that there is no loss in compression effectiveness that is caused by encryption, and encryption is faster if there is less data to encrypt.

The following example shows how to exclude objects that are already compressed or encrypted, by using **exclude.compression** statements:

```
exclude.compression ?:\...\*.gif
exclude.compression ?:\...\*.jpg
exclude.compression ?:\...\*.zip
exclude.compression ?:\...\*.mp3
exclude.compression ?:\...\*.cab
exclude.compression ?:\...\*.aes
exclude.compression ?:\...\*.rsa
```

The preferred setting is **compressalways yes**, and then use **exclude.compression** statements to omit files that cannot be compressed.

Tuning client-side data deduplication

The performance of client-side data deduplication can be affected by processor requirements and deduplication configuration.

About this task

Data deduplication is a method of reducing storage needs by eliminating redundant data. Client-side data deduplication is the process of removing the redundant data during a backup operation on the client system. Client-side data deduplication is especially effective when you want to conserve bandwidth between the IBM Storage Protect client and server.

Procedure

- To help you enhance the performance of client-side data deduplication, take the following actions based on the task that you want to complete.

Table 25. Actions for tuning client-side data deduplication performance	
Action	Explanation
Ensure that the client system meets the minimum hardware requirements for client-side data deduplication.	<p>Before you decide to use client-side data deduplication, verify that the client system has adequate resources available during the backup window to run the deduplication processing.</p> <p>The preferred minimum processor requirement is the equivalent of one 2.2 GHz processor core per backup process with client-side data deduplication. For example, a system with a single-socket, quad-core, 2.2-GHz processor that is used 75% or less during the backup window is a good candidate for client-side data deduplication.</p>
Use a combination of deduplication and compression to obtain significant data reduction.	<p>When data is compressed after it is already deduplicated, it can give you more savings in data reduction as compared to running data deduplication alone. When data deduplication and compression are both enabled during a backup operation on the backup-archive client, the operations are sequenced in the preferred order (data deduplication followed by compression).</p>
Avoid running client compression in combination with server-side data deduplication.	<p>When you use client compression in combination with server-side data deduplication, it is typically slower and reduces data volume less than the preferred alternatives of server-side data deduplication alone, or the combination of client-side data deduplication and client-side compression.</p>

Table 25. Actions for tuning client-side data deduplication performance (continued)

Action	Explanation
<p>Increase the number of parallel sessions as an effective way to improve overall throughput when you are using client-side deduplication. This action applies to client systems that have sufficient processor resources, and when the client application is configured to perform parallel backups.</p>	<p>For example, when you use IBM Storage Protect for Virtual Environments, it might be possible to use up to 30 parallel VMware backup sessions before a 1 Gb network becomes saturated. Rather than immediately configuring numerous parallel sessions to improve throughput, increment the number of sessions gradually, and stop when you no longer see improvements in throughput.</p> <p>For information about optimizing parallel backups, see “Optimizing parallel backups of virtual machines” on page 230.</p>
<p>Configure the client data deduplication cache with the enablededupcache option.</p>	<p>The client must query the server for each extent of data that is processed. You can reduce the processor usage that is associated with this query process by configuring the cache on the client. With the data deduplication cache, the client can identify previously discovered extents during a backup session without querying the IBM Storage Protect server.</p> <p>The following guidelines apply when you configure the client data deduplication cache:</p> <ul style="list-style-type: none"> – For the backup-archive client, including VMware virtual machine backups, always configure the cache for client-side data deduplication. – For IBM Storage Protect for Virtual Environments operations, if you configure multiple client sessions to back up a vStorage backup server, you must configure a separate cache for each session. – For networks with low latency that process a large amount of deduplicated data daily, disable the client deduplication cache for faster performance. <p>Restriction:</p> <ul style="list-style-type: none"> – For applications that use the IBM Storage Protect API, do not use the client data deduplication cache because backup failures can occur if the cache becomes out of sync with the IBM Storage Protect server. This restriction applies to the IBM Storage Protect Data Protection applications. Do not configure the client data deduplication cache when you are using the data protection products. – If you use image backups, do not configure the client data deduplication cache.

Table 25. Actions for tuning client-side data deduplication performance (continued)	
Action	Explanation
Decide whether to use client-side data deduplication or server-side data deduplication.	<p>Whether you choose to use client-side data deduplication depends on your system environment. In a network-constrained environment, you can run data deduplication on the client to improve the elapsed time for backup operations. If the environment is not network-constrained and you run data deduplication on the client, it can result in longer elapsed backup times.</p> <p>To evaluate whether to use client-side data or server-side data deduplication, see the information in Table 26 on page 212.</p>

Use the following checklist to help you choose whether to implement client-side or server-side data deduplication.

Table 26. Checklist for choosing client-side versus server-side data deduplication	
Question	Response
Does the speed of your backup network result in long backup times?	<p>Yes</p> <p>Use client-side data deduplication to obtain both faster backups and increased storage savings on the IBM Storage Protect server.</p> <p>No</p> <p>Determine the importance of storage savings versus faster backup process.</p>
What is more important to your business: The amount of storage savings that you achieve through data reduction technologies, or how quickly backups complete?	<p>Consider the trade-offs between having the fastest elapsed backup times and gaining the maximum amount of storage pool savings:</p> <ul style="list-style-type: none"> – For the fastest backups in an unconstrained network, choose server-side data deduplication. – For the largest storage savings, choose client-side data deduplication that is combined with compression.

What to do next

For more information about using IBM Storage Protect deduplication, see [Configuring storage](#).

Related concepts

Checklist for [data deduplication](#)

Data deduplication requires more processing resources on the server or client. Use the checklist to verify that hardware and your IBM Storage Protect configuration have characteristics that are key to good performance.

Related tasks

[Evaluating data deduplication results](#)

You can evaluate the effectiveness of IBM Storage Protect data deduplication by examining various queries or reports. Actual data reduction results can show whether the expected storage savings are

achieved. You can also evaluate other key operational factors, such as database utilization, to ensure that they are consistent with expectations.

Tuning server-side data deduplication

Tune settings and configuration for different operations to ensure that the performance of server-side data deduplication is efficient.

Reducing client data flow with include and exclude options

You can use client **include** and **exclude** options to reduce the amount of data that is backed up, which shortens the backup and restore windows.

Almost every file system collects files that are not critical to your applications and users. Examples of such files include operating system files or updates that can be easily downloaded and be reinstalled if you do not have a local copy, core files, log files, and cached data. Use the IBM Storage Protect client **include** and **exclude** options to ensure that the client operations are protecting only what is important.

Adjusting the I/O buffer size of the client

The **diskbuffsize** client option specifies the maximum disk input/output (I/O) buffer size (in KB) that the client uses when it is reading files.

Optimal client performance during backup, archive, or HSM migration processes might be achieved if the value for this option is equal to or smaller than the amount of file read ahead that is provided by the client file system. A larger buffer requires more memory and might not improve performance.

The default value is 32 KB for all clients except clients that run on AIX. For clients that run on AIX operating systems, the default value is 256 KB, except when `enablelanfree yes` is specified. When `enablelanfree yes` is specified on AIX, the default value is 32 KB. API client applications have a default value of 1023 KB, except for Windows API client applications, version 5.3.7 and later, which have a default value of 32 KB.

The default value is the preferred value for the client operating system.

If the performance of restore operations seems slow, consider resizing the **diskbuffsize** option:

1. Stop the current restore operation.
2. Set this option in the appropriate server stanza in the `dsm.sys` file: `diskbuffsize 32`.
3. Restart the restore operation.

Optimizing the transaction size

A transaction is a unit of work that is exchanged between a client and the server.

A client program can transfer more than one file or directory between the client and server before the client commits the data in the transaction to server storage. If a transaction contains more than one file or directory, it is called a transaction group.

Data in a transaction is sent from the client to the server during backup operations. Data in a transaction is sent from the server to the client during a restore operation.

You can control the amount of data that is sent between the client and server by setting the **txnbytelimit** client option. The server administrator can also limit the number of files or directories that are contained within a transaction group by setting the **TXNGROUPMAX** option.

These two options work together such that the size of any transaction is based on which of these parameter values is reached first. For example, if the **TXNGROUPMAX** option is set to 4096 and the **txnbytelimit** option is set to 25600 KB (25 MB), then up to 4096 small files can be included in a transaction if the sum of their sizes does not exceed 25600 KB. If objects are 25 MB in size, or larger, they are sent as one file in one transaction.

Changing how much data can be sent in a transaction affects the speed at which the client performs work. The default values are sufficient in most environments, except when the data is being written directly to a

tape device. At the end of each transaction, the tape buffers must be written to physical media, which is a slow process; transferring more data with each transaction can improve performance when the server is writing data directly to tape.

Consider the following tips when you set a value for **txnbytelimit**:

- Increasing the amount of data per transaction increases the size of logs and log pools space on the server. Verify that you have enough free disk space to contain larger transaction logs and log pool space. Increasing log sizes can also cause longer server start times.
- Increasing the amount of data per transaction causes more data to be retransmitted if an error occurs. Resending data reduces performance, and resending larger transactions reduces performance even more.
- The benefits of changing the setting for the **txnbytelimit** option depend on the configuration and the type of workloads. In particular, increasing this value benefits tape storage pool backup more so than disk storage pool backup, especially if you are protecting many small files.

Consider setting a smaller **txnbytelimit** value if error conditions cause repeated retransmissions of transactions, when you specify static, shared static, or shared dynamic as the copy serialization attribute in the standard management class. A smaller **txnbytelimit** value applies to static and shared attributes. If a file changes during a backup operation and the client does not send it, the client still must resend the other files in that transaction.

To enhance performance, set the **txnbytelimit** option to 2 GB, and on the server, set the **TXNGROUPMAX** option to 256 KB. Additionally, for small file workloads, stage the backups to a disk storage pool and then migrate the files to tape.

For the **txnbytelimit** option, you can specify a value in the range 300 KB - 32 GB. The default value is 25600 KB.

Suggested settings for **txnbytelimit**

txnbytelimit setting when backing up objects to disk before migrating them to tape:

```
txnbytelimit 25600K
```

txnbytelimit setting when backing up objects directly to tape:

```
txnbytelimit 10G
```

If you are using IBM Storage Protect with the IBM Content Manager application and you notice slow server data movement operations, see the support article at [technote 85953](#) for information about using the CM VOL_AGGREGATESIZE setting in the CM RMVOLUMES table to improve transaction performance.

Effects of management classes on transactions

Each copy of a file or directory that is backed up by IBM Storage Protect is bound to (associated with) a management class.

A management class includes a backup copy group. A backup copy group defines how IBM Storage Protect manages objects that were backed up. Management class attributes include such things as the storage pool where the object is stored, how many versions of each object are created, and how long the versions are retained.

During backup operations, IBM Storage Protect bundles file and directory backup copies into transactions. That is, the client opens a transaction with the server database, backs up one or more objects, and then closes the transaction. If the IBM Storage Protect server database successfully commits the transaction, the client repeats the process until all eligible objects are backed up.

Transaction-based processing provides for reliable backups, but each commit operation also increases processing time. In general, the best performance is achieved by grouping as many objects as possible into a single transaction.

The maximum size of a transaction is governed by the following two parameters:

TXNGROUPMAX

This option is set on the server. It specifies the maximum number of objects that can comprise a transaction

txnbytelimit

This option is set on each client. It specifies the maximum size of a transaction, in KB.

