



IBM Software Group

# Introducing z/TPF!

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## AIM Core and Enterprise Solutions

IBM z/Transaction Processing Facility Enterprise Edition 1.1.0

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# Opening

- New application model
  - Large memory
  - Large programs
  - Subroutines
  - Stacks

## Opening

- Other z/TPF presentations
  - Applications: Debuggers
  - Applications: Coding today for z/TPF tomorrow - Assembler
  - Languages: GNU Compiler Collection (GCC) Overview
  - Languages: Coding today for z/TPF tomorrow - C/C++
  - Languages: C/C++ single source APARs
  - Operations: Positive feedback enhancements
  - SCP: Scheduler enhancements
  - Database: TPFDF status
- Other demonstrations and education
  - TPF Toolkit
  - Make environment

## Agenda

- z/Architecture Overview
- Development Environment
- Application Productivity Enhancements
  - Debug and Test Enhancements
  - Design and Coding Enhancements - z/TPF Globals Support
  - Design and Coding Enhancements
- Availability Enhancements
- New Tuning Capabilities
- Additional Enhancements
- Migration to z/TPF: Concepts and Tools
- Closing

# z/Architecture Overview

## Architecture Definition

- z900, z990, z800, and z890 machines have 2 architecture modes
  - ESA/390 - TPF 4.1 runs in this mode
    - Limited to 2 GB of storage
  - z/Architecture - z/TPF runs in this mode
    - Designed to exploit larger storage spaces
      - For example, 32 GB, 64 GB, 128 GB, and so on

## What is z/Architecture?

- 64-bit addresses
  - Real
  - Virtual
    - Region tables
    - 64-bit entries in page - segment - region table entries
- 64-bit registers
  - 16 general registers
  - 16 control registers
  - 16 floating point registers
- 128-bit PSW
- 8-KB prefix area
- Format 2 IDAWs - 64-bit entries



## What is z/Architecture (continued)?

- z/Architecture instruction set is a superset of ESA/390
  - ESA/390 instructions work in z/Architecture
    - Except: No vector instructions
- Addressing modes
  - 64-bit addressing mode
  - 31-bit addressing mode
  - 24-bit addressing mode (not supported in z/TPF)
- Reference...
  - z/Architecture Principles of Operation (SA22-7832)
  - z/Architecture Reference Summary (SA22-7871)



## Use of z/Architecture Enables Use of New Features

- Relative Instructions - BRC, LARL
  - Program expansion is easier
- Immediate Instructions
  - Halfword Immediate - AHI, CHI, LHI, and MHI
  - For portions of 64-bit register - I1xx, N1xx, O1xx, TMxx
    - For example TMHH, TMHL, TMLH, TMLL
  - Programs easier to write and more efficient
- Storage-protection override
- C/C++ benefits as well as assembler because compiler takes advantage of these instructions

## z/TPF Memory Model

- z/TPF has ability to exploit > 2 GB of memory
  - Several existing large system tables moved above 2 GB
  - Most new system tables reside above 2 GB
  - Customer tables can reside above 2 GB
  - Programs can reside and execute above 2 GB
  - Additional diagnostics
  - New tuning capabilities

## z/TPF Memory Model (continued)

- Program execution
  - C programs execute in 64-bit addressing mode
    - Mostly recompile of TPF 4.1 source code
  - Assembler programs can execute in 31-bit or 64-bit addressing mode
    - Assembler programs brought from TPF 4.1 with no changes will execute in 31-bit addressing mode
    - Most IBM realtime assembler system programs did not require change and still run in 31-bit addressing mode
    - **Expect most assembler programs will not require source changes**

# z/TPF SVM < 4 GB

...see next page

4 GB	Not valid
2 GB	<b>31-bit System Heap*</b>
	4KB Frames
	ECBs
	SWBs
	IOBs
	Common Blocks
	<b>31-bit CRPA*</b>
	<b>31-bit Virtual User Areas*</b>
	31-bit System Tables
16 MB	IPLB/CCIO/Keypts/CIMR
	Control Program
0	Prefix Page

\* - Use 1-MB frames

# z/TPF SVM > 4 GB

Other I-Streams SVM region/segment/page
VFA
<b>64-bit System Heap*</b>
<b>System Heap Control*</b>
Dump Buffer Area
EVM region/segment/page
1-MB Frames
ECB Trace Tables
ECB Preallocated Areas
<b>64-bit CRPA*</b>
<b>64-bit Virtual User Areas*</b>
<b>TCP/IP Areas*</b>
64-bit System Tables
LDEV Trace
Main I-Stream SVM region/segment/page

\* - Use 1-MB frames

4 GB  
0

...see preceding page

# z/TPF EVM

	EVM	SVM
	Same as SVM	
	64-bit ECB Heap	1MB Frames
	Same as SVM	
	Preallocated ECB Storage	Preallocated ECB Storage
4 GB	Same as SVM	
2 GB	Not Valid	
	Same as SVM	
	31-bit ECB Heap	4KB Frames
	Preallocated 31-bit ECB Heap	
	Preallocated ECB Stack	
	ECB Stack Area	
	Thread Stack Area	
	ECB Private Area	ECBs
0	Same as SVM	

## z/TPF Memory Model - Real Address Handling

- SVM address does not equal real address
  - In TPF 4.1 most SVM addresses equaled real addresses
  - In z/TPF SVM is truly virtual
- Do not need > 2 GB of real storage to run z/TPF
  - Important for VM test systems
- z/TPF SVM and EVM do not map 2 GB - 4 GB
  - Real storage is not discarded
  - Intent is to find problems where programs in 31-bit addressing mode pass an incorrect pointer to programs in 64-bit addressing mode
    - Pointer created using BAL, BAS, BALR, BASR instructions.

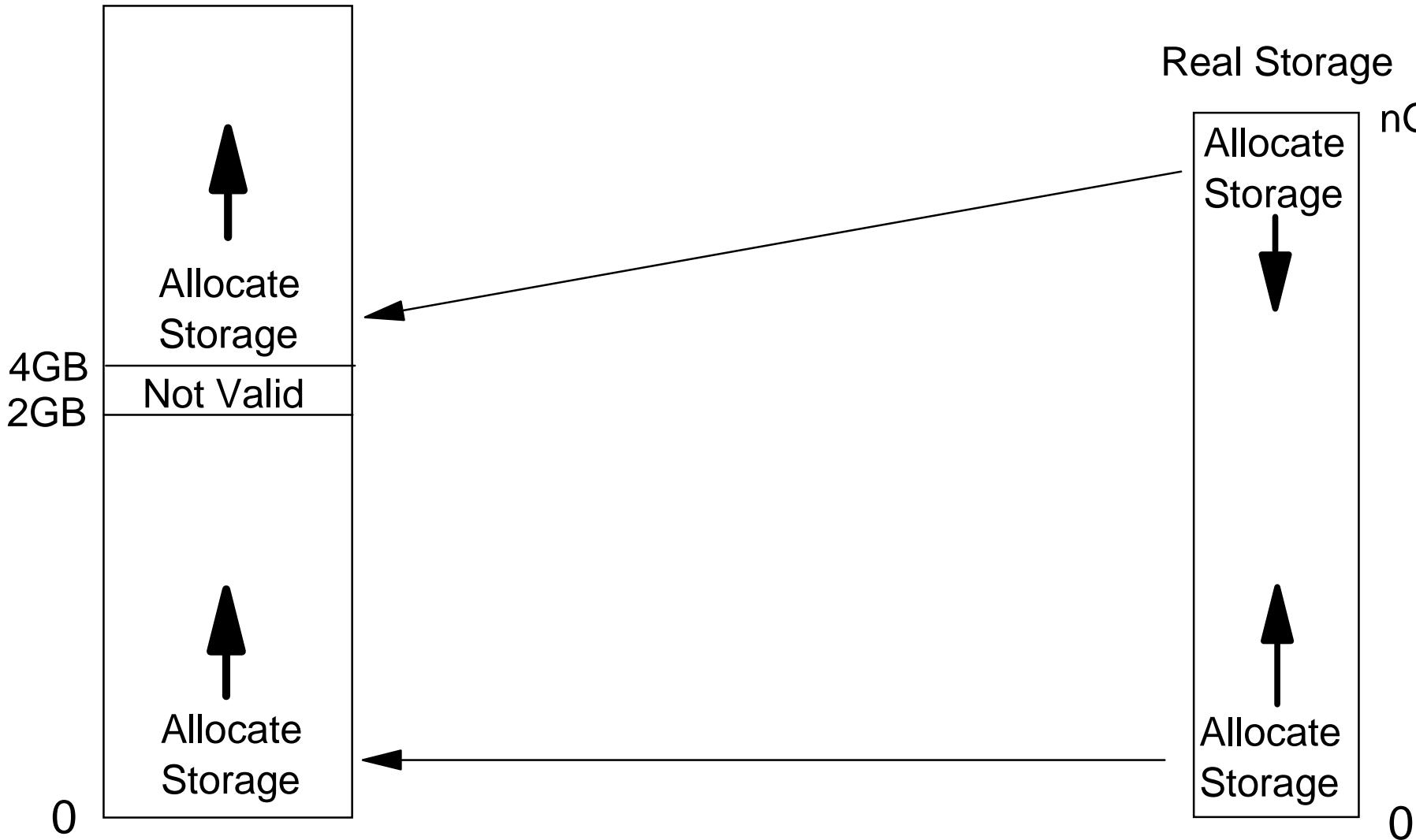


# z/TPF Real Address Mapping

SVM

Real Storage

nGB



# Development Environment and Program Management

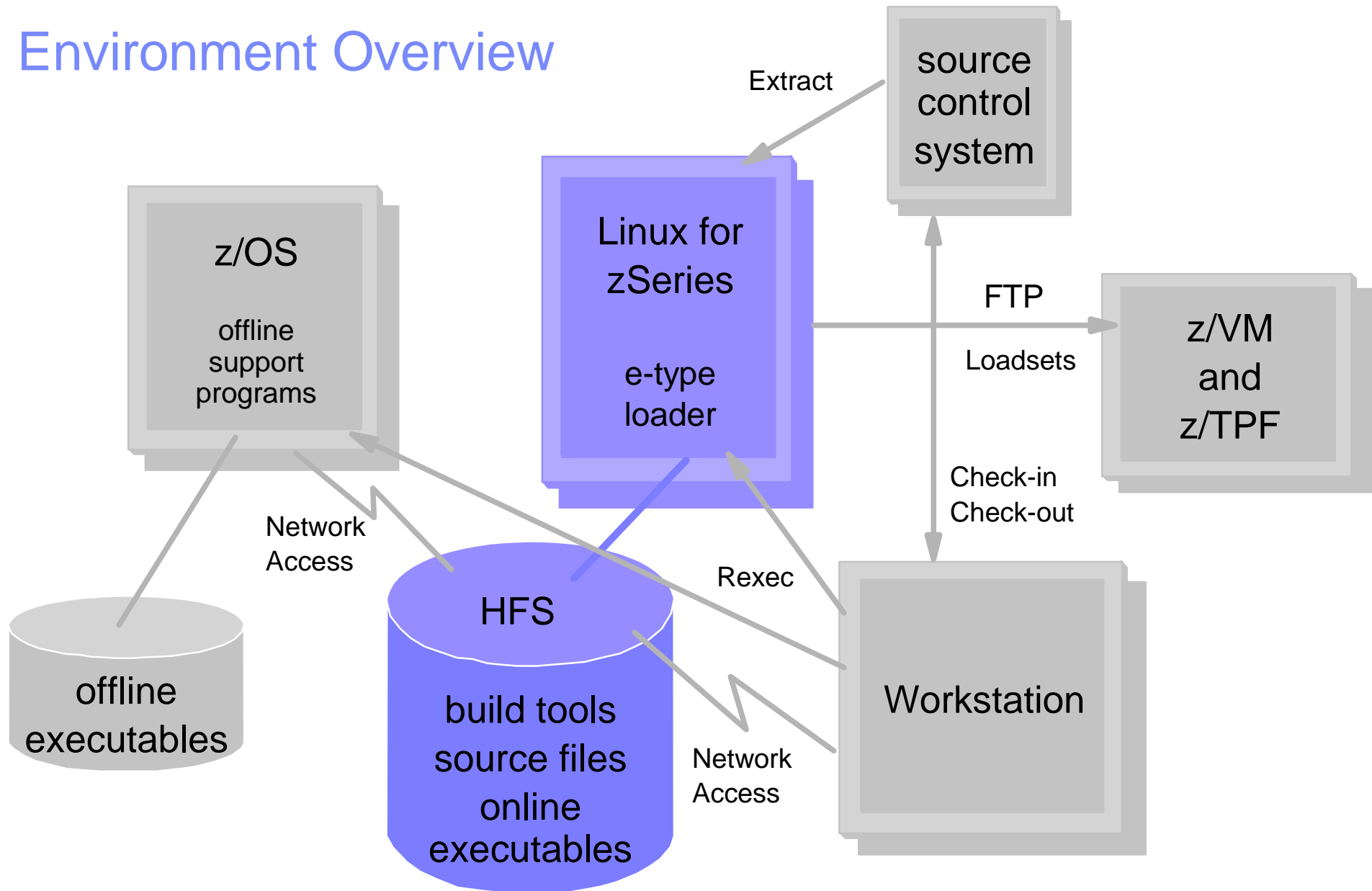
## Useful Terms

- **FSF** : Free Software Foundation
  - Organization that sponsors and controls GNU software
  
- **GNU** : GNU's Not Unix
  - Collaborative effort to provide free software
  
- **GPL** : GNU Public License
  - License used to ensure that free software remains free
  
- **GCC** : GNU Compiler Collection
  - Set of tools used to build programs
  - Supports a number of languages
  
- **ELF** : Executable and Linking Format
  
- **SO** : Shared Object

# Development Environment Objectives

- Open
  - Based on open tools
    - gnumake, korn shell
  - Easy to modify or enhance
  
- Consistent
  - System and development builds use the same tools
  - MakeTPF integrated with TPF Toolkit for WebSphere Studio
  
- Configurable
  - Different TPF configurations
  - Layers for development
  
- Easy to use
  - HFS based

# Environment Overview



# System and Application Build

- Driven from Linux
  - Compiler, Assembler, Linker, and Loader\*
- Uses maketpf makefiles
  - All compile, assemble, and link options
  - Used to build entire system
- Managed by a control file
  - Central point of information regarding all programs
  - Replaces sppgml.mac, ibmpal.cpy, and usrtpf.cpy
- Generates load deck
- Easy to build and maintain the system

# C/C++ Compiler Support

- GNU Compiler Collection (GCC)
  - C/C++, assembler, linker, and other utilities
  - Open Source
    - All changes incorporated by the FSF
  - Cross compiler
  - Open standards
    - Executable and Link Format (ELF)
    - Standard Application Binary Interface (ABI)
  - Extensions
    - Imbedded assembler



## z/TPF Transformation

- z/TPF modified to match open standards
  - All programs are loaded as ELF
    - C/C++ and TPF Assembler
    - DLMS, DLLs, LLMs, BAL Segments, and CIMR Components
  - Enter/Back modified to be a dynamic loader
  - All program linkages modified to match the ABI
  - All programs use stack as linkage mechanism
- Better integration of C/C++ and BAL
- Improved performance of program linkages
- Easier to manage

## Executable and Linking Format (ELF)

- Describes binary file format used by LINUX and documented as part of System V Application Binary Interface (ABI)
  - Tool Interface Standard (TIS) Portable Formats Specification
    - <http://www.linuxbase.org/spec/refspecs/elf/TIS1.1.pdf>
  - System V ABI Edition with 64-bit ELF update
    - <http://www.linuxbase.org/spec/refspecs/elf/gabi41.pdf>
- Must be used in conjunction with a processor-specific ABI supplement
  - LINUX for zSeries
    - [oss.software.ibm.com/developerworks/opensource/linux390/docu/lzsabi0.pdf](http://oss.software.ibm.com/developerworks/opensource/linux390/docu/lzsabi0.pdf)

# Program Linkage

- Dynamic Linker / Loader
  - Transforms Enter/Back to standard processing
    - Retrieves programs from file
    - Performs relocation
  - Resolves external references to a specific stub at *system level*
    - Handles any interface transformations
    - Provides dynamic hooks
      - E-type loader
      - User exits
      - Trace tools
- Standard processing with z/TPF efficiency

## PAT : Changing Concept of Program Allocation

- No offline allocator
  - SALTBL or TABLE40 are obsolete
  - No more timestamp mismatches
- No need to maintain program order
  - System detects new programs
  - Adds new entries to end
- Dynamically change via command (ZAPAT)
- Maintenance handled via positive feedback controls
  - Online capture of IPAT via command (ZDECK)
  - Offline tool to merge captured information into control file
- Easier to maintain and manage program base

