z/TPF EE V1.1
z/TPFDF V1.1
TPF Toolkit for WebSphere® Studio V3
TPF Operations Server V1.2



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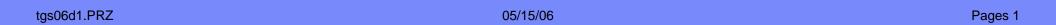
# TPF Users Group Spring 2006

Storage Allocation with z/TPF





IBM z/Transaction Processing Facility Enterprise Edition 1.1.0 © IBM Corporation 2006





## Introduction

- General storage recommendations for z/TPF systems
  - test
  - production
- Preallocation methods to optimize performance
  - models
  - distributions
  - data collection reports



#### Don't be afraid of this

- TPF customers have solved similar problems for years
  - Utilization to run
    - CPU
    - DASD
    - CH/CU
    - SAPs
  - VFA candidates
  - ► File resident programs



## Storage for test systems

- In the lab test systems grew from 250M to 1G
- For test systems use factor of 4
- z/TPF does't touch large storage until used
  - VFA blocks
  - ► CRPA
  - 1M frames
- Experiments
  - z/TPF touched roughly 1/4 of its storage
    - depends on how memory was allocated
  - TPF4.1 touched all its storage
  - z/TPF vs TPF 4.1 working sets about equal (4/4=1)
- 3 memory sizes for test systems
  - Often test systems have mean ECBs < 10</p>



## Storage for production systems

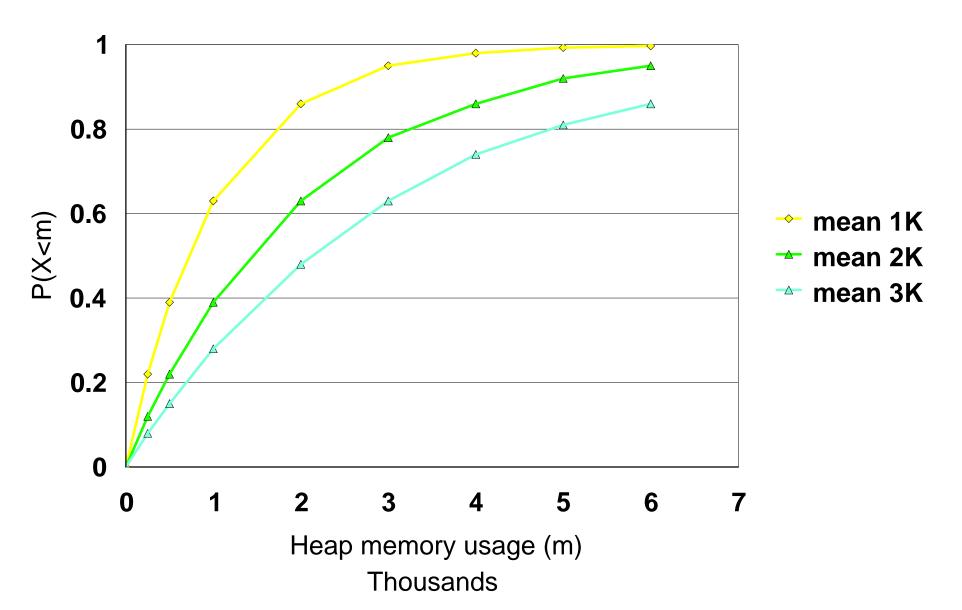
- Assume production system of 2G
  - ► .5G for VFA
  - ► 1.5G for not-VFA
- Not-VFA multiply by 3
  - $\triangleright$  1.5G x 3 = 4.5G
- VFA 20% increase
  - $\triangleright$  .5G x 1.2 = .6G
- Sum is 5.1G we round to 6G
- Production systems use memory factor of 3



## Initial z/TPF configuration

- When customers convert to z/TPF there will be extra storage
  - ► TPF4.1 < 2G and late model minimum CEC > 8G
- Be generous in initial z/TPF memory allocations
- Optimize using z/TPF data collection
  - ► TPF4.1 does not have some required reports
- Key point: have a huge number of 1M frames
  - Much more important than putting memory to VFA
  - System will be protected even with suboptimal preallocation

## **CDFs** with different means





## Determine 31-bit heap preallocation

- Each ECB gets fixed amount of heap call it P
- If ECB exceeds P then gets 1M frame
- Memory vs memory and memory vs performance
- Assume mean heap of 20K
  - ▶ If P=0 then a.s. each ECB gets 1M
  - ► If P=1M then wasted memory--few ECBs need > 100K
- Use TPF ECB 31-BIT HEAP AREA USAGE REPORT
  - Pmf of max heap use by ECBs over their life
    - from this create cdf F(x)
  - choose x such that .93 < F(x) < .98</p>



## 31-bit heap calculations

- define r = mean active ECB / defined ECBs
- let r = 300/1000
- cost is 1000x + (1-F(x)) 300 (1M)
  - g(x) = x + (1-F(x))r M
- for exponential 1-F(x) = exp(-lx)
- minimizing g(x) we get x=2.7 times mean
- F(2.7 mean) = .933



#### Determination of A1 to A4 sizes

- Use the ECB HEAP REQUEST SIZE REPORT
  - ► Set A4 = 3 times mean heap
  - Pmf shows 0 to 2A4 in increments 2A4/64
- Set
  - ► F(A1)=.2
  - ► F(A2)=.4
  - ► F(A3)=.6
  - ► F(A4)=.8



#### Determination of A1 to A4 counts

- Set buffer count = 0 so all ECBs miss
- Compute m= mean heap per ECB
- Counts =  $(m + 3 \operatorname{sqrt}(m))/4$ 
  - counts will vary from the mean
  - using Poisson estimate
- Key point not concerned about single ECB greatly exceeding counts
  - not a ruin problem like insurance/dam construction
  - inherently a ratio problem
  - system performance problem but not allocation problem



## **Application Stack**

- If ECB exceeds preallocation gets 1M frame
  - Memory vs memory has been solved
  - ightharpoonup x such that .92 < F(x) < .96
- Virtual limit
  - too small then too high % of ECBs perish
  - too large then steal virtual space < 2G</p>
  - choose y such that F(y) > .995
  - Use PREALLOCATED APPLICATION STACK STORAGE USAGE REPORT



#### **Thread Stack**

- No preallocated amount backed by real as needed
- Virtual storage issues as Application Stack
- Virtual = Tstack x Maxthread
  - If > Tstack then death
  - If > maxthread get negative returns so can control
- Use THREAD STACK STORAGE USAGE REPORT
- Use THREAD NUMBER PER PROCESS REPORT
- for both y such that F(y) > .995
- weak upper bounds control tail of process to get tighter



#### **ECB** Private Area

- If exceed prealloc by even 4K then build page/segment tables up to virtual limit
  - ▶ limit in 4.1 is 1M so good start point
- More accuracy/change use ECB PRIVATE AREA STORAGE USAGE REPORT
  - set prealloc at 6 times mean ECB private usage
  - if have cdf then x such that F(x)=.99
  - examine long tail--bad ECBs or error conditions



### **2G Virtual limit**

- 31-bit system heap
- ECB 31-bit heap
- Thread stack
- Application stack
- ECB private area(<32M)</p>
- 31-bit CRPA



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