The size of any transaction is based on which of these parameter values is reached first. For example, if the **TXNGROUPMAX** option is set to 4096 and the **txnbytelimit** option is set to 25600 KB (25 MB), then up to 4096 small files can be included in a transaction if the sum of their sizes does not exceed 25600 KB. If objects are 25 MB in size, or larger, they are sent as one file in one transaction.

Another factor that can influence the transaction size is the destination storage pool for the backed up objects. The objects within a transaction must all be directed to the same storage pool. As the transaction is processed, if one of the objects is directed to a different storage pool, then the current transaction is committed, and a new transaction is opened for the object that is destined for another storage pool.

If the destination storage pool is frequently changed during a backup operation, performance is reduced because new transactions must be created. For example, assume that you have a directory structure that contains many media files of different media file formats, such as the following files:

```
/media/vid001.jpg
/media/vid001.wmv
/media/vid002.jpg
/media/vid002.wmv
.
.
.
/media/vid9999.wmv
```

Assume also that you have **include** statements that bind these files types to different management classes, such as the following examples:

```
include /media/*.jpg diskclass
include /media/*.wmv tapeclass
```

The management classes named DISKCLASS and TAPECLASS each specify a different storage pool: one writes to disk and the other writes to tape. When the media files are backed up, the `/media/vid001.jpg` file is backed up in one transaction and is directed to the disk storage pool. The next object, `/media/vid001.wmv`, is backed up in another transaction that is directed to the tape storage pool. Then, the `/media/vid002.jpg` file is backed up in yet another new transaction and is directed to the disk storage pool. This behavior adversely affects the performance of backup operations. In addition to the additional processing time from the inefficient transactions, more delays can occur if you must wait for tapes to be mounted.

Consider revising your management classes, or management class bindings, to reduce or eliminate the number of different storage pools that the client uses for backed up objects.

A similar situation can occur when directory objects are backed up. By default, directory objects are bound to the management class with the longest **RETONLY** (retain only version) value. If multiple management classes in the active policy set have the same value for **RETONLY**, then the management class that is last, when sorted alphabetically, is used. For example, if the management classes named DISKCLASS and TAPECLASS both have the same **RETONLY** setting, and they are both in the active policy set, then the default management class for directory objects is TAPECLASS.

If file objects are directed to a disk storage pool and directory objects are directed to a different storage pool, such as tape, that also reduces transaction efficiency and slows performance. One way to avoid the inefficiencies of using a different management class for directory objects is to use the **dirmc** option and specify the same management class that you use to back up files. Using the example management classes named DISKCLASS and TAPECLASS, set **DIRMC** DISKCLASS to bind directory objects to the management class, and storage pool, that you use for file objects.

Setting options to minimize processor usage

You can set several client options to reduce the time that is required for client processing tasks and improve performance. The client options to consider are **quiet**, **virtualnodename**, **ifnewer**, **incrbydate**, and **tapeprompt**.

For Mac OS X file systems, limiting the length of extended attributes can help improve the client performance.

For clients on any operating system, turning off antivirus programs, or other programs that compete with the client for system resources, can also improve client performance.

quiet client option

Two client options determine whether messages are displayed during backup operations: **quiet** and **verbose**. The **verbose** client option is the default option; it causes messages to be displayed in the output during client operations. The **quiet** client option can be set to suppress messages from being displayed.

When you set the **quiet** option, messages and summary information are still written to the log files, but they are not displayed in GUI or command-line output. The **quiet** option provides two main benefits that can improve client performance:

- For tape backups, the first transaction group of data is always resent. To avoid the transaction from being resent, use the **quiet** option to reduce retransmissions at the client.
- If you are using the client scheduler to schedule backups, the **quiet** option reduces entries in the schedule log, which might improve client throughput.

Although the **quiet** option can offer some modest performance improvements, consider using the default (**verbose**) option. The benefits of having messages displayed and logged can outweigh the performance gains that the **quiet** option provides.

virtualnodename client option

When you restore, retrieve, or query objects that are owned by another node, consider using the client **virtualnodename** option instead of the **fromnode** option.

The **fromnode** option uses more system resources than the **virtualnodename** option. By using the **virtualnodename** option instead of the **fromnode** option, you might improve client performance.

ifnewer client option

The **ifnewer** client option is used only with restore commands. This option can reduce network traffic during restore operations. This option ensures that files are restored only if the date of the file that is stored on the server is newer than the date of the same file that is stored on the client node.

The **ifnewer** option can be set only on the command line, and, like all options that are specified on the command line, it must be preceded by a hyphen (-) character. For example:

```
dsmc restore "/home/grover/*" -sub=y -rep=y -ifnewer
```

incrbydate client option

The **incrbydate** client option can shorten backup windows because the option causes the client to back up objects only if they are new, or changed, since the last incremental backup was run. As its name implies, this option can be used only to perform incremental backups.

Incremental backups that use the **incrbydate** option, have limitations that regular incremental backups do not have. You must understand these limitations to correctly use this option. For incremental by date backups, consider the following limitations:

- Files that are created or modified after a directory was processed by the IBM Storage Protect client, but before the backup completes, are skipped at the next **incrbydate** backup.

- **incrbydate** backups do not cause expired files to be deleted from the server.
- If a management class changes for a file or a directory after an **incrbydate** is run, the stored objects are not rebound to the new management class.
- After an **incrbydate** backup runs, if only the attributes of an object are changed, the file is not included in the next **incrbydate** backup.

During an incremental backup operation, where the **incrbydate** option is not used, the server reads the attributes of all the client files that are in the server file system and passes this information to the client. The client then compares the server attribute list to all of the files in the client file system. This comparison can be time-consuming, especially on clients that have limited memory.

With an incremental-by-date backup, the server passes only the date of the last successful incremental backup operation to the client and the client backs up only files that are new or changed since the previous incremental backup. The time savings can be significant. However, regular, periodic incremental backups are still needed to back up files workstation files that are excluded from backups, by the limitations of incremental by date backups.

For example, if a new file in your file system has a creation date that is earlier than the last successful backup date, future incremental-by-date backups do not back up this file because the client assumes that the file has already been backed up. Also, files that were deleted are not detected by an incremental-by-date backup and these deleted files are restored if you perform a full system restore.

For a full list of restrictions about the **incrbydate** option, see [incrbydate option](#).

The **incrbydate** option can be specified only on the command line with the **incremental** command. Like all options that are specified on the command line, it must be preceded by a hyphen (-) character. For example:

```
dsmc incremental -incrbydate
```

Consider journal-based backups as an alternative to incremental by date backups. Journal-based backups perform a traditional incremental backup of the file system when the first backup occurs. A journal file records which file system objects are changed after the initial backup, and the journal is used to determine which objects to include in subsequent backups. Journal-based backup is most appropriate for file systems that do not change many objects, often. For more information about journal-based backups, see [“Journal-based backup” on page 191](#).

tapeprompt client option

The **tapeprompt** client option specifies whether you want to be prompted to wait for a tape to be mounted, if a tape is required to back up or restore objects.

Backup and restore, and archive and retrieve operations can be processed with fewer interactive input delays if you specify **tapeprompt no**. With **tapeprompt** no specified, the client still waits for tapes to be mounted if they are required, but the prompts that ask whether you want to wait for a tape, or skip objects that require a tape, are suppressed.

Improving client performance by using multiple sessions

You can set up the client to use multiple backup sessions to improve client performance.

Running concurrent client sessions

Running two or more client program instances at the same time on the same system might provide better overall throughput than a single client instance, depending on the available resources.

You can schedule backups for multiple file systems concurrently on one IBM Storage Protect client system with any of the following methods:

- By using one node name, and running one client scheduler, and setting the **resourceutilization** client option to 5 or greater, and including multiple file systems in the schedule, or in the domain specification. This method is the simplest way to run concurrent sessions.

- By using one node name, running one client scheduler, and scheduling a command that runs a script on the client system, where the script includes multiple client (**dsmc**) commands.
- By using multiple node names and running one client scheduler for each node name, and where each scheduler uses its own client options file.

Multiple session backup and restore

A multiple session restore operation allows backup-archive clients to start multiple sessions that use no-query restore operations, thus increasing the speed of restore operations. A multiple session restore operation is similar to a multiple session backup operation.

Multiple session restores can be used under the following conditions:

- The data to be restored is stored on several tape volumes or file device class volumes.
- Sufficient mount points are available.
- The restore is done by using the no-query restore protocol.

When you request a backup or archive, the client can establish more than one session with the server. The default is to use two sessions: one to query the server and one to send file data.

Parallel (concurrent) backup and restore operations that work with sequential file or tape storage pools require multiple mount points. A mount point is a tape or a file device class volume. The **resourceutilization** client option governs the maximum number of concurrent backup or restore sessions that the client can use. The **MAXNUMMP** server parameter, on the **UPDATE NODE** or **REGISTER NODE** commands, and the **MOUNTLIMIT** setting in the **DEFINE DEVCLASS** and **UPDATE DEVCLASS** commands, determines how many mount points a client node can use, at one time.

Configure these settings according to your requirements and available hardware. Take into account the number of mount points that all nodes might need, at any one time. For example, if you have four client nodes and only eight tape drives, if you configure all four nodes with **MAXNUMMP 8**, one node can seize all of the tape drives, leaving no tape drives for other nodes to use.

If all the files are on random disk, only one session is used. There is no multiple session restore for a random-access disk-only storage pool restore. However, if you are restoring files and the files are on four sequential disk volumes (or on four tape volumes) and other files are on random access disk, you can use up to five sessions during the restore.

Server settings take precedence over client settings. If the client **resourceutilization** option value exceeds the value of the server **MAXNUMMP** setting for a node, you are limited to the number of sessions that are specified by the **MAXNUMMP** parameter.

Multiple restore sessions are allowed only for *no-query* restore operations. A no-query restore is started by using an unrestricted wildcard in the file specification on the **restore** command. The following is an example of a no-query restore.

```
dsmc restore /home/*
```

The wildcard character (*) is unrestricted because it does not filter on object names or extensions. For example, `dsmc restore /home/????.*` is unrestricted. No-query restores also cannot use any of the object-filtering options. Specifically, you cannot use the **inactive**, **latest**, **pick**, **fromdate**, or **todate** options. For details about running a no-query restore, see the [Restore command](#).

The server sends the **MAXNUMMP** value to the client during sign-on. During a no-query restore operation, if the client receives a notification from the server that another volume that contains data to be restored was found, the client checks the **MAXNUMMP** value. If another session would exceed the **MAXNUMMP** value, the client does not start the session.

Backup considerations

Only one producer session per file system compares attributes for incremental backup. Incremental backup throughput does not improve for a single file system with a small amount of changed data.

Data transfer sessions do not have file system affinity; each consumer session could send files from multiple file systems, which helps balance the workload. Sending files from multiple file systems is not beneficial if you are backing up directly to a tape storage pool that is collocated by file space. Do not use multiple sessions to back up objects directly to a storage pool collocated by file space. Use multiple commands, one per file space.

The setting of the **resourceutilization** option and internal heuristics determine whether new consumer sessions are started.

When you backup objects directly to tape, you can prevent multiple sessions, so that data is not spread across multiple volumes, by setting the **resourceutilization** option to 2.

Restore considerations

Only one session is used when files are restored from random access disk storage pools.

Only one file system can be restored at a time with the command line, but multiple sessions can be used on a single file system.

