



# IBM Directory Server Version 5.1 Performance Tuning Guide





# IBM Directory Server Version 5.1 Performance Tuning Guide

**First Edition (November, 2002)**

This edition applies to version 5, release 1, of The IBM® Directory and to all subsequent releases and modifications until otherwise indicated in new editions.

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## Preface

Welcome to the IBM Directory Server Performance Tuning Guide. The purpose of this document is to provide performance tuning information for IBM Directory Server. The document is divided into sections dealing with LDAP server, IBM Universal Database 2 (DB2®), AIX® operating system, and hardware tuning issues.

For the most current and accurate tuning information, see the Web version of the Performance Tuning Guide on the IBM Directory Web site:

<http://www-4.ibm.com/software/network/directory/library>

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## Who should read this guide

The target audience for this guide includes:

- System installation and deployment administrators
- Network system administrators
- Information Technology architects
- Application developers

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## Typeface conventions

This guide uses several typeface conventions for special terms and actions. These conventions have the following meaning:

<b>Bold</b>	Command names and options, keywords, and other information that you must type as shown.
<i>Italics</i>	Variables and values you must provide appear in italics.
Monospace	Code examples, command lines, screen output, file names, programming keywords, message text or prompts addressed to the user, and text that the user must enter appear in monospace font.

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## Acronyms used in this document

- ACL – Access Control List
- DB2 – DB2 Universal Database™
- LDAP – Lightweight Directory Access Protocol
- SMP – Symmetric Multi-Processor





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## Chapter 1. IBM Directory tuning general overview

This guide provides tuning information for the IBM Directory Server and related IBM Database 2 (DB2) database. The IBM Directory Server is a Lightweight Directory Access Protocol (LDAP) directory that provides a layer on top of DB2, allowing users to efficiently organize, manipulate and retrieve data stored in the DB2 database. Tuning for optimal performance is primarily a matter of adjusting the relationships between the LDAP server and DB2 according to the nature of your workload.

Because each workload is different, instead of providing exact values for tuning settings, we have provided, where appropriate, guidelines for how to determine the best settings for your system. The measurements in this guide use the DirectoryMark (DirMark) 1.2 industry-standard benchmark provided by Mindcraft Inc., but the processes outlined can be applied to any workload. See Appendix A, "DirMark 1.2" on page 33 and Appendix B, "Platform configurations" on page 35 for more information about DirMark 1.2 and the platform configurations used in the guide's examples.

For the most current and accurate tuning information, see the Web version of the Performance Tuning Guide on the IBM Directory Server Web site:

<http://www-4.ibm.com/software/network/directory/library>

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### IBM Directory Server 5.1 application components

The following figure illustrates how IBM Directory Server components interact with each other. Tuning these components can result in improved performance.

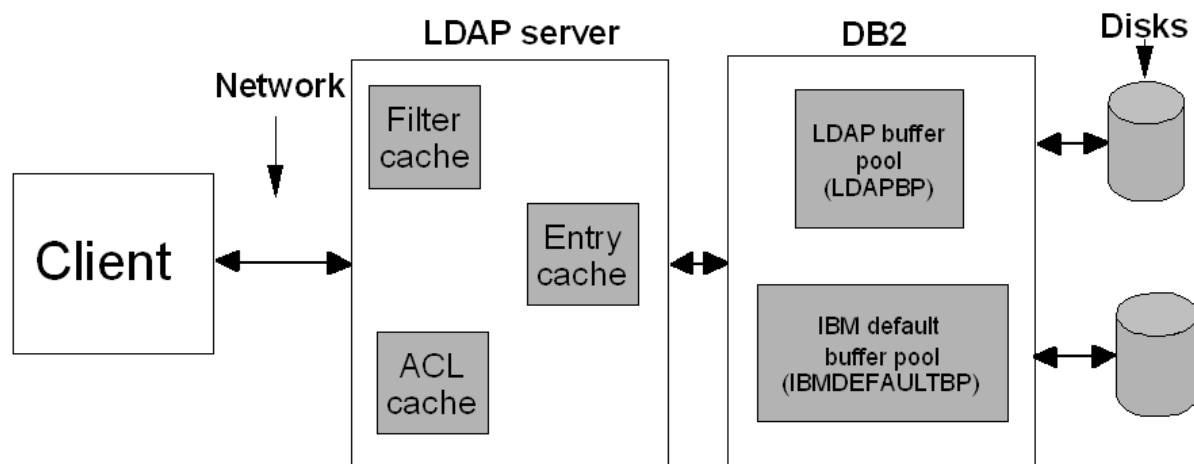


Figure 1. IBM Directory Server 5.1

The arrows in Figure 1 represent the path of a query issued from a client machine. The query follows a path from the IBM Directory Server client to the LDAP server, to DB2, to the physical disks in search of entries that match the query's search filter settings. The shorter the path to matching entries, the better overall performance you can expect from your system.

For example, if a query locates all the matching entries in the LDAP Server, access to DB2 and the disks is not necessary. If the matching entries are not found in the LDAP Server, the query continues on to DB2, and, if necessary to the physical disks as a last resort. Because of the time and resources it takes to retrieve data from disk, it is better from a performance standpoint to allocate a significant amount of memory to the LDAP server caches and DB2 buffer pools.

## LDAP caches and DB2 buffer pools

Caches and buffer pools store previously retrieved data and can significantly improve performance by reducing disk access. When requested data is found within a cache or buffer pool, it is called a cache hit. A cache miss occurs when requested data is not located in a cache or buffer pool.

Because the type of information in each cache and buffer pool differs, it is useful to understand why and when each cache is accessed.

### LDAP caches

There are three LDAP caches:

#### Filter cache

The filter cache contains cached entry IDs. When the client issues a query for some data, the first place the query goes is the filter cache. There are two things that can happen when a query arrives at the filter cache:

- **The IDs that match the filter settings used in the query are located in the filter cache.** If this is the case, the list of the matching entry IDs is sent to the entry cache.
- **The filter cache does not contain the matching entry IDs.** In this case, the query must access DB2 in search of the desired data.

See “LDAP filter cache” on page 5 for more information.

#### Entry cache

The entry cache contains cached entry data. Entry IDs are sent to the entry cache. If the entries that match the entry IDs are in the entry cache, then the results are returned to the client. If the entry cache does not contain the entries that correspond to the entry IDs, the query goes to DB2 in search of the matching entries.

See “Entry cache” on page 8 for more information.

#### ACL cache

The Access Control List (ACL) cache contains information about the access permissions of recently queried entries, such as the entry owner and whether the entry’s permissions are explicit or inherited. Having this information cached in memory can speed up the process of determining whether the user who submitted the query is authorized to see all, some, or none of its results.

### DB2 buffer pools

There are two DB2 buffer pools:

#### LDAPBP

LDAPBP contains cached entry data (ldap\_entry) and all of the associated indexes. LDAPBP is similar to the entry cache, except that LDAPBP uses different algorithms in determining what entries are cached. It is possible that an entry that is not cached in the entry cache is located in LDAPBP. If the requested data is not found in the entry cache or LDAPBP, the query must access the physical disks.

See “LDAPBP buffer pool size” on page 16 for more information.

#### **IBMDEFAULTBP**

DB2 system information, including system tables and other LDAP information, is cached in the IBMDEFAULTBP. You might need to adjust the IBMDEFAULTBP cache settings for better performance. See “IBMDEFAULTBP buffer pool size” on page 17 for more information.

### **Memory allocation between LDAP caches and buffer pools**

The LDAP caches are generally more efficient as a means of caching LDAP searches; however, parts of the LDAP cache get invalidated on updates and must be reloaded before performance benefits return. Some experimentation between the two caching schemes is required to find the best memory allocation for your workload. See “Filter cache bypass limits” on page 7 for more information.

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## **IBM Directory Server tuning overview**

Tuning the LDAP server can significantly improve performance by storing useful data in the caches. It is important to remember, however, that tuning the LDAP server alone is insufficient. Some tuning of DB2 is also required for optimal performance.

The most significant performance tunings related to the IBM Directory Server involve the LDAP caches. LDAP caches are fast storage buffers in memory used to store LDAP information such as queries, answers, and user authentication for future use. While LDAP caches are mostly useful for applications that frequently retrieve repeated cached information, they can greatly improve performance by avoiding calls to the database. See “LDAP caches” on page 5 for information about how to tune the LDAP caches.

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## **DB2 tuning overview**

DB2 serves as the data storage component of the IBM Directory Server. Tuning DB2 results in overall improved performance.

This guide contains several recommendations for tuning DB2, but the most commonly tuned items are:

- **DB2 buffer pools** – Buffer pools are DB2 data caches. Each buffer pool is a data cache between the applications and the physical database files. Adjusting the size of the DB2 buffer pools can result in improved performance. See “DB2 buffer pool tuning” on page 15 for information about buffer pool tuning.
- **Optimization and organization** – After initially loading a directory, or after a number of updates have been performed, it is very important to update database statistics and organization for DB2 to perform optimally. See “Optimization and organization (reorgchk and reorg)” on page 18 for more information.
- **Indexes** –Indexes can make locating data on disk very fast, providing a significant boost to performance. For information about how to create indexes, see “Indexes” on page 23.

**Attention:** You must place the DB2 log on a physical disk drive separate from the data. While there might be some performance benefit to having the DB2 log and data on different drives, data-integrity concerns require the separation. Use the following command to set the path to the DB2 log file directory:

```
UPDATE DATABASE CONFIGURATION FOR database_alias USING NEWLOGPATH path
```

Be sure the database instance owner has write access to the specified path or the command fails.

---

## Generic LDAP application tips

The following are some generic tips that can help improve performance:

- Perform searches on indexed attributes only. See “Indexes” on page 23 for instructions on how to define and verify indexes for IBM Directory.
- Open a connection only once and reuse it for many operations if possible.
- Minimize the number of searches by retrieving multiple attribute values at one time.
- Retrieve only the attributes you need. Do not use ALL by default. For example, when you search for the groups a user belongs to, ask for just the Distinguished Names (DNs), and not the entire group.
- Minimize and batch updates (add, modify, modrdn, delete) when possible.

