

z/TPF V1.1

2013 TPF Users Group

Title: Performance Topics

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AIM Enterprise Platform Software IBM z/Transaction Processing Facility Enterprise Edition 1.1

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HW on floor can have variable MIPS

- Without an IPL TPF will adjust to the increased/decreased MIPS
- As an example with 15 CPs(current maximum EC12) we get 4 different MIPS ratings with no IPL

scaling factor ((n+1)/n)

- 2817-415 77.9 ITRR ------
- 2817-515 186.7 ITRR 2.38
- 2817-615 238.5 ITRR 1.28
- 2817-715 365.9 ITRR 1.53
- Talk to IBM account team
 - Conditions and HW/SW charges will apply
 - Designed to handle seasonal peaks
 - Infrequent adjustment

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2G to 4G for C programs

- Typically C programs (in 64-bit CRPAs) reside at addresses > 4G
- High branch frequency from >4G down to below 2G
 - Program linkage
 - Function Trace
- The z196 is not thrilled with these large jumps
 - Branch prediction affected
- z/TPF now permits 64-bit CRPAs to reside at 2G to 4G
 - Lab measurements showed large gains with artificially intense driver
- Customer measurement/inference showing 6 to 11% performance gain
- Joint work with lab and STG
 - EC12 much better at large jumps
 - APAR still has some value even on the EC12



Existence time calculations

• TPF Internal Existence time defined as

- Message ECB input time into TPF to output time from TPF
- Little's Law time honored method to calculate response time
 - Using msg rate and active ECBs
- Things would be simpler if the message came in, created no additional ECBs, waited until all work was completed, response sent and then exited
 - This doesn't happen often enough
- The major source of divergence is the listener/demon ECB
- For increased accuracy and partitioning by message group need to use Owners introduced in z/TPF
 - Use Little's Law on each partition

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Example with different estimators

Data

- Msg/sec=2150
- Active ECB =680
- Active IOB=147, IO/sec=89718, IO resp=1.6 mills
- Pysical IO per msg = 41.9
- IO wait predictor of msg response time = 41.9 x 1.6 = 67 mills
 - Crude lower bound
- Little's law with raw ECB count => 680/2150 = 316 mills
 - upper bound
- Little's law with adjusted ECB count => 280/2150 = 130 mills
 - assuming 400 ECBs were listeners
 - TPF data reduction will include a long running ECB count

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HyperPAV – Increased Device Throughput

• Assume for a CU

- .5 mills for CU cache hit
- 10 mills for CU cache miss
- P(cache hit) = .8; P(cache miss) = .2
- We have sufficient aliases and host adapter capacity
- Expected service time = E(s)=.8(.5) +.2(10)=2.4 mills
- Thus device rate at full utilization = 1/(.0024) = 416.6 IO/sec
 - Without HyperPAV- doing one IO at a time

Rate to Disk = 416.6(.2) = 83.4 IO/sec

- Disk is busy 2 mills out of the E(s)=> 2/2.4 = .83 of total device busy time
- So if the Disk could get more IO it could do more IO/sec; specifically 100-83.4 = 16.6 IO/sec
 - To send 16.6 IO/sec to the disk we need to send 83.4 IO/sec to the device
- Now with HyperPAV permitting the CU cache hits to 'MP' with the cache misses we do a total of 416.6 + 83.4 = 500 IO/sec to the device



HyperPAV Data Collection

• TPF will report IOBs in use by SSID

- Increase or decrease Aliases
 - Use Little's Law to determine IO response time by SSID

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Release performance analysis

- Increasing trend in mills/msg so only have about 10 days before system change and 10 days after
- Assume Tuesday msg rate was 20% higher than other days and hence performance is more critical on Tuesday
 - On the surface it seems reasonable to do a performance comparisons with 'Tuesday to Tuesday'
 - Problem is we are giving up too much precision because we only have 2 points before and after
- Suggest do differences of Mon to Mon, Tues to Tues etc



Sample Mean

- Well before the date of the TPF Release or PUT change
 - compute the sample variance of the daily TPF ITRs
 - With many observations this sample variance will be extremely close to the real underlying variance
 - For Release/PUT change the mean will change but the real variance will not

$$\mu_{\overline{x}} = \mu$$

$$\sigma_{\overline{x}} = \frac{\sigma}{\sqrt{n}}$$

$$\overline{x} = \frac{1}{N} \sum_{i=1}^{N} x_i$$

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Few sample points will have significant variation





Variance reduction with the sample mean



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Sample Mean convergence to true mean



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Notes

- Performance is in Internal Throughput Rate (ITR) ratio based on measurements and projections using standard IBM benchmarks in a
 controlled environment. The actual throughput that any user will experience will vary depending upon considerations such as the
 amount of multiprogramming in the user's job stream, the I/O configuration, the storage configuration, and the workload processed.
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