Even small clients can realize improved throughput of restore operations if the data to be restored is on multiple tapes. One session can be restoring data while another might be waiting for tapes to be mounted, or be delayed while it is reading a tape, looking for the data to restore.

Tape cartridge contention might occur, especially when files are not restored from a collocated pool. Collocating files reduces the likelihood of tape cartridge contention.

Optimizing the number of multiple sessions to run

IBM Storage Protect clients can establish concurrent sessions to back up and restore data. The creation of concurrent sessions is controlled by an algorithm within the client software; you cannot directly control this algorithm. The default behavior is to use two sessions: one to query the server and one to send file data. You can set the **resourceutilization** option to cause the client to use additional concurrent sessions to query and send data.

Multiple sessions are used when you specify multiple file specifications on a backup command, **restore** command, **archive** command, or a **retrieve** command. For example, if you enter the following command and you specify **resourceutilization 5**, the client might start a second session to query the server for a list of files that were backed up on file space B:

```
inc /Volumes/filespaceA /Volumes/filespaceB
```

Whether the second session starts depends on how long it takes to query the server about files that are backed up on file space A. The client might also try to read data from the file system and send it to the server on multiple sessions.

The value that you specify for the **resourceutilization** option is an integer in the range 1 - 100. The value that you specify does not correlate directly to the number of sessions the client can create. For example, setting **resourceutilization 5** does not mean that the client can run only five concurrent sessions. What it does indicate is that this client might create more concurrent sessions than a client that has **resourceutilization** set to 1, but fewer concurrent sessions than a client that has **resourceutilization** set to 10. The **resourceutilization** option setting increases or decreases the ability of each client to create multiple sessions.

The following factors affect the performance of concurrent sessions:

Available server resources and processing capacity

The hardware that the IBM Storage Protect server runs on must have sufficient memory, storage, and processor capacity to efficiently support multiple sessions.

Available client resources and processing capacity

The hardware that the IBM Storage Protect client runs on must also have sufficient memory, storage, and processor capacity to efficiently support multiple sessions.

Configuration of the client storage subsystem

File systems that are spread across multiple disks, either by software striping, RAID-0, or RAID-5, can accommodate the increases in random read requests that concurrent sessions generate with better efficiency than a single-drive file system can. In fact, a single-drive file system might not show any performance improvement if you set the **resourceutilization** option.

For file systems spread across multiple physical disks, setting the **resourceutilization** to 5 or higher can produce optimal performance in configurations where the server has sufficient processing capacity and memory to handle the load.

Network bandwidth

Concurrent sessions increase the amount of data that flows over the network. In particular, LANs might be adversely affected by the increased data traffic.

If you set the **resourceutilization** option and you are backing up client files directly to a sequential device, update the server **MAXNUMMP** setting to accommodate the additional mount points that concurrent sessions might need.

Before you change any settings, consider the potential disadvantages of concurrent sessions:

- Concurrent sessions might produce multiple accounting reports.
- The server might not be configured to support all of the potential concurrent sessions. Review the server **MAXSESSIONS** setting and change it if the client-initiated sessions can exceed its current value.
- A **QUERY NODE** command might not accurately summarize client activity.

During restore operations, the default client behavior is to use a single session, unless the **resourceutilization** option was specified to a value greater than 2. When you are restoring files for a critical client system from tape, and the files are on many tape volumes, set the **RESOURCEUTILIZATION** value to 10. If four tape drives are available, and you want to restore operations to use all four tape volumes concurrently, set the **MAXNUMMP** value for the node to 4. If all of the client files that are being restored are in random-access disk storage pools, only one restore session is used, regardless of the **resourceutilization** option value.

The default value for the **resourceutilization** option is 2, and the maximum value is 100.

For example, if the data to be restored is on five different tape volumes, and the maximum number of mount points for the node that is requesting the restore is 5, and **resourceutilization** option is set to 3, three sessions are used for the restore. If the **resourceutilization** setting is increased to 5, five sessions are used for the restore. A one-to-one relationship exists between the number of restore sessions that are allowed and the **resourceutilization** setting.

The following values are the preferred settings:

For workstations

```
resourceutilization 1
```

For a small server

```
resourceutilization 5
```

For a large server

```
resourceutilization 10
```

The following table shows the maximum number of concurrent sessions that are possible, for each of the values in the **resourceutilization** range. A producer thread is a session that scans the client system for eligible files. The remaining sessions are consumer threads and are used to transfer data. Subtract the producer sessions that are listed in the table from the maximum number of sessions to determine the number of consumer threads. In the table, the threshold column shows how soon a subsequent thread can be started after the previous thread is started, for each of the values that are specified for the **resourceutilization** option.

resourceutilization value	Maximum number of sessions	Unique number of producer sessions	Threshold (seconds)
1	1	0	45
2 (default)	2	1	45
3	3	1	45
4	3	1	30
5	4	2	30
6	4	2	20
7	5	2	20
8	6	2	20
9	7	3	20
10	8	4	10
n > 10	n	n/2, rounded down to the nearest whole number	10

Tuning journal-based backups

To help improve the performance of incremental backups, you can run journal-based backups.

About this task

Journal-based backups have the following advantages over standard incremental backups:

- Journal-based backups can complete faster than standard incremental backups because they do not compare file system object attributes with information that is stored on the server. Instead, on a file system that supports journaling, changes to a file system are recorded in a locally stored journal database. The locally stored journal database entries are used to determine which objects to include in backup operations.

The benefits of using a journal-based backup do diminish if the file systems have many file changes. Journal-based backups perform best on large file systems in which many of the files do not often change.

- Journal-based backups require less memory and less client disk I/O than full incremental backups.

Procedure

- Use the information in the following table to help you tune journal-based backups.

Action	Explanation
Ensure that there is enough disk space on the client system to contain the journal database.	The amount of disk space that is required for the journal database depends on the number of files and directories that change between each successive journal-based backup operation.
Use default settings.	The default settings for journal size, log names and locations, file system check intervals, and other journal settings work well in most environments.

Action	Explanation
Change the default settings. Edit the <code>tsmjbbd.ini.smp</code> file to include or exclude file systems to monitor for changes, set the size of the journal database, and specify notification options and other settings. Save the changes to a file named <code>tsmjbbd.ini</code> (without the <code>smp</code> extension).	<p>If the default settings do not work well in your system environment, change them. For example, by excluding file systems, you can limit the amount of data to monitor for journal-based backups. This action can improve the performance of the backups.</p> <p>Configuration settings for the journal service (on Windows) or journal daemon (on Linux and AIX) are copied to the client disk when you install the backup-archive client. The default settings are in the <code>tsmjbbd.ini.smp</code> file. Comments in the <code>tsmjbbd.ini.smp</code> file provide the documentation for the journal settings.</p> <p>The journal service or journal daemon uses the <code>tsmjbbd.ini</code> file when the journal service is started.</p> <p>Tips for the Windows client:</p> <ul style="list-style-type: none"> – You can use the configuration wizard to edit the default settings. – Changes that are made to the <code>tsmjbbd.ini</code> file are applied dynamically. When changes are made to the settings in the file, the journal service applies the changes automatically without needing to restart the service.

What to do next

For more information about when to use journal-based backups, see [“Journal-based backup” on page 191](#).

Optimizing restore operations for clients

Standard IBM Storage Protect progressive incremental backup operations are optimized to restore individual files or small numbers of files.

Progressive incremental backup minimizes tape usage, reduces network traffic during backup operations, and eliminates the storage and tracking of multiple copies of the same data. Progressive incremental backup might reduce the impact to client applications during backup. For a balanced level of backup and restore performance, try running progressive incremental backup with collocation set on, in the storage pool.

If restore performance is more important than a balance between backup and restore operations, you can optimize based on your goals for restore performance. When you optimize for restore operations, there are often costs in tape usage and backup performance.

When you optimize restore operations, the performance depends on the type of media that you use. For more information about the media that you can use to restore data, see [Table 27 on page 223](#).

Table 27. Advantages and disadvantages of the different device types for restore operations

Device type	Advantages	Disadvantages
Random access disk	<ul style="list-style-type: none"> • Quick access to files • No mount point needed 	<ul style="list-style-type: none"> • No reclamation of unused space in aggregates • No deduplication of data
Sequential access disk (FILE)	<ul style="list-style-type: none"> • Reclamation of unused space in aggregates • Quick access to files (disk based) • Allows deduplication of data 	Requires mount point but not as severe an impact as real tape
Virtual tape library	<ul style="list-style-type: none"> • Quick access to files because of disk-based media • Existing applications that were written for real tape do not have to be rewritten 	<ul style="list-style-type: none"> • Requires mount point but not as severe an impact as real tape • No deduplication of data
Active data pools	<ul style="list-style-type: none"> • No sorting through inactive files to get to active data • Can be defined on any type of storage pool • Tapes can be taken offsite for disaster recovery 	Cannot be used with random access disk pools
Tape	<ul style="list-style-type: none"> • A large amount of data can be stored on a tape • Tapes can be taken offsite for disaster recovery 	<ul style="list-style-type: none"> • Requires mount point and physical tape mounting/dismounting • No deduplication of data • Slower access to files because of sequential access of tapes

The following tasks can help you balance the costs against the need for optimized restore operations:

- Identify systems that are most critical to your business. Consider where your most important data is, what is most critical to restore, and what needs the fastest restore. Identify which systems and applications you want to focus on, optimizing for restore.
- Identify your goals and order the goals by priority. The following list has some goals to consider:
 - Disaster recovery or recovery from hardware crashes, requiring file system restores
 - Recovery from loss or deletion of individual files or groups of files
 - Recovery for database applications (specific to the API)
 - Point-in-time recovery of groups of files

The importance of each goal can vary for the different client systems that you identified as being most critical.

For more information about restore operations for clients, see [“Concepts for client restore operations”](#) on page 226.

Environment considerations

IBM Storage Protect performance depends upon the environment.

The environment includes network characteristics, storage hardware, and time constraints for backup and restore operations. Consider the following items when deciding on the storage hardware:

- Types of tape drives used
- The availability of snapshot functions
- The availability of disk drives
- The availability of fiber-channel adapters

Consider sequential-access disk (FILE) to store data that requires quick restoration. For data that is less critical, store the data to random access disk, then allow or force the data to migrate to tape.

You can also use active-data pools to store active versions of client backup data. Archive and space-managed data is not allowed in active-data pools. Inactive files are removed from the active-data pool during expiration processing. Active-data pools that are associated with a FILE device class do not require tape mounts, and the server does not have to position past inactive files. In addition, FILE volumes can be accessed concurrently by multiple client sessions or server processes. You can also create active-data pools that use tape media, which can be moved off-site, but which require tape mounts.

If you do not use FILE or active-data pools, consider how restore performance is affected by the layout of data across single or multiple tape volumes. You can have multiple simultaneous sessions when you use FILE to restore, and mount overhead is skipped with FILE volumes. Major causes of performance problems are excessive tape mounts and needing to skip over expired or inactive data on a tape. After a long series of incremental backups, perhaps over years, the active data for a single file space can be spread across many tape volumes. A single tape volume can have active data that is mixed with inactive and expired data.

Restoring entire file systems

Using a file system image backup optimizes restore operations when an entire file system must be restored. For example, in disaster recovery or recovery from a hardware failure.

Restoring from an image backup minimizes concurrent mounts of tapes and positioning within a tape during the restore operation.

