May 22-26 • Colorado Convention Center • Denver, Colorado, USA

Platform: z/OS

New & Cool SQL: Version 8

William Favero

IT Software Specialist IBM North American Lab Services DB2 for z/OS IBM Software Group

Session: G10 Wednesday, May 25, 2005 /12:30 PM - 1:40 PM.



Where

Data Converge

Business &

What's new in V8?

- Multi-row FETCH and INSERT
- Scalar Fullselect
- Multiple DISTINCT Clauses
- Dynamic Scrollable Cursors
- INSERT within SELECT Statement
- GET DIAGNOSTICS Statement
- Current Package Path Special Register



Where

Data Converge

Business &



Where Business & Data Converge

Multi-Row FETCH and INSERT

- Benefits
 - Enhances usability and power of SQL
 - Facilitates Portability
 - Performance improved by eliminating multiple trips between application and DB engine; for distributed, reduced network traffic
 - Combined with scrollable cursors important for browse applications
- What is it?
 - Multi-row Fetch:
 - A single fetch statement can retrieve multiple rows of data from the result table of a query as a rowset
 - A rowset is a group of rows of data that are grouped together and operated on as a set
 - Multi-row Insert:
 - A single SQL statement can insert one or more rows into a table or view
 - Multi-row insert can be implemented as either static or dynamic SQL



Where Business & Data Converge

Multi-row FETCH topics

- DECLARE CURSOR Statement
- Host Variable Arrays
- FETCH Statement
- POSITIONED UPDATE & DELETE Statement



DECLARE CURSOR example

- Declare C1 as the cursor of a query to retrieve a rowset from the table DEPT.
- The prepared statement is MYCURSOR EXEC SQL

DECLARE CURSOR C1 CURSOR

WITH ROWSET POSITIONING

FOR MYCURSOR;

- Rowset positioning specifies whether multiple rows of data can
- be accessed as a rowset on a single FETCH statement -default is WITHOUT ROWSET POSITIONING



Host Variable Arrays

- Host variable array is an array in which each element of the array contains a value for the same column
- Changes to allow host variable arrays
 - COBOL
 - PL/1
 - C++
 - NOTE: Assembler support is limited to cases where USING DESCRIPTOR is allowed. Assembler pre-compiler does not recognize declaration of host variable arrays. Programmer responsible for allocating storage correctly, etc..
- Can only be referenced in multi-row fetch or insert
- In general, arrays may not be arrays of structures



Where Business & Data Converge

COBOL

Example 1: Declare CURSOR C1 and fetch 10 rows using multi-row FETCH

- 01 OUTPUT-VARS.
 - 05 NAME OCCURS 10 TIMES.
 - 49 NAME-LE PIC S9(4)COMP-4 SY C.
 - 49 NAME-DATA PIC X(40).
 - 05 SERIAL-NUMBER PIC S9(9)COMP-4 OCCURS 10 TIMES.

PROCEDURE DIVISION.

EXEC SQL

DECLARE C1 CURSOR WITH ROWSET POSITIONING FOR SELECT NAME, SERIAL# FROM CORPORATE.EMPLOYEE END-EXEC.

EXEC SQL

OPEN C1 END-EXEC.

EXEC SQL

FETCH FIRST ROWSET FROM C1 FOR 10 ROWS INTO :NAME, :SERIAL-NUMBER END-EXEC.



Where Business & Data Converge

FETCH example

EXAMPLE 1:

Fetch the previous rowset and have the cursor positioned on that rowset

EXEC SQL FETCH PRIOR ROWSET FROM C1 FOR 3 ROWS INTO... -- OR --EXEC SQL FETCH ROWSET STARTING AT RELATIVE -3 FROM C1 FOR 3 ROWS INTO...

EXAMPLE 2:

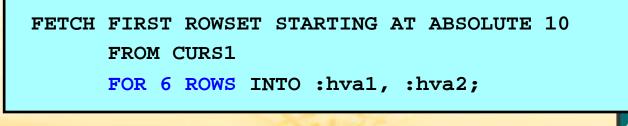
Fetch 3 rows starting with row 20 regardless of the current position of the cursor

EXEC SQL FETCH ROWSET STARTING AT ABSOLUTE 20 FROM C1 FOR 3 ROWS INTO...