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## Chapter 2. IBM Directory Server tuning

This chapter discusses the following performance tuning tasks for the IBM Directory Server:

- Tuning LDAP caches
- Determining how directory size affects performance

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### LDAP caches

LDAP caches are fast storage buffers in memory used to store LDAP information such as queries, answers, and user authentication for future use. Tuning the LDAP caches is crucial to improving performance.

An LDAP search that accesses the LDAP cache can be faster than one that requires a connection to DB2, even if the information is cached in DB2. For this reason, tuning LDAP caches can improve performance by avoiding calls to the database. The LDAP caches are especially useful for applications which frequently retrieve repeated cached information. See Figure 1 on page 1 for an illustration of the LDAP caches.

The following sections discuss each of the LDAP caches and demonstrate how to determine and set the best cache settings for your system. The examples used in the tables and measurements are based on DirMark 1.2 data. Keep in mind that every workload is different, and some experimentation will likely be required in order to find the best settings for your workload.

#### LDAP filter cache

When the client issues a query for some data, the first place the query goes is the filter cache. This cache contains cached entry IDs. There are two things that can happen when a query arrives at the filter cache:

- **The IDs that match the filter settings used in the query are located in the filter cache.** If this is the case, the list of the matching entry IDs is sent to the entry cache.
- **The matching entry IDs are not cached in the filter cache.** In this case, the query must access DB2 in search of the desired data.

#### Filter cache size

To determine how big your filter cache should be, run your workload with the filter cache set to different values and measure the differences in operations per second. For example, Figure 2 on page 6 shows the varying operations per second based on different filter cache sizes:

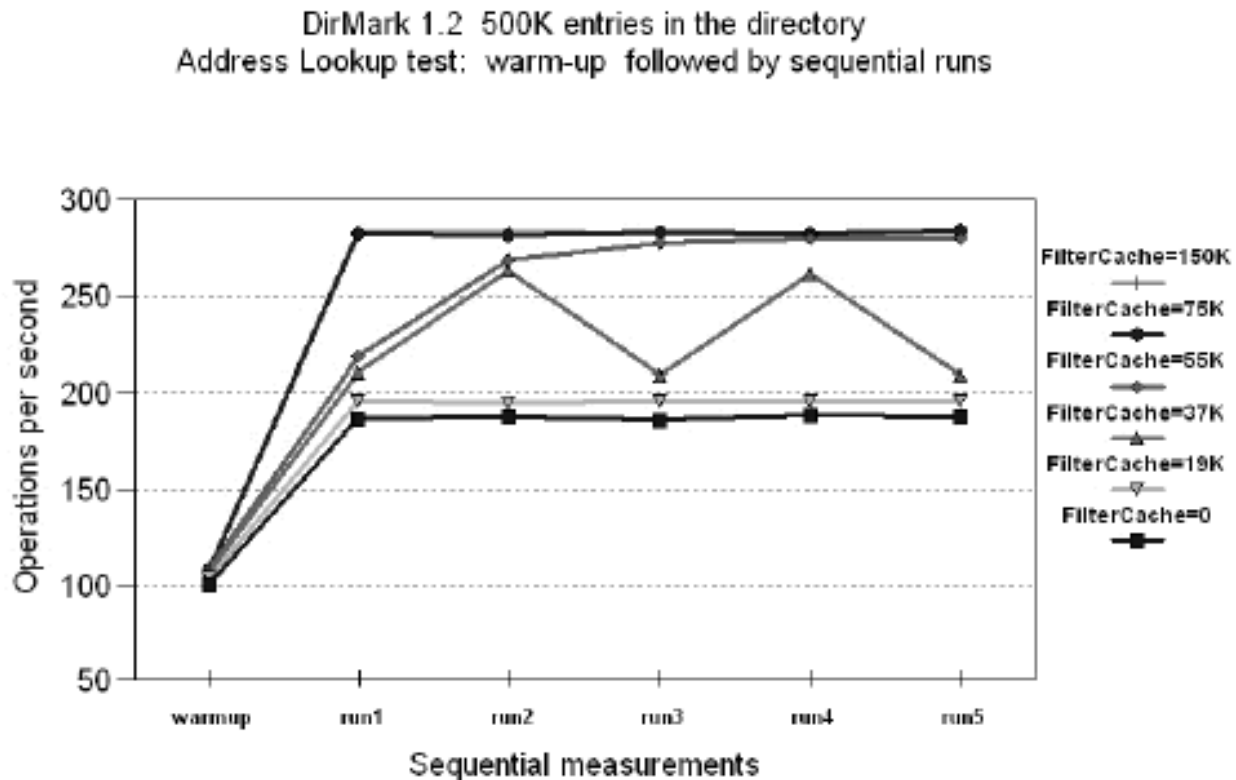


Figure 2. Varying the size of the filter cache

From the results in Figure 2, it can be concluded that for this workload, there is little benefit to a filter cache size of 20,000 over no filter cache at all. Sizes in the 20,000–50,000 range result in unpredictable varying behaviors as some queries hit in the cache while others miss.

For this workload it appears that a filter cache large enough to hold 75,000 entries results in the best performance. There is no benefit making the filter cache any larger than this. The following table shows that setting the filter cache at 75,000 results in 50% more operations per second than if the filter cache was set to zero:

Table 1. Effects of varying the size of the filter cache

Filter cache size	75,000	0
Filter cache bypass limit	100	100
Entry cache size	460,000	460,000
Address Lookup (operations per second)	283	188
Address Lookup (operations per second)	108	101

See “LDAP cache configuration variables” on page 10 to set the filter cache size.

### Filter cache size with updates

Figure 3 on page 7 shows that there is no performance benefit in allocating any memory to the filter cache if even a small fraction of the operations in the workload are updates.

If this proves to be the case for your workload, the only way to retain the performance advantage of a filter cache when updates are involved is to batch your updates. This allows long intervals during which there are only searches. If you cannot batch updates, specify the filter cache size of zero and allot more memory to other caches and buffer pools. See “LDAP cache configuration variables” on page 10 for instructions on how to set configuration variables such as filter cache size.

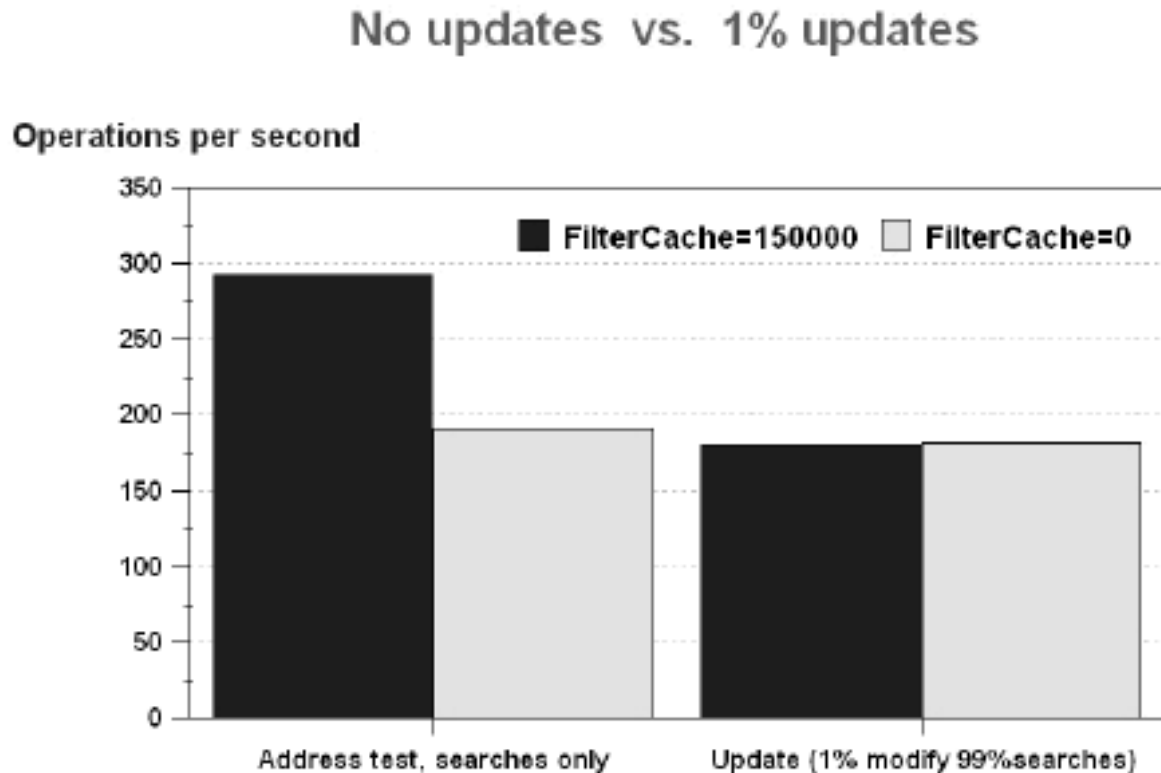


Figure 3. Effects of updates on the performance of the filter cache

#### Filter cache bypass limits

The filter cache bypass limit configuration variable limits the number of entries that can be added to the filter cache. For example, if the bypass limit variable is set to 1,000, search filters that match more than 1,000 entries are not added to the filter cache. This prevents large, uncommon searches from overwriting useful cache entries. To determine the best filter cache bypass limit for your workload, run your workload repeatedly and measure the throughput.

For example, Figure 4 on page 8 shows the operations per second based on varying cache bypass limit sizes.