## IPAT : Program Attribute Table

- Defined in the control file
  - Fetch type: pre-load, on demand, and default
  - Load type: BSS only, SS only, ALL, SS Shared
  - Debug
  - Timeout value: caller, default, interval
  - Affinity
  - Dump group name
  - Trace group name
  - Authorizations: restrict, montc, key0, cmb, bypass
  - User attributes
  
- Built from control file using a makefile
  
- Maintenance can be automated

## Load Deck Changes

- No fixed column restrictions
- More freedom in ordering cards
  - using control statements
- More options for search paths
  - Multiple paths for a given component type
  - Override path with program lists
- Output report
  - shows directory in which components were found
  - more helpful diagnostic messages
- New patch format
  - allows validation of data
  - consistent with ZAPGM format
- Option to verify programs against IPAT

# Sample TLDR Load Deck

```
@DEFINE
SYSID=BSS
&SYS1=/u/ztpf/linuxbin/cur/bss/load
&APP1=/u/cmast/proj1

@CIMR &SYS1
CPS0.so (my patches to fix displacement problem)
@@CCNUCL B2A D0C8 VALDATA-D0C5

@KEYPOINT &SYS1
* Processor Shared Keypoints
CTK6.kpt
* Processor Unique Keypoints for Processor C
CTKA.kpt%C

@VERIFY &SYS1
IPAT.so

@APPLICATION &APP1
@INCLUDE /u/cmast/base/PROJ.BASE
MYP1.so /* fix for defect 12345 */
MYP2.so
```



# Sample TLDR Output

```
@DEFINE
SYSID=BSS
&SYS1 = /u/ztpf/linuxbin/cur/bss/load/
&APP1 = /u/cmast/proj1/
@CIMR &SYS1
CPS0.so (my patches to fix displacement problem)
@@CCNUCL b2a D0C8 VALDATA-D0C5
@KEYPOINT &SYS1
CTK6.kpt
CTKA.kpt %C
@VERIFY &SYS1
IPAT.so
@APPLICATION &APP1
@INCLUDE /u/cmast/base/PROJ.BASE
>QZZ1.so
>QZZ2.so
>QZZ3.so
MYP1.so
MYP2.so
```

```
LOAD0007I The default of ELDRCLEAR=NO was accepted. Any loadsets that were
           loaded previously to the image will remain unchanged.
```

```
LOAD0010I Parsing step ended with return code 0.
```

# Sample TLDR Output

```
*****  
                END PARSING PHASE                *****
```

```
LOAD8070E MYP1.so is not present in the specified IPAT  
          and will not be loaded.
```

```
LOAD8070E MYP2.so is not present in the specified IPAT  
          and will not be loaded.
```

```
*****  
*****  
                SUMMARY REPORT                *****  
*****
```

Status of DEFINE Flags:

SYSID=BSS

OVERLAY\_IPAT=NO

PROGCLEAR=NO

ELDRCLEAR=NO

DEBUGFILES=YES

WARN\_BSO\_RELO=NO

&APP1 = /u/cmast/proj1/

&SYS1 = /u/ztpf/linuxbin/cur/bss/load/

THE FOLLOWING PAT WAS USED FOR VERIFICATION

/u/ztpf/linuxbin/cur/bss/load/IPAT.so

# Sample TLDR Output

```
CIMR Compnts      Path      Found In
  CPS0.so         &SYS1    /u/ztpf/linuxbin/cur/bss/load/
                (my patches to fix displacement problem)
                @@CCNUCL b2a D0C8
```

```
Keypoints        P Path      Found In
  CTKA.kpt       %C &SYS1    /u/ztpf/linuxbin/cur/bss/load/
  CTK6.kpt       &SYS1    /u/ztpf/linuxbin/cur/bss/load/
```

```
App. Pgms        Path      Found In
  MYP1.so        &APP1    /u/cmast/proj1/
  MYP2.so        &APP1    /u/cmast/proj1/
  QZZ1.so        &APP1    /u/cmast/proj1/
  QZZ2.so        &APP1    /u/cmast/proj1/
  QZZ3.so        &APP1    /u/cmast/proj1/
```

THE FOLLOWING COMPONENTS WERE NOT LOADED.

APP programs:

```
MYP1.so          - NOT PRESENT IN THE SPECIFIED IPAT
MYP2.so          - NOT PRESENT IN THE SPECIFIED IPAT
```

LOAD0020I Load step ended with return code 8.

## Sample OLDR Load Deck

```
@DEFINE
  SYSID=BSS
  &MYPGMS=/u/cmast/proj1      /* define search path */
  &SYS1=/u/ztpf/linuxbin/cur/bss/load /* path that contains IPAT.so */

@VERIFY &SYS1
  IPAT.so

@LOADSET BASE1 &MYPGMS
  MYP1.so, MYP2.so
@INCLUDE /u/cmast/base/PROJ.BASE
```

# Sample OLDR Output

```
@DEFINE
SYSID=BSS
&MYPGMS = /u/cmast/proj1/
&SYS1 = /u/ztpf/linuxbin/cur/bss/load/
@VERIFY &SYS1
IPAT.so
@LOADSET BASE1 &MYPGMS
MYP1.so
MYP2.so
@INCLUDE /u/cmast/base/PROJ.BASE
>QZZ1.so
>QZZ2.so
>QZZ3.so
LOAD0010I Parsing step ended with return code 0.
```

```
*****                               END PARSING PHASE                               *****
```

```
LOAD4070W MYP1.so is not present in the specified IPAT.
           The program will be loaded anyway.
LOAD4070W MYP2.so is not present in the specified IPAT.
           The program will be loaded anyway.
```

# Sample OLDR Output

```
*****
***** SUMMARY REPORT *****
*****
```

## Status of DEFINE Values:

```
SYSID=BSS
WARN_BSO_RELO=NO
DEBUGFILES=YES
&MYPGMS = /u/cmast/proj1/
&SYS1 = /u/ztpf/linuxbin/cur/bss/load/
```

## LOADSETS SUCCESSFULLY PROCESSED:

### Loadset BASE1:

```
MYP1.so      &MYPGMS      /u/cmast/proj1/
MYP2.so      &MYPGMS      /u/cmast/proj1/
QZZ1.so      &MYPGMS      /u/cmast/proj1/
QZZ2.so      &MYPGMS      /u/cmast/proj1/
QZZ3.so      &MYPGMS      /u/cmast/proj1/
```

## LOADSETS SUCCESSFULLY PROCESSED WITH WARNINGS:

### Loadset BASE1:

```
MYP1.so      - NOT PRESENT IN THE SPECIFIED IPAT
MYP2.so      - NOT PRESENT IN THE SPECIFIED IPAT
```

LOAD0020I Load step ended with return code 4.

# Applications Productivity Enhancements Debug and Test Enhancements



## ECB Trace Enhancements - C Function Trace

- C function trace is always available
  - Can be used in both test and production environments
  - Enabled by default compiler option
- Function entry and exit can be traced
  - Function entry and macro trace are equivalent
  - Function exit is a separate option
- Formatted macro and C function trace are collated together
- Function trace entries include:
  - Prototype
  - Parameter values
- Ability exists to add free-form text into function trace.

# ECB Trace Example

## ECB TRACE

```

TRACE      PROG  IS  LOADMOD  MACRO      PARAMETERS
GROUP      MODE          OR          OR
          OFFSET  PROGRAM INFORMATION
IBM_DEFT   CTAL
          01  +001E1D0 void __serrc(t_serrc status=00000000,int number=0092EA6C,const char* msg=
          0000000000000000,void ** slist=0000000000000000,char prefix=0xE4,+0001
          MORE PARAMETERS)
          CTIS
          01  +000C124 return(void * 0000000480000000) errno=00000
          from void * malloc64(size_t)
          64PU1 01  +000046C QZZ8  MALOC  S-000000000000F0F78 B-0000000480000000
R0  00000000 00000000 00000000 00000000  R1   R2  00000000 00000000 00000000 00000500  R3
R4  00000003 8DBF1110 00000000 000F0F78  R5   R6  00000000 00000082 00000003 8C5E2CF4  R7
R8  00000003 8C5E2DF0 00000000 CC022800  R14  R15 00000000 0AA0039A

          01  +000C110 void * malloc64(size_t size=000000000000F0F78)
          QZZ8
          01  +0000B6C return(long int 000000000000F0F78) errno=00000
          from long int get_size_of_table(char*)
          01  +0000B62 long int get_size_of_table(char* name=000000038C5E2EEA)
          CTAL
          01  +0023452 return(int 00000049) errno=00000
          from int tpf_trace_info(char*)
          01  +00233AC This is a test to show that free form text can be added to function trace

```

## ECB Trace Enhancements - More Trace Entries

- Number of trace entries per ECB can be set by the user
  - Maximum number of trace entries 65,535
- Provides more diagnostic information for each ECB
  - Can see more of the flow through an application
- Makes debugging easier.

## ECB Trace Enhancements - Trace Groups

- Problem: middleware programs can dominate ECB trace
- Solution: trace groups
  - Each ECB has multiple trace buffers
    - Trace buffers per ECB is user settable (min = 4, max = 254)
  - Trace Group = set of programs which share an ECB trace buffer
    - User settable by load module
      - Control File (offline) PAT (online)
    - IBM use
      - MQ has its own trace group = IMQSGRP
      - z/TPFDF has its own trace group = DATABASE
- Makes debugging easier.

# ECB Trace Group Example

## ECB TRACE

TRACE GROUP	PROG MODE	IS	LOADMOD OR	MACRO OR	PARAMETERS
IBM_DEFT			OFFSET	PROGRAM INFORMATION	
			QXDH		
		01	+00000C6	QXDH ENTRC	ENTER TO UAC0
R0	00000000	00000000	00000000	00000000	R1 R2 00000000 00000000 00000000 00000004 R3
R4	00000000	00000001	00000000	0A308000	R5 R6 00000000 0C6001C0 00000000 0A306EA1 R7
R8	00000004	00000064	00000000	0A308150	R14 R15 00000000 83745CCA
			QXCE		
		01	+0000172	QXCE ENTRC	ENTER TO QXDH
R0	00000000	00000000	00000000	00000000	R1 R2 00000000 00000000 00000000 00000004 R3
R4	00000000	00000001	00000000	0A308000	R5 R6 00000000 0C6001C0 00000000 0A306EA1 R7
R8	00000004	00000064	00000000	0A308150	R14 R15 00000000 83745CCA
DATABASE			UTDF		
		01	+0000E98	UAA0 BACKC	RETURN TO QXCE
R0	00000000	00000000	00000000	00000000	R1 R2 00000000 00000000 00000000 00000004 R3
R4	00000000	00000001	00000000	0A308000	R5 R6 00000000 0C6001C0 00000000 0A306EA1 R7
R8	00000004	00000064	00000000	0A308150	R14 R15 00000000 83745CCA
			UFTH		
		01	+0009582	return(void) errno=00000	from void ufth_btree_interface(void)

## ECB Trace Enhancements - New ECB Trace Capabilities

- ECB heap trace
  - Records all ECB heap requests (malloc, calloc, realloc, free)
  - Formatted in dumps
  - Number of entries is user settable (max = 32,767)
  - Ability to track heap usage even if ECB trace has wrapped
- ECB data level trace
  - Records name of load module
    - Obtained 4K frame in ECB private area
    - Released 4K frame in ECB private area
  - Formatted in dumps
  - Ability to track ECB private area usage even if ECB trace has wrapped



# New ECB Trace Examples

## ECB heap trace:

### ECB HEAP TRACE

```

LOADMOD   LOADSET           OBJECT NAME
OR
OBJ DSP   IS ECB SVM  THRDID  HEAP BUFFER ADDR API
CTS7      LOADSET-BASE     OBJECT-cts7.goff
+000012F4 01 0ABE4000 00000000                FREEC BLOCK=000000000CE00FB0
+000011BA 01 0ABE4000 00000000 000000000CE00FB0 MALOC SIZE=0000000000000044
+00000500 01 0ABE4000 00000000                FREEC BLOCK=000000000CE00FF0
+000012F4 01 0ABE4000 00000000                FREEC BLOCK=000000000CE0CFB0
+000011BA 01 0ABE4000 00000000 000000000CE0CFB0 MALOC SIZE=0000000000000044
CMQO      LOADSET-BASE     OBJECT-cxcsmw.cpp
+00000046 01 0ABE4000 00000000 000000000CE02FF8 malloc(size=0000000000007FFE)
CTSC      LOADSET-BASE     OBJECT-ctsc.goff
+0000042C 01 0ABE4000 00000000 000000000CE00FF0 MALOC SIZE=0000000000000008

```

## ECB data level trace:

\*DATA BLK, LEVEL 5 OBTAINED BY PROGRAM-BXAM RELEASED BY PROGRAM-N/A

\*AVAIL L1 BLK OBTAINED BY PROGRAM-CVAA RELEASED BY PROGRAM-CVAU



## ECB Trace Enhancements - New ECB Trace Capabilities

- Trace log
  - Ability to record all macro and C function trace items for an ECB
  - Started by API call - TLOGC, tpf\_tracelog\_on()
  - Stopped by API call or ECB exit - tpf\_tracelog\_off()
  - Data written to tape or HFS file
  - Post processed on Linux
    - Output to HFS file

## Socket API Trace

- Trace socket API calls
- Entries contain detailed information, including:
  - API input parameters
    - Includes implied parameters like time out values
  - Output, including the API return code
    - *sockerrno()* value included if error return code
  - How long it took the API to be completed
  - Data is displayed in user friendly (readable) format
- Two trace entries created if ECB becomes blocked while processing the API:
  - First entry shows the input and when the API was issued
    - Allows you to determine if an API is pending
  - Second entry shows the output and when the API was completed

## There are Actually Two Socket API Traces

- Trace at a per-ECB level:
  - Contains socket APIs issued by this ECB
  - Independent of C function trace
  - Useful in debugging socket application programs
  - Included and formatted in dumps
- Trace at a per-socket level:
  - Contains APIs issued for a particular socket
  - Useful because multiple ECBs can share a socket
  - A socket can be passed from one ECB to another via *activate\_on\_receipt* (AOR)
  - Display online via new ZSOCK API command
  - Allows you to view recent history of the socket for faster online debugging
    - Can be used in conjunction with individual IP trace

# Socket API Trace Example - Compact Format

ZSOCK TRACE SOCK-C000002

SOCK0035I 15.26.53 BEGIN SOCKET TRACE FOR 00C00002

ECB	API	RC	PROGRAM	COMPLETION TIME (SEC)	TIME STAMP
07650000	socket	00C00002	QZZQ	0.003	May 15 15:25:39
07650000	bind	0	QZZQ	0.011	May 15 15:25:39
07650000	listen	0	QZZQ	0.002	May 15 15:25:39
07650000	accept		QZZQ		May 15 15:25:39
07650000	accept	00C00008	QZZQ	13.634	May 15 15:25:52
07650000	accept		QZZQ		May 15 15:25:52
07650000	accept	00C00013	QZZQ	5.213	May 15 15:25:57
07650000	accept	00C00014	QZZQ	0.016	May 15 14:25:57
07650000	accept		QZZQ		May 15 15:25:57

END OF DISPLAY

## Socket API Trace Example - Detailed Format

```
ECB-07650000  API-socket          PROG-QZZQ  OFFSET-000058  IS-02  May 15 15:25:39
  type-SOCK_STREAM  prot-IPPROTO_TCP
  RC-00C00002      COMPLTIME-0.003sec

ECB-07650000  API-bind          PROG-QZZQ  OFFSET-000148  IS-02  May 15 15:25:39
  port-5004  ip-9.117.241.1  addrLen-16
  RC-0      COMPLTIME-0.011sec

ECB-07650000  API-listen       PROG-QZZQ  OFFSET-000334  IS-02  May 15 15:25:39
  backlog-15
  RC-0      COMPLTIME-0.002sec

ECB-07650000  API-accept       PROG-QZZQ  OFFSET-000520  IS-02  May 15 15:25:39
  addrLen-16  timeout-0
  BLOCKED

ECB-07650000  API-accept       PROG-QZZQ  OFFSET-000520  IS-02  May 15 15:25:52
  port-1027  ip-9.117.232.167
  RC-00C00008  COMPLTIME-13.634sec
```

## Network Management - Start with SNMP

- Simple Network Management Protocol (SNMP) tracks exception conditions at the system-wide level for TCP/IP nodes (hosts and routers). For example:
  - MIB variable *ipInDiscards* contains the number of datagrams received by this node that contained no errors, but were discarded anyway (for example, due to lack of buffer space)
    - Useful for determining which routers are overloaded
- SNMP can tell you if a server is having problems. For example:
  - MIB variable *tcpRetransSegs* contains the number of segments retransmitted by this node and *ipReasmReqds* contains the number of IP fragments received by this node
- However, SNMP will not tell you which sockets are causing the most retransmits, or have the most IP fragments flowing
- In addition, SNMP does not track all exception conditions
  - No MIB variable for TCP data received out of order