ROWSETs

- A group of rows for the result table of a query which are returned by a single FETCH statement
- Program controls how many rows are returned (i.e., size of the rowset)
 - Can be specified on the FETCH statement (maximum rowset size is 32767)
- Each group of rows are operated on as a rowset
- Ability to intertwine single row and multiple row fetches for a multi-fetch cursor



Determining rowset size

- If FOR n ROWS is NOT specified and cursor is declared for rowset positioning..
- Size of rowset will be the same as the previous rowset fetch as long as
 - It was the previous fetch for this cursor
 - Or the previous fetch was a FETCH BEFORE or FETCH AFTER and the fetch before that was a rowset fetch
- Else rowset is 1

FETCH FIRST ROWSET FOR 5 ROWS	Returns 5 rows
FETCH NEXT ROWSET	Returns the next 5 rows
FETCH NEXT	Returns a single row
FETCH NEXT ROWSET	Return a single row
FETCH NEXT ROWSET FOR 3 ROWS	Returns 3 rows
FETCH BEFORE	Returns 0 rows
FETCH NEXT ROWSET	Returns 3 rows



Fetching beyond the result set

- If you try to fetch beyond the result set you will receive end of data condition
 - i.e., When there are only 5 rows left in result table and you request FETCH NEXT ROWSET FOR 10 ROWS, 5 rows will be returned with an SQLCODE +100
- This includes where FETCH FIRST n ROWS ONLY has been specified