Consider the following information when you run file system restore operations:

- Combine image backups with progressive incremental backups for the file system to allow for full restore to an arbitrary point-in-time.
- To minimize disruption to the client during backup, use either hardware-based or software-based snapshot techniques for the file system.
- Perform image backups infrequently. More frequent image backups give better point-in-time granularity, but there is a cost. The frequent backups affect the tape usage, there is an interruption of the client system during backup, and there is greater network bandwidth needed.

As a guideline you can run an image backup after a percentage of data is changed in the file system, since the last image backup.

Image backup is not available for all clients. If image backup is not available for your client, use file-level restore as an alternative.

Restoring parts of file systems

Progressive incremental backups optimize restore operations for small numbers of files or groups of files. These backups also make optimal use of network bandwidth for backup operations, and can minimize elapsed backup time and tape usage.

To optimize for restoring a file or a group of files, or for a system on which an image backup cannot be made, consider the following methods:

- Use collocation by group, by a single client node, or by client file space for primary sequential pools that clients back up to. For large file spaces for which restore performance is critical, consider creating mount points on the client system. The mount points would allow collocation of data under the file space level.
- Specify the client option COLLOCATEBYFILESPEC. This option helps limit the number of tapes that are written to by objects from one file specification. For more information about this option, see [Collocatebyfilespec](#).
- Create backup sets that can be taken to the client system and can be used to restore from, directly. This method is effective if there is sufficient lead time before the restore, and can save network bandwidth.

Backup set creation can also be done periodically when resources are available, for example, on weekends.

- Use progressive incremental backups, but periodically force a backup of all files.

Some users reported it effective to define multiple IBM Storage Protect client nodes on a system. One client node runs the incremental backups and uses policies which retain multiple versions. Another client node runs either full backups or incremental backups with collocation, but uses policies that retain a single version. One node can be used for restoring older versions of individual files. You can use the other client node for restoring a complete file system or directory tree to the latest version.

Another effective way to optimize restore operations is to occasionally create a backup image.

- Create multiple storage pool hierarchies for clients with different priorities. For the most critical data, the best choice might be to use only disk storage. You can use different storage hierarchies to set collocation differently in the hierarchies.
- Although it might affect server performance, issue the **MOVE NODEDATA** command to consolidate critical data in tape storage pools. You can even issue the command in storage pools that have collocation that is turned on. It might be important to consolidate data for certain nodes, file spaces, and data types more often than for others. If you do not use collocation or are limited by tape quantity, you can consolidate data more often. You can also consider the rate of data turnover.

For more information about collocation, see [Optimizing operations by enabling collocation of client files](#).

Restoring databases for applications

Doing more frequent full backups leads to faster restores for databases. For some database products, you can use multiple sessions to restore, you can restore just the database, or restore just the database log files.

For information about data protection for databases, see [IBM Storage Protect for Databases](#)

Restoring files to a point-in-time

Keeping many versions is not essential for restoring to a point-in-time. But by increasing the number of versions that you keep, you might restore from an earlier point-in-time and still find the versions corresponding to that time.

If you also schedule incremental backups regularly, you might have greater granularity in restoring to a discrete point-in-time. However, keeping many versions can degrade restore operation performance. Setting policy to keep many versions also has costs, in terms of database space and storage pool space. Your policies might have overall performance implications.

If you cannot afford the resource costs of keeping the large numbers of file versions and must restore to a point-in-time, consider the following options:

- Use retention sets
- Use backup sets
- Export the client data
- Use an archive
- Take a volume image, including virtual machine backups

You can restore to the point-in-time when the backup set was generated, the export was run, or the archive was created. Remember, when you restore the data, your selection is limited to the time at which you created the backup set, export, or archive.

Tip: If you use the archive function, create a monthly or yearly archive. Do not use archive as a primary backup method because frequent archives with large amounts of data can affect server and client performance.

See [“Restoring parts of file systems” on page 224](#).

Concepts for client restore operations

The client restore includes the following operations:

[“No-query restore operations” on page 226](#)

[“Running multiple commands with backup and restore” on page 227](#)

[“Running multiple sessions on clients for a restore” on page 227](#)

[“Controlling resource utilization by a client” on page 227](#)

No-query restore operations

The client uses two different methods for restore operations: Standard restore (also called classic restore), and no-query restore.

The no-query restore requires less interaction between the client and the server, and the client can use multiple sessions for the restore operation. The no-query restore operation is useful when you restore large file systems on a client with limited memory. The advantage is that no-query restore avoids some processing that can affect the performance of other client applications. In addition, it can achieve a high degree of parallelism by restoring with multiple sessions from the server and storage agent simultaneously.

With no-query restore operations, the client sends a single restore request to the server instead of querying the server for each object to be restored. The server returns the files and directories to the client without further action by the client. The client accepts the data that comes from the server and restores it to the destination named on the restore command.

The no-query restore operation is used by the client only when the restore request meets both of the following criteria:

- You enter the restore command with a source file specification that has an unrestricted wildcard.

An example of a source file specification with an unrestricted wildcard is:

```
/home/mydocs/2002/*
```

An example of a source file specification with a restricted wildcard is:

```
/home/mydocs/2002/sales.*
```

- You do not specify any of the following client options:

- inactive
- latest
- pick
- fromdate
- todate

To force classic restore operations, use `?*` in the source file specification rather than `*`. For example:

```
/home/mydocs/2002/?*
```

For more information about restore processes, see the [Restore command](#).

Running multiple commands with backup and restore

You can run multiple commands instead of multiple sessions to speed up the backup and restore of client nodes with critical data.

When you use multi-sessions to back up data, the sessions might be contending for the same underlying hard disk. The contention for resources can cause delays in processing.

An alternative method is to manage backups by starting multiple client commands, where each command backs up a predetermined number of file systems. Using this method, with collocation at the file space level, can improve backup throughput and allow for parallel restore processes across the same hard disk drives.

You must issue multiple commands when you are restoring more than one file space. For example, when you are restoring both a C drive and a D drive on a Windows system you must issue multiple commands.

You can issue the commands one after another in a single session or window, or issue them at the same time from different command windows.

When you enter multiple commands to restore files from a single file space, specify a unique part of the file space in each restore command. Be sure that you do not use any overlapping file specifications in the commands. To display a list of the directories in a file space, issue the **QUERY BACKUP** command on the client. For example:

```
dsmc query backup -dironly -subdir=no /usr/
```

Running multiple sessions on clients for a restore

To use multiple sessions, data for the client must be on multiple sequential access volumes in a file or tape storage pool. Or the data can be contained in a random disk storage pool (with a device class with a device type of DISK). The data for a client usually becomes spread out over some number of volumes over time.

To potentially benefit from multi-session restores, consider collocating client data by group. Collocation by group can cause data for a node to be distributed on more than one volume. The distribution happens while the group's total data is kept on as few volumes as possible.

Restore operations can be restricted on mount points. The **MAXNUMMP** parameter in the **REGISTER NODE** or **UPDATE NODE** command applies to restore operations. The client can restrict the number of sessions, which are based on the combination of the **MAXNUMMP** value and the client **RESOURCEUTILIZATION** value. Unlike tape, you can mount FILE volumes in multiple sessions simultaneously, for restore or retrieve operations.

Set the client option for resource utilization to one greater than the number of sessions that you want. Use the number of drives that you want that single client to use. The client option can be included in a client option set.

Issue the restore command so that it results in a no-query restore process.

Controlling resource utilization by a client

You can control the number of mount points (equivalent to drives) allowed to a client by setting the **MAXNUMMP** parameter on either the **UPDATE NODE** or **REGISTER NODE** command.

At the client, the option for resource utilization also affects how many drives (sessions) the client can use. The client option, resource utilization, can be included in a client option set. If the number specified in the **MAXNUMMP** parameter is too low and there are not enough mount points for each of the sessions, it might not be possible to achieve the benefit of the multiple sessions that are specified in the resource utilization client option.

- For backup operations, prevent multiple sessions if the client is backing up directly to tape so that data is not spread among multiple volumes. Multiple sessions can be prevented at the client by using a value of 2 for the resource utilization option on the client.

- For restore operations, set the resource utilization option to one greater than the number of sessions that you want. Use the number of drives that you want that single client to use.
- With file spaces, a session is limited to processing a single file space. You cannot have multiple backup or restore sessions processing a single file space. However, if you have multiple file spaces on a client, you can have multiple sessions processing those file spaces.

File space tuning

Using IBM Storage Protect virtual mount points can enhance the performance of backup and restore operations on file systems that contain millions of files.

On many operating systems that IBM Storage Protect supports, you might be able to use file system or operating system tools to divide the file systems into manageable units, such that each file system can be protected in an acceptable backup or restore window.

On AIX, Linux, and Solaris operating systems, the IBM Storage Protect **virtualmountpoint** option can be used to logically divide a large file system into smaller increments. *Virtual mount points* are an IBM Storage Protect construct. Virtual mount points are not recognized by the operating system as mount points; they are only recognized and used by IBM Storage Protect.

When it protects objects that are contained under a virtual mount point, IBM Storage Protect treats each virtual mount point as a separate file space. Creating virtual mount points can enhance performance in the following ways:

- Less memory is needed for client operations because the virtual mount points divide a large file system into smaller increments, and processing fewer objects uses less memory.
- IBM Storage Protect can do more work in parallel, by running concurrent backup or restore operations on objects that are under two or more of the virtual mount points.

Using virtual mount points to improve performance works best if each of the virtual mount points has approximately the same number of files. If you cannot divide your file system in this way, virtual mount points might not be a suitable means to improve performance.

To illustrate how to use virtual mount points, assume that a client has a large file system called `/data`. Also, assume that the `/data` file system has multiple subdirectories that you want to frequently protect.

You can use the **virtualmountpoint** option to create virtual mount points that divide the `/data` file system into manageable logical units, such as shown in the following example:

```
virtualmountpoint /data/dir1
virtualmountpoint /data/dir2
.
.
virtualmountpoint /data/dir19
virtualmountpoint /data/dir20
```

These example **virtualmountpoint** options create 20 virtual mount points for the `/data` file system. The objects (`dir1`, `dir2`, and so on) that are shown on the example **virtualmountpoint** statements are directory objects on the file system. When objects in these directories are stored on the server, they are stored in a file space that matches name of the objects that are included on each **virtualmountpoint** statement. That is, objects in `dir1` are stored in file space named `dir1`, and so on.

You can back up and restore objects in each virtual mount point, independently of the others, and independently of other objects that are not in a virtual mount point. Any objects that get added to the `/data` file system, but that are not under a virtual mount point, are protected when you back up the objects in the `/data` file system. Objects that are under a virtual mount point are protected when you back up the virtual mount point.

If you use **virtualmountpoint** options, monitor the growth of the physical file system. If many new objects are added to the physical file system in locations that are not defined as logical mount points, it might eventually be easier to abandon the use of the virtual mount points and just backup the entire file system.

If you intend to use virtual mount points to divide the contents of a large file system, be aware that adding virtual mounts after a file system has been backed up can change the command syntax that is needed to restore objects.

For example, assume that you back up the `/data/dir1/file1` object before you create any virtual mount points. The `/data/dir1/file1` object is stored on the server in the `/data` file space. Assume that you later create a virtual file space by setting `virtualmountpoint /data/dir1` and that you create and back up a `file1` object in it. This new `file1` object is stored on the server in the `/dir1` file space (the file space matches the virtual mount point name).