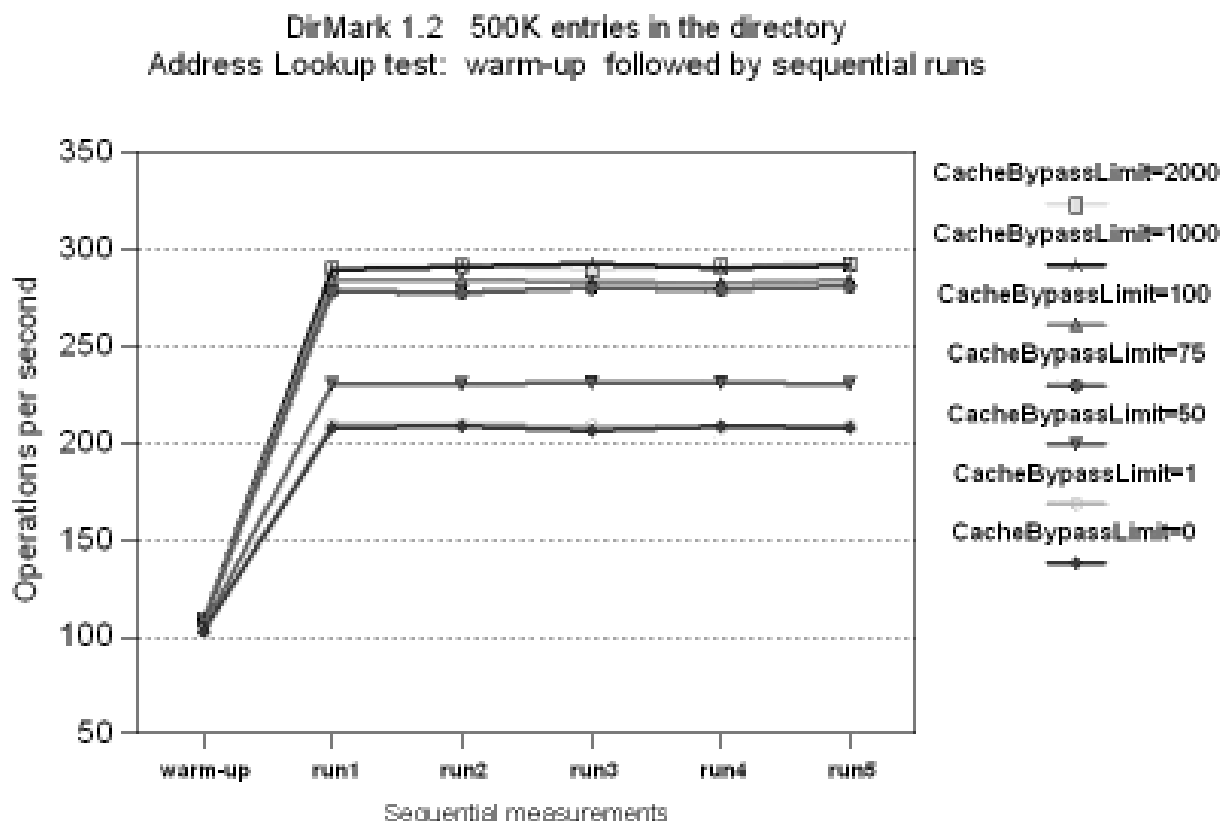


Figure 4. Varying the filter cache bypass limit

For the workload in Figure 4, setting the limit too low hurts performance by preventing valuable filters from being cached. Setting the filter bypass limit to approximately 100 appears to be the best size for this workload. Setting it any larger benefits performance only slightly.

See “LDAP cache configuration variables” on page 10 to set the filter cache bypass limit.

## Entry cache

The entry cache contains cached entry data. Entry IDs are sent to the entry cache. If the entries that match the entry IDs are in the entry cache, then the results are returned to the client. If the entry cache does not contain the entries that correspond to the entry IDs, the query goes to DB2 in search of the matching entries.

### Entry cache size

To determine how big your entry cache should be, run your workload with the entry cache set to different sizes and measure the differences in operations per second. For example, Figure 5 on page 9 shows the varying operation per second based on different entry cache sizes:



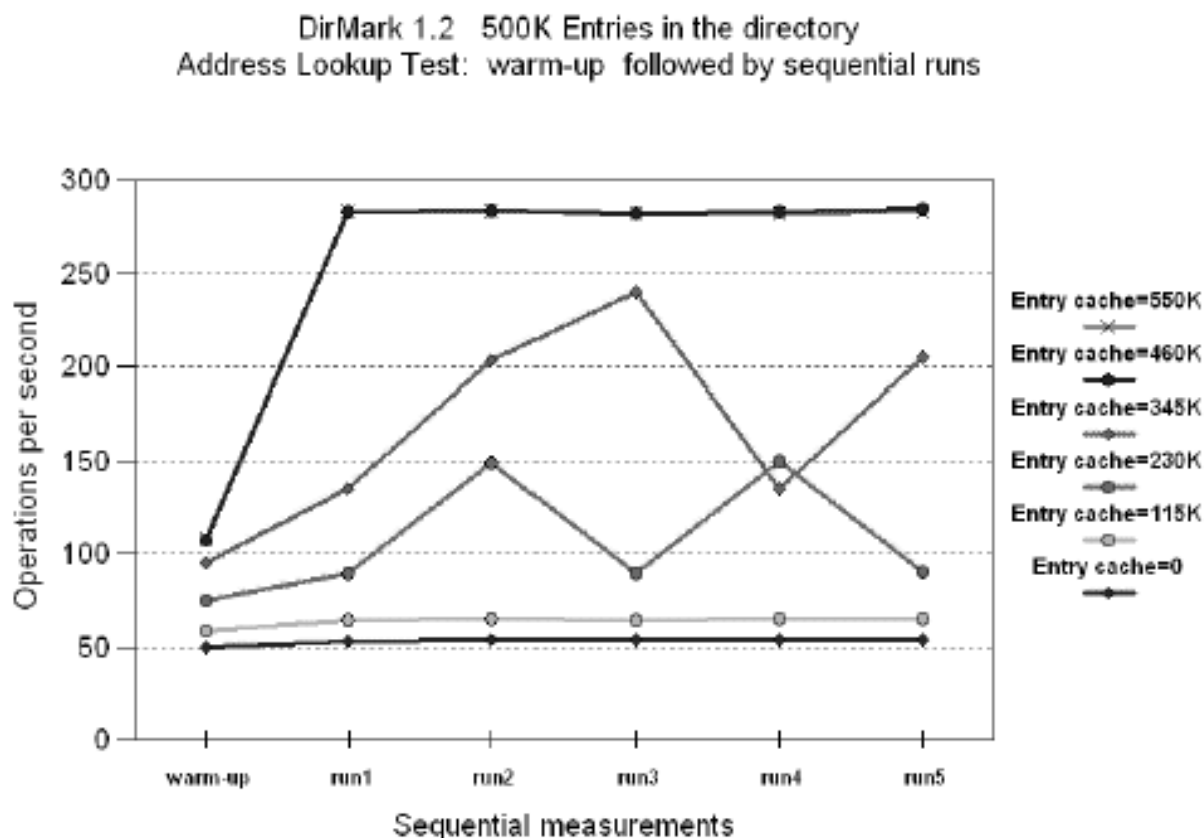


Figure 5. Varying the size of the entry cache

From the results in Figure 5, it can be concluded that for this workload, there is little benefit to an entry cache size of 100,000 over no entry cache at all. Sizes in the 200,000–400,000 range result in unpredictable varying behaviors as some queries hit in the cache while others miss.

For this workload it appears that an entry cache large enough to hold 460,000 entries results in the best performance. There is no benefit to making the entry cache any larger than this. Setting the entry cache at 460,000 results in 5 times as many operations per second than if entry cache was set to zero:

Table 2. Effects of varying the size of the entry cache

Entry cache size	460,000	0
Filter cache size	100	100
Filter cache bypass limit	150,000	150,000
Address Lookup (operations per second)	284	54
Address Lookup warmup (operations per second)	107	50

To find the best cache size for your workload, you must run your workload with different cache sizes. See “LDAP cache configuration variables” on page 10 to set the filter cache size.

## Entry cache bypass limits

When the entry cache bypass limit is set to any value, the filter cache bypass limit applies to the entry cache as well.

---

## Cache entry sizes in megabytes

Cache sizes are measured by number of entries. When determining how much memory to allocate to your LDAP caches, it can be useful to know how big the entries in your cache are in megabytes.

The following example shows how to measure cached entries in megabytes:

**Note:** This example calculates the average size in megabytes of an entry in a sample entry cache, but the average filter cache entry size can be calculated similarly.

1. From the LDAP server:
  - a. Set the filter cache size to zero.
  - b. Set the entry cache size to a small value; for example, 200.
  - c. Start **ibmslapd**.
2. From the client:
  - a. Run your application
  - b. Find the entry cache population (call this *population1*) using the following command:

```
ldapsearch -h servername -s base -b cn=monitor objectclass=* | grep entry_cache_current
```
3. From the LDAP Server:
  - a. Find the memory used by **ibmslapd** (call this *ibmslapd1*):
    - On AIX operating systems, use `ps v`
    - On Windows® operating systems, use the **VM size** column in the **Task Manager**.
  - b. Stop **ibmslapd**.
  - c. Increase the size of the entry cache but keep it smaller than your working set.
  - d. Start **ibmslapd**.
4. Run your application again and find the entry cache population (call this *population2*). See 2b for the command syntax.
5. Find the memory used by **ibmslapd** (call this *ibmslapd2*). See 3a for the command syntax.
6. Calculate the size of an entry cache entry using the following formula:
$$\frac{(\text{ibmslapd size2} - \text{ibmslapd size1})}{(\text{entry cache population2} - \text{entry cache population1})}$$

For example, using this formula with the 500,000-entry database used by DirMark 1.2 results in the following measurement:

$$(192084\text{KB} - 51736\text{KB}) / (48485 - 10003) = 3.65\text{KB per entry}$$

---

## LDAP cache configuration variables

LDAP cache configuration variables allow you to set the LDAP cache sizes, bypass limits, and other variables that affect performance.

## Setting LDAP configuration variables

You can set LDAP configuration variables using the Web Administration Tool or the command line.