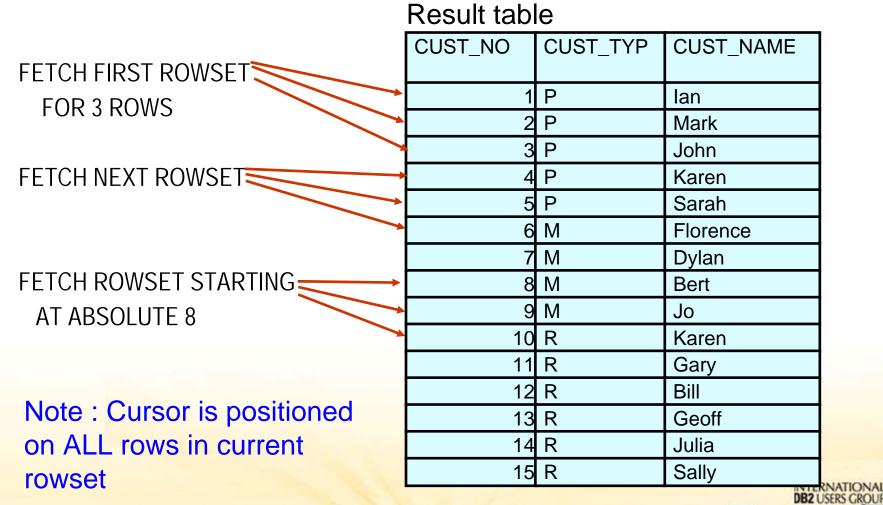


Cursor Positioning : Rowset positioned fetches

Where

Data Converge

Business &

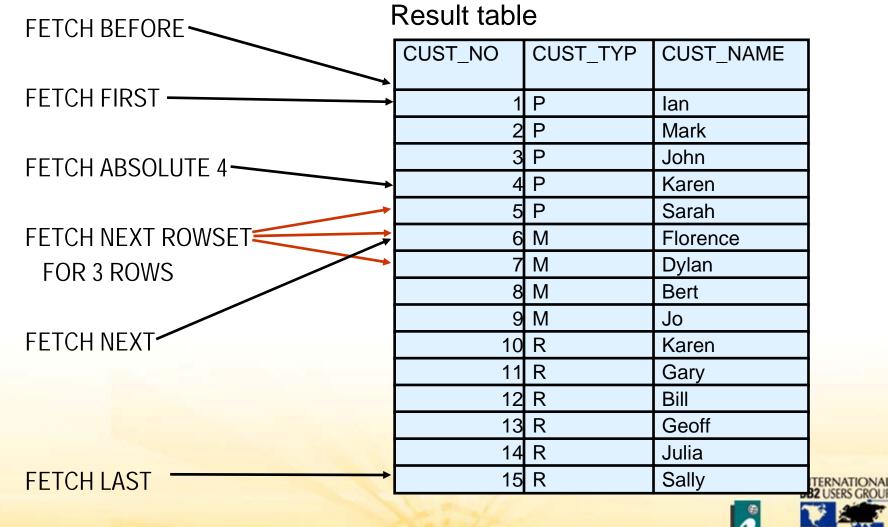




Slide 12 of 36

Where Business & Data Converge

Cursor Positioning : Row positioned fetches



2

Where Business & Data Converge

Positioned UPDATE of Multi-row FETCH

 Allows positioned UPDATE or DELETE to be used on a "wide" cursor

> UPDATE T1 SET COL1='ABC' FOR CURSOR C1 FOR ROW :hv OF ROWSET;



Multi-row insert

- New third form of insert
 - INSERT via VALUES is used to insert a single row into the table or view using values provided or referenced
 - INSERT via SELECT is used to insert one or more rows into table or view using values from other tables or views
 - INSERT via FOR "n" ROWS form is used to insert multiple rows into table or view using values provided in host variable array
- FOR "n" ROWS
 - For static, valid to specify FOR "n" ROWS on INSERT statement (for dynamic INSERT, specify FOR "n" ROWS on EXECUTE statement)
 - Maximum value of n is 32767
 - Input provided with host variable array -- each array represents cells for multiple rows of a single column
- VALUES clause allows specification of multiple rows of data
 - Host variable arrays used to provide values for a column on INSERT
 - Example: VALUES (:hva1, :hva2)



Where Business & Data Converge

Multi-row insert - ATOMIC vs. NOT ATOMIC

- ATOMIC (default) -- if the insert for any row fails, all changes made to database by any inserts are undone
- NOT ATOMIC -- inserts are processed independently
 If errors occur during execution of INSERT, processing continues
 - Diagnostics are available for each failed row through GET DIAGNOSTICS
 - SQLCODE will indicate if all failed, all were successful or at least one failed



GET DIAGNOSTICS statement

- Enables more diagnostic information to be returned than can be contained in SQLCA
- Returns SQL error information
 - for overall statement
 - for each condition (when multiple conditions occur)
- Supports SQL error message tokens greater than 70 bytes (SQLDA Limitation)

```
INSERT INTO T1 FOR 5 ROWS VALUES (:array);
GET DIAGNOSTICS :errcount = NUMBER;
DO || = 1 TO ERR_COUNT;
GET DIAGNOSTICS FOR CONDITION :||
:rc = RETURNED_SQLSTATE;
END;
Moved from end of presentation
```

Some examples for GET DIAGNOSTICS

- To determine how many rows were updated in an UPDATE statement
 - GET DIAGNOSTICS :rcount = ROW_COUNT;
- To handle multiple SQL errors during a NOT ATOMIC multi-row insert
 - GET DIAGNOSTICS :numerrors = NUMBER;
 - Then code a loop to execute the following for the number of errors
 - GET DIAGNOSTICS CONDITION :i :retstate = RETURNED_SQLSTATE
- To see all diagnostic information for an SQL statement
 - GET DIAGNOSTICS : diags = ALL STATEMENT
 - Sample output in :diags
 - Number=1; Returned_SQLSTATE=02000; DB2_RETURNED_SQLCODE=+100;
 - Would continue for all applicable items and for all conditions
 - Items are delimited by semicolons