## Network Management - Add Socket Exceptions

- z/TPF tracks exception conditions on a per-socket basis
- Exceptions for a given socket are shown when that socket is displayed (via ZSOCK DISPLAY with FORMAT)
- Use the new ZSOCK EXCEPTION command to determine the sockets with the most total exceptions or most exceptions of a specific type
  - Allows you to identify troubled sockets in realtime rather than having to examine trace data after the fact
  - Helps identify the scope of the problem. For example:
    - One socket
    - Sockets with one local or remote application
    - Sockets with one business partner
    - Many/all sockets
- Turns SNMP into a Superior Network Management Protocol



## Problem Determination Example Using Socket Exceptions

1. Large amount of retransmitted packets on z/TPF is detected by:
  - An SNMP manager monitoring z/TPF
  - ZSNMP or ZTTCP DISPLAY STATS command
2. Issue ZSOCK EXCEPTION to gather further information

### ZSOCK EXCEPTION RETRANS TOP-5

```
CSMP0097I 12.48.57 CPU-B SS-BSS SSU-HPN IS-01
SOCK0042I 12.48.57 BEGIN SOCKET EXCEPTIONS DISPLAY
```

RANK	FD	TOTAL	RETRANS	OUT ORDER	FRAG OUT	FRAG IN
----	-----	-----	-----	-----	-----	-----
1	00C00011	376	366	10	0	0
2	00C00256	25	23	2	0	0
3	00C0014B	7	7	0	0	0
4	00C00033	935	5	0	0	930
5	00C001CD	8	5	3	0	0

END OF DISPLAY

## Dump Formatting Enhancements

- Storage formatting always has:
  - 8-byte core address on left
  - 16-byte hex display in middle
  - Translated display on right
- Translation can be
  - EBCDIC
  - ASCII
  - User defined code page
  - Specified on IDATG or LISTC macro

# Dump Formatting Translation Example

```

*AREA REFERENCED BY REGISTER 9 00000000ED37D14  PROTECT KEY      9
00000000ED37000      00000000      00000000      00000000      0ED39000      .....L..
00000000ED37010      00000000      00000000      C3C8D3F1      40404040      .....CHL1
00000000ED37020      40404040      40404040      40404040      C3C8D3F1      CHL1
00000000ED37030      40404040      40404040      40404040      40404040
00000000ED37040      00000004      00000003      00000002      00000000      .....
    
```

```

*DATA BLK      OBTAINED BY PROGRAM-BXDY  RELEASED BY PROGRAM-N/A
*BLOCK IN USE
0000000010A08000      C1F277C0      00000000      00000000      00060000      A2.....
0000000010A08010      00000000      00000000      00000000      00000000      .....
    
```

```

*MSG BLK, LEVEL 0
00000000AB06E80      00000040      00212850      00000000      00000000      ... ..
00000000AB06E90      00410100      00E9C9D5      C5E340C1      C4C440E2      .....ZINET ADD S
00000000AB06EA0      60D4E8E2      C5D9E540      D7D9D6C3      60E740D7      -MYSERV PROC-X P
00000000AB06EB0      C7D460D4      E8D7C740      D4D6C4C5      D360E6C1      GM-MYPG MODEL-WA
00000000AB06EC0      C9E340D7      D6D9E360      F9F9F9F9      F9F9F9F9      IT PORT-99999999
    
```

## Dump Formatting Enhancements - Extensions

- Dump processing has ability to call dump formatting extension
- Extension allows customization
  - What is to be dumped
  - Actual formatting of the data
- User extensions are available
- IBM uses extensions
  - MQ uses to dump queue and channel information

# Dump Formatting Extensions Example

\*MQ Profile (QMGR) in Memory:

000000000ECF6000	00000064	E3D7C6D8	D4404040	40404040	....TPFQM
000000000ECF6010	40404040	40404040	40404040	40404040	
000000000ECF6030	40404040	415DB9D6	00000001	00000000	...O.....
000000000ECF6040	00000000	000000B4	0000000A	0000001E	.....
000000000ECF6050	00000005	00001388	00000100	00000032	.....h.....
000000000ECF6060	C4C5C1C4	4BD3C5E3	E3C5D94B	D8E4C5E4	DEAD.LETTER.QUEU
000000000ECF6070	C5404040	40404040	40404040	40404040	E

Local Normal Queue DEAD.LETTER.QUEUE, Common=No:

\*

000000000ED32048	C4C5C1C4	4BD3C5E3	E3C5D94B	D8E4C5E4	DEAD.LETTER.QUEU
000000000ED32058	C5404040	40404040	40404040	40404040	E
000000000ED32068	40404040	40404040	40404040	40404040	
000000000ED32078	00000000	D4D8D8C4	00000001	00000000	....MQQD.....

## Dump Formatting Enhancements - Branch Trace

- Branch Trace is formatted for last 10 entries
  - Gives name of program and displacement into program
  - For CP, gives name of CP CSECT and displacement into CSECT

```
*LAST 10 BRANCH TABLE ENTRIES          CSECT      OFFSET*
      00000000 801C70BE          CCTAPE      000060BE
      00000000 8102C060          CIO         000014E8
      00000000 801C4852          CCTAPE      00003852
      00000000 8102BD98          CIO         00001220
      00000000 8020D528          CCOSAE      00004528
      00000000 8004C3E8          CCCLHR      000003E8
      00000000 8004CB44          CCCLHR      00000B44
      00000000 8004C330          CCCLHR      00000330
      00000000 80002780          CCNUCL      00002780
      00000000 8011A16C LAST CCDBAF      0000116C
```

## Dump Formatting Enhancements - Application Stack

- Each in use stack frame is formatted in a dump
- Formatting contains:
  - Name of the function using the stack
    - C function name includes prototype
  - Tags to identify fields in the stack, for example:
    - BCH = Back Chain
    - PAT = PAT address
    - ...many more...



# Dump Application Stack Formatting Example - C Function

```
*STACK FRAME FUNCTION- void __serrc(t_serrc status,int number,const char* msg,
                               void ** slist,char prefix,+0001 MORE PARAMETERS)
```

```
000000000CC0F490 BCH 00000000      0CC0FC48      00000000      0CD00400      .....
000000000CC0F4A0 R2 00000000      0CC0F710 R3 00000000      0CC0F430      .....7.....4.
000000000CC0F4B0 R4 00000000      0CC0F788 R5 00000000      0013306C      .....7h.....
000000000CC0F4C0 R6 00000000      00000002 R7 00000000      0CC0FBF0      .....0
000000000CC0F4D0 R8 00000000      00000001 R9 00000000      0AA00000      .....
000000000CC0F4E0 R10 00000000     00133012 R11 00000000     0092EA6C      .....k..
000000000CC0F4F0 R12 00000000     00002000 R13 00000003     8D8BB888      .....h
000000000CC0F500 R14 00000000     0012990A R15 00000000     0CC0F490      .....r.....4.
000000000CC0F510 F0 00000000     00000001 F2 00000003     8D8BBA10      .....
000000000CC0F520 F4 00000003     8D8B8DE4 F6 00000000     0CC0F4B0      .....U.....4.
000000000CC0F530 PAT 00000000     0A201420 CRE 00000003     8D8B7B40      .....
000000000CC0F540 BAS 00000000     0013306C TRN D8E9E9F8     0CD00500      .....QZZ8....
000000000CC0F550      00000000     00002000      0CC0FC48     8D8BB7F8      .....8
```

# Dump Application Stack Formatting Example - Assembler

```

*STACK FRAME FUNCTION- BXEY
0000000012C139C0 BCH 00000000      12C13C00      00000000      00000001      .....A.....
0000000012C139D0 R2  00000000      10A00000 R3  00000000      00135012      .....
0000000012C139E0 R4  00000000      00000020 R5  00000000      000000F8      .....8
0000000012C139F0 R6  00000000      12C13BC8 R7  0A320010      07F10008      .....A.H....1..
0000000012C13A00 R8  00000000      10A00008 R9  00000000      12C13BC8      .....A.H
0000000012C13A10 R10 00000000      000000F2 R11 00000000      0013506C      .....2.....
0000000012C13A20 R12 00000000      00000008 R13 00000000      00000001      .....
0000000012C13A30 R14 00000000      8D3C6B2A R15 00000000      0FA00000      .....
0000000012C13A40 F0  00000000      12C13C00 F2  00000000      000000F8      .....A.....8
0000000012C13A50 F4  00000000      10A00000 F6  00000000      001351FE      .....
0000000012C13A60 PAT 00000000      0FE030A0 CRE 00000000      00002000      .....
0000000012C13A70 BAS 00000000      0D3C6AD0 TRN C2E7C5E8      80000000      .....BXEY....
0000000012C13A80      00010000      00C13A08      00000000      0000E5F0      .....A.....V0

```

## Dump Formatting Enhancements - ECB Heap

- ECB Heap formatting includes:
  - Buffer address
  - Size of buffer
  - Obtaining load module
  - Displacement into load module where malloc() request was done
  - Buffer contents

### \*IN USE HEAP STORAGE

```

BUFFADDR-00000000CE02FF8  SIZE-0000000000008008  ECB SVA-0ABE4000  THRD ID-00000000
GET PROGRAM-CMQO  OFFSET-000000000002F36E  REL PROGRAM-N/A  OFFSET-N/A

```

```

00000000CE02FF8      E3E2C840      00000240      01043000      BBE78D09      TSH ... ..X..
00000000CE03008      285644E2      00000111      01F40000      D4E2C840      ...S.....4..MSH
00000000CE03018      000297DE      00000064      00000000      00000210      ..p.....
00000000CE03028      E7D8C840      00000001      D9D8F140      40404040      XQH ....RQ1

```

## Dump Formatting Enhancements - z/TPFDF

- Current SW00SR block is formatted
- Pointer to Database Interface Block (DBIFB) in ECB is labeled
  - DBS is the tag name
  - DBIFB is in ECB Heap - included in the dump

# Dump Formatting Enhancements - z/TPFDF SW00SR

## \*CURRENT SW00SR BLOCK

```

000000000C6001C0 BID E2E60000 PGM D8E7C4C8 KLS 00000000 KLS 036152AC SW..QXDH.....
000000000C6001D0 DET 00000000 WID B0750001 INB 001A00BF ILT 000100A5 .....v
000000000C6001E0 EOR 0000004D      00320000 RBV 0001D980 OP2 06040000 .....R.....
000000000C6001F0      00000000 ALG 0A30816C      D1000000      00000000 .....a.J.....
000000000C600200      00000000      00000000 FAD 00000000 FAD 435801EE .....
000000000C600210 IPT 00000000 IPT 0A308150      0000001A      00000000 .....a.....
000000000C600220      00000000      00000000      00000000      00000000 .....
000000000C600230      00000000      00000000 UKY 00000000 ID1 08082600 .....
000000000C600240 ID5 00010000 ID9 00000000 IDF 00000000      00000050 .....
000000000C600250 RTN 00000400 PNB 001A0005      00000000 REC 0A30A61A .....w.
000000000C600260 IPA 00000000      001AA000 NLR 00000000 ORD 0000003D .....
000000000C600270      00000000      00000000 SRV 00000000 BWF 00000000 .....
000000000C600280 BPT 00000000      00000000      00000000      00000000 .....
000000000C6002A0      00000000      00000000      0000B075      00000CE7 .....X
000000000C6002B0      C4C88001      FDAE0000      00000000      00000000 DH.....
000000000C6002C0      00000000      00000000      00000000      00000000 .....
000000000C600300      00000000      00000000      00000000      C4C6E3C4 .....DFTD
000000000C600310 CAL 00000004 RET 00000000 PM1 00000002 PM2 800200BE .....
000000000C600320 PM3 0A308150      00000000      00000000      00000000 ..a.....

```

## Dump Formatting Enhancements - Program Link Map

- Failing load module's link edit map is dumped for:
  - CTL-1, CTL-2, CTL-3, and OPR-4 dumps
  - SYSDUMP=YES

### \*\*\*\*\* ASSOCIATED LINK MAP DATA

```
000000008450E30 OBJECT FILE: /ztpf/bld/bas/obj/uaa0.o
                COMPILED ON: 2004/10/04 AT 19.12.26
                COMPILER: HLASM V10 R52
```

OFFSET	ADDRESS	FUNCTION NAME
00000000	000000008450E30	UAA0_BSO
00000008	000000008450E38	UAA9_BSO
00000010	000000008450E40	UAA2_BSO
00000040	000000008450E70	UAA0_INT31
00000066	000000008450E96	UAA0_INTN31
000000BE	000000008450EEE	UAA9_INT31
000000E4	000000008450F14	UAA9_INTN31
0000013C	000000008450F6C	UAA2_INT31
00000162	000000008450F92	UAA2_INTN31



## Dump Formatting Enhancements - Architectural Changes

- 64-bit Registers
- 128-bit PSWs
- 64-bit core addresses

### \*GENERAL REGISTERS

R0	00000000	00000004	00000000	00000001	R1	R2	00000000	0A308000	00000000	0C6001C0	R3
R4	00000000	0A306EA1	00000004	00000064	R5	R6	00000000	0A308150	00000000	83745CCA	R7
R8	00000000	08461408	00000000	0A300000	R9	R10	00000000	00000010	00000000	00001000	R11
R12	00000000	00002000	00000000	FFFFFFFF	R13	R14	00000000	00000000	00000000	0C50F4C0	R15

### \*CONTROL REGISTERS

C0	00000000	1CB4BEF0	00000000	33261007	C1	C2	00000000	00000000	00000000	70000000	C3
C4	00000000	00000000	00000000	00000000	C5	C6	00000000	81000000	00000000	63FF1007	C7
C8	00000000	0000FFFF	00000000	C0000000	C9	C10	00000000	03600000	00000003	92BFFFFFFF	C11
C12	80000000	010FA760	00000000	63FF1007	C13	C14	00000000	D7000000	00000000	00000000	C15



## Dump Formatting Architectural Example (continued)

### \*FLOATING POINT REGISTERS

```

Y0  43277D94 AE415D24 43277DB5 128336B6 Y1  Y2  43277DB5 128336B6 43277DB5 128336B6 Y3
Y4  00000000 00000000 43277DB5 128336B6 Y5  Y6  43277D94 AD7DEDE0 00000000 00000000 Y7
Y8  00000000 00000000 00000000 00000000 Y9  Y10 00000000 00000000 00000000 00000000 Y11
Y12 00000000 00000000 00000000 00000000 Y13 Y14 00000000 00000000 00000000 00000000 Y15

```

\*FLOATING POINT CONTROL REGISTER C0000000

\*CURRENT PSW AT TIME OF DUMP - 07150000 80000000 00000000 08461C54

### \*PSWS

	OLD				NEW			
* RESTART	00000000	00000000	00000000	00000000	0004C001	80000000	00000000	00000F80
* EXTERNL	07150000	80000000	00000000	08461C50	0404C001	80000000	00000000	0001A488
* SVC	07150000	80000000	00000000	08482BD4	05040001	80000000	00000000	00094098
* PROGRAM	07150000	80000000	00000000	08461C54	0004C001	80000000	00000000	000A4008
* MACHINE	0706C001	80000000	00000000	0004E71C	0000C001	80000000	00000000	000A0000
* I/O	07150000	80000000	00000000	08461C50	0404C001	80000000	00000000	0102C1C0

## Owners

- A new mechanism to track users of system resources
  - 32-byte name is given to users of system resources
    - CMBs, SWBs, 4-KB Frames, 1-MB Frames
    - Name associated when resource obtained (GETBC, GSWBC)
  - The name is comprised of
    - 8-byte high level qualifier
    - 8-byte mid level qualifier
    - 16-byte low level qualifier
    - For example ITAPE.CCWTRANS.REAL\_CCW\_BLOCK.

## Owners (continued)

- ECBs can be registered to an owner
  - New API will register the ECB - EOWNRC, tpf\_eownrc()
  - All resources obtained by the ECB after registration will use the registered owner name
- Online displays provided
  - ZSTAT OWNER.

# ZSTAT OWNER Example

ZSTAT OWNER BLOCK-FRM1MB

CSMP0097I 21.49.05 CPU-B SS-BSS SSU-HPN IS-01

STAT0023I 21.49.05 BLOCK OWNER DISPLAY

	IOB	FRAME	COMMON	SWB	ECB	FRM1MB
ALLOCATED	2704	5000	250	1416	300	300
AVAILABLE	2704	4952	247	1376	293	233

—

	IOB	FRAME	COMMON	SWB	ECB	FRM1MB
ICRPA						53
ISYSHEAP						13
IECB						1
ITCPIP						1

END OF DISPLAY+

## Program Trace Name Enhancements

- Standardized mechanism to modify program name used for tracing
  - No longer need to modify program name at 4(,R8)
  - PNAMC, `tpf_pnamc()` can be used to modify trace name
    - Set your own trace name
    - Use the caller of this routine
    - Use the name in the PAT
  - Automatically setup on enter
- Used:
  - Macro trace
  - Dump messages
  - File a record
  - Logging pool gets and releases.