Running `dsmc restore /data/dir1/file1` restores the `file1` object from the copy that is stored on the server in the virtual mount point (`dir1`) file space.

To restore the `file1` object that was saved in the `/data` file space, you must use the following syntax:

```
dsmc restore {/data}/dir1/file1
```

The brace (`{` and `}`) characters force the server to search the `/data` file space for the `file1` object.

Consider the following items if you use virtual mount points to create addition file spaces on the server:

- For applications that use the IBM Storage Protect API, limit the number of file spaces to no more than 100 per client. Examples of programs that use the API are IBM Storage Protect for Virtual Environments, IBM Storage Protect for Mail, IBM Storage Protect for Enterprise Resource Planning, and IBM Storage Protect for Databases.
- For sequential-accessed storage pool volumes, collocate files by node or group, rather than by file space. For example, 100 small file systems require 100 volumes if they are collocated by file space, but fewer volumes are needed if the files are collocated by node or group.

Windows system state backups

The most recently released versions of the IBM Storage Protect backup-archive client and IBM Storage Protect server software include updates that improve the performance of Windows system state backup and restore operations.

There are no user-configurable options or settings that can be adjusted to improve the efficiency of Windows system state protection. Backing up or restoring the Windows system state is a resource-intensive and time consuming operation. If you decide that you must back up the Windows system state, consider whether you can define server policy settings that retain fewer versions of system state backups. For example, your organization might require you to retain data files for 60 days, but require only 10 days of retention for system state information. The client **include.systemstate** option can be used to specify a different management class to be used for system state backups.

The decision to explicitly back up the Windows system state depends on how you plan to restore a node, after a node fails. The following factors can influence your decision to back up the Windows system state data:

- If you plan to restore a node by reinstalling the operating system from the Windows installation media, or from a repair disk and image backup, you do not need to back up the Windows system state data.
- If you plan to restore a physical machine from an image backup or from a snapshot image backup, back up all volumes because system state data might exist on other disks, and not just on the C drive.
- If you plan to restore a Windows virtual machine, system state objects are backed up when you perform a full backup of the virtual machine. A separate backup of the system state data is not required to restore a Windows virtual machine from a full virtual machine backup.
- If you plan to perform a bare metal restore of a client node, you must explicitly back up the system state files so they are available to restore the system state objects to the same, or a different, system. To reduce storage requirements, associate system state backups with policies that limit the number of backup copies that are retained on the server or in storage pools.

Restriction: Bare metal restore of Microsoft Windows servers and workstations that adhere to the Unified Extensible Firmware Interface (UEFI) specification is possible only from IBM Storage Protect backup-archive clients that are version 7.1, or later.

Tuning virtual machine backup operations

You can improve the performance of backup operations for virtual machines by adjusting client options.

About this task

To help improve the performance of backup operations for virtual machines with the backup-archive client or IBM Storage Protect for Virtual Environments, adjust the settings for the following options:

- Options for optimizing parallel backups of virtual machines
- Options for the transport mode for VMware backups
- Options for adjusting the scalability of virtual machine backup operations (applicable only to Tivoli Storage Manager for Virtual Environments 6.4 or later, or IBM Storage Protect for Virtual Environments 7.1.3 or later)

Optimizing parallel backups of virtual machines

IBM Storage Protect backup-archive client 6.4 and later provides parallel backup processing for backing up multiple virtual machines at the same time with one IBM Storage Protect data mover node.

About this task

The data mover node is the node that represents a specific IBM Storage Protect backup-archive client that moves data from one system to another.

With IBM Storage Protect for Virtual Environments, you can run parallel backup processing of incremental-forever full and incremental-forever incremental virtual machine backups.

To help optimize parallel virtual machine backups for IBM Storage Protect for Virtual Environments: Data Protection for VMware, adjust the settings for the **vmmaxparallel**, **vmlimitperhost**, and **vmlimitperdatastore** options. These options can also help reduce the processor load that parallel backups can create on a host in the vSphere infrastructure.

To optimize parallel virtual machine backups for IBM Storage Protect for Virtual Environments: Data Protection for Hyper-V, adjust the settings for the **vmmaxparallel** option.

For more information option, see [Parallel backups of virtual machines](#).

vmmaxparallel client option

The **vmmaxparallel** option specifies the maximum number of virtual machines that can be backed up to an IBM Storage Protect server at one time per client process.

Before you set a value for the **vmmaxparallel** option, review the following information:

Preferred setting

The preferred value for the **vmmaxparallel** option depends on the following factors:

- Resource availability on the IBM Storage Protect server and client
- Network bandwidth between the server and the client
- The load tolerance on the participating VMware infrastructure

To determine the preferred setting for this option, experiment with parallel backups of the virtual machines. In that way, you can select a setting that is appropriate for the backup window and the hardware and system configuration in the environment.

The default value is 1, which might be too restrictive. The maximum value is 50, which might be ineffective.

In general, set this option to the highest value that still maintains a tolerable processor load on the vSphere hosts and a tolerable I/O load on the data stores for a virtual machine backup workload.

Effects on performance

Increasing this value can result in more parallelism in terms of backup processing and can improve aggregate throughput to the IBM Storage Protect server from a backup-archive client.

Potential trade-offs for this setting

Setting a value too low might limit the potential of a backup environment by restricting the aggregate throughput to the IBM Storage Protect server. However, a low value might be necessary to throttle the amount of data that is sent from the client to the IBM Storage Protect server, or to reduce the processor or I/O load on the vSphere hosts and data stores.

Setting a value too high might result in oversaturating the link between the client and the IBM Storage Protect server, or elevating the processor load on certain vSphere hosts. Beyond a certain value, you might not experience improvements in aggregate throughput performance, depending on the network bandwidth and proxy or vSphere host processor resources that are available.

vm`limitperhost` client option

The **vm`limitperhost`** option specifies the maximum number of virtual machines on an ESX server that can be included in a parallel backup operation.

Before you set a value for the **vm`limitperhost`** option, review the following information:

Preferred setting

The preferred value is the default value of 0. By using the default value, you set no limits on the maximum number of virtual machines on an ESX server that can be included in a parallel backup operation. Verify that the value for the **vm`limitperhost`** option is compatible with the value for the **vm`maxparallel`** option.

Set the **vm`limitperhost`** option to the highest value that can still maintain a tolerable processor load on any single vSphere host for a virtual backup workload. Ensure that backup sessions are distributed equally across the affected vSphere hosts.

When you are specifying the setting, consider the set of virtual machines that are being backed up.

For example, if a set of 10 virtual machine guests is hosted on 5 vSphere hosts, and **vm`maxparallel`** is set to 10, set the **vm`limitperhost`** option to 2. In this way, you can distribute the parallel backup sessions across the hosts during a 10-guest parallel backup operation.

Effects on performance

The **vm`limitperhost`**, **vm`maxparallel`**, and **vm`limitperdatastore`** options limit the number of parallel backup operations that occur overall and for any single vSphere host. You can use these options to reduce the processor load that parallel backups can create on a vSphere host.

For a set of virtual machine guests that you back up, the order that IBM Storage Protect uses for creating backup sessions is random. Depending on the setting for the **vm`maxparallel`** option, it might be possible that too many backup sessions would involve too few vSphere hosts at any one time during a backup operation.

The **vm`limitperhost`** option can be used to ensure that no more than the number of backup sessions that are indicated by the **vm`limitperhost`** option value affect any one host.

Potential trade-offs for this setting

Setting a value too low might artificially limit the maximum number of concurrent virtual machine backups in an environment to less than what is feasible. However, a low value might be necessary to throttle the amount of data that is sent to the IBM Storage Protect server, or to reduce the processor load on the involved vSphere hosts.

Setting a value too high might result in elevated processor loads on certain vSphere hosts.

vm-limit-per-datastore client option

The **vm-limit-per-datastore** option specifies the maximum number of virtual machines in a data store that can be included in a parallel backup operation.

Before you set a value for the **vm-limit-per-datastore** option, review the following information:

Preferred setting

The preferred value is the default value of 0. By using this value, you set no limit on the maximum number of virtual machines in a data store that can be included in a parallel backup operation. However, ensure that the selected value is compatible with the value that is being used for the **vm-max-parallel** option.

Set the **vm-limit-per-datastore** option to the highest value that can still maintain a tolerable processor load on any single vSphere host for a virtual machine backup workload. In addition, adjust this value so that the backup workload is spread across as many vSphere data stores as possible.

When you are specifying the setting, consider the set of virtual machines that are being backed up.

For example, if a set of 10 virtual machine guests is hosted on 5 vSphere data stores, and **vm-max-parallel** is set to 10, set the **vm-limit-per-datastore** option to 2. In this way, you can distribute the parallel backup sessions across the data stores during a 10-guest parallel backup operation.

Effects on performance

The **vm-limit-per-datastore**, **vm-max-parallel**, and **vm-limit-per-host** options limit the number of parallel backups that occur overall and for any single vSphere datastore. You can set these options to reduce the processor load that parallel backups can create on a vSphere host or hot spots on vSphere datastore LUNs.

For a set of virtual machine guests that you back up, the order that IBM Storage Protect uses to create backup sessions is random. Depending on the setting for the **vm-max-parallel** option, it might be possible that too many backup sessions would involve too few vSphere datastores.

The **vm-limit-per-datastore** option can be used to ensure that no more than the number of backup sessions that are indicated by the **vm-limit-per-datastore** option affect any one datastore.

Potential trade-offs for this setting

Setting a value too low might artificially limit the maximum number of concurrent virtual machine backups in an environment to less than what is feasible. However, a low value might be necessary to throttle the amount of data that is sent to the IBM Storage Protect server, or to reduce the processor load on the vSphere hosts or the I/O load on vSphere data stores.

Setting a value too high might result in elevated processor loads on certain vSphere hosts, depending on the mapping of VMware data stores to hosts. Setting too high a value might also result in elevated loads on certain vSphere data stores. This outcome might lead to inefficiencies because the underlying LUNs of those data stores must handle excessive I/O processes compared to others.

Selecting a transport mode for VMware backups

To set the preferred transport order or hierarchy for backup or restore operations on VMware virtual machines, specify the **vm-transport** option.

Procedure

Before you set the **vm-transport** option, review the following information:

- In most cases, set the **vm-transport** option to default (san:hotadd:nbdssl:nbd). If you do not specify a value for the **vm-transport** option, the default value is used.

The optimal transport mode depends on the composition of the backup environment. [Table 28 on page 233](#) specifies the transport mode to use for specific backup environments.

Table 28. Preferred transport modes for specific backup environments	
Backup environment	Value for the <code>vmvstortransport</code> option
You want to offload backup traffic from the LAN, and instead, move it over a SAN.	<code>san</code>
You are using a data mover node that is installed on a virtual machine to back up other virtual machines. The backups can be moved over a SAN or a LAN.	<code>hotadd</code>
You are backing up virtual machines over an Ethernet LAN; you either do not have, or do not want to use, a SAN to offload backup traffic from the LAN.	<code>nbd</code>
You are backing up virtual machines over an Ethernet LAN and you want to use SSL to encrypt the data. Note that encrypting data can degrade backup performance.	<code>nbdssl</code>

- The values that are listed in [Table 28 on page 233](#) are the preferred transport methods, but you might not want to specify a single transport method as the value. You might specify multiple transport methods so that you can fail over to another transport method when the first one fails; otherwise, the operation might fail. However, you might want to restrict this option so that only a certain set of transport methods is used. If you leave an entry out of the colon-separated value list, that entry is no longer available and is skipped.
- Consider the effects of the setting on performance.
It is typically preferable to use the fastest available transport mode. However, in some environments, it might be necessary to avoid certain transport modes or emphasize others for resource management purposes.
- Consider the potential trade-offs for this setting.
Specifying a slower transport mode might reduce the aggregate throughput of the backup environment.