Where Business & Data Converge

Scalar Fullselect

- Benefits:
 - Enhances usability and power of SQL
 - Facilitates Portability
 - Conforms with SQL Standards
- What is it?
 - Scalar fullselect is a fullselect, in parentheses, that returns single value
 - Scalar fullselect can be used in an expression

• Example: SELECT PRODUCT, PRICE FROM PRODUCTS WHERE PRICE <= 0.7 * (SELECT AVG(PRICE) FROM PRODUCTS);



Scalar Fullselect -- Restrictions

- Scalar Fullselect Not Supported in...
 - a CHECK constraint
 - a grouping expression
 - a view that has a WITH CHECK OPTION
 - a CREATE FUNCTION (SQL scalar)
 - a column function
 - ORDER BY clause
 - join-condition of the ON clause for INNER/OUTER JOINS



Where Business & Data Converge

Multiple DISTINCT clauses

- Benefits
 - Enhances usability and power of SQL
 - ► DB2 Family Compatibility
- What is it?
 - Allows more than one DISTINCT keyword on the SELECT or HAVING clause for a query
 - DB2 can now evaluate standard deviation & variance column functions



Where Business & Data Converge

Multiple DISTINCT clauses

Prior to Version 8

- ► SELECT DISTINCT C1, C2 FROM T1;
- SELECT COUNT(DISTINCT C1) FROM T1;
- SELECT C1, COUNT(DISTINCT C2) FROM T1 GROUP BY C1;

•With Version 8

- SELECT DISTINCT COUNT(DISTINCT C1), SUM(DISTINCTC2) FROM T1;
- SELECT COUNT(DISTINCT C1), AVG(DISTINCT C2) FROM T1 GROUP BY C1;
- SELECT SUM(DISTINCT C1), COUNT(DISTINCT C1), AVG(DISTINCT C2) FROM T1; GROUP BY C1 HAVING SUM(DISTINCT C1) = 1;

Not Supported in Version 8

- SELECT COUNT(DISTINCT A1,A2) FROM T1 GROUP BY A2;
- SELECT COUNT(DISTINCT(A1,A2)) FROM T1 GROUP BY A2;





Where Business & Data Converge

Dynamic Scrollable Cursors

- Benefits:
 - Enhances usability and power of SQL
 - Facilitates Portability
 - Conforms to SQL Standards
 - Performance improved by sort elimination
 - Elimination of work file (temporary table)
 - Scrolls directly on base table (TS scan, IX scan & DPSI)
 - No result table materialized at any time
 - Immediate visibility of committed updates, inserts, deletes
 - Positioned updates and deletes allowed when possible
 - Avoid sort by using backward index scan for ORDER BY

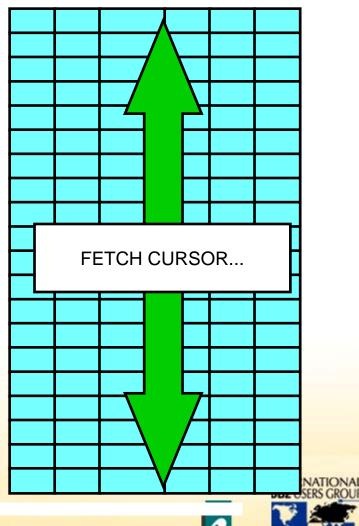




Where Business & Data Converge

Static Scrollable Cursors -- Version 7 Review

- •Cursors can be scrolled
 - -backwards
 - -forwards
 - -to an absolute position
 - -to a position relative to the current cursor
 - -before/after position
 - -Sensitive and Insensitive Cursors
- •Result table in TEMP DB
- Rows refreshed on demand