## Dump Message Enhancements

- Message includes:
  - Trace name
  - Load module name
  - Object name
  - Displacement into the object where dump happened
  - If dump is in CP, CP CSECT name
  - If CTL-1, CTL-2, CTL-3, or OPR-4
    - Instruction length code + Program-interruption code
    - Failing instruction
  - If no core available dump
    - Summary of Owners information
- Dump Tape VSN where dump started is displayed.

# Dump Message Examples

## Standard OPR dump message:

```
CPSE0152E 12.07.10 IS-0001 SS-BSS  SSU-BSS  SE-000499 OPR-IDBC161
010000A TRC-QXEN
CTDF      OBJ-ufgq.goff          000003E6  LOADSET-BASE
DFOPN -- REFERENCE NAME SHORTER THAN 8 BYTES +
```

## OPR-4 dump message:

```
CSMP0097I 09.06.59 CPU-B SS-BSS  SSU-HPN  IS-01
CPSE0156I 09.06.59 DUMP SEQUENCE NUMBER 2380 STARTED ON VSN 000230+
CSMP0097I 09.07.56 CPU-B SS-BSS  SSU-HPN  IS-01
CPSE0152E 09.06.59 IS-0001 SS-BSS  SSU-HPN  SE-002380 OPR-I000004
010000B TRC-CVXS
CVXS      OBJ-cvxsf4.goff        00000164  LOADSET-BASE
PSW      07150001 80000000  00000003 8C5EF4CC  PIC 003B ILC 0004 I-50F02070
R0-2     00000000 00000008  00000003 8C5EF3BE  00000000 8CD04780
R3-5     00000000 0CD07028  BBF0067F ED5CAA2A  00000000 0AA0F000
R6-8     00000000 0AA101D8  00000000 0AA00390  00000000 0CC0E3E0
R9-11    00000000 0AA00000  00000000 00014BB0  00000000 00001000
R12-14   00000000 00002000  00000003 8C5EF4AC  00000000 00000000  _
R15      00000000 00000003
```



## Dump Message Examples (continued)

### No core dump message:

```
CPSE0151T 09.21.21 IS-0001 SS-BSS  SSU-HPN  SE-002579 CTL-I064C02 CATASTROPHIC
                                CSECT-CCNUCL  0000CA12
```

NO FRAMES AVAILABLE

	IOB	FRAME	COMMON	SWB	ECB	FRM1MB
ALLOCATED	0002504	0001000	0000100	0001092	0000050	0000000220
AVAILABLE	0002502	0000000	0000096	0000065	0000047	0000000153
	IOB	FRAME	COMMON	SWB	ECB	FRM1MB
ITAPE		0000993				
IECB		0000004				
ISYSTEM		0000003				
END OF OWNERS DISPLAY						

## Dump Enhancements - OPR dumps

- Enhanced ability to select core areas in an OPR dump
  - Ability to dump 4-KB of core around each register
  - Ability to include collated macro trace and I/O trace
  - Ability to include prefix pages of all I-streams
  - ZIDOT extended to OPR dumps
    - Ability to include core areas in OPR dumps
      - System tables
      - User tables.

## Dump Enhancements - Dump Groups

- Problem: If a core area is to be included in OPR dumps for 50 programs, doing 50 ZIDOTs is time consuming and error prone.
- Solution: Dump Groups
  - Set of programs which share ZIDOT dump characteristics
    - One ZIDOT command includes tables in OPR dumps for many programs
  - Defined in control file
  - Online maintained in PAT
    - ZDPAT to display
    - ZAPAT to change.

## Dump Enhancements - Select by CP CSECT

- Problem: If a core area is to be included in all dumps in a functional area of the CP, ZIDOTs for each dump is time consuming and error prone.
- Solution: Select core areas for dumps in the CP by CP CSECT
  - ZIDOT enhanced to support additions and alters
  - IDOTC macro created
  - Some examples...
    - CCTAPE dumps include ITAPE and ISWBUSE
    - CCSONS dumps include IMFST and IIOBS
    - CCSONP dumps include IPOOL and IRIAT
  - Selection used for CTL-1, CTL-2, CTL-3, OPR-4, and SYSDUMP=YES dumps.

## Dump Enhancements - Dump by Owner

- Ability to dump blocks in a resource by Owner name
- Example, if ISWBUSE is selected...
  - A tape dump only needs SWBs which are owned by ITAPE
  - A MPIF dump only needs SWBs which are owned by MPIF
- ZIDOT provides ability to select by Owner
  - IDOTB provides same capability.

## Dump Enhancements - Named Manual Dumps

- Problem: ZDUMP does not give enough information and ZDUMP ALL gives too much information
- Solution: Named manual dumps
  - ZDUMP has list of manual dumps
    - Predefined list targets specific functions
      - Example: Tape or SNA
    - ZIDOT can alter existing named manual dumps
    - IDOTM macro defines named manual dumps
  - Create your own manual dump
    - ZIDOT can create new named manual dumps
    - Or include IDOTM definitions in CUDP.CPY
    - You target the data that is dumped.

# Named Manual Dump Example

## Predefined Manual dump:

```
ZDUMP IBMSNA
CSMP0097I 12.47.19 CPU-B SS-BSS  SSU-HPN  IS-01
CPSE0156I 12.47.19 DUMP SEQUENCE NUMBER 2778 STARTED ON VSN 000230+
CSMP0097I 12.47.19 CPU-B SS-BSS  SSU-HPN  IS-01
DUMP0000I 12.47.19 ZDUMP-OK+
CSMP0097I 12.47.19 CPU-B SS-BSS  SSU-HPN  IS-01
CPSE0115I 12.47.19 IS-0001 SS-BSS  SSU-HPN  SE-002778 MANUAL DUMP
010000B TRC-CVFM
CVFM      OBJ-cvfm.goff          0000024E  LOADSET-BASE
NAMED MANUAL DUMP-      IBMSNA    +
```



## Named Manual Dump Example (continued)

### Create my own Named Manual dump:

```
ZIDOT INCLUDE MYMANUAL IMFST
CSMP0097I 12.54.45 CPU-B SS-BSS SSU-HPN IS-01
IDOT0002I 12.54.45 KEYWORD IMFST INCLUDED FOR NAMED MANUAL DUMP - MYMANUAL
IDOT0001I 12.54.45 COMPLETED+
```

```
ZDUMP MYMANUAL
CSMP0097I 12.55.44 CPU-B SS-BSS SSU-HPN IS-01
CPSE0156I 12.55.44 DUMP SEQUENCE NUMBER 2779 STARTED ON VSN 000230+
CSMP0097I 12.55.44 CPU-B SS-BSS SSU-HPN IS-01
DUMP0000I 12.55.44 ZDUMP-OK+
CSMP0097I 12.55.44 CPU-B SS-BSS SSU-HPN IS-01
CPSE0115I 12.55.44 IS-0001 SS-BSS SSU-HPN SE-002779 MANUAL DUMP
010000B TRC-CVFM
CVFM OBJ-cvfm.goff 0000024E LOADSET-BASE
NAMED MANUAL DUMP- MYMANUAL +
```

## Dump Enhancements - No Core Dumps

- CTL-C dump number no longer used
- Unique dump number for each resource
  - 064C00 - System out of SWBs
  - 064C01 - System out of Common Blocks
  - 064C02 - System out of 4-KB frames
  - 064C03 - System out of 1-MB frames
  - 064C04 - System out of ECBs
- Allows customization
  - In TPF 4.1 need to dump all possible storage areas on a CTL-C
  - Ability to dump only what is needed
    - Example ... 064C00 - do not need to dump common blocks.

## Dump Enhancements - Additional ZIDOT Enhancements

- Provide ability to take a dump if SERRC is in NODUMP table
  - ZIDOT FORCE
- Provide ability to take a dump if SNAPC is in NODUMP table
  - ZIDOT SNAP FORCE

## LDEV Trace Enhancements

- Ability to define the number of trace entries for all LDEVs
  - Does not require a reassembly of IPLB
- Ability to dynamically increase the number of trace entries for a range of SDAs
- Ability to display LDEV trace online
- ZIOTR provides all of these new capabilities.

```
ZIOTR SET COUNT 420 TO 422 575
CSMP0097I 13.34.33 CPU-B SS-BSS SSU-HPN IS-01
IOTR0002I 13.34.33 TRACE COUNT SET DISPLAY
THE NUMBER OF TRACE ENTRIES FOR SDA 0420 IS SET TO 575
THE NUMBER OF TRACE ENTRIES FOR SDA 0421 IS SET TO 575
THE NUMBER OF TRACE ENTRIES FOR SDA 0422 IS SET TO 575
END OF DISPLAY+
```

## LDEV Online Trace Display Example

```
ZIOTR DISPLAY TRACE 420 15
CSMP0097I 13.37.47 CPU-B SS-BSS SSU-HPN IS-01
IOTR0010I 13.37.47 TRACE ENTRIES, SUBCHANNEL 0001000F, LDEV 00000000012E5E00
ISDAC BBF04310 E2D68C17 01 00 00013994 00000004 01D2AC00
ISDAC BBF042C4 C914C41A 01 00 00013994 00000004 01B3CC00
INT BBF0429A 58E8A6DE 01 00C04007 01F5F020 0C000000 001C44E8
SIOSC BBF0429A 58E5A79E 01 00 001C428A 01F5F018 00 80
ISDAC BBF0429A 58E5999E 01 00 001C6176 00000000 001D1720
INT BBF0429A 58E57A5E 01 00C04007 01F5F020 0C000000 001C44E8
SIOSC BBF0429A 58E2109E 01 00 001C428A 01F5F018 00 80
ISDAC BBF0429A 58E2065E 01 00 001C6176 00000000 001D1720
INT BBF0429A 58E1E15E 01 00C04007 01F5F020 0C000000 001C44E8
SIOSC BBF0429A 58DE6C9E 01 00 001C428A 01F5F018 00 80
ISDAC BBF0429A 58DE609E 01 00 001C6176 00000000 001D1720
INT BBF0429A 58DE355E 01 00C04007 01F5F020 0C000000 001C44E8
SIOSC BBF0429A 58D9421E 01 00 001C428A 01F5F018 00 80
ISDAC BBF0429A 58D934DE 01 00 001C6176 00000000 001D1720
INT BBF0429A 58D9145E 01 00C04007 01F5F020 0C000000 001C44E8
END OF DISPLAY+
```

## SNAPC Enhancement

- ECB trace can be included in a SNAPC
  - Collated macro and C function trace
  - Last 40 trace items.



## Added Protection in ECB Private Area

- 4-KB frame allocation in the ECB Private Area alternates:
  - Valid 4-KB frame
  - Unmapped 4-KB frame
- If a program references beyond the end of a 4-KB block, an OPR-4 will happen
- This is one capability of TPF 4.1 BLKCHK mode.
  - Running BLKCHK mode is processor and memory intensive
- Using this allocation method z/TPF gains some capabilities of BLKCHK mode without the overhead
- Improved ability to detect logic errors as they happen
  - Prevents core corruption.

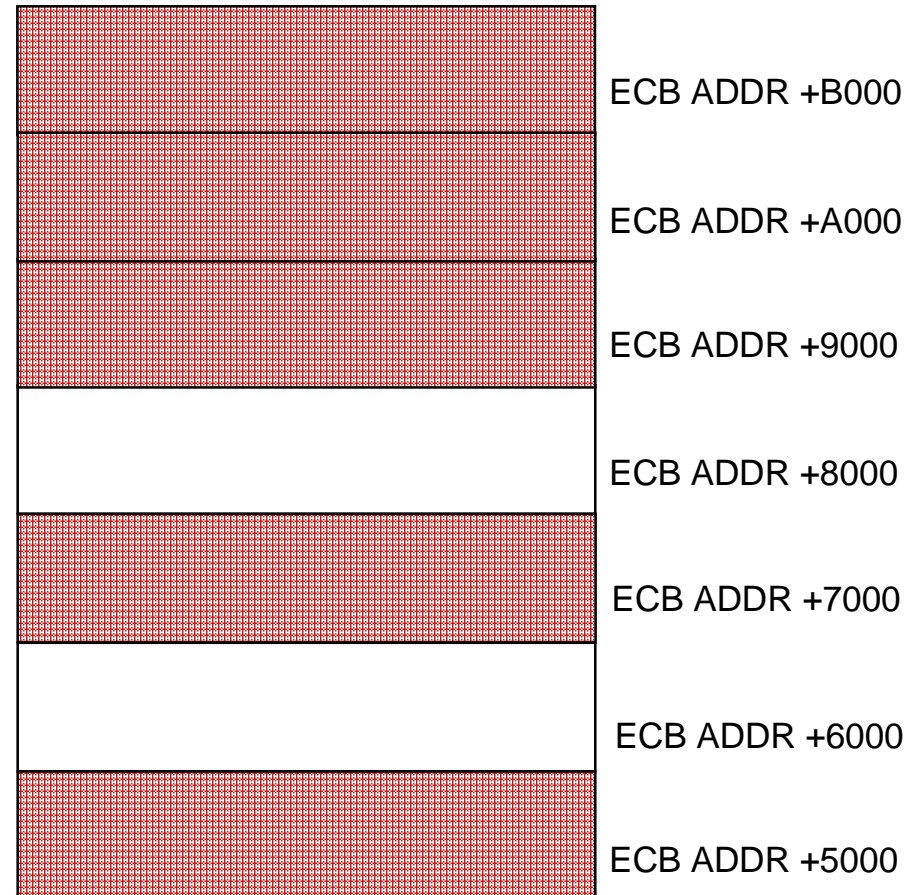


# ECB Private Area Representation

Will be 3rd 4-KB Frame →

2nd 4-KB Frame →

1st 4-KB Frame →



## New Diagnostic Tool - ECB Heap HEAPCHECK Mode

- System-wide debugging mode intended for test system use only
- Similar to block check mode in the ECB private area (EPA)
- Intended to detect common programming errors in the usage of the ECB heap *before* applications are released to production environments
  - Applications that write beyond the end of allocated ECB heap buffers
  - Applications that continue to access ECB heap buffers after they have been released
- Allows ECB heap corruption and unauthorized accesses to be detected at the exact moment they occur
- ZSTRC ALTER (NO)HEAPCHECK
  - When turned on/off, the change takes effect immediately for all newly created ECBs. No IPL necessary!
  - All existing ECBs continue to operate with the HEAPCHECK setting that was in effect when they were created.