To set LDAP configuration variables using the Web Administration Tool:

1. Expand the **Manage server properties** category in the navigation area of the Web Administration tool.
2. Click **Performance**.
3. You can modify any of the following configuration variables
  - **Cache ACL information** — This option must be selected in order for the **Maximum number of elements in ACL cache** settings to take effect.
  - **Maximum number of elements in ACL cache** (ACL cache size) — The default is 25,000.
  - **Maximum number of elements in entry cache** (entry cache size) — Specify the maximum number of elements in the entry cache. The default is 25,000. See “Entry cache” on page 8 for more information about the entry cache.
  - **Maximum number of elements in search filter cache** (filter cache size)— The search filter cache consists of the requested search filters and resulting entry identifiers that matched. On an update operation, all filter cache entries are invalidated. The default is 25,000. See “LDAP filter cache” on page 5 for more information about the filter cache.
  - **Maximum number of elements from a single search added to search filter cache** (filter cache bypass limit) — If you select **Elements**, you must enter a number. The default is 100. Otherwise select **Unlimited**. Search filters that match more entries than the number specified here are not added to the search filter cache. See “Filter cache bypass limits” on page 7 for more information about bypass limits.
4. When you are finished, click **OK** to apply your changes, or click **Cancel** to exit the panel.

To set LDAP configuration variables using the command line, issue the following command:

```
ldapmodify -DAdminDN -wAdminpassword  
-ifilename
```

where the file *filename* contains:

```
dn: cn=Directory,cn=RDBM Backends,cn=IBM Directory,cn=Schemas,cn=Configuration  
changetype: modify  
replace: ibm-slapdDbConnections  
ibm-slapdDbConnections: 15  
  
dn: cn=Front End, cn=Configuration  
changetype: modify  
replace: ibm-slapdACLCache  
ibm-slapdACLCache: TRUE  
-  
replace: ibm-slapdACLCacheSize  
ibm-slapdACLCacheSize: 25000  
-  
replace: ibm-slapdEntryCacheSize  
ibm-slapdEntryCacheSize: 25000  
-
```

```
replace: ibm-slapdFilterCacheSize
ibm-slapdFilterCacheSize: 25000
-
replace: ibm-slapdFilterCacheBypassLimit
ibm-slapdFilterCacheBypassLimit: 100
```

There are several additional settings that affect performance by putting limits on client activity, minimizing the impact to server throughput and resource usage, such as:

- ibm-slapdSizeLimit: 500
- ibm-slapdTimeLimit: 900
- ibm-slapdIdleTimeOut: 300
- ibm-slapdMaxEventsPerConnection: 100
- ibm-slapdMaxEventsTotal: 0
- ibm-slapdMaxNumOfTransactions: 20
- ibm-slapdMaxOpPerTransaction: 5
- ibm-slapdMaxTimeLimitOfTransactions: 300
- ibm-slapdPagedResAllowNonAdmin: TRUE
- ibm-slapdPagedResLmt: 3
- ibm-slapdPageSizeLmt: 201
- ibm-slapdSortKeyLimit: 3
- ibm-slapdSortSrchAllowNonAdmin: TRUE

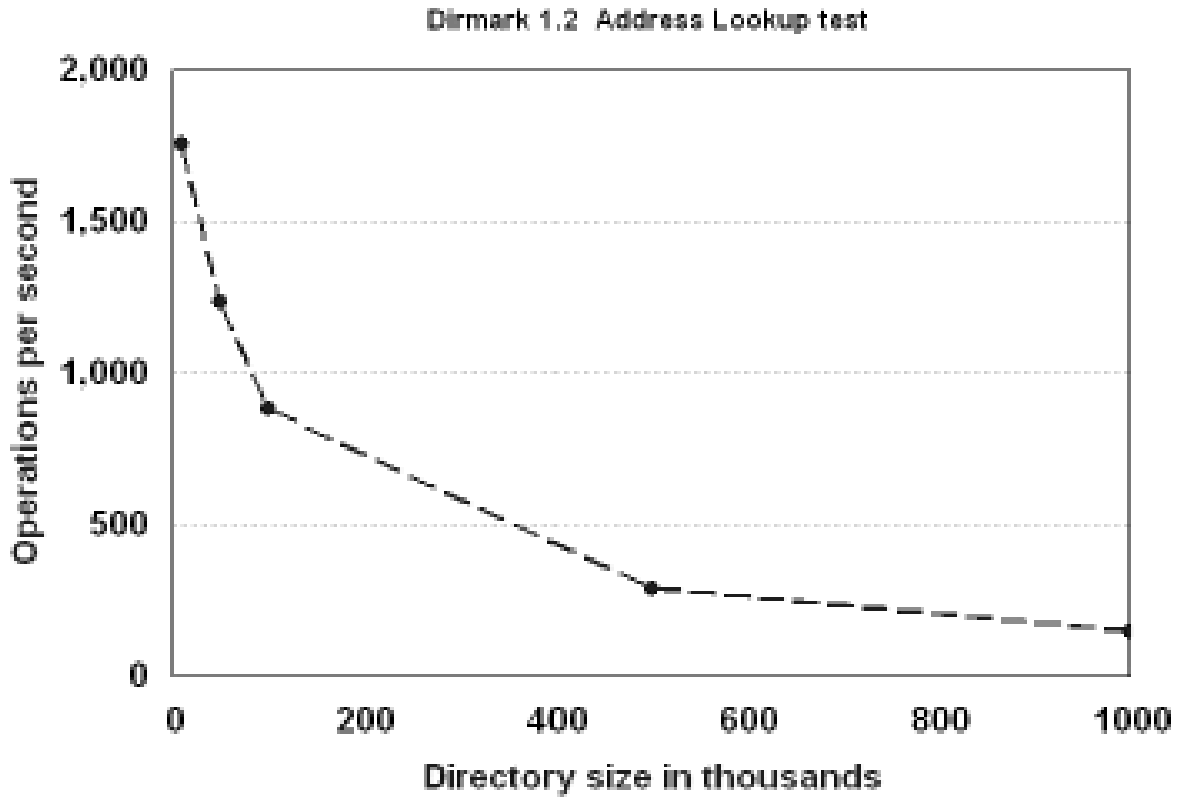
**Note:** Default values are shown.

---

## Directory size

It is possible to measure how the size of your directory impacts performance by running your workload with different directory sizes. For example, if your current directory contains 500,000 entries, try tuning your workload with a 100,000 entry directory and a 1,000,000 entry directory.

It is important when you run your workload that you consider several measurements. For example, measuring the number of operations per second as shown in Figure 6 on page 13, it appears that performance degrades significantly as the the database size grows.



*Figure 6. Operations per second*

However, the DirMark Address Lookup test includes a large fraction of wildcard searches and exact-match searches, such as "(sn=Smith)" that return all entries where the sn value is "Smith". Both of these types of searches typically return multiple entries in response to a single search request. As Figure 7 on page 14 shows, as the size of the directory grows, so does the number of entries returned in response to wildcard and exact-match search requests.

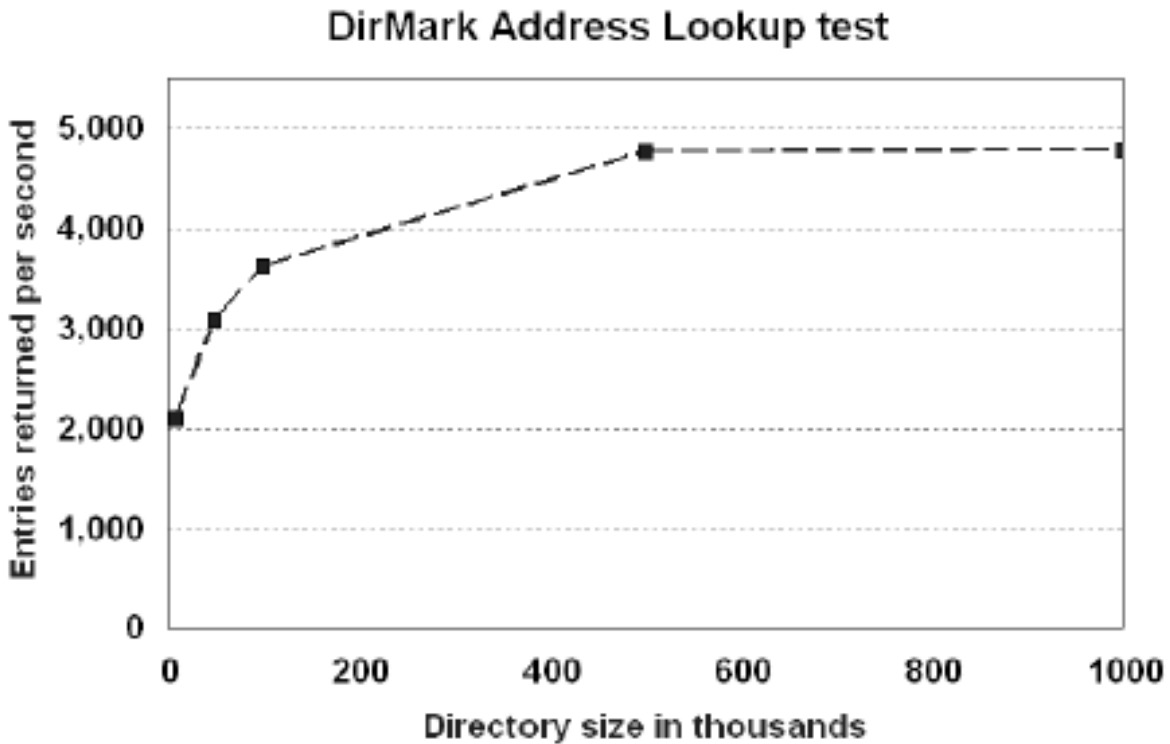


Figure 7. Entries returned per second

In this situation, the number of entries returned per second is a truer measure of throughput than operations per second, because each operation requires more work to be performed as the size of the database grows.

**Note:** As your directory grows, it might become necessary to readjust the sizes of the LDAP caches and DB2 buffer pools. You can determine the optimal sizes for your caches and buffer pools using the guidelines in “LDAP caches” on page 5 and “DB2 buffer pool tuning” on page 15.

---

## Chapter 3. DB2 tuning

The IBM Directory uses DB2 as the data store and Structured Query Language (SQL) as the query retrieval mechanism. While the LDAP server caches LDAP queries, answers, and authentication information, DB2 caches tables, indexes, and statements.