What to do next

For more information about the `vmvstortransport` option, see the [Vmvstortransport client option](#).

Adjusting the scalability of virtual machine backup operations

You can improve performance by adjusting the scalability within IBM Storage Protect for incremental-forever backups of virtual machines.

Before you begin

You must have a license to use IBM Storage Protect for Virtual Environments.

About this task

Virtual machine disk files are stored on the IBM Storage Protect server as data blocks called megablocks. Each of these megablocks contains 128 MB of data. When a change occurs on a disk in an area that is represented by a megablock, an IBM Storage Protect object is created. For every subsequent incremental backup, if a change is detected, an extra IBM Storage Protect object is created on the server. When a large amount of objects exists for the same virtual machine data, excessive demands are placed on the IBM Storage Protect server database.

Procedure

- To fine-tune these IBM Storage Protect server scalability conditions, use either the **mbobjrefreshthresh** or **mbpctrefreshthresh** option, but not both.

mbobjrefreshthresh client option

Use this option when you estimate the IBM Storage Protect objects that represent production data for each virtual machine backup.

For example, when the number of IBM Storage Protect objects exceeds this value, the megablock is refreshed. This action means that the entire 128 MB block is backed up to the IBM Storage Protect server and is represented as a single IBM Storage Protect object.

Before you set a value for the **mbobjrefreshthresh** option, review the following information:

Preferred setting

The preferred value is the default, 50. When the number of IBM Storage Protect objects that are needed to describe a 128 MB megablock for a virtual machine guest disk exceeds this value, the entire megablock is refreshed.

If you are backing up to a tape storage pool on the server, you might decrease this value so that megablock refreshes occur more often. This way, the data that was backed up for each virtual machine guest disk is more likely to be collocated on tape volumes. This setting might improve restore performance in such cases.

Effects on performance

When a megablock is refreshed, the objects that were used to represent the megablock area in previous backups are expired. This option can affect the quantity of data that is copied to the IBM Storage Protect server and the server database-related processor utilization during incremental forever backups.

Potential trade-offs for this setting

Setting this option close to its maximum value of 8192 can result in sending less data to the IBM Storage Protect server during an average incremental forever backup operation. However, the number of database entities that the IBM Storage Protect server must track increases. This result can increase server processor utilization during incremental forever backups by a small degree.

Setting this option close to its minimum value of 2 can result in marginal database processing savings during incremental forever backups. However, the quantity of data that is copied to the IBM Storage Protect server might be higher and might approach the size of a full backup.

mbpctrefreshthresh client option

The **mbpctrefreshthresh** option defines a threshold for the percentage of a megablock that can change before a full refresh is initiated. Use this option when you estimate the amount of extra data that is backed up for each virtual machine.

For example, when a 128 MB block of a production disk changes more than the percentage specified by the **mbpctrefreshthresh** option, the entire 128 MB block is copied to the IBM Storage Protect server. The block is represented as a single IBM Storage Protect object.

Before you set a value for the **mbpctrefreshthresh** option, review the following information:

Preferred setting

The preferred value is the default, 50. When a 128 MB megablock changes by a percentage that exceeds this value since its last refresh (a full copy to the IBM Storage Protect server), the entire megablock is refreshed.

If you are backing up to a tape storage pool on the server, you might decrease this value so that megablock refreshes occur more often. This way, the data that was backed up for each virtual machine guest disk is more likely to be collocated on tape volumes. This setting might improve restore performance in such cases.

Effects on performance

When a megablock is refreshed, the objects that were used to represent the megablock area in previous backups are expired. This option can affect the quantity of data that is copied to the IBM Storage Protect server and the server database-related processor utilization during incremental forever backups.

Potential trade-offs for this setting

Setting this option close to its maximum value of 100 can result in sending less data to the IBM Storage Protect server during an average incremental forever backup operation. However, the number of database entities that the IBM Storage Protect server must track increases. This result can increase server processor utilization during incremental forever backups by a small degree.

Setting this option close to its minimum value of 1 can result in marginal database processing savings during incremental forever backups. However, the quantity of data that is copied to the IBM Storage Protect server might be higher and might approach the size of a full backup.

Performance tuning for LAN-free environments

LAN-free backup can improve performance because the backup traffic can be routed over the SAN instead of the LAN. LAN-free data movement can make LAN bandwidth available for other uses and decrease the load on the IBM Storage Protect server, allowing it to support a greater number of concurrent client connections.

Backing up data to tape or disk over the SAN, or restoring data from tape or disk over the SAN, has the following advantages over equivalent operations that are performed only over the LAN:

- Metadata is sent to the server over the LAN; sending metadata over the LAN has negligible impact on LAN performance. Client data bypasses the potentially busy and slower LAN and is sent over the faster SAN. Backing up or restoring data over a SAN is generally faster than the same operation over a LAN.
- Sending client data over the SAN frees the IBM Storage Protect server from the task of handling data, which leads to more efficient use of server resources because the data goes directly to storage.
- Using a SAN is more efficient than a LAN when you are protecting large files, or databases; IBM Storage Protect Data Protection products generally benefit from SAN efficiencies.

When you configure IBM Storage Protect in a SAN environment, consider the following points:

- Ensure that you provide a sufficient number of data paths to tape drives.
- Backing up many small files directly to a real tape device can be inefficient. For file systems that have many small files, consider sending the files over the LAN to a disk storage pool, and migrate the files to tape later.
- Optimize the transaction size for writing files to tape or disk; for information, see [“Optimizing the transaction size”](#) on page 213.
- To improve backup and restore performance, include `lanfreecommmethod shardemem` in the client options file if the storage agent and client are on the same system. Setting this option allows the IBM Storage Protect client and IBM Storage Protect storage agent to communicate by using RAM, instead of using TCP/IP.
- Set the `tcpnodeLAY` option to YES, in the server or client options. This setting allows packets that are smaller than the maximum transmission unit (MTU) to be sent immediately.

Do not use LAN-free backup and restore if you are using IBM Storage Protect server-side data deduplication. You cannot bypass the server if server-side data deduplication is used to reduce redundant objects processing.

Chapter 14. Tuning network performance

If possible, use a dedicated local area network (LAN) or a storage area network (SAN) for backup operations. Keep device drivers for network components updated to take advantage of the latest fixes and improvements. Consider the IBM Storage Protect options that can help you tune how the clients and the server use the network. Ensure that you understand how TCP/IP works and review the information about TCP flow control and the sliding window.

Tuning TCP/IP settings for clients and servers

Typically, the default values for the client and server options for TCP/IP work well. In some cases, however, you can tune the settings to improve communication performance.

Before you begin

Review the information in [“TCP flow control and the sliding window”](#) on page 239. Ensure that you observe system performance before and after you change settings.

Procedure

- Typically, the default values for the **TCPWINDOWSIZE** options on the client and server are preferred. However, a larger window might improve communication performance, especially on fast networks with high latency, such as a long-distance wide area network (WAN) connection.

If you decide to tune TCP/IP window sizes, apply the following guidelines:

- If you increase the size of the **TCP/IP** window, do so in increments. For example, try doubling the value of the **TCPWINDOWSIZE** option and observing the results before you increase the value again. A larger value does not always improve performance.

Tip: Do not set the sliding window size to be larger than the buffer space on the network adapter. The window acts as a buffer on the network. A window size that is larger than the buffer space on the network adapter might cause packets to be lost on the network adapter. Because packets must be sent again when they are lost, throughput might degrade.

- Some Microsoft Windows and Linux systems have a TCP auto-tuning feature that monitors session transfer statistics and then adjusts the send and receive window sizes to optimize performance. For IBM Storage Protect servers and clients that run on these systems, the IBM Storage Protect **TCPWINDOWSIZE** option is set to the default value of 0 to use auto-tuning. When the option is set to 0, server sessions use the send and receive windows sizes that are specified by the operating system.

Tip: The TCP auto-tuning feature is enabled by default on some versions of Windows, and disabled by default on others. If you intend to use auto-tuning, ensure that it is enabled for the Windows system.

Always monitor operations after changes in the **TCPWINDOWSIZE** value to verify that the performance is not degraded.

- If the operating system cannot automatically tune TCP window size, change the **TCPWINDOWSIZE** option default value of 0 to a value in the range 16 - 2048. If you specify a value in the range 16 - 2048, the window size is in the range 1 KB - 2 MB.
- The window size that you set with the **TCPWINDOWSIZE** client option might have to be a compromise for different operations on the system. For example, you might have to determine the optimal value for backup-archive client operations and for IBM Storage Protect for Virtual Environments operations, and then specify a compromise value that is suitable for both.
- For the **TCPNODELAY** option, use the default value of YES.
This setting disables the Nagle algorithm and allows packets that are smaller than the MTU size to be immediately sent.

Controlling network traffic from client schedules

You can control the network traffic from scheduled client operations by using certain IBM Storage Protect server **SET** commands. The commands control how frequently clients contact the server, and how the sessions are spread across the schedule window.

Procedure

- Use the **SET RANDOMIZE** command to randomize start times within the startup window of each client schedule.

Communication errors can happen when many clients contacts the server simultaneously. If concurrent scheduled operations for clients have communication errors, you can increase the randomization percentage so that client contact is spread out. Increased randomization decreases the chance for communication overload and failure. The randomized start times apply only to clients that use the client-polling scheduling mode.

- Set how frequently a client can contact the server to obtain scheduled work by using the **SET QUERYSCHEDPERIOD** command. This command overrides the client setting, and applies when the client-polling mode is used for schedules.

A shorter time period means more network traffic because of client polling. Use longer settings (6 - 12 hours) to reduce network traffic. Alternately, use the server-prompted scheduling mode to eliminate network traffic because of client polling.

- Set a global limit on the number of times that a scheduled command is tried again on a client by using the **SET MAXCMDRETRIES** command. This command overrides the client setting. A smaller number reduces network traffic that is caused by scheduled commands that are tried again.

If you use the **SET MAXCMDRETRIES** command, be sure to consider schedule startup windows. If a retry is attempted outside of the schedule startup window, it fails.

- Set the number of minutes between retries of a scheduled command after a failed attempt to contact the server. Use the **SET RETRYPERIOD** command. This command overrides the client setting. A larger value reduces the network traffic that is caused by retries and increases the chance of a successful retry.

If you use the **SET RETRYPERIOD** command, be sure to consider schedule startup windows. If a retry is attempted outside of the schedule startup window, it fails.

Setting network options for IBM Storage Protect on AIX systems

The default values for network options on the IBM AIX operating system can be used for most IBM Storage Protect configurations. Review the preferred network-related settings.