## ZDMAP - Locate a Program

- Displays where a program resides in memory and program size
- Displays the objects in that program, object location, object size, and compile date and time

### ZDMAP CTS7

CSMP0097I 13.05.38 CPU-B SS-BSS SSU-HPN IS-01

DMAP0003I 13.05.38 LINK MAP DATA DISPLAY

CTS7 ACTIVE IN LOADSET BASE IN SUBSYSTEM BSS

PROGRAM ADDRESS 000000038CBEE000

PROGRAM SIZE 00007000

cts7 - OBJ FILE AT ADDR 000000038CBEF550

OBJECT FILE SIZE 00001B7C

COMPILED ON 2004/09/27 AT 08.56.21

END OF DISPLAY

# ZDMAP - Locate an Object Within a Program

## ZDMAP COMX OBJ-cgthbn

CSMP0097I 13.05.21 CPU-B SS-BSS SSU-HPN IS-01

DMAP0003I 13.05.21 LINK MAP DATA DISPLAY

COMX ACTIVE IN LOADSET BASE IN SUBSYSTEM BSS

PROGRAM ADDRESS 000000038D86A000

PROGRAM SIZE 00018000

cgthbn

- OBJ FILE AT ADDR 000000038D870710

OBJECT FILE SIZE

00000C10

COMPILED ON 2004/09/27 AT 08.45.24

END OF DISPLAY

# ZDMAP - Locate the Object Where a Function Resides

## ZDMAP COMX FUNC-codepage\_conv

CSMP0097I 13.07.58 CPU-B SS-BSS SSU-HPN IS-01

DMAP0003I 13.07.58 LINK MAP DATA DISPLAY

COMX ACTIVE IN LOADSET BASE IN SUBSYSTEM BSS

PROGRAM ADDRESS 000000038D86A000

PROGRAM SIZE 00018000

cgthba

- OBJ FILE AT ADDR 000000038D86F25C

OBJECT FILE SIZE

00001120

COMPILED ON 2004/09/27 AT 08.45.25

OFFSET	ADDRESS	FUNCTION NAME
--------	---------	---------------

00000000	000000038D86F25C	codepage_conv
----------	------------------	---------------

END OF DISPLAY

## ZDMAP - Determine the Object and Function Given an Address

ZDMAP ADDR-38D86F333

CSMP0097I 13.03.08 CPU-B SS-BSS SSU-HPN IS-01

DMAP0003I 13.03.08 LINK MAP DATA DISPLAY

COMX ACTIVE IN LOADSET BASE IN SUBSYSTEM BSS

PROGRAM ADDRESS 000000038D86A000

PROGRAM SIZE 00018000

cgthba - OBJ FILE AT ADDR 000000038D86F25C

OBJECT FILE SIZE 00001120

COMPILED ON 2004/09/27 AT 08.45.25

OFFSET	ADDRESS	FUNCTION NAME
--------	---------	---------------

00000000	000000038D86F25C	codepage_conv
----------	------------------	---------------

000000038D86F333 IS AT OFFSET 000000D7 INTO OBJECT FILE cgthba

END OF DISPLAY

# ZDMAP - Display Object File Source Information

## ZDMAP CTSR INFO-spath

CSMP0097I 13.09.53 CPU-B SS-BSS SSU-HPN IS-01

DMAP0003I 13.09.53 LINK MAP DATA DISPLAY

CTSR ACTIVE IN LOADSET BASE IN SUBSYSTEM BSS

PROGRAM ADDRESS 000000038CB95000

PROGRAM SIZE 00006000

/tpf51/bld/base/obj

- OBJ FILE AT ADDR 000000038CB96690

OBJECT FILE SIZE

000012FC

COMPILER INFO - HLASM V10 R52

COMPILED ON 2004/09/27 AT 08.56.34

END OF DISPLAY



# ZDMAP - Display Object File Source Information

## ZDMAP CNSD INFO-SPATH

CSMP0097I 11.51.29 CPU-B SS-BSS SSU-HPN IS-01

DMAP0003I 11.51.29 LINK MAP DATA DISPLAY

CNSD ACTIVE IN LOADSET BASE IN SUBSYSTEM BSS

PROGRAM ADDRESS 000000038C77B000

PROGRAM SIZE 0002E000

/tpf51/bld/base/obj

- OBJ FILE AT ADDR 000000038C77DABC

OBJECT FILE SIZE

000000C0

COMPILER INFO - GCC: (GNU) 3.4-gnupro-04r1-3

COMPILED ON 2004/09/27 AT 10.17.14

## Enhanced Disassembler

- Ability to disassemble code in core storage using:
  - ZDCOR
  - ZDPGM
  - TPF Debugger
- Updated to support the z/Architecture instruction set
- Macro parameters are disassembled as well now.

# Disassembler Example

ZDCOR 38CB96D96.44 INST

```
CSMP0097I 12.59.54 CPU-B SS-BSS SSU-HPN IS-01
DCOR0010I 12.59.54 BEGIN DISPLAY
000000038CB96D96 0A32 0008 RELCC D1
000000038CB96D9A 0A2C 0851 GETCC D1,L4
000000038CB96D9E E360 1250 0016 LLGF R6,592(,R1)
000000038CB96DA4 A779 0FFF LGHI R7,X'FFF'
000000038CB96DA8 A7F9 0FFF LGHI R15,X'FFF'
000000038CB96DAC 0EE6 MVCL R14,R6
000000038CB96DAE D703 9080 9080 XC 128(4,R9),128(R9)
000000038CB96DB4 D201 9080 89D4 MVC 128(2,R9),2516(R8)
000000038CB96DBA 0A16 0800 FILEC D1,TAG=Y,GDS=Y
000000038CB96DBE E3E0 0C28 0017 LLGT R14,3112
000000038CB96DC4 A7EB 0780 AGHI R14,X'780'
000000038CB96DC8 E3E0 E000 0004 LG R14,0(,R14)
000000038CB96DCE B904 001E LGR R1,R14
000000038CB96DD2 A775 00DC JAS R7,X'DC'
000000038CB96DD6 D601 1290 1290 OC 656(2,R1),656(R1)
END OF DISPLAY
```

## Ability to View Program Listings Online

### ZDPGM XODD 78.20 LIST

```
CSMP0097I 09.18.06 CPU-B SS-BSS  SSU-HPN  IS-01
DPGM0010I 09.18.06 BEGIN DISPLAY OF FILE COPY FOR XODD.XODD
                VERSION TA IN PROGRAM BASE 2 ACTIVE IN LOADSET TESTZ
0078 58F0 9298          3364          GLOBZ REGR=R15
007C 9015 9014          3896          STM   R1,R5,EBW012      SAVE REGS
0080 58E0 9058          3897          L     R14,EBW080      LOAD XTAT TERML
0084 50E0 9064          3898          ST    R14,EBW092      STORE XTAT ENTRY
0088 4860 82C4          3899          LH    R6,XGAH45      LOAD MAX MSG COU
008C 9516 E012          3900          CLI   XZ1IN1,X'16'   IS THIS 23RD XOLD
0090 4770 8028          3901          BNE   XGNLST         NO, BRANCH
0094 4160 000A          3902          LA    R6,10          LOAD MAX OF 10 O
END OF DISPLAY - ZEROED LINES NOT DISPLAYED
```

## View Program Listing - Example 2

- Can specify the starting column of the listing to display
- Enables you to view the desired sections of the listing.

```
ZDPGM XODD 78.20 LIST STCOL-30
```

```
CSMP0097I 11.50.26 CPU-B SS-BSS SSU-HPN IS-01
```

```
DPGM0010I 11.50.26 BEGIN DISPLAY OF FILE COPY FOR XODD.XODD
```

```
VERSION TA IN PROGRAM BASE 2 ACTIVE IN LOADSET TESTZ
```

```
  GLOBZ REGR=R15
```

```
  STM   R1,R5,EBW012
```

```
        SAVE REGS
```

```
  L     R14,EBW080
```

```
        LOAD XTAT TERML FLD CORE ADDRESS
```

```
  ST    R14,EBW092
```

```
        STORE XTAT ENTRY ADDR
```

```
  LH    R6,XGAH45
```

```
        LOAD MAX MSG COUNT OF XOLD
```

```
  CLI   XZ1IN1,X'16'
```

```
        IS THIS 23RD XOLD RCD
```

```
  BNE   XGNLST
```

```
        NO, BRANCH
```

```
  LA    R6,10
```

```
        LOAD MAX OF 10 ON RCD 23
```

```
END OF DISPLAY - ZEROED LINES NOT DISPLAYED
```

## New Debugger Capabilities

- Can now debug programs built with optimization
  - Can debug same code loaded to production system
- Linkage into DF library standardized
  - Can trace calls into the DF library
- Can use debugger against a captured file
  - Dumps and snapshots of running ECBs can be written to HFS file
  - Debugger makes it easier to locate program variables
    - And other data using expression evaluator
- Can use the trace log utility in conjunction with debugger
  - Can turn trace log on and off from the debugger
  - Allows you to create a trace log between two breakpoints.

# Applications Productivity Enhancements Design and Coding Enhancements z/TPF Globals Support



## Overview of Existing Globals Support

- In TPF 4.1, globals are a portion of fixed main storage that are designed to contain application records; globals permit fast and efficient communication among application programs and between application programs and the control program
- In z/TPF, globals support as it existed in TPF 4.1, remains essentially the same, with a few modifications and enhancements
- Modifications:
  - No longer support 24-bit globals
    - Existing TPF 4.1 global layout is still supported
      - Separate primary and extended global areas, both containing 31-bit globals
  - Global attribute table (GAT) required even if extended global areas are not present

## Enhancements to Existing Globals Support in z/TPF

- Easier for application programs to update globals
  - C/C++ applications now have the ability to modify protected globals directly
  - Existing TPF 4.1 methods to modify globals are still supported
    - `glob_modify( )`, `glob_update( )`, and `global( )`
- Faster global restart
  - Reduces the time required to IPL the z/TPF system
- New synchronization option
  - Provides capability for applications to receive notification when global synchronization has completed across all processors in the z/TPF complex
    - Enhances data integrity and system availability
  - New assembler and C/C++ API options
    - ASM: `SYNCC` macro with `WAIT=YES|NO`
    - C/C++: `tpf_glob_sync_wait( )` function

## Globals Support in z/TPF

- **Format-1 globals:** The same globals support that exists in TPF 4.1
  - 4 global areas (GL1, GL2, GL3, and GL4)
  - Global storage allocator (GOA) records on DASD
  - Global directories (GL0BA and GL0BY) in main storage
  - Application macros: GLOBZ, GLOUC, FILKW, SYNC, etc.
- **Format-2 globals:** New globals support for z/TPF
  - Used and managed independently of format-1 globals
  - Enhanced features unique to z/TPF

## Format-2 Globals Design Initiatives

- Exploit 64-bit architecture
- No impacts to existing applications
- Make globals easy to use
- Provide robust functionality
- Enable globals to be highly extensible
- Offer an easy method for migrating existing applications to the new globals support

## Similar Attributes Between Format-1 and Format-2 Globals

- SSU-unique or SSU-shared
- IS-unique or IS-shared
- Processor-unique or processor-shared
- Protected or unprotected
- Keypointable or non-keypointable
- Synchronizable or non-synchronizable

## Unique Features of Format-2 Globals Support

- Global records only; no global fields
- Reside in dynamically allocated main storage
- Can reside above or below 2-GB bar
- Managed by a new set of z/TPF commands
- Accessed by new z/TPF-unique APIs
- Keypointable globals may be > 1055 bytes
- Synchronizable globals may be > 1055 bytes
- No limit on the size of global data
- No limit on the number of global definitions
- Access format-1 global data with new format-2 globals APIs



## Much Easier to Manage Format-2 Globals

- New set of z/TPF commands that are used to manage format-2 globals
  - No SIP requirements
  - No PILOT tapes
  - No macros or header files that must be kept current
- ZGLBL command
  - Define, alter, and delete format-2 globals definitions
  - Initialize format-2 globals data
  - Define and display z/TPF system control information
- ZDGBL command
  - Display attributes for a format-2 global
  - Display data for a format-2 global
- ZAGBL command
  - Alter data for a format-2 global



## Defining Format-2 Globals

- ZGLBL GLOBAL DEFINE
  - Create a format-2 global definition
  - Global names are 8-characters long
    - Valid characters: A-Z, a-z, 0-9, and '\_'
    - Padded on the right with blanks if less than 8 characters
  - Attributes
    - Uniqueness: SSU, IS, processor
    - Load: restart, cycle-up, on-demand
    - Protected or not protected
    - Location: 31-bit or 64-bit
    - Keypointable, synchronizable, or neither
  - LIKE parameter: define a new global with same attributes as an existing global

## Defining Format-2 Globals Example

- Only need to issue the ZGLBL GLOBAL DEFINE command once
  - The global definition is viewable by all SSUs within the SS, all processors in the L/C complex, and all I-streams
  - The global *cannot* be accessed by an application until it has been initialized with data
- ZGLBL GLOBAL DEFINE example:

```
ZGLBL GLOBAL DEFINE _Glob2 LOC-64 SSU-N IS-N PROC-Y KEY-Y SYNC-N PROT-Y  
LOAD-DEMAND
```

```
CSMP0097I 13.35.22 CPU-B SS-BSS SSU-HPN IS-01  
GLBL0050I 13.35.22 FORMAT-2 GLOBAL RECORD _Glob2 DEFINED SUCCESSFULLY
```

```
ZGLBL GL DEF _myglob LIKE-_Glob2
```

```
CSMP0097I 13.35.22 CPU-B SS-BSS SSU-HPN IS-01  
GLBL0050I 13.35.22 FORMAT-2 GLOBAL RECORD _myglob DEFINED SUCCESSFULLY
```

## Initializing Format-2 Globals

- ZGLBL GLOBAL INITIALIZE
  - Initialize the data for a format-2 global
  - Must initialize global before it can be used
  - Specify the source from which the global is to be initialized
    - ASDEFINED: reserves specified amount of storage only
    - ZERO: reserves specified amount of storage and clears it
    - GLOBAL: copies data from another format-2 global
    - INPUTDECK: uses global data input deck created online or offline
  - z/TPF system reserves storage, both in memory and on DASD, for the new global data

## Creating Format-2 Globals Data Input Deck

- Use a global data input deck with ZGLBL GLOBAL INITIALIZE
- User can use any mechanism to create the input deck
  - Create the input deck on any offline platform and transfer it to the online z/TPF system
    - GLINIT: sample tool available to build global data input deck
      - z/OS: JCL interface creates a general data set (GDS), HFS, tape, or VRDR
      - LINUX: command line interface creates HFS
  - Create the input deck on the online z/TPF system itself

## Initializing Format-2 Globals Example

- Create global data input deck on LINUX, and FTP it to the z/TPF file system into file /temp/globaldata/\_myglob.data
- Define a temporary data definition name (DDNAME) on z/TPF

```
ZDSMG DEF MYGLOBDDNAME HFS-'/temp/globaldata/_myglob.data'  
CSMP0097I 13.35.22 CPU-B SS-BSS SSU-HPN IS-01  
DSMG0001I 13.35.22 DDNAME MYGLOBDDNAME DEFINED
```

- Use the DDNAME as input to ZGLBL GLOBAL INITIALIZE command

```
ZGLBL GLOBAL INITIALIZE _myglob SOURCE-INPUT INPUT-MYGLOBDDNAME PROC-ALL  
CSMP0097I 13.35.22 CPU-B SS-BSS SSU-HPN IS-01  
GLBL0200I 13.35.22 FORMAT-2 GLOBAL RECORD _myglob SUCCESSFULLY INITIALIZED
```

## Reinitializing Format-2 Globals

- ZGLBL GLOBAL INITIALIZE can also be used to dynamically reinitialize the contents of a format-2 global while the z/TPF system is running (even in NORM state)
- Allows global data to be reloaded multiple times in a month, week, or even a single day, without disrupting the flow of the z/TPF system
- Backup copy of the original data is stored on DASD
- Backup copy of the original data remains in memory until it is no longer being used
- After the reinitialization, all new accesses of the format-2 global will occur for the new data



## Miscellaneous Format-2 Globals Commands

- ZGLBL GLOBAL ALTER
  - Alter attributes for an existing format-2 global definition
- ZGLBL GLOBAL DELETE
  - Delete an unused format-2 global definition
- ZGLBL GLOBAL UNDO
  - Reverse an unwanted or unintentional initialization or deletion
  - Undo of initialization or deletion can be processed up until the user-specified expiration time limit has elapsed
- ZGLBL GLOBAL RELEASE
  - Force release of backup copies of global data on DASD



## More Miscellaneous Format-2 Globals Commands

- ZGLBL CANCEL and ZGLBL CONTINUE
  - Operator confirmation for certain actions
- ZGLBL SET
  - Set subsystem-wide format-2 globals control values (eg, keypoint interval, expiration time limit, and load preference)
- ZGLBL DISPLAY
  - Display subsystem-wide control values
  - Monitor resource utilization

## Displaying Format-2 Globals

- ZDGBL
  - Display an individual format-2 global's attributes
  - Display all format-2 globals defined to a subsystem
  - Display the data for an individual format-2 global
  - Redirect output of display to an HFS file
    - Create log of global definitions and attributes
    - Capture an individual global's data on-demand

# Displaying a Format-2 Global's Attributes Example

ZDGBL \_glob1

CSMP0097I 13.35.22 CPU-B SS-BSS SSU-HPN IS-01

DGBL0025I 13.35.22 GLOBAL RECORD DISPLAY

GLOBAL NAME: \_glob1  
SUBSYSTEM USER: HPN  
ISTREAM NUMBER: 1  
ADDRESS: 0000000880184000  
SIZE (BYTES): 0000000000001000  
RECORD TYPE: #IF2G  
GORD ORDINAL: 0000000000000002

LOCATION: 64BIT  
PROTECTED: NO  
SSU: SHARED  
PROC: SHARED  
IS: MEMORY  
LOAD: RESTART  
KEYPOINTABLE: NO  
SYNCHRONIZABLE: YES  
CONTROLLED BY: SYSTEM

END OF DISPLAY+

## Displaying All Format-2 Globals Example

ZDGBL ALL

CSMP0097I 13.35.22 CPU-B SS-BSS SSU-HPN IS-01

DGBL0026I 13.35.22 BEGIN DISPLAY ALL FORMAT-2 GLOBAL RECORDS IN BSS SUBSYSTEM

GLOBAL	PROC	SSU	IS	PROT	LOC	LOAD	KEYPT	SYNC	CONTROL
_glob1	N	N	M	N	64	RES	N	Y	SYSTEM
_myglob	Y	Y	Y	Y	64	DEM	Y	N	SYSTEM
_Glob2	Y	N	N	Y	31	CYC	Y	N	SYSTEM