Many DB2 configuration parameters affect either the memory (buffer pools) or disk resources. Since disk access is usually much slower than memory access, the key database performance tuning objective is to decrease the amount of disk activity.

This chapter covers the following DB2 tunings:

- DB2 buffer pool tuning
- Optimization and organization (**reorgchk** and **reorg**)
- Other DB2 configuration parameters
- **backup** and **restore**

**Attention:** Only users listed as database administrators can execute the DB2 commands. Make sure the ID running the DB2 commands is a user in the `dbsysadm` group (UNIX® operating systems) or a member of the Administrator group (Windows operating systems.) This includes the DB2 instance owner and root.

If you have any trouble running the DB2 commands, check to ensure the DB2 environment variables have been established by running **db2profile** (if not, the **db2 get** and **db2 update** commands will not work). The script file **db2profile** is located in the `sqllib` subdirectory under the instance owner's home directory. If you need to tailor this file, follow the comments inside the file to set your instance name, user paths, and default database name (the default path is `/home/ldapdb2/sqllib/db2profile`.) It is assumed that the user is logged in as **ibm-slapdDbUserId**. If logged in as the root user on a UNIX operating system, it is possible to switch to the instance owner as follows:

```
su - instance owner
```

where *instance owner* is the defined owner of the LDAP database.

To log on as the database administrator on a Windows 2000 operating system, run the following command:

```
runas /user:instance owner db2cm
```

where *instance owner* is the defined owner of the LDAP database.

For additional stability and performance enhancements, upgrade to the latest version of DB2.

---

### DB2 buffer pool tuning

A buffer pool is a data cache between LDAP and the physical DB2 database files for both tables and indexes. DB2 buffer pools are searched when entries and their attributes are not found in the entry cache.

There are several considerations to keep in mind when tuning the DB2 buffer pools, for example:

- If there are no buffer pools, all database activity results in disk access.
- If the size of each buffer pool is too small, LDAP must wait for DB2 disk activity to satisfy DB2 SQL requests.
- If one or more buffer pools is too large, memory on the LDAP server might be wasted.
- If the total amount of space used by both buffer pools is larger than physical memory available on the server, operating system paging (disk activity) will occur.

Determining the settings of your DB2 buffer pools is one of the most significant DB2 performance tunings. Buffer pool tuning typically needs to be done only once.

To get the current DB2 buffer pool sizes, run the following commands:

```
db2 connect to instance owner
db2 "select bpname,npages,pagesize from SYSIBM.SYSbufferpools"
```

where *instance owner* is the defined owner of the database.

The following example output shows the default settings for the example above:

BPNAME	NPAGES	PAGESIZE
IBMDEFAULTBP	29500	4096
LDAPBP	1230	32768

2 record(s) selected.

## Buffer pool sizes

In IBM Directory Server 5.1, the LDAP directory database (DB2) has two buffer pools, LDAPBP and IBMDEFAULTBP. The size of each buffer pool needs to be set separately, but the method for determining how big each should be is the same: Run your workload with the buffer pool sizes set to different values and measure the differences in operations per second.

**Note:** DB2 does not allow buffer pools to be set to zero.

### LDAPBP buffer pool size

This buffer pool contains cached entry data (ldap\_entry) and all of the associated indexes. LDAPBP is similar to the entry cache, except that LDAPBP uses different algorithms in determining what entries are cached. It is possible that an entry that is not cached in the entry cache is located in LDAPBP.

To determine the best size for your LDAPBP buffer pool, run your workload with LDAPBP buffer pool size set to different values and measure the differences in operations per second. For example, Figure 8 on page 17 shows the varying operations per second based on different LDAPBP buffer pool sizes:



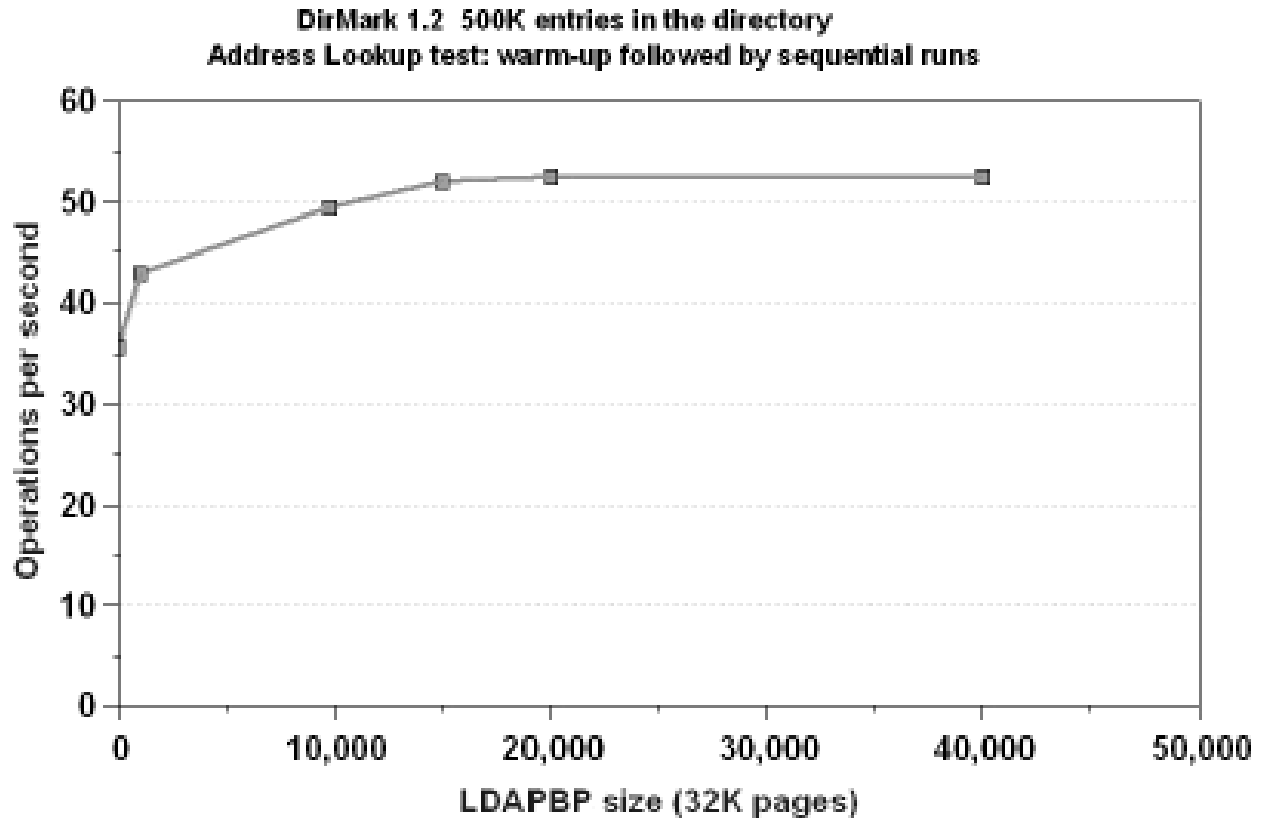


Figure 8. Varying the size of LDAPBP

For the workload in the above example, the best performance results from a LDAPBP size of approximately 15,000 32K pages. However, the performance gain of 15,000 over a size of 9,800 is slight. In a memory-constrained environment, setting the LDAPBP size to 9,800 saves approximately 166MB of memory.

#### **IBMDEFAULTBP buffer pool size**

DB2 system information, including system tables and other LDAP information, is cached in the IBMDEFAULTBP. Although no LDAP data is contained in the IBMDEFAULTBP, you may need to adjust the IBMDEFAULTBP cache settings for better performance in the LDAPBP.

To determine the best size for your IBMDEFAULTBP buffer pool, run your workload with the buffer pool sizes set to different values and measure the differences in operations per second. For example, Figure 9 on page 18 shows the varying operations per second based on different IBMDEFAULTBP buffer pool sizes:

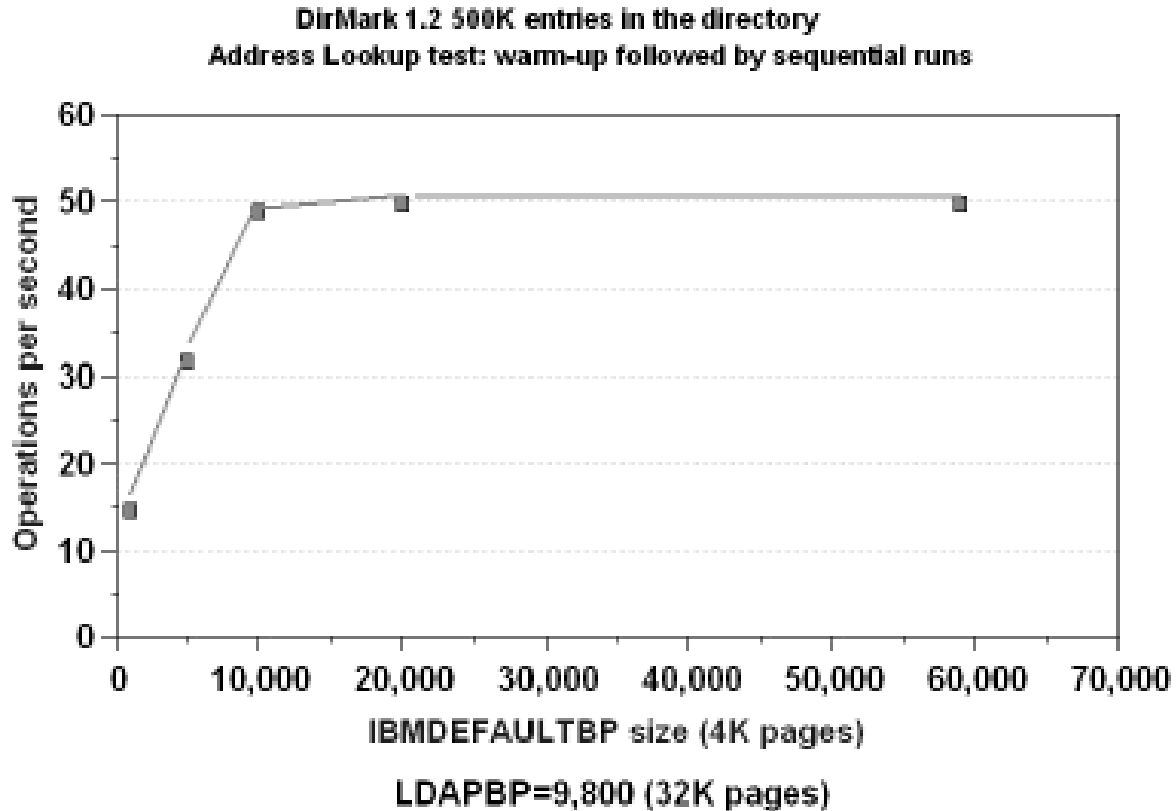


Figure 9. Varying the size of IBMDEFAULTBP

For the workload in the above example, setting the IBMDEFAULTBP large enough to hold the working set improves throughput more than three times over a small buffer pool size. There is little additional benefit to setting IBMDEFAULTBP larger than 20,000 4K pages.