Procedure

- If you use 10 Gbit Ethernet adapters, enable flow control on the switch or router port that the AIX system is connected to. For instructions about enabling flow control, consult your network administrator or see the manuals for your router or switch.
- Set the `rfc1323` network option to 1. To ensure that the new setting is used whenever the system is restarted, use the `-p` option on the command. For example, issue the following command:

```
no -o rfc1323=1 -p
```

- If you see nonzero values for the `no mbuf errors` field in the output from the **entstat**, **fddistat**, or **atmstat** commands, increase the value for the `thewall` option. Set the `thewall` option to at least 131072 and the `sb_max` option to at least 1310720.

Some versions of the AIX operating system have larger default values for these options, so that modifying the values might not be necessary.

TCP/IP and network concepts for advanced tuning

If you plan to tune TCP/IP settings for the IBM Storage Protect client or server, first ensure that you understand key concepts.

TCP/IP sends and receives data for applications on a system. TCP/IP is composed of two protocols: Transmission Control Protocol (TCP) and Internet Protocol (IP).

Applications such as the IBM Storage Protect client and server interact with TCP. By changing the **TCPWINDOWSIZE** client and server options, you affect the flow control function in TCP.

Applications do not interact with IP or lower-level protocols that control how one system communicates its receive window size to another, retransmission of lost data, or receipt of data from a sending system.

The following factors can affect network operations:

- System resources, such as memory and processors.
- Communications adapters. Link utilizations and the limitations of various communication layer implementations affect the use of resources.
- Data sizes and load on the network.

TCP flow control and the sliding window

Transmission Control Protocol (TCP) uses a *sliding window* for flow control. Before you tune any TCP/IP settings, first understand how the TCP sliding window works.

The TCP sliding window determines the number of unacknowledged bytes, x , that one system can send to another. Two factors determine the value of x :

- The size of the send buffer on the sending system
- The size and available space in the receive buffer on the receiving system

The sending system cannot send more bytes than space that is available in the receive buffer on the receiving system. TCP on the sending system must wait to send more data until all bytes in the current send buffer are acknowledged by TCP on the receiving system.

On the receiving system, TCP stores received data in a receive buffer. TCP acknowledges receipt of the data, and *advertises* (communicates) a new *receive window* to the sending system. The receive window represents the number of bytes that are available in the receive buffer. If the receive buffer is full, the receiving system advertises a receive window size of zero, and the sending system must wait to send more data. After the receiving application retrieves data from the receive buffer, the receiving system can then advertise a receive window size that is equal to the amount of data that was read. Then, TCP on the sending system can resume sending data.

The available space in the receive buffer depends on how quickly data is read from the buffer by the receiving application. TCP keeps the data in its receive buffer until the receiving application reads it from that buffer. After the receiving application reads the data, that space in the buffer is available for new data. The amount of free space in the buffer is advertised to the sending system, as described in the previous paragraph.

Ensure that you understand the TCP window size when you use sliding window for flow control. The window size is the amount of data that can be managed. You might need to adjust the window size if the receive buffer receives more data than it can communicate. For more information about optimizing the TCP window size, see [“Optimization of window size for different operations on the same system” on page 241](#).

How the send and receive buffers interact has the following consequences:

- The maximum number of unacknowledged bytes that a system can send is the smaller of two numbers:
 - The send buffer size on the sending system
 - The receive window size that the receiving system advertises to the sending system

- When the receiving application reads data as fast as the sending system can send it, the receive window stays at or near the size of the receive buffer. The result is that data flows smoothly across the network. If the receiving application can read the data fast enough, a larger receive window can improve performance.
- When the receive buffer is full, the receiving system advertises a receive window size of zero. The sending system must pause and temporarily cannot send any more data.
- In general, more frequent occurrences of zero size for the receive window results in overall slower data transmission across the network. Every time the receive window is zero, the sending system must wait before sending more data.

Typically, you set the send window and the receive window sizes separately for an operating system. In AIX, for example, the `tcp_sendspace` and `tcp_recvspace` parameters of the **no** command can be used to set the send and receive window sizes.

The sliding window that is used by IBM Storage Protect operations is controlled with the **TCPWINDOWSIZE** option.

Related concepts

TCP window controls in IBM Storage Protect

The **TCPWINDOWSIZE** options for IBM Storage Protect server and clients override the operating system settings for the size of send and receive windows for TCP/IP sessions. The **TCPWINDOWSIZE** option is available as a server option and a client option. With each option, you specify one value, which is used as the size for both the send and receive windows.

TCP window controls in IBM Storage Protect

The **TCPWINDOWSIZE** options for IBM Storage Protect server and clients override the operating system settings for the size of send and receive windows for TCP/IP sessions. The **TCPWINDOWSIZE** option is available as a server option and a client option. With each option, you specify one value, which is used as the size for both the send and receive windows.

During incremental backup operations for files, both client and server act as receivers of data:

- The server sends metadata about the inventory of active backup versions to the client. The metadata consists of file names and attributes. For file systems that contain millions of files, this data can be a substantial amount, as much as hundreds of megabytes or even gigabytes.
- The client sends backup copies of new and changed files to the server.

Typically the default value for the **TCPWINDOWSIZE** option works well. A larger window might improve communication performance, especially on fast networks with high latency, such as a long-distance wide area network (WAN) connection.

Specifying the **TCPWINDOWSIZE** option with a value of 0 causes IBM Storage Protect to use the operating system default for the TCP window size. If the operating system does not automatically tune the TCP window size, avoid using the operating system default. The operating system default might be optimized for other applications, which might not be the optimal setting for IBM Storage Protect.

If IBM Storage Protect clients and server are on the same subnet, a larger TCP window size is unlikely to improve throughput. Also, you might need more kernel memory if you set a large TCP receive window size. The risk of the increased memory requirements might be greater than the benefit of a larger TCP window size.

Modern operating systems provide TCP/IP stacks that commit the requested memory as it is needed. Therefore, these systems have less risk of increased kernel memory for the send and receive buffers. These operating systems also automatically tune the receive buffer size by observing the session transfer statistics, and either increasing or decreasing the receive window, as appropriate. For these operating systems only, you might set the IBM Storage Protect server **TCPWINDOWSIZE** option to 0 and use the automatic tuning feature. These settings are especially useful when clients that connect to the server are remote.

The **TCPWINDOWSIZE** option is not related to the **TCPBUFSIZE** server option or the **tcpbuffsize** client option. The **TCPWINDOWSIZE** option is also not related to the send and receive buffers that are allocated in client or server memory.

Related concepts

TCP flow control and the sliding window

Transmission Control Protocol (TCP) uses a *sliding window* for flow control. Before you tune any TCP/IP settings, first understand how the TCP sliding window works.

Optimization of window size for different operations on the same system

The sizes of TCP send and receive windows that work well for one application might not work well for another application, even for another IBM Storage Protect application.

Finding the correct balance for the window sizes between the server and the client is also important. For example, if you reduce the **TCPWINDOWSIZE** option on the client from 2000 to 63 and set the option on the server to 1024, slower backup performance is likely the result for the following reasons:

- IBM Storage Protect uses the **TCPWINDOWSIZE** to configure both the send buffer size and the receive buffer size. Thus on the client, the send buffer size is 63 KB.
- The maximum number of unacknowledged bytes that can be sent is limited by the smaller of the send buffer size and the receive window size that is communicated by the receiving system. Therefore, although the receiver (IBM Storage Protect server) has a window size of up to 1024 KB, the effective window size is 63 KB.

For backup operations to the IBM Storage Protect server, typically the server takes in data fast enough so that larger TCP receive windows do not cause the window size to become zero. If the server takes in data fast enough, larger window sizes can improve performance, and smaller window sizes can decrease performance.

Restore operations by a backup-archive client have different characteristics than restore operations by an IBM Storage Protect for Virtual Environments client. The backup-archive client performs the following actions for a restore operation:

1. The backup-archive client reads data that is sent by the IBM Storage Protect server from the TCP receive buffer.
2. The backup-archive client writes the data directly to files on disk.

If this write operation is slow and the IBM Storage Protect server sends data faster than the client can write it, the TCP receive buffer becomes full. When the TCP receive buffer becomes full, the receiver advertises a zero window size more frequently and the operation slows.

3. The client repeats steps 1 and 2 until all data is restored.

Typically, the restore operation by an IBM Storage Protect for Virtual Environments client can be slower than a restore operation by a backup-archive client because of the operations that occur for writing the data. The IBM Storage Protect for Virtual Environments client performs the following actions for a restore operation:

1. The IBM Storage Protect for Virtual Environments client reads data that is sent by the IBM Storage Protect server from the TCP receive buffer.
2. The IBM Storage Protect for Virtual Environments client writes the data to the vStorage API. Then, more operations and resources might be required, including communications with VMware, VMware processing of the data, and allocation of new blocks as the virtual machine disk (VMDK) grows.

If this write operation is slow and the IBM Storage Protect server sends data faster than the client can write it, the TCP receive buffer becomes full. When the TCP receive buffer becomes full, the receiver advertises a zero window size more frequently and the operation slows.

3. The IBM Storage Protect for Virtual Environments client repeats steps 1 and 2 until all data is restored.

In step “2” on page 241, the IBM Storage Protect for Virtual Environments restore operation has more possible operations and might need more resources than a restore operation by a backup-archive client.

Therefore, a larger receive window size is more likely to cause the TCP receive window size to become zero for the IBM Storage Protect for Virtual Environments restore operation. When both IBM Storage Protect for Virtual Environments and the backup-archive client are in use on a system, you must find a window size that balances the needs of the two types of operations. For example, in one case a window size of 1008 gave better overall throughput for such a system.

Chapter 15. Performance tuning for products that are used with IBM Storage Protect

Performance tuning information is available for products that are used with IBM Storage Protect products and components.

IBM Storage Protect Snapshot

Look for the most recent performance tuning information in the online product information at <https://www.ibm.com/docs/en/sps/>.

IBM Storage Protect for Space Management

See “Tuning for IBM Storage Protect for Space Management” on page 243.

IBM Storage Protect for Virtual Environments

Look for the most recent performance tuning information in the online product information at <https://www.ibm.com/docs/en/SSERB6>.

See also “Tuning virtual machine backup operations” on page 230.

Content Management

See “Content Management environments” on page 244.

Tuning for IBM Storage Protect for Space Management

IBM Storage Protect for Space Management provides techniques that you can use to optimize object migration to, and recall from, tape storage.

Optimizing migration to tape

If you must migrate many small files to the server, performance is better if the data moves to disk rather than to tape. After the files are migrated to disk, you can use storage pool migration to move the files to tape.

If you attempt to migrate many files directly to tape, performance of the migration operation might be unsatisfactory, particularly if each file is relatively small (<100 MB). By default, IBM Storage Protect for Space Management operates on one file at a time, so there is one transaction for each file that is migrated (or recalled). Setting the **hsmgroupedmigrate** option to YES can improve performance because it causes the selected files to be grouped into a single transaction. This option can be especially helpful, if you use a wildcard pattern on the file specification parameter on the **dsmmigrate** command. You might not know, in advance, how many files might match the wildcard pattern.

The number of objects that are grouped in a single transaction is limited by the IBM Storage Protect server option named **txngroupmax**. The number of bytes that are sent in a single transaction is limited by the IBM Storage Protect client option named **txnbytelimit**.

Optimizing recall from tape

Tape processing is automatically optimized when you use the **dsmrecall** command and include both the **-filelist** option, which specifies the file that contains the list of files to recall, and the destination file system. If you do not specify a file system, the recall process does not optimize tape processing.