END OF DISPLAY +

## Displaying Format-2 Global Data Example

ZDGBL \_glob1 0.30 COPY-B

CSMP0097I 13.35.22 CPU-B SS-BSS SSU-HPN IS-01

DGBL0010I 13.35.22 BEGIN DISPLAY

FILE COPY FOR \_glob1 PROC-B SSU-HPN IS-1

0000000000000000- A9E3D7C6 89A2C799 8581A3A9 E3D7C689 zTPFisGr eatzTPFi

0000000000000010- A2C79985 81A3A9E3 D7C689A2 C7998581 sGreatzT PFisGrea

0000000000000020- A3A9E3D7 C689A2C7 998581A3 A9E3D7C6 tzTPFisG reatzTPF

CORE COPY FOR \_glob1 PROC-B SSU-HPN IS-1

0000000880184000- A9E3D7C6 89A2C799 8581A3A9 E3D7C689 zTPFisGr eatzTPFi

0000000880184010- A2C79985 81A3A9E3 D7C689A2 C7998581 sGreatzT PFisGrea

0000000880184020- A3A9E3D7 C689A2C7 998581A3 A9E3D7C6 tzTPFisG reatzTPF

END OF DISPLAY - ZEROED LINES NOT DISPLAYED+

## Modifying Format-2 Global Data Example

- ZAGBL
  - Change the data for an individual format-2 global
  - Change the data in memory, on DASD, or both

```
ZAGBL _glob1 0 C1C2C3C4 VALDATA-A9E3D7C689A2C799 COPY-c
CSMP0097I 13.35.22 CPU-B SS-BSS SSU-HPN IS-01
AGBL0010I 13.35.22 BEGIN DISPLAY
                CORE COPY FOR _glob1 PROC-B SSU-HPN IS-1
0000000880184000- A9E3D7C6 89A2C799 8581A3A9 E3D7C689 zTPFisGr eatzTPFi
ALTERED TO - _
0000000880184000- C1C2C3C4 89A2C799 8581A3A9 E3D7C689 ABCDisGr eatzTPFi
END OF DISPLAY - ZEROED LINES NOT DISPLAYED+
```

## Format-2 Globals Subsystem User (SSU) Groups

- Allow SSU-unique globals to be shared among a subset of SSUs
- Data resides in the "owning" SSU
- Manage SSU groups with z/TPF commands
  - ZGLBL SSUGROUP DEFINE
  - ZGLBL SSUGROUP ALTER
  - ZGLBL SSUGROUP DELETE
  - ZGLBL DISPLAY SSUGROUP
- SSU-unique globals may include as many predefined SSU groups as necessary
  - ADDGROUP parameter on ZGLBL GLOBAL DEFINE and ZGLBL GLOBAL ALTER
  - REMGROUP parameter on ZGLBL GLOBAL ALTER



## Using Format-2 Globals SSU Groups

```
WP/ZGLBL SSUGROUP DEFINE group1 OWNER-WP2 ADDSSU-WP1.WP2
```

```
CSMP0097I 12.25.50 CPU-B SS-WP SSU-WP1 IS-01
```

```
GLBL0700I 12.25.50 FORMAT-2 GLOBAL SSU GROUP group1 DEFINED SUCCESSFULLY
```

```
WP/ZGLBL GLOBAL DEFINE _globwp LOC-31 SSU-Y PROC-N IS-N KEY-N SYNC-N
```

```
ADDGROUP-group1
```

```
CSMP0097I 12.25.50 CPU-B SS-WP SSU-WP1 IS-01
```

```
GLBL0050I 12.25.50 FORMAT-2 GLOBAL RECORD _globwp DEFINED SUCCESSFULLY
```

```
WP/ZGLBL GLOBAL INIT _globwp SO-ZERO SIZE-5000 SSUGROUP-group1
```

```
CSMP0097I 12.25.50 CPU-B SS-WP SSU-WP1 IS-01
```

```
GLBL0200I 12.25.50 FORMAT-2 GLOBAL RECORD _globwp SUCCESSFULLY INITIALIZED
```

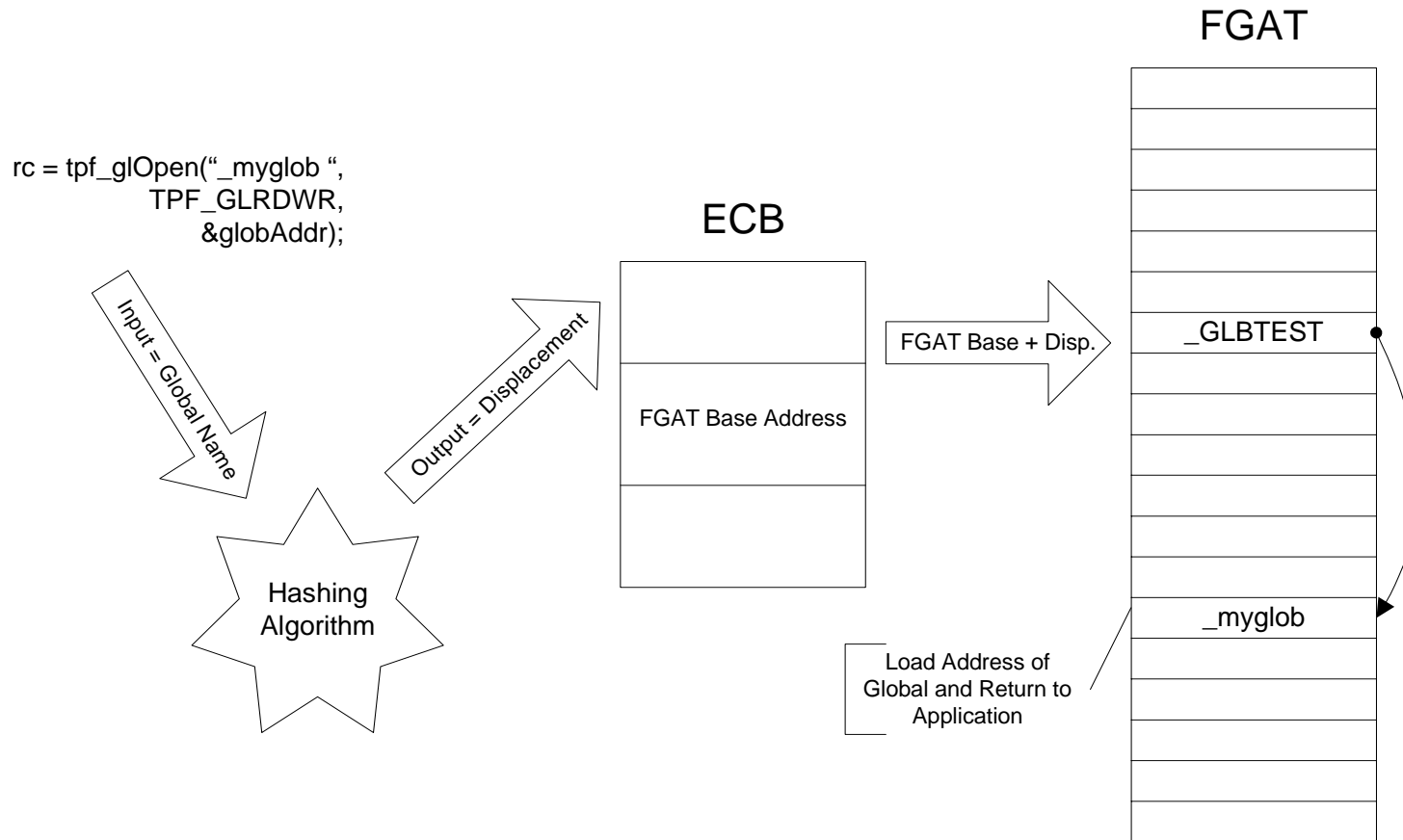
## New z/TPF APIs for Format-2 Globals

- Abstract APIs determine behavior based on the defined attributes of the global
- C/C++ language APIs modeled after standard file system APIs
- Perform various actions on globals
  - Open a global:  
tpf\_glOpen( ) and GLOBLC FUNC=OPEN
  - Request updates made to a global to be written to the database:  
tpf\_glWrite( ) and GLOBLC FUNC=WRITE
  - Change options on how a global was opened:  
tpf\_glCntl( ) and GLOBLC FUNC=CNTRL
  - Obtain information about an open global:  
tpf\_glStat( ) and GLOBLC FUNC=STAT
  - Close a global:  
tpf\_glClose( ) and GLOBLC FUNC=CLOSE

## Opening Format-2 Globals

- `tpf_glOpen(char *globalname, enum t_glopt options, void **globaladdr);`
  - Returns main storage address of requested global for the current SSU and IS
  - Specify options for opening a global
    - `TPF_GLRD`: open for read-only access
    - `TPF_GLRDWR`: open for read/write access (ie, lock global for exclusive use)
    - `TPF_GLRDFST`: open for read-only access using fast path
  - Creates a global descriptor (non fast path only)
    - Keeps track of all globals currently open for an ECB
    - Used as input for subsequent `tpf_gl`-type API calls
- Assembler equivalent: `GLOBLC FUNC=OPEN`
  - `OPT=RDONLY`
  - `OPT=RDWRITE`
  - `OPT=RDFAST`

# Opening Format-2 Globals Example



## Opening Format-2 Globals for Read-Only or Read/Write

- If the global is opened for read/write access, it will be locked for exclusive use by this ECB
  - Synchronizable: perform equivalent of SYNC LOCK operation to lock the global on DASD and refresh main storage copy
  - Non-synchronizable: perform CORHC on global's main storage address
- If the global is opened for read-only access, application can select one of two methods:
  - Fast Path
    - Best performance for accessing read-only data that will not be updated by the application
  - Non-Fast Path
    - Provides capability of efficiently changing a read-only access to read/write access
    - Provides diagnostic capabilities to keep track of all globals that have been accessed by this application

## Updating Format-2 Global Data

- `tpf_glWrite(int globaldesc, enum t_glopt options);`  
`tpf_glWrite(int globaldesc, enum t_glopt options, long offset, long length);`
  - File out updates made to a global that was previously opened for read/write access
  - Only allowed for keypointable or synchronizable globals
  - Global remains open for read/write access for this ECB
  - Specify options for writing a global
    - `TPF_GLALL`: update the entire global
    - `TPF_GLPART`: update only the specified part of the global (using "offset" and "length" parameters)
- Assembler equivalent: `GLOBLC FUNC=WRITE`
  - Optionally specify: `OFFSET=, LENGTH=`



## Changing the Open Options for Format-2 Globals

- `tpf_glCntl(int globaldesc, enum t_glopt options);`
  - Change open options for a previously opened global
    - Read/write access to read-only access
    - Read-only access to read/write access
  - Specify options for changing global access
    - `TPF_GLRDWR`: change access of global to read/write
    - `TPF_GLUNLK`: unlock a global previously opened for read/write access
    - `TPF_GLUNLKWT`: unlock a global previously opened for read/write access, and wait for any previous updates to be filed out
- Assembler equivalent: `GLOBLC FUNC=CNTL`
  - `OPT=RDWRITE`
  - `OPT=UNLOCK`
  - `OPT=UNLOCKWT`



## Obtaining Information About Format-2 Globals

- `tpf_gStat(int globaldesc, struct iglst *statinfo);`
  - Obtain information about a previously opened global
  - Get attributes of global from format-2 global attribute table (FGAT)
    - Global's main storage address and size
    - SSU-unique or SSU-shared
    - IS-unique or IS-shared
    - Processor-unique or processor-shared
    - Keypointable or non-keypointable
    - Synchronizable or non-synchronizable
    - Protected or not protected
    - Opened for read-only or read/write access
- Assembler equivalent: `GLOBLC FUNC=STAT`

## Closing Format-2 Globals

- `tpf_glClose(int globaldesc, enum t_glopt options);`  
`tpf_glClose(int globaldesc, enum t_glopt options, long offset, long length);`
  - Close a global and optionally write updates to the global to DASD
  - Specify options for closing a previously opened global
    - TPF\_GLUPD: file out updates made to the global (ie, keypoint or synchronize)
    - TPF\_GLUPDWT: file out updates made to the global and wait for completion (only for synchronizable)
    - TPF\_GLNOU: close global without updating it
    - TPF\_GLPART: only update the specified portion of the global (using "offset" and "length" parameters)
  - Remove format-2 global descriptor from ECB
- Assembler equivalent: `GLOBLC FUNC=CLOSE`
  - `OPT=UPDATE`
  - `OPT=UPDATEWT`
  - `OPT=NOUPDATE`

## Considerations for Updating Format-2 Globals

- Updating a keypointable format-2 global
  - Action taken depends on state of subsystem
    - 1052 state: file the updated global to the database before returning to the application
    - Above 1052 state: turn on keypoint request indicator. Time-initiated keypoint process will scan through all keypointable format-2 globals; those that have requested a keypoint will be written to the database
  - Size limit on keypointable format-2 globals is determined by the user
- Updating a synchronizable format-2 global
  - Synchronize tightly-coupled and loosely-coupled based on IS-uniqueness and processor-uniqueness attributes of global
  - Size limit on synchronizable format-2 globals is determined by the user
  - Ability to update a portion of the global
  - Ability to have issuing ECB wait for all other processors in loosely coupled complex to acknowledge completion of synchronization

## Modifying Format-2 Globals is Easier

- Both unprotected *and* protected format-2 globals can be updated in main storage directly by both assembler and C/C++ applications
- No longer need to call intermediate function to perform the update on behalf of the application
- `tpf_glmod(enum t_glmod key);`
  - Sets the correct storage protection access for the ECB to be able to alter protected format-2 globals
  - Use `keyrc( )` to return the ECB to working storage protection when done modifying protected format-2 globals

## Migrating Format-1 Global Data to Format-2 Globals

- Applications can be migrated one at a time to use new format-2 globals APIs to access format-1 globals data
- ZGLBL GLOBAL DEFINE with CONTROL-FORMAT1
- Allows user to define a format-2 global record that points at existing format-1 global data (ie, *format-1 global in migration state*)
  - All format-2 globals APIs can be used on format-1 globals in migration state.
  - New UGLM user exit to help the migration process
- ZGLBL GLOBAL MIGRATE
  - Migrate data from format-1 to format-2
  - Issue after all applications have been updated to use format-2 globals APIs
  - Global will become a system-controlled format-2 global, and will no longer reference the format-1 global data

## Customizing Format-2 Globals Support

- When the requirements for a particular global record do not fit into the definition of a system-controlled global record, it can be defined as user-controlled
  - Large tables that have unique loading, keypointing, or synchronization requirements
  - Data that has unique storage requirements or may not even physically reside on the z/TPF system
- ZGLBL GLOBAL DEFINE with CONTROL-USER
- Easy customization for user-controlled format-2 globals with various user exits
- User-controlled or system-controlled is transparent to the application
  - Same format-2 globals APIs are used by the application



# Operators Can Customize Format-2 Global Data in Dumps

- Using format-2 global names as dump keywords
  - 2 ways to create a dump keyword that references a format-2 global
    - ZIDOT CREATE command
    - IDATG macro
  - Use the ZIDOT INCLUDE command to specify that this global should be included for specific system errors. The copy of the global for the failing SSU and IS is dumped
  - "IF2GLOPN": pre-defined dump keyword that will dump all format-2 globals currently opened by the failing ECB



## Programs Can Customize Format-2 Global Data in Dumps

- Specifying format-2 globals to dump at system error time
  - LISTC macro
    - Specify name of a format-2 global and maximum length to dump
    - Can be used for SERRC and SNAPC dumps
  - serrc\_list and snapc\_list structures updated for C/C++
    - SERRC\_F2GLOBAL and SNAPC\_F2GLOBAL indicators tell the z/TPF system that the "tag" field is a format-2 global name
  - "IF2GLOPN" can be used as the global name, as noted for dump keywords

## Tuning Format-2 Globals Usage and Characteristics

- Data collection and data reduction
  - Data is collected and reported for the usage of synchronizable format-2 globals
  - Rate of locks, unlocks, synchronizations, and synchronizations with wait option

FORMAT-2 GLOBALS SUMMARY  
79 OBSERVATIONS

BSS SUBSYSTEM

14 JUN

09:18:32

GLOBAL NAME ADDED/DELETED	LOCKS/SEC	UNLOCKS/SEC	SYNCS/SEC	SYNC WAIT S/SEC	
_gl obs	0.124	0.051	0.057	0.016	
_myGl ob	0.115	0.056	0.043	0.016	
_newGl ob	0.009	0.008	0.000	0.001	Y
_Del ete	0.000	0.000	0.000	0.000	Y
	-----	-----	-----	-----	
TOTAL	0.248	0.114	0.100	0.034	

# Applications Productivity Enhancements Design and Coding Enhancements

## Expanded ECB Private Area

- Private area above 16-MB and below 2-GB
- Size increased
  - Minimum size is 4 MB
  - Absolute maximum size is 16 MB
  - Usable maximum size is user settable from 4 MB to 16 MB
    - For all ECBs
    - For individual ECBs
      - EBMAXC, tpf\_ebmaxc() can override usable maximum
- Easier management of traditional TPF I/O.