### Setting buffer pool sizes

Use the `alter bufferpool` command to set the IBMDEFAULTBP and LDAPBP buffer pool sizes. The following example shows the IBMDEFAULTBP and LDAPBP buffer pools being set:

```
db2 alter bufferpool ibmdefaultbp size 20000
db2 alter bufferpool ldapbp size 9800
db2 force applications all
db2stop
db2start
```

## Optimization and organization (reorgchk and reorg)

DB2 uses a sophisticated set of algorithms to optimize the access to data stored in a database. These algorithms depend upon many factors, including the organization of the data in the database, and the distribution of that data in each table. Distribution of data is represented by a set of statistics maintained by the database manager.

In addition, IBM Directory Server creates a number of indexes for tables in the database. These indexes are used to minimize the data accessed in order to locate a particular row in a table.

In a read-only environment, the distribution of the data changes very little. However, with updates and additions to the database, it is not uncommon for the distribution of the data to change significantly. Similarly, it is quite possible for data in tables to become ordered in an inefficient manner.

To remedy these situations, DB2 provides tools to help optimize the access to data by updating the statistics and to reorganize the data within the tables of the database.

## Optimization

Optimizing the database updates statistics related to the data tables, which improves performance and query speed. Optimize the database periodically or after heavy database updates, for example, after importing database entries. The IBM Directory Server **Optimize** button uses DB2 runstats to update statistical information used by the query optimizer for all the LDAP tables.

**Note:** **reorgchk** also updates statistics. If you are planning to do a **reorgchk**, optimizing the database is unnecessary. See “Database organization (reorgchk and reorg)” for more information about **reorgchk**.

To optimize the database using the Configuration Tool:

1. Start the Configuration Tool by typing **ldapxcfg** on the command line.
2. Click **Optimize database**.

After a message displays indicating the database was successfully optimized, you must restart the server for the changes to take effect.

To optimize the database using the command line, run the following commands:

```
DB2 RUNSTATS ON TABLE table-name AND DETAILED INDEXES ALL SHRLEVEL REFERENCE
```

Run the following commands for more detailed lists of runstats that improve performance:

```
DB2 RUNSTATS ON TABLE table-name WITH DISTRIBUTION AND DETAILED INDEXES ALL SHRLEVEL REFERENCE
```

```
DB2 RUNSTATS ON TABLE ldapdb2.objectclass WITH DISTRIBUTION AND DETAILED INDEXES ALL SHRLEVEL REFERENCE
```

## Database organization (reorgchk and reorg)

Tuning the organization of the data in DB2 using the **reorgchk** and **reorg** commands is important for optimal performance.

**reorgchk** updates statistical information to the DB2 optimizer to improve performance, and reports statistics on the organization of the database tables.

**reorg**, using the data generated by **reorgchk**, reorganizes table spaces to improve access performance and reorganizes indexes so that they are more efficiently clustered. The **reorgchk** and **reorg** commands can improve both search and update operation performance.

**Note:** Tuning organizes the data on disk in a sorted order. Sorting the data on disk is beneficial only when accesses occur in a sorted order, which is not typically the case. For this reason, organizing the table data on disk typically yields little change in performance.

## Performing a reorgchk

After a number of updates have been performed against DB2, table indexes can become sub-optimal and performance can degrade. Correct this situation by performing a DB2 **reorgchk** as follows:

```
db2 connect to ldapdb2
db2 reorgchk update statistics on table all
```

Where *ldapdb2* is the name of your database.

To generate a **reorgchk** output file (recommended if you plan to run the **reorg** command) add the name of the file to the end of the command, for example:

```
db2 reorgchk update statistics on table all >reorgchk.out
```

A sample **reorgchk** report looks something like this:

```
db2 => reorgchk current statistics on table all
```

Table statistics:

```
F1: 100 * OVERFLOW / CARD < 5
F2: 100 * TSIZE / ((FPAGES-1) * (TABLEPAGESIZE-76)) > 70
F3: 100 * NPAGES / FPAGES > 80
```

CREATOR	NAME	CARD	OV	NP	FP	TSIZE	F1	F2	F3	REORG
-----										
LDAPDB2	ACLPERM	2	0	1	1	138	0	-	100	---
LDAPDB2	ACLPROP	2	0	1	1	40	0	-	100	---
LDAPDB2	ALIASEDOBJECT	-	-	-	-	-	-	-	-	---
LDAPDB2	AUDIT	1	0	1	1	18	0	-	100	---
LDAPDB2	AUDITADD	1	0	1	1	18	0	-	100	---
LDAPDB2	AUDITBIND	1	0	1	1	18	0	-	100	---
LDAPDB2	AUDITDELETE	1	0	1	1	18	0	-	100	---
LDAPDB2	AUDITEXTOPEVENT	1	0	1	1	18	0	-	100	---
LDAPDB2	AUDITFAILEDOPONLY	1	0	1	1	18	0	-	100	---
LDAPDB2	AUDITLOG	1	0	1	1	77	0	-	100	---
...										
SYSIBM	SYSINDEXCOLUSE	480	0	6	6	22560	0	100	100	---
SYSIBM	SYSINDEXES	216	114	14	28	162216	52	100	50	***
...										
SYSIBM	SYSPLAN	79	0	6	6	41554	0	100	100	---
SYSIBM	SYSPLANAUTH	157	0	3	3	9106	0	100	100	---
SYSIBM	SYSPLANDEP	35	0	1	2	5985	0	100	50	---
-----										

Index statistics:

F4: CLUSTERRATIO or normalized CLUSTERFACTOR > 80

F5:  $100 * (KEYS * (ISIZE+8) + (CARD-KEYS) * 4) / (NLEAF * INDEXPAGESIZE) > 50$

F6:  $(100-PCTFREE) * (INDEXPAGESIZE-96) / (ISIZE+12) ** (NLEVELS-2) * (INDEXPAGESIZE-96) / (KEYS * (ISIZE+8) + (CARD-KEYS) * 4) < 100$

CREATOR	NAME	CARD	LEAF	LVLS	ISIZE	KEYS	F4	F5	F6	REORG
-----										
Table: LDAPDB2.ACLPERM										
LDAPDB2	ACLPERM_INDEX	2	1	1	6	2	100	-	-	---
Table: LDAPDB2.ACLPROP										
LDAPDB2	ACLPROP_INDEX	2	1	1	6	2	100	-	-	---
Table: LDAPDB2.ALIASEDOBJECT										
LDAPDB2	ALIASEDOBJECT	-	-	-	-	-	-	-	-	---
LDAPDB2	ALIASEDOBJECTI	-	-	-	-	-	-	-	-	---
LDAPDB2	RALIASEDOBJECT	-	-	-	-	-	-	-	-	---
Table: LDAPDB2.AUDIT										
LDAPDB2	AUDITI	1	1	1	4	1	100	-	-	---
Table: LDAPDB2.AUDITADD										
LDAPDB2	AUDITADDI	1	1	1	4	1	100	-	-	---
Table: LDAPDB2.AUDITBIND										
LDAPDB2	AUDITBINDI	1	1	1	4	1	100	-	-	---
Table: LDAPDB2.AUDITDELETE										
LDAPDB2	AUDITDELETEI	1	1	1	4	1	100	-	-	---
Table: LDAPDB2.AUDITEXTOPEVENT										
...										
Table: LDAPDB2.SN										
LDAPDB2	RSN	25012	148	2	14	25012	99	90	0	---
LDAPDB2	SN	25012	200	3	12	25012	99	61	119	--*
LDAPDB2	SNI	25012	84	2	4	25012	99	87	1	---
...										
Table: LDAPDB2.TITLE										
LDAPDB2	TITLEI	-	-	-	-	-	-	-	-	---
Table: LDAPDB2.UID										
LDAPDB2	RUID	25013	243	3	17	25013	0	62	79	*--
LDAPDB2	UID	25013	273	3	17	25013	100	55	79	---
LDAPDB2	UIDI	25013	84	2	4	25012	100	87	1	---
Table: LDAPDB2.UNIQUEMEMBER										
LDAPDB2	RUNIQUEMEMBER	10015	224	3	47	10015	1	60	44	*--
LDAPDB2	UNIQUEMEMBER	10015	284	3	47	10015	100	47	44	--*
LDAPDB2	UNIQUEMEMBERI	10015	14	2	4	7	100	69	8	---
...										
Table: SYSIBM.SYSFUNCTIONS										
SYSIBM	IBM127	141	1	1	13	141	65	-	-	*--
SYSIBM	IBM25	141	2	2	34	141	100	72	60	---
SYSIBM	IBM26	141	2	2	32	141	78	68	63	*--
SYSIBM	IBM27	141	1	1	23	68	80	-	-	*--
SYSIBM	IBM28	141	1	1	12	2	99	-	-	---
SYSIBM	IBM29	141	1	1	4	141	100	-	-	---
SYSIBM	IBM30	141	3	2	59	141	78	76	38	*--
SYSIBM	IBM55	141	2	2	34	141	99	72	60	---
...										
-----										

CLUSTERRATIO or normalized CLUSTERFACTOR (F4) will indicate REORG is necessary for indexes that are not in the same sequence as the base table. When multiple indexes are defined on a table, one or more indexes may be flagged as needing REORG. Specify the most important index for REORG sequencing.