For example, the following command optimizes tape recall because it specifies both the **-filelist** option and a file system:

```
dsmrecall -filelist=myFileList myFileSystem
```

The **-preview** option on the **dsmrecall** command produces file lists that list the files that are in the file that is specified by the **-filelist** option and orders the file lists by the tapes that contain the files. You

can specify each of these cartridge-centric file lists, on a separate **dsmrecall** command, to recall the files from tapes in a particular order.

If the list file that is specified on the **filelist** option is a collection file that was created by using the **dsmrecall** command with the **-preview** option, the recall begins immediately. The files are recalled in the order that they are specified in the collection file. To change the order in which files are recalled, you can edit the entries in the collection file.

If the list file is not a collection file, and the list file is correctly formatted, the file entries are sorted for optimized tape processing and they are then recalled.

In the following example, the recall begins immediately because the file that is identified by the **-filelist** parameter is a collection file:

```
dsmrecall -filelist=/HsmManagedFS/.SpaceMan/tapeOptimizedRecall/node_ID/PID/  
FileList.ordered.collection myFileSystem
```

For more information about optimizing tape recall processing, see [Optimized tape recall processing](#).

Content Management environments

Server performance can be affected by Content Management applications that interface with IBM Storage Protect through the IBM Storage Protect client application programming interface (API).

Minimizing time for each transaction

While most IBM Storage Protect operations process many files per transaction, Content Management applications tend to process few, or just one, file per transaction. With as few as one file per transaction, the time for each such transaction becomes critical. Key to Content Management application performance is the time that it takes to write to the storage pool and the active log.

- To minimize the time for writing to the storage pool and the active log, use disk systems that use write cache, which hides the latency of writing to physical disk.
- Consider avoiding the use of IBM Storage Protect features such as simultaneous write or active log mirroring. When you use these features, the server must perform more write operations at the end of each transaction. The additional write operations might cause slow performance in Content Management environments.
- Be careful when you are mirroring to storage over long distances. The time that is involved in the I/O process grows with the distance.

Reducing wasted space in FILE storage pools

If the average file that Content Management sends to IBM Storage Protect for backup is smaller than 256 KB, and you use FILE device classes for the storage pools, a substantial amount of space might be wasted in the storage pools.

Transactions that are 256 KB or less can waste space because the server writes a minimum of one block, or 256 KB, to a volume in a storage pool that uses a FILE device class. For example, if a transaction is only 64 KB, the space that is used on disk for the transaction is still 256 KB.

You can consider using the NONBLOCK data format for FILE storage pools that are used for Content Management data. Using the NONBLOCK data format instead of the NATIVE data format might save space under these conditions.

The data format for an existing storage pool cannot be changed. If your storage pools use the NATIVE data format and you want to try the NONBLOCK data format, you must define new storage pools.

Appendix A. Server instrumentation reference

You can use server instrumentation to track operations, such as backup and restore, and to help identify where performance problems originate.

Use the servermon component that is automatically installed and configured as part of the server installation to collect data at regular intervals.

Selecting a server instrumentation strategy

Follow the usage strategies to get the best results when you use server instrumentation.

Procedure

You can select any of the following server instrumentation strategies:

- Start and stop server instrumentation around the operation. An operation can be any procedure that affects performance, such as backup or restore operations.
 1. Start server instrumentation, and start the operation that you want to monitor.
 2. End server instrumentation immediately after the operation completes. If a thread is started when instrumentation is active, session and process statistics are included in the output. A thread is a sequence of actions that are managed by an operating system scheduler. A process might require more than one thread. For example, a backup operation uses at least two threads. You can issue an IBM Storage Protect administrative client macro command to start server instrumentation before you start the operation.
- Set a time limit when you run server instrumentation.
 - The optimum length of time to run server instrumentation for most cases is 5 - 15 minutes. You can run server instrumentation for up to 30 minutes.
 - If server instrumentation is active for 30 minutes, hundreds of threads are traced and the output can be overwhelming. Reports with that many threads can make it more difficult to diagnose a problem.
 - Do not run server instrumentation on a busy server for the maximum amount of time. When possible, limit instrumentation on the server. If the system workload is the problem, instrumentation results might not help you resolve the source of system performance problems.
- Find a match for the multiple threads of a particular session or process. Look for the parent and child relationships between the threads. In the instrumentation output per thread, use the thread ID and the parent thread ID to find the other threads that are associated with the operation.
 - Find the thread in the instrumentation data. For example, look in the IBM Storage Protect activity log file for a session ID that matches a particular client session in the instrumentation data.
 - During the operation, take the output from the **SHOW THREADS** command to see the session or process ID that a particular thread is working on. Use the thread ID from the output to find that same thread ID in the instrumentation.
 - Look for related threads, which are based on the amount of data that is moved.

Related reference

INSTRUMENTATION BEGIN

Use this command to start server instrumentation.

INSTRUMENTATION END

Use this command to stop server instrumentation and save the output.

Starting and stopping server instrumentation

You can start server instrumentation from an administrative command line or from an administrative client. After you stop server instrumentation, you can use the results to determine where performance problems are occurring.

About this task

You must have system privilege to start or stop server instrumentation.

Procedure

Complete the following steps to use server instrumentation:

1. Issue the **INSTRUMENTATION BEGIN** command to start instrumentation from the administrative client:

```
dsmadm -credentialsfile=secretpwdfile instrumentation begin
```

The *id* that is specified in the credentials file must be an IBM Storage Protect administrator ID that has system privilege.

2. Start an operation for which you want to analyze performance.
3. Issue the **INSTRUMENTATION END** command to stop server instrumentation. Specify an output file for the data. If you do not specify an output file, data is written only to the screen. If you issue any remote commands from an administrative client and redirect the output to a file, that file is saved to the administrative client.

The following command can be issued from the administrative client:

```
dsmadm -credentialsfile=secretpwdfile instrumentation end > filename
```

Related concepts

Server instrumentation categories

IBM Storage Protect server instrumentation can report on the elapsed times for the process categories that are documented in the table. Server instrumentation tracks all input and output on a thread-by-thread basis for the categories.

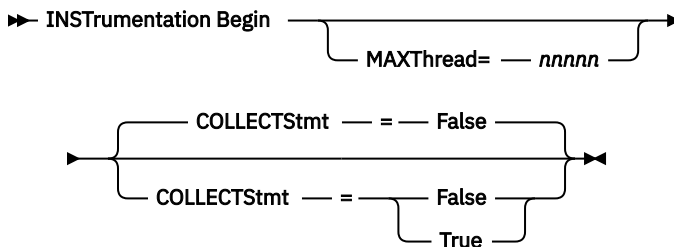
INSTRUMENTATION BEGIN

Use this command to start server instrumentation.

Privilege class

You must have system privilege to start server instrumentation.

Syntax



Parameters

MAXThread

Specifies the maximum number of threads that you want to trace. The default is 4096. If more than 4096 threads might be running during the instrumentation interval, raise this value. The maximum number of threads that you can run is 100,000. The minimum number of threads that you can run is 512.

COLLECTStmt

Collects detailed information on SQL commands, especially SQL commands that take a long time to run. The default value is **FALSE**. If you suspect that the Db2 program that is the database manager for the server is causing the problem, change this value to **TRUE**.

Example: Start server instrumentation at the administrative command-line client

Start an administrative client session in command-line mode, and begin collecting data. Use a credentials file that is named "secretpwdfile".

```
dsmadm -credentialsfile=secretpwdfile instrumentation begin
```

Example: Use command redirection to start server instrumentation for a storage agent

Start server instrumentation on storage agent StgAgnt_375.

```
dsmadm -credentialsfile=secretpwdfile  
StgAgnt_375:instrumentation begin
```


INSTRUMENTATION END

Use this command to stop server instrumentation and save the output.

Privilege class

You must have system privilege to stop server instrumentation.

Syntax

►► INSTRumentation End 

Parameters

filename

Specifies the file name for the output. Instrumentation output is generated when instrumentation is stopped. The file that you specify for instrumentation output is saved to the server directory. If you are running the administrative client from another system to issue this command, the output is not stored on the local system, but on the server system. You can specify a path to a different location if you cannot write to the server directory.

The preferred approach is to instead redirect the output to a file. See the examples that follow.

If you do not either specify a file name or redirect the output to a file, the output is only displayed on the screen and is not saved.

Example: Stop server instrumentation and redirect the output to a file

Stop server instrumentation and send the output to the instr_041413.ods file.

```
dsmadm -credentialsfile=secretpwdfile instrumentation end > instr_041413.ods
```

Example: Use command redirection to stop server instrumentation for a storage agent

Stop server instrumentation on storage agent StgAgnt_375 and send the output to the instr_041413.ods file.

```
dsmadm -credentialsfile=secretpwdfile StgAgnt_375:instrumentation end > instr_041413.ods
```

Related concepts

Server instrumentation categories

IBM Storage Protect server instrumentation can report on the elapsed times for the process categories that are documented in the table. Server instrumentation tracks all input and output on a thread-by-thread basis for the categories.

Server instrumentation for different operating platforms

Server instrumentation differs on the various IBM Storage Protect server operating systems.

The operating systems differ for server instrumentation in these respects:

- On operating systems such as AIX and Linux, only one thread does I/O to any disk storage-pool volume (called **DiskServerThread**). This thread provides a disk volume-centric view and can be difficult to get complete operation disk statistics.
- On Windows servers, the following processes occur:
 - Any thread can do I/O on a disk storage-pool volume (called **SsAuxThread** for backup)
 - These threads provide a process or session-oriented view
 - It can be more difficult to identify disk contention issues
 - Windows timing statistics have only about a 15-millisecond granularity

Appendix B. Accessibility features for the IBM Storage Protect product family

Accessibility features assist users who have a disability, such as restricted mobility or limited vision, to use information technology content successfully.

Overview

The IBM Storage Protect family of products includes the following major accessibility features:

- Keyboard-only operation
- Operations that use a screen reader

The IBM Storage Protect family of products uses the latest W3C Standard, WAI-ARIA 1.0 (www.w3.org/TR/wai-aria/), to ensure compliance with US Section 508 and Web Content Accessibility Guidelines (WCAG) 2.0 (www.w3.org/TR/WCAG20/). To take advantage of accessibility features, use the latest release of your screen reader and the latest web browser that is supported by the product.

The product documentation in IBM Documentation is enabled for accessibility.

Keyboard navigation

This product uses standard navigation keys.

Interface information

User interfaces do not have content that flashes 2 - 55 times per second.

Web user interfaces rely on cascading style sheets to render content properly and to provide a usable experience. The application provides an equivalent way for low-vision users to use system display settings, including high-contrast mode. You can control font size by using the device or web browser settings.

Web user interfaces include WAI-ARIA navigational landmarks that you can use to quickly navigate to functional areas in the application.

Vendor software

The IBM Storage Protect product family includes certain vendor software that is not covered under the IBM license agreement. IBM makes no representation about the accessibility features of these products. Contact the vendor for accessibility information about its products.

Related accessibility information

In addition to standard IBM help desk and support websites, IBM has a TTY telephone service for use by deaf or hard of hearing customers to access sales and support services:

TTY service
800-IBM-3383 (800-426-3383)
(within North America)

For more information about the commitment that IBM has to accessibility, see [IBM Accessibility](http://www.ibm.com/able) (www.ibm.com/able).

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Glossary

A glossary is available with terms and definitions for the IBM Storage Protect family of products.

See the [IBM Storage Protect glossary](#).

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