## Expanded ECB Heap

- 31-bit ECB Heap - below 2-GB
  - Default heap area - malloc() uses this area
  - With z/TPF memory model , 31-bit ECB Heap size can be large
    - 100 MB ... 500 MB
- 64-bit ECB Heap - above 2-GB
  - New API - malloc64() uses this area
  - Same virtual size as 1-MB frame area
  - Intended as very large work area
    - Large ECB unique in core tables
    - Large sort area
    - ECB unique cache.

## Expanded System Heap

- Dynamically allocate memory available to all ECBs
- Infrastructure which other function is built upon
  - Format-2 Globals
  - Memory File System
- 31-bit System Heap - below 2 GB
  - Virtual area
  - Uses from top of ECB virtual area to 2 GB
- 64-bit System Heap - above 2 GB
  - Virtual area - same size as 1 MB frames
  - Preallocated - backed with real memory
    - Intended for known large tables.

## Assembler Program Enhancements

- Size no longer limited to 4 K
- BEGIN macro allows
  - Multiple base registers ... R1 - R8
  - No base register (baseless)
- Architecture enablements
  - Relative instructions
  - Long displacement instructions
- Macros updated to have baseless capabilities
- New standard linkage macros for E-Type programs
  - Based on CP linkage macros
  - CLINKC, SLINKC, RLINKC, ELINKC
  - Easier subroutine coding.



## Assembler Program Enhancements - Register Handling

- R8 saved / restored across all SVCs
  - Does not need to be the program base register
  - Enables baseless coding
- R10, R11, R12, R13 available for use in E-type programs
  - Not guaranteed across macro calls
- Extended Register Save
  - Will save R10, R11, R12, R13 over general macro calls
  - Can be enabled for all general macros in a program
  - Can be enabled for general macros in a specific section of code.

## C Environment Enhancements

- Every ECB is initialized with a C Environment
  - Stack is always allocated
- Dynamically allocate stack storage - ALLOC and alloca()
  - Efficient way to allocate temporary storage
  - Valid during scope of function
  - Accessible from both TPF Assembler and C/C++
- TPF Assembler programs can define space on stack frame upon entry
  - DSECT defined in program

```
BEGIN NAME=ABCD,DSECT=WORKAREA,USEREG=Rz
REG_SAVE DS 8FD
TOD_SAVE DS 2FD
APSTKC
```

## C Environment Enhancements (continued)

- Ability to call C functions from TPF Assembler
  - New API - CALLC
  - Write function once and call it from C and Assembler
    - New IBM functions using this capability
  - Enhance existing assembler applications using C functions
    - Don't need to rewrite application in C
    - Gain benefits of using higher level language.

# New File Systems

- Processor and Subsystem Unique
  
- Memory File System (MFS)
  - Based on system heap storage
  - Used for temporary files
  - Deleted on IPL or file system dismount
  - Operates at memory speeds
  
- Fixed and Pool File Systems (FFS and PFS)
  - Based on pool and fixed files
  - Used for persistent files
  - Dramatically different performance characteristics
    - More than 100x faster
    - Creating and deleting files
    - Appending to existing files

## New Device Support

- Can use standard APIs to access data independent of the device type
  
- Device drivers under the TPF File System (TFS)
  - General Data Sets
  - General Tapes
  - Virtual Readers
  - Sockets
  - User written driver
  
- Standard Streams (stdin, stdout, stderr)
  - Special character files
  - Eliminated initialization overhead

## New File Systems APIs

- All file system APIs are now supported including the following:
  - Full file and byte locks
  - File attributes
    - Record ID to assign to data records
    - File Service Level
    - User assigned

## File System Check Utility

- Scandisk-like function with fix capability for all file systems
  - TFS, MFS, FFS, PFS
- Ability to correct a file system without requiring a re-initialization of the file system or a system IPL
- Performs checking of specified file system while the file system is in use
- Invoked via the zfile fsck comand
  - Use can be automated



## Additional New Functionality - Storage-Protection Override

- Problem: When updating a protected global in TPF 4.1 the program can only write into the global area. It cannot write into the ECB or Stack
- Solution: Use hardware feature called storage-protection override
  - When on, write to storage in key 9 as well as the PSW key
  - Allows programs to:
    - Write into protected globals
    - And write into EVM work areas (ECB, ECB private area, ECB Heap, and Application Stack)
  - New parameters on GLMOD, KEYCC allow use of storage-protection override
  - New C API - `tpf_stpoc()`

## Additional New Functionality

- Do not need to load a RIAT to use a new record ID
  - Simply input ==> ZRD TM MOD RECID-xxxx,....
  - Reduces bureaucracy for new development
  
- Improved addressability to system Heap areas
  - Unique tokens can be assigned to system heap areas
    - Assigned when system heap is obtained
      - GSYSC or tpf\_gsysc()
    - Address of in use system heap can be obtained
      - New API - FSYSC or tpf\_fsysc()
    - System heap can be released using a unique token
      - RSYSC or tpf\_rsysc()
  - Simplifies inter-program and inter-ECB communication via system heap storage.
  - Eliminates need to allocate space to save system heap address

## Additional NEW APIs

- `mallinfo()` and `EHEAPC FUNC=INFO`
  - Gives total storage being used for 31-bit ECB heap
  - Gives maximum size 31-bit ECB heap buffer that can currently be allocated
  
- `tpf_eheap_tag()`, `tpf_eheap_locate()` and `EHEAPC FUNC=TAG|LOCATE`
  - Applications can assign names to ECB heap buffers
    - Alternative to passing ECB heap addresses between programs or routines.
    - Simplifies inter-program and inter-routine communication
  
- `malloc64()` and `MALOC HEAP=64`
  - Allocate storage from 64-bit ECB Heap
  - Similar APIs for `calloc64()` and `realloc64()`

## Additional New APIs (continued)

- SWISC IMMEDIATE and `swisc_immediate()`
  - Move to another I-stream and start work where the SWISC is coded
  
- SYNCC SYNC, WAIT=YES
  - On return global update is in core on all processors
  
- TAMCC
  - Test Addressing mode
  
- `tpf_tgetc()`, `tpf_tputc` and TGETC, TPUTC
  - Ability for tape to read into or write out of discontinuous areas of memory
  - Ability to use any storage area such as ECB heap

## Additional New APIs (continued)

- tpf\_STCKE()
  - Obtain 128-bit extended clock
  - Provides C programs access to STCKE instruction
  
- ERRNOC
  - Provides assembler programs access to C variable errno
  
- DEBUGC
  - Debug facility for assembler programs
  - Can be used in the CP
  - Example ... take a dump if a routine is executed in 31-bit addressing mode

# Availability Improvements

## TCP/IP Tables - They're Movin' On Up

- Large tables reside above 2-GB bar, including:
  - Socket block table
  - IP message table (IPMT)
- Experience has shown there are many factors in determining what size IPMT is needed
  - Can make IPMT very large and not a concern
  - Larger IPMT helps enable OSA polling changes
- Can dynamically increase the size of these tables via ZNKEY command
  - Change takes effect immediately
  - No longer requires reloading CTK2 and an IPL



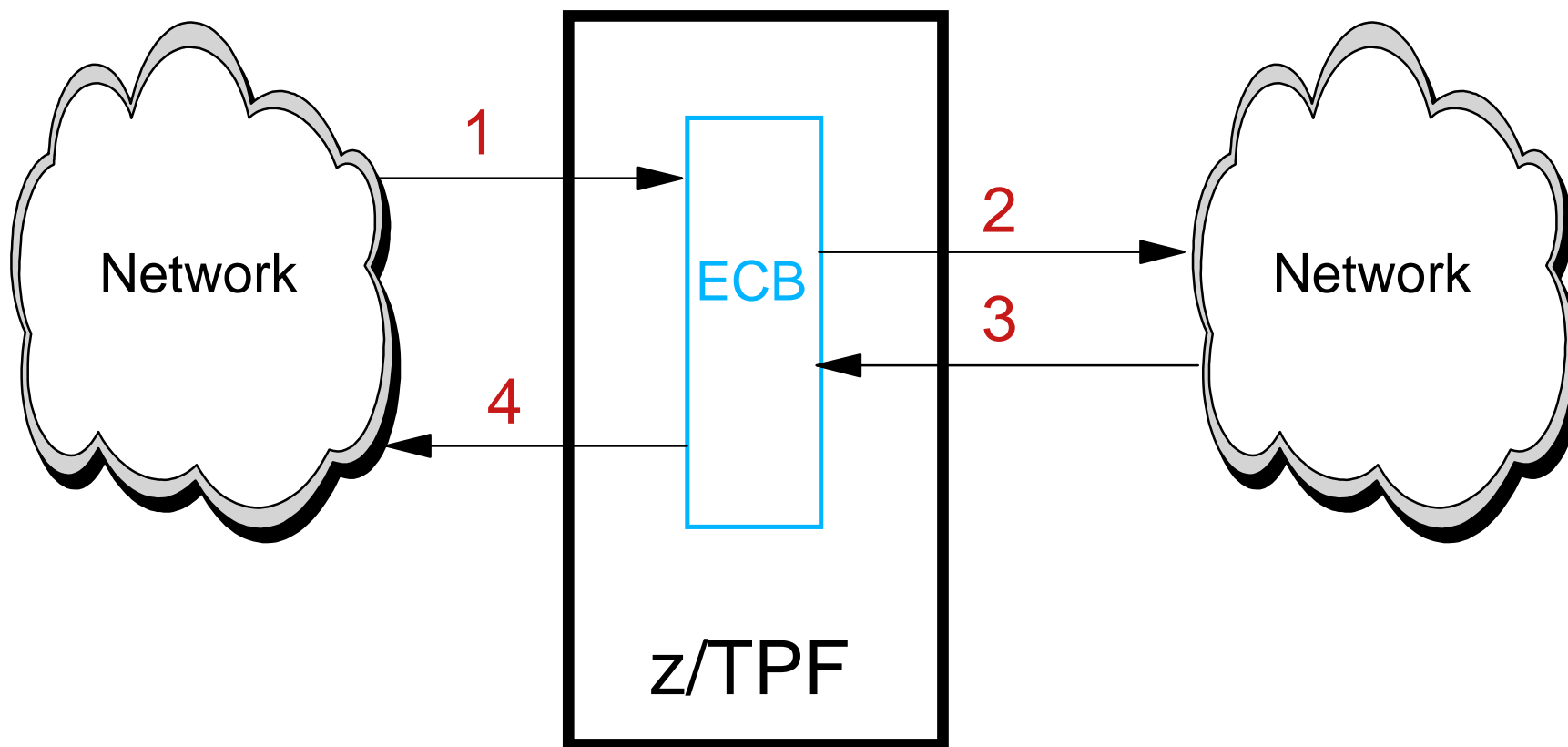
## TCP/IP Input Message Priority

- Allows you to define the priority of TCP/IP input messages
- Determines how a message is processed by z/TPF
  - Independent of network priority
- Priority value is defined at the application level
- Application program can:
  - Override the priority value for a given socket
  - Change the priority value during the life of a socket
- "High Priority" input messages
  - Processed via the ready list rather than the input list
  - These messages are read in from the network and processed even when in input list shutdown
- Other priority values exist
  - Determines what messages are discarded first if the system runs low on TCP/IP block resources

## OSA Polling Changes

- TPF polls (reads input messages from) OSA even when in an input list shutdown condition
  - High priority input messages continue to be processed
  - Other input messages are queued in memory (in the IP message table) and do not begin processing until system resources return to an acceptable level
- TPF polls OSA during dump processing
  - All messages are queued in memory and do not begin processing until after the dump is completed
- Reading input messages from the network during shutdown and dump processing reduces the likelihood of lost messages
  - Not polling would result in routers discarding messages that are destined for z/TPF
    - Lost messages would result in exception handling overhead following the shutdown or dump

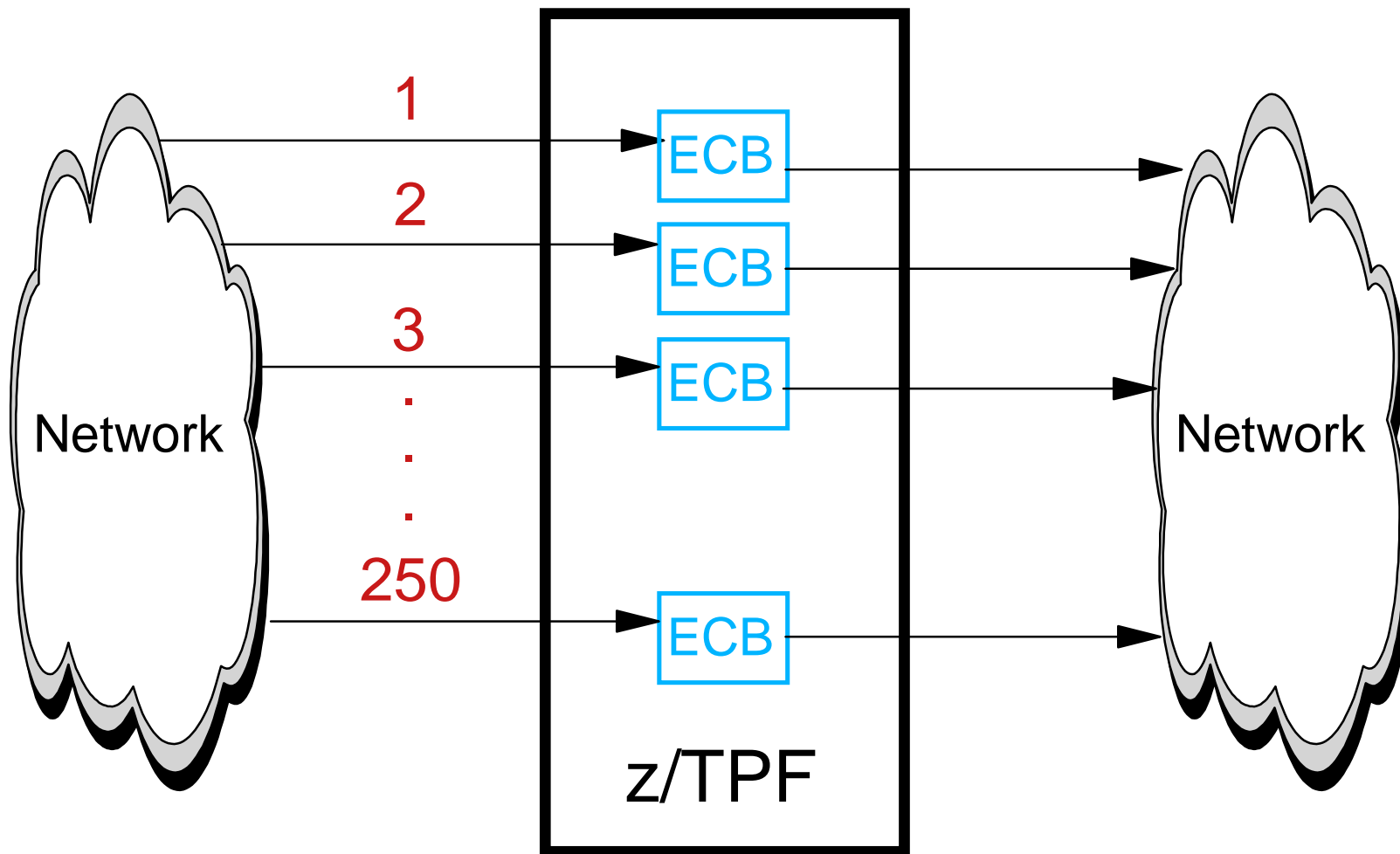
## Distributed Transaction Processing



## Distributed Transaction Details

1. z/TPF receives an input message from a remote client and an ECB is created to process the message.
  - How the message was received from the network (over TCP/IP, SNA, and so on) does not matter.
2. As part of the transaction, the z/TPF application sends a request to a remote server over TCP/IP. The application then issues a socket read() API causing the ECB to be suspended while waiting for the reply from the remote server.
3. The remote server sends the reply back over TCP/IP causing the application ECB to be dispatched again.
4. The application sends the output message to the remote client and the ECB exits.

## Example of When to Use High Priority Input Messages



## High Priority Input Message Example Details

1. Input message #1 is received, ECB #1 is created and sends a request to a remote server over TCP/IP. ECB #1 is then suspended waiting for the reply from the remote server.
  - The socket between z/TPF and the remote server is defined as high priority.
2. Step 1 is repeated 249 more times. There are now 250 active ECBs all waiting for replies from the remote server.
  - In this example, there are enough active ECBs to drive z/TPF into input list shutdown.
  - These ECBs will not go away until replies are received over the TCP/IP network from the remote server.
  - If the sockets between z/TPF and the remote server were not high priority, there would be a deadlock condition.
  - However, because these sockets are high priority, replies from the remote server are read in and processed, enabling the ECBs to exit. This in turn reduces the number of active ECBs and gets the system out of input list shutdown.