Using the statistics generated by **reorgchk**, run **reorg** to update database table organization. See "Performing a reorg" on page 22.

Keep in mind that **reorgchk** needs to be run periodically. For example, **reorgchk** might need to be run after a large number of updates have been performed. Note that LDAP tools such as **ldapadd**, **ldif2db** and **bulkload** can potentially do large numbers of updates that require a **reorgchk**. The performance of the database should be monitored and a **reorgchk** performed when performance starts to degrade. See “Monitoring performance” on page 29 for more information.

**reorgchk** must be performed on all LDAP replicas, since each uses a separate database. The LDAP replication process does not include the propagation of database optimizations.

In general, reorganizing a table takes more time than running statistics. Therefore, performance might be improved significantly by running statistics first.

Because LDAP caches prepared DB2 statements, you must stop and restart **ibmslapd** in order for DB2 changes to take effect.

### Performing a reorg

After you have generated organizational information about the database using **reorgchk**, the next step in reorganization is finding the tables and indexes that need reorganizing and attempting to reorganize them. This can take a long time. The time it takes to perform the reorganization process increases as the DB2 database size increases.

To reorganize database table information:

1. If you have not done so already, run **reorgchk**:

```
db2 reorgchk update statistics on table all >reorgchk.out
```

The **reorgchk** update statistics report has two sections; the first section is the table information and the second section is the indexes. An asterisk in the last column indicates a need for reorganization.

2. To reorganize the tables with an asterisk in the last column:

```
db2 reorg table table name
```

where *table name* is the name of the table to be reorganized, for example, LDAPDB2.LDAP\_ENTRY.

Generally speaking, since most data in LDAP is accessed by index, reorganizing tables is usually not as beneficial as reorganizing indexes.

3. To reorganize the indexes with an asterisk in the last column:

```
db2 reorg table table name index index name
```

where *table name* is the name of the table, for example, LDAPDB2.LDAP\_ENTRY. And where *index name* is the name of the index, for example, SYSIBM.SQL000414155358130.

4. Run **reorgchk** again. The output from **reorgchk** can then be used to determine whether the reorganization worked and whether it introduced other tables and indexes that need reorganizing.

Some guidelines for performing a reorganization are:

- If the number on the column that has an asterisk is close to the recommended value described in the header of each section and one reorganization attempt has already been done, you can probably skip a reorganization on that table or index.

- In the table LDAPDB2.LDAP\_ENTRY there exists a LDAP\_ENTRY\_TRUNC index and a SYSIBM.SQL index. Preference should be given to the SYSIBM.SQL index if attempts to reorganize them seem to alternate between one or the other needing reorganization.
- Reorganize all the attributes that you want to use in searches. In most cases you will want to reorganize to the forward index, but in cases with searches beginning with '\*', reorganize to the reverse index.

For example:

**Table: LDAPDB2.SECUUUID**

```
LDAPDB2 RSECUUID <- This is a reverse index
LDAPDB2 SECUUID <- This is a forward index
LDAPDB2 SECUUIDI <- This is an update index
```

## Indexes

Indexing results in a considerable reduction in the amount of time it takes to locate requested data. For this reason, it can be very beneficial from a performance standpoint to index all attributes used in searches.

Use the following DB2 commands to verify that a particular index is defined. In the following example, the index being checked is for the attribute **principalName**:

```
db2 connect to database name
db2 list tables for all | grep -i principalName
db2 describe indexes for table database name.principalName
```

Where *database name* is the name of your database.

If the second command fails or the last command does not return three entries, the index is not properly defined. The last command should return the following results:

IndexSchema	Index Name	Unique Rule	Number of Columns
-----	-----	-----	-----
LDAPDB2	PRINCIPALNAMEI	D	1
LDAPDB2	PRINCIPALNAME	D	2
LDAPDB2	RPRINCIPALNAME	D	2

3 record (s) selected.

To have IBM Directory Server create an index for an attribute the next time IBM Directory Server is started, do one of the following:

- To create an index using the Web Administration Tool:
  1. Expand **Schema management** in the navigation area, and click **Manage attributes**.
  2. Click **Edit attribute**.
  3. On the **IBM extensions** tab, select the **Equality** checkbox under **Indexing rules**.
- To create an index from the command line, issue the following command:

```
ldapmodify -f /ldap/etc/addindex.ldif
```

The addindex.ldif file should look like this:

```
dn: cn=schema
changetype: modify
replace: attributetypes
```

```

attributetypes: ( 1.3.18.0.2.4.318 NAME ( 'principalName' 'principal' ) DESC
'A naming attribute that may be used to identify eUser object entries.' EQUALITY
1.3.6.1.4.1.1466.109.114.2 ORDERING 2.5.13.3 SUBSTR 2.5.13.4 SYNTAX
1.3.6.1.4.1.1466.115.121.1.15 USAGE userApplications )
-
replace: ibmattributetypes
ibmattributetypes: ( 1.3.18.0.2.4.318 DBNAME( 'principalName' 'principalName' )
ACCESS-CLASS normal LENGTH 256 EQUALITY ORDERING SUBSTR APPROX )

```

---

## Other DB2 configuration parameters

Performance benefits can come from setting other DB2 configuration parameters, such as APPLHEAPZ and LOGFILSIZ. The current setting of parameters can be obtained by issuing the following command:

```
db2 get database configuration for database name
```

where *database name* is the name of your database.

This command returns the settings of other DB2 configuration parameters as well.

The following command also shows the DB2 configuration parameters for the entire database instance:

```
db2 get database manager configuration
```

To set the DB2 configuration parameters use the following syntax:

```

db2 update database configuration for database name using \
parm name parm value
db2stop
db2start

```

where *database name* is the name of your database and where *parm name* is the parameter to change and *parm value* is the value it is to be assigned.

Changes to DB2 configuration parameters do not take effect until the database is restarted with **db2stop** and **db2start**.

For a list of DB2 parameters that affect performance, visit the DB2 Web site:  
<http://www.ibm.com/software/data/db2>

**Note:** If DB2 recognizes that a parameter is configured insufficiently, the problem is posted to the diagnostic log (db2diag.log). For example, if the DB2 buffer pools are too large, DB2 overrides the buffer pool settings and uses a minimal configuration. No notice of the change in buffer pool sizes is given except in the diagnostic log, so it is important to view the log if you are experiencing poor performance.

---

## Database backup and restore considerations

When using the database **backup** and **restore** commands it is important to keep in mind that when you restore over an existing database, any tuning that has been done on that existing database is lost.



---

## Chapter 4. AIX Operating system tuning

This chapter discusses the following performance tuning tasks for the AIX operating system:

- Enabling large files
- Setting MALLOCMULTIHEAP
- Viewing **ibmslapd** environment variables (AIX operating system only)

---

### Enabling large files

The underlying AIX operating system files that hold the contents of a large directory can grow beyond the default size limits imposed by the AIX operating system. If the size limits are reached, the directory ceases to function correctly. The following steps make it possible for files to grow beyond default limits on an AIX operating system:

1. When you create the file systems that are expected to hold the directory's underlying files, you should create them as Large File Enabled Journaled File Systems. The file system containing the DB2 instance's home directory, and, if **bulkload** is to be used, are file systems that can be created this way. The file system containing the **bulkload** temporary directory can be specified via the LDAPIMPORT environment variable.
2. Set the soft file size limit for the root, ldap, and the DB2 instance owner users to -1. A soft file size limit of -1 for a user specifies the maximum file size for that user as unlimited. The soft file size limit can be changed using the **smitty chuser** command. It is necessary for each user to log off and log back in for the new soft file size limit to take effect. You must also restart DB2.

---

### Setting MALLOCMULTIHEAP

The MALLOCMULTIHEAP environment variable can improve LDAP performance on SMP systems. To set this variable, run the following command just before starting **ibmslapd**:

```
export MALLOCMULTIHEAP=1
```

The disadvantage to using MALLOCMULTIHEAP is increased memory usage.

It might take less memory, yet have less of a performance benefit, if the variable is set as follows:

```
export MALLOCMULTIHEAP=heaps: numprocs+1
```

where *numprocs* is the number of processors in the multiprocessor system.

More information on MALLOCMULTIHEAP can be found in the AIX documentation.

---

## Viewing ibmslapd environment variables (AIX operating system only)

To view the environment settings and variables for your **ibmslapd** process, run the following command:

```
ps ewww PID
```

where *PID* is the **ibmslapd** process ID.