## New and Improved Scheduler

- TPF scheduler
  - Also known as the CPU loop or task dispatcher
- Now uses a combination of shared lists along with the existing I-stream unique lists
  - Replaces the TPF 4.1 routing weight based scheduler
- Better suited for z/TPF processors running a larger number of I-streams
- Enables more efficient use of compute resources within a processor, especially at very high utilizations
  - May result in better response time



## How to Take Advantage of the New Scheduler

- Program affinity attribute
  - Determines whether an ECB can dynamically switch I-streams while processing in a given program/application
    - For example, after a DASD I/O operation is completed
    - Default is not to dynamically switch
      - Behavior is consistent with TPF 4.1 to make migration easier
- Certain middleware and system services will be defined to take advantage of dynamic balancing, including:
  - MQSeries, mail server
  - TCP/IP native stack, file systems
- Application programs that are most frequently used are prime candidates to exploit the dynamic balancing capability
  - Assuming the application is tightly-coupled capable

## Dynamic Program Allocation

- New ADD option on ZAPAT command
  - Allocates a new program in the PAT
  - Replaces PATU option on ZOLDR LOAD
- Different ways to define the attributes (like KEY0, MONTC, RESTRICT) of the new program:
  - Specify the attributes on the ZAPAT ADD
  - Use the same attributes as an existing program
    - Use the LIKE option on ZAPAT ADD
  - Use the attributes that are assigned to unallocated programs
    - Do not specify any attributes or the LIKE option on the ZAPAT ADD command

# ZAPAT ADD Example

**ZAPAT ADD CLDW LIKE CLDB**

**APAT0200I 16.54.42 PROGRAM CLDW ADDED**

**BEGIN DISPLAY OF FILE COPY FOR IMAGE CUR51B**

<b>PROGRAM</b>	<b>CLDW</b>
<b>VERSION</b>	
<b>LINKAGE TYPE</b>	<b>---</b>
<b>BASE PAT SLOT</b>	<b>00000000</b>
<b>FILE ADDRESS</b>	<b>00000000D43014D5</b>
<b>CORE ADDRESS</b>	<b>0000000000000000</b>
<b>AUTHORIZATION</b>	<b>KEY0 MONTC RESTRICT</b>
<b>FETCH</b>	<b>DEFAULT</b>
<b>QUALIFIERS</b>	<b>NONE</b>
<b>TIMEOUT</b>	<b>50</b>
<b>DUMP GROUP</b>	<b>NONE</b>
<b>TRACE GROUP</b>	<b>IBM_DEFT</b>
<b>AFFINITY</b>	<b>PROGRAM</b>

## Enable More Utilities to Run in 1052 State

- More system services can be enabled in 1052 state:
  - File systems
  - Pools (GFS)
  - TCP/IP native stack, including Internet Daemon (INETD)
- Allows certain utilities to run in 1052 state, including:
  - FTP
    - Upload configuration files needed by applications before cycling to NORM and starting those applications
  - TPF Debugger
    - Can debug code in 1052 state now

## 1052 State Enabled System Services Details

- GFS does not automatically start until cycle up above 1052 state (same behavior as TPF 4.1)
  - Can manually start pools in 1052 state
    - New ZPOOL 1052 UP command
- An OSA-Express connection can be defined as automatically starting in 1052 state (default is still CRAS state)
  - Do not need additional OSA-Express adapters
    - Can define multiple logical connections to the same physical OSA-Express adapter
      - For example, one connection that starts in 1052 state and another connection that does not start until CRAS state
- Sockets can be started in 1052 state using new *SetTCP1052()* API
  - Default is still that sockets cannot be started until CRAS state
- INETD can be set up to automatically start in 1052 state and start certain applications

## Reduce Time Spent Dumping

- When a dump is in progress, all I-streams are stopped
  - No work can be done; the system is unavailable
- In TPF 4.1 data is written to tape during dump processing
- The amount of time that the system is unavailable is a function of:
  - Tape speed
  - Amount of storage being dumped.
  
- z/TPF addresses both tape speed and amount of storage dumped.



## Reduced Time Spent Dumping - Tape Speed

- In z/TPF data is written to a dump buffer area during the dump
  - Memory to memory move during dump processing
  - Write to tape is started after the dump has completed
  - Size of the dump buffer area is set by users
    - Make it as big as you need - 2 GB, 5 GB, 10 GB.
  
- If dump buffer area is full, dump processing will write to tape
  - RTL tape has been changed to increase speed
    - Must be blocked
    - Dump blocks are 128-KB
      - Increased from 32,760 bytes
      - 128-KB is optimal size of writing to tape devices.



## Reduce Time Spent Dumping - Reduce Amount of Data

- Ability to dump percentage of frames and SWBs
  - When dumping frames or SWBs, rather than dumping all blocks, dump a sample of each block
  - Specify from 0% to 100%
    - ZASER MAXBLKS.
  
- Ability to limit number of bytes dumped for an ECB heap buffer
  - If an ECB has a 100-MB ECB heap buffer, an OPR dump could take a long time
    - Is all 100 MB needed?
  - Specify number of bytes to dump for each ECB heap buffer
    - ZASER MAXEHEAP.

## Reduce Dump Impacts

- Prevent streaming dumps causing an outage
  - Ability to limit the number of CTL dumps
    - If more than a specified number of CTL dumps are taken in 1 minute, NODUMP remaining CTL dumps in that minute
      - ZASER MAXCTL.
  - Ability to limit the number of dump messages sent to prime cras
    - If more than a specified number of dump messages are sent to prime cras in 1 minute, do not send any more messages in that minute
      - Dumps will still happen
      - Dump message will be written to tape
      - ZASER MAXCPSE.

## Identify Run Away ECBs

- ECB Resource Manager
  - Based on ECB Resource Policeman
- Identify ECBs which:
  - Do too many FINDs, FILEs, SERRCs, or create ECB-type macros
  - Get too many SWBs, Common Blocks, pool file addresses, or system heap
- Actions taken can be:
  - Send a message
  - Take a SERRC with return
  - Take a SERRC and exit the ECB.

# ECB Resource Manager Example

```
ZECBM DISPLAY
CSMP0097I 18.34.59 CPU-B SS-BSS  SSU-HPN  IS-01
ECBM0010I 18.34.59 ECB RESOURCE LIMIT TABLE DISPLAY FOR CORE COPY

          1ST LIMIT  2ND LIMIT
RESOURCE WARNING  PERCENTAGE
FIND             5000           0
FILE             5000           0
SERR              10           0
SWBK             200           0
GRFS            300000         300
CMBK              50           0
SYSH              0            0
CRET             100           0
FILT= 1000 FILR= 1000
ecbRM MONITORING IS: OFF
1ST LIMIT ACTION IS: MESSAGE
2ND LIMIT ACTION IS: DUMPEXIT
END OF DISPLAY+
```

## Additional Enhancements

- Ability to include registers in macro trace without an IPL
  - Provides ability to get more diagnostics information quickly and easily.
  
- New ECB Heap Control Table (EHCT).
  - Each ECB has its own unique EHCT used to store information required in the management of the ECB heap.
  - Separates system control information from user data.
  - Applications overwriting ECB heap buffers will not overwrite system control information.

# New Tuning Capabilities

## BAL Packaging

- Assembler programs are linked
- Ability to package many assembler programs into one load module
  - For example, the following programs are link edited into CVAA
    - CVAA, CVAD, CVAH, CVAI, CVAO, CVAU
- Reduces linkage costs
- No need to update or reassemble programs
- Loading is done with ZOLDR
  - No need to IPL
- Further linkage cost reductions can be done by modifying ENTRC to programs that are in the same load module
  - ENTRC aaaa,TYPE=INT
  - Most efficient linkage



## Tuning Performance of the ECB Heap

- Most ECB heap requests are for small amounts of storage
  - New algorithm enables users to optimize allocation performance for these smaller buffer requests.
  - 4 fixed-size available lists are used to fulfill these smaller buffer requests.
    - Virtually eliminates the search time for small buffer requests.
    - Reduces amount of fragmented, or "unusable", ECB heap storage
- Larger requests are fulfilled from a 5th variable-sized available list.
- Users can customize
  - Size of each fixed-size buffer
  - Number of fixed-sized buffers allocated at ECB initialization.



## Sample ECB Heap Request Size Report from Data Collection

- ECB Heap Request Size data collection report provides information for setting sizes of buffers on the fixed-size available lists.

### TPF ECB HEAP REQUEST SIZE REPORT

CURRENT BUFFER SIZES IN BYTES:           AVL1 -       64   AVL2 -       256   AVL3 -       1024   AVL4 -       4096

CLASS UPPER LIMIT	FREQUENCY OBSERVED	PERCENT OF TOTAL	FREQUENCY DIAGRAM (SCALE = 1/1)
128	33	17.37%	*****
256	0	0.00%	
384	38	20.00%	*****
512	0	0.00%	
640	1	0.53%	*
768	30	15.79%	*****
896	0	0.00%	
1024	0	0.00%	
1152	0	0.00%	
1280	0	0.00%	
1408	0	0.00%	
1536	0	0.00%	
1664	0	0.00%	

## ECB Preallocated Storage

- Storage permanently assigned to an ECB
  - Private Area
  - ECB Heap
  - Application Stack
- Reduces overhead of updating EVM when added storage is needed
  - Also, reduces cleanup at exit
- Size can be set by the user
  - Intent is to set size of preallocated areas so that storage requests will be satisfied with preallocated for the majority of your ECBs
- Efficient use of storage based on application usage patterns
- Optimizes performance based on your unique applications.

# Additional Enhancements

## IP Scan

- Performs many TCP/IP stack functions on a time-initiated basis, including:
  - Sending TCP delayed acknowledgments
  - Retransmitting messages lost by the network
  - Sends output messages from partially full output buffers
- Now can run on any I-stream
  - Used to run on the main I-stream only
  - Enables a more balanced load on z/TPF
- Can be invoked more frequently
  - May improve response time

## Reduce IPL Time and Reduce VM Impact

- z/TPF does not test block (TB) every 4-KB in IPLB
  - TB 4K pages when they are needed
    - VFA buffers (if VFA is rebuilt)
    - Resources: SWBs, Common Blocks, 4-KB frames, 1-MB frames
  - Use of Available lists and Allocated lists for resources
    - All resources are on the allocated list initially
    - No resources on available list initially
    - Resource moved to available list on get when list is empty
      - TB done when resource moved to available list
    - All releases go to available list
    - All gets look at available list first
- Both combined minimizes VM working set size.



## Reduce IPL Time

- Enhancement to Format-1 Globals restart
  - Global restart runs in parallel with
    - TCP/IP restart
    - SNA restart
    - TPFDF restart
    - CCP restart

## Enhanced Memory Management

- Ability to define up to 8 memory configurations in CTKA
  - Memory configuration consists of:
    - Number of IOBs, CMBs, SWBs, ECBs, 4-KB and 1-MB frames
    - Size of dump buffer area, preallocated ECB areas, and more
    - Minimize size needed for 31-bit system heap and VFA
  - On IPL CCCTIN selects the configuration to use based on:
    - Defined preference
    - Best fit based on available memory
- Allows predefined memory configurations for
  - New faster processors with more memory
  - Various sizes of test systems
- Easier management of CTKA.

## Improved Timeout Processing

- Each load module has its own timeout value
  - Defined offline in control file, online in PAT
  - Ability to use the calling program's timeout value
    - C libraries such as CTAL uses this option
    - Ability to set a default timeout value for user programs
- Internal CP table which allows unlimited processing has been eliminated
- Timeout processing enabled in all states
  - Including restart and 1052.

## Constraint Relief

- Support up to 255 SSUs
- Support SDA addresses to x'FFFF'
  - Including tape
- Support up to 40,000 DASD devices

## Improved Coordination of Online Changes to Offline Source

- Ability to put online tables into offline source format
  - ZDECK command
    - RIAT
    - PAT
  
- Enhancements to Positive Feedback
  - Improved log management
  - Ability to select any command to be added to positive feedback.

## New Definable Options Requested by Customers

- Prime / Dupe module pairing
- DASD least queuing option
  - When DASD queue length is equal, always go to dupe rather than flip flop
- Set in SIP.

## New User Exits Requested by Customers

- Online DBR
- Midnight processing
- Copy member of user fields in ECB page 1
- Dump error message processing (CPSA)
- Dump data for OPR dumps
- Duplicate dump processing
- ZSTAT command



## Various Command Changes Requested by Customers

- New command to display CINFC information
  - ZDCNF / ZACNF
- Display state of all subsystems - ZDSYS ALL
- Display TOD clock sync information - ZPSMS DISPLAY TOD
- Ability to get macro name from SVC number - ZDSVC

## Various Enhancements Requested by Customers

- Eliminate assembly and compile warnings
- Improved lock release routine in dump processing
- Prevent looping catastrophic dumps
- Increase shutdown levels to 4 bytes
- Counts of DASD finds from DASD and VFA in the ECB
- Repeat tape mount request in restart
- Ability to request keypoint of in core user keypoint
- C function for FLVFC - tpf\_flvfc()
- Improve CXFRC Parent processing
- RLCHA program RLCH to use 4-KB block

## Enhancements to LODIC Support

- Additional 4 user classes supported
- Provide ability to specify a maximum amount of time that an ECB will lose control on initial LODIC call due to low resources
- When marked ECBs create child ECBs, allow the child ECBs to use a different class
- Supports shutdown for 1-MB frames
- Enhanced lose of control support
  - Save and restore new floating point registers.

## Larger Tape Blocking Support

- Support tape blocking up to 128-KB
- Any tape can use this
  - Command ZTLBL updates tape label mask
  - Command ZTMNT can override tape label mask
- 128-KB is most efficient size to write to tape devices
- RTL tape handling
  - Always write dump blocks at 128-KB
  - Blocking for other writes (TOURC) is user settable
- RTA tape handling
  - Always blocks at 32,760 bytes - dumps and user data
  - Most traces go to RTA.

## More VFA

- VFA gets remaining storage after all tables are allocated
  - Could be a large amount of storage - gigs and gigs
- New function added to support large number of delayed file records
  - Do not file delayed filed short term record if it has not been accessed in a specified period of time
    - Likely file address is not used anymore, a release file was not done for it
    - User settable by record ID using command
      - ZRTDM MOD RECID-xxxx,PDELAY-yyy

## Larger Keypoints

- Support 48-KB keypoints
- IBM keypoints moved to 48-KB are: CTKA, CTK2
- User keypoints moved to 48-KB are: CTKD, CTK3
- New APIs to access keypoints
  - GETKC / tpf\_getkc() and UPDKC / tpf\_updkc()

# Migration to z/TPF: Concepts and Tools



## Single Source Concept

- There are some changes which have to be made to application programs in order to migrate from TPF 4.1 to z/TPF
- Single Source support allows these changes to be made to applications while they are running in TPF 4.1
  - Don't need two copies of application source
  - Provides capability to reassemble / recompile TPF 4.1 applications for z/TPF.

## Architectural Concepts

- For assembler programs, 31-bit addressing mode is supported
- Storage areas which are shared among C and assembler programs are below 2-GB
  - ECB private area
  - Default ECB heap
  - Application stack
  - Format-1 Globals
  - Format-2 Globals are supported below 2-GB
- 31-bit pointer is available to C programs
  - Provides access to addresses stored in 4-byte fields

# Closing

## Comments

- Performance
  - Impact is estimated to be around 10% though individual results may vary.
  
- Migration to z/TPF from TPF 4.1
  - PUT 15 or higher installed on TPF 4.1
    - 32LC conversion complete (ZPMIG & ZMIGR done)
    - FARF6 conversion complete (ZMODE 6 done)
  - Coexistence supported for loosely coupled complexes

## Software not supported in z/TPF

- Target(TPF) C
- 24-bit addressing mode
- 24-bit Global support
- Fallback keypoint extents
- File resident programs
- PTV
- RTT
- VEQR mode
- Old comms: SLC and bisync
- CLAW offload support for TCP/IP
- SNA PU 2.1 Logon Manager

## Hardware not supported in z/TPF

- 3350, 3375, 3380, 3390, 9345, RAMAC DASD
- 3880, 3990-2, 3990-3, 3990-6 DASD CU
- 3480, 3495 tape
- 3705, 3725 CCU
- 3172, 3174 Terminal CU
- 3088 CTC
- 3505 card reader
- LLF, ELLF
- TOD RPQ
- STR subset
- 1403 printer

## Closing

➤ Questions???



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