Example output for a PID of 27594:

```
$ ps ewww 27594
  PID  TTY  STAT  TIME  COMMAND
  27594 pts/2  A      0:02 ibmslapd -n -f /home/gwllmz/Test/Unit/Repl/
taliesin.ibmslapd.conf _=/usr/bin/ibmslapd LANG=en_US LOGIN=root IMQCONFIGCL
/etc/IMNSearch/dbcshep PATH=/usr/bin:/etc:/usr/sbin:/usr/local/bin:/usr/
contrib/bin:/usr/prod/bin:/usr/ucb:/bin:/usr/bin/X11:/sbin:/
usr/lib/netls/conf/nodeLock:/usr/bin:/etc:/usr/sbin:/usr/ucb:/usr/bin/
X11:/sbin:/home/ldapdb2/sqllib/bin:/home/ldapdb2/sqllib/adm:/home/ldapdb2/
sqllib/misc ID=root LC_FASTMSG=true IMQCONFIGSRV=/
etc/IMNSearch CGI_DIRECTORY=/var/docsearch/cgi-bin CLASSPATH=/home/ldapdb2/
sqllib/java/db2java.zip:/home/ldapdb2/sqllib/java/db2jcc.jar:/home/ldapdb2/
sqllib/java/sqlj.zip:/home/ldapdb2/sqllib/function:
. LOGNAME=root MAIL=/usr/spool/mail/root LOCPATH=/usr/lib/nls/
loc HNAME=taliesin PS1=[ID@HNAME: $PWD ]?=> OS=AIX VWSPATH=/home/
ldapdb2/sqllib DOCUMENT_SERVER_MACHINE_NAME=localhost
USER=root AUTHSTATE=compat _EUC_SVC_DBG_LOGGING=1 DCE_USE_WCHAR_NAMES=
1 _EUC_SVC_DBG_FILENAME=/tmp/gss.out SHELL=/bin/ksh ODMDIR=/etc/objrepos JAVA_HOME=/
usr/jdk_base DOCUMENT_SERVER_PORT=49213
_EUC_SVC_DBG=*.8 HOME=/ DB2INSTANCE=ldapdb2 VWS_TEMPLATES=/home/
ldapdb2/sqllib/templates LD_LIBRARY_PATH=/home/ldapdb2/sqllib/lib TERM=vt220 MAILMSG=
[YOU HAVE NEW MAIL] PWD=/home/gwllmz/Defects/77260/aus51ldap.20021106a DOCUMENT_DIRECTORY=/
usr/docsearch/html TZ=CST6CDT
TRY_PE_SITE=1 VWS_LOGGING=/home/ldapdb2/sqllib/logging A_z=! LOGNAME LIBPATH=/
usr/lib:/usr/ldap/java/bin:/usr/ldap/java/bin/classic:/home/ldapdb2/sqllib/
lib NLSPATH=/usr/lib/nls/msg/%L/%N:/usr/lib/nls/msg/%L/%N.cat KRB5_KTNAME=/home/
gwllmz/Test/Unit/Repl/taliesin.keytab
ODBCCONN=15
```

---

## Chapter 5. Hardware tuning

This chapter contains some suggestions for improving disk drive performance.

---

### Disk speed improvements

With millions of entries in LDAP server, it can become impossible to cache all of them in memory. Even if a smaller directory size is cacheable, update operations must go to disk. The speed of disk operations is important. Here are some considerations for helping to improve disk drive performance:

- Use fast disk drives
- Use a hardware write cache
- Spread data across multiple disk drives
- Spread the disk drives across multiple I/O controllers
- Put log files and data on separate physical disk drives



---

## Chapter 6. IBM Directory Server features

The sections in this chapter briefly describe the following additional performance-related IBM Directory features.

- Bulk loading (**bulkload**)
- Replication
- Monitoring Performance
- When to configure LDAP change log

---

### Bulk loading (**bulkload**)

The **bulkload** utility loads directory data to the LDAP database using an LDIF file. **bulkload** usually is significantly faster than **ldif2db** and **ldapadd** when loading approximately 100,000 to a million entries. Read the **bulkload** documentation in the *IBM Directory Server Version 5.1 Administration Guide* for information about using **bulkload**.

---

### Replication

IBM Directory Server supports replication. Through replication, multiple, synchronized copies of directory data are maintained on multiple directory servers. Using replication can improve performance related to search throughput because the workload can be divided across several servers. This allows for an enterprise's LDAP search activity to be distributed among several servers.

In the configuration of replication, you can specify that updates be replicated either immediately, or at scheduled times. If your application environment does not depend on the updates occurring immediately, it might be beneficial to schedule replication for off-peak hours, or on some more frequent, periodic basis. This will cause the updates to be batched together, reducing the cost of cache invalidation that results from updates on the replica servers.

Replication also adds to the workload on the master server where the updates are first applied. In addition to updating its copy of the directory data, the master server must send the changes to all replica servers. Careful scheduling of replication to avoid peak activity times will help minimize the impact to throughput on the master server.

---

### Monitoring performance

The following **ldapsearch** command can be used to monitor performance.

```
ldapsearch -h ldap_host -s base -b cn=monitor objectclass=*
```

where *ldap\_host* is the name of the LDAP host.

The monitor search returns some of the following attributes of the server:

**cn=monitor**

**version=IBM Directory, Version 5.1**

**total connections**

The total number of connections since the server was started.

**current connections**  
The number of active connections.

**maxconnections**  
The maximum number of active connections allowed.

**writewaiters**  
The number of threads sending data back to the client.

**readwaiters**  
The number of threads reading data from the client.

**opsinitiated**  
The number of initiated requests since the server was started.

**livethreads**  
The number of worker threads being used by the server.

**opscompleted**  
The number of completed requests since the server was started.

**entriessent**  
The number of entries sent by the server since the server was started.

**searchesrequested**  
The number of initiated searches since the server was started.

**searchescompleted**  
The number of completed searches since the server was started.

**filter\_cache\_size**  
The maximum number of filters allowed in the cache.

**filter\_cache\_current**  
The number of filters currently in the cache.

**filter\_cache\_hit**  
The number of filters found in the cache.

**filter\_cache\_miss**  
The number of filters not found in the cache.

**filter\_cache\_bypass\_limit**  
Search filters that return more entries than this limit are not cached.

**entry\_cache\_size**  
The maximum number of entries allowed in the cache.

**entry\_cache\_current**  
The number of entries currently in the cache.

**entry\_cache\_hit**  
The number of entries found in the cache.

**entry\_cache\_miss**  
The number of entries not found in the cache.

**acl\_cache**  
A Boolean value indicating that the ACL cache is active (TRUE) or inactive (FALSE).

**acl\_cache\_size**  
The maximum number of entries in the ACL cache.

**currenttime**  
The current time on the server. The current time is in the format:

year month day hour:minutes:seconds GMT

**Note:** If expressed in local time the format is

day month date hour:minutes:seconds timezone year

**starttime**

The time the server was started. The start time is in the format:

year month day hour:minutes:seconds GMT

**Note:** If expressed in local time the format is

day month date hour:minutes:seconds timezone year

**en\_currentregs**

The current number of client registrations for event notification.

**en\_notificationssent**

The total number of event notifications sent to clients since the server was started.

## Example

The following example shows how to calculate the throughput of the server by monitoring the server statistic called **opscompleted**, which is the number of operations completed since the LDAP server started.

Suppose the values for the **opscompleted** attribute obtained by issuing two **ldapsearch** commands to monitor the performance statistics, one at time **t1** and the other at a later time **t2**, were **opscompleted(t1)** and **opscompleted(t2)**. Then, the average throughput at the server during the interval between **t1** and **t2** can be calculated as:

$$(\text{opscompleted}(t2) - \text{opscompleted}(t1) - 3) / (t2 - t1)$$

(3 is subtracted to account for the number of operations performed by the **ldapsearch** command itself.)

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## When to configure LDAP change log

IBM Directory 5.1 has a function called **change log** that results in a significantly slower LDAP update performance. The **change log** function should be configured only if needed.

The **change log** function causes all updates to LDAP to be recorded in a separate change log DB2 database (that is, a different database from the one used to hold the LDAP server Directory Information Tree). The change log database can be used by other applications to query and track LDAP updates. The change log function is disabled by default.

One way to check for existence of the **change log** function is to look for the suffix **CN=CHANGELOG**. If it exists, the **change log** function is enabled.





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## Appendix A. DirMark 1.2

DirectoryMark 1.2 is an industry-standard benchmark provided by Mindcraft Inc. as a performance test for LDAP implementations.

<http://www.mindcraft.com/directorymark/>

DirMark 1.2 consists of a loading phase and two scenarios, Messaging and Address Lookup. Both Messaging and Address Lookup consist entirely of searches; there are no updates.

Each scenario consists of two phases, a warmup phase and a run phase. During the warmup phase, the searches primarily request entries that are not in the LDAP caches; most of these requests require interaction with the DB2 backing store. For all the measurements reported in this document, warmup consisted of running all queries at least once; consequently, during the run phase all entries requested are potentially already in LDAP caches in memory if the caches are large enough to hold all of them. Thus the warmup phase and the run phase comprise two distinctly different workloads.

During the run phase of Address Lookup, a number of client threads issue search requests to the IBM Directory Server from predetermined scripts. The scripts include a number of different kinds of searches, including wildcard and other searches that return multiple entries per request. The client threads run through their scripts continuously for three minutes. Throughput is measured on the server for each three-minute interval, and then each client starts over at the beginning of its script. Each three-minute interval is referred to as a run. The server is not restarted between runs.



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## Appendix B. Platform configurations

The examples in this guide use the following platform configurations:

- Clients
  - Two 866 MHz Intel PIII with 512 MB RAM; IBM 10/100 EtherJet™ PCI Adapter (1 processor active)
  - Two 4 CPU 1.8 GHz Intel PIV with 1GB RAM; Intel PRO/100 VE Desktop Connection (when 4 processors active)
  - Windows 2000 Professional with SP2®
- Server
  - 4-Way 600 MHz, pSeries™ eServer Model 660-6H1 with 1 or 4 processor active, 4 GB RAM
  - IBM 10/100 Ethernet Adapter
  - AIX 5.1
  - IBM Directory Server Version 4.1 GA
  - AIXTHREAD\_SCOPE=S
  - MALLOCTYPE=3.1 (1way) or MALLOCMULTIHEAP=1 (4 way)
  - NODISCLAIM=true (1way)
  - ACLCACHE\_SIZE=150000
  - RDBM\_CACHE\_SIZE=460000 except where noted
  - RDBM\_FCACHE\_SIZE=75000 except where noted
  - RDBM\_CACHE\_BYPASS\_LIMIT=100 except where noted.
  - RDBM\_ENTRY\_CACHE\_BYPASS=YES
  - Necessary Indexes created (for attribute seeAlso)
  - No ACLs were set. By default, anyone can search and compare. Directory Administrator can update
- DB2 v 7.2
  - maxlocks 100 sortheap 2500 dbheap 5000 ibmdefaultbp 236000 (4K pages) ldapbp 9800 (32K pages)
  - logfilsiz 2048, logprimary 6
- Miscellaneous
  - Caches were warmed up by running all scripts once
  - Measurements were taken using 50 client threads except where noted
  - DB2 log files are on the same disk as the containers
  - DirectoryMark Version. 1.2



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## Appendix C. Notices

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