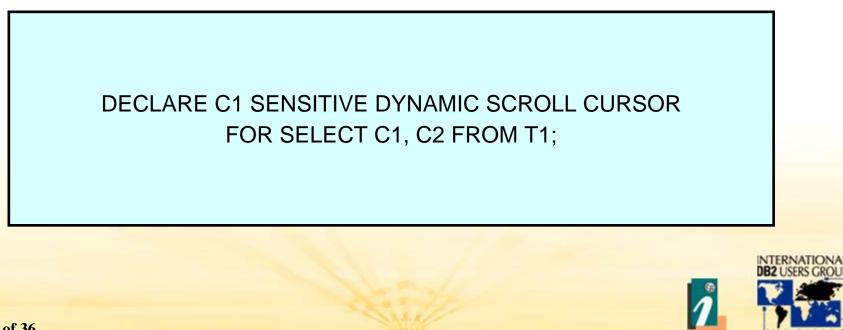


Where Business & Data Converge

Dynamic Scrollable Cursors –V8 new function

 Scrollable cursor that provides access to base table rather than workfile -- allows immediate visibility of updates, deletes, and *inserts*

Supported for index and tablespace scan access paths (type 2 indexes)



Where Business & Data Converge

New DECLARE CURSOR statement attributes

ASENSITIVE

- DB2 determines sensitivity of cursor
- If read-only...
 - -It behaves as an insensitive cursor
 - -Cursor is INSENSITIVE if SELECT statement does not allow it to be SENSITIVE (Union, Union All, For Fetch Only, For Read Only)
- ► If not read only, SENSITIVE DYNAMIC is used for maximum sensitivity
- Mainly for Client applications that do not care whether or not the server supports the sensitivity or scrollability

SENSITIVE DYNAMIC

- Specifies that size of result table is not fixed at OPEN cursor time
- Cursor has complete visibility to changes
 - -All committed inserts, updates, deletes by other application processes
 - -All positioned updates and deletes within cursor
 - -All inserts, updates, deletes by same application processes, but outside cursor
- FETCH executed against base table since no temporary result table created



Where Business & Data Converge

Implications on FETCH

- INSENSITIVE not allowed with FETCH statement
 - If the associated cursor is either declared as SENSITIVE DYNAMIC SCROLL or
 - If the cursor is declared ASENSITIVE and DB2 chooses the maximum allowable sensitivity of SENSITIVE DYNAMIC SCROLL
- There are no "holes" as there is no temporary result table
 - Special case: If FETCH CURRENT or FETCH RELATIVE + 0 requested but row on which cursor is positioned was deleted or updated so that it no longer meets selection criterion (returns +222)
- Inserts by the application itself are immediately visible inserts by others are visible after committed
- Order is always maintained
 - If current row deleted, the cursor is positioned before the next row of the original location and there is no current row





Cursor position and scrolling

- At OPEN CURSOR, cursor is positioned before first row
- After FETCH, fetched row becomes current row and cursor positioned on current row
- When FETCH reaches end of file, cursor positioned after last row if scroll number positive and before first row if scroll number negative
- Ability to scroll backwards and forwards through result set which is ever changing
 - Note that scroll quantity counts new inserts and has no way of counting deleted rows
 - Usage example: airline reservation or credit card processing



Some considerations

- Dynamic scrollable cursors are supported with stored procedures
 - SP itself can update via dynamic scrollable cursor but program calling SP is restricted from updating using allocated cursor
- Scalar functions/arithmetic expressions in SELECT list are reevaluated every fetch
- Column functions (AVG, MIN, MAX, etc.) are calculated once at open cursor
 - Functions may not be meaningful because size of result set can change
- Use of non-deterministic function (built-in or UDF) in WHERE clause of select statement or statement name of scrollable cursor can cause misleading results
 - Result of function can vary from one FETCH to subsequent FETCH of same row
- Parallelism is not supported with dynamic scrollable cursors



Backward index scan enabled

- DB2 will now select an ascending index and use a backward scan to avoid the sort for the descending order
- DB2 will use the descending index to avoid the sort and scan the descending index backwards to provide the ascending order
- To be able to use an index for backward scan,
 - Index must be defined on the same columns as ORDER BY and
 - Ordering must be exactly opposite of what is requested in ORDER BY.
 - i.e., If index defined as DATE DESC, TIME ASC, can do:
 - Forward scan for ORDER BY DATE DESC, TIME ASC
 - Backward scan for ORDER BY DATE ASC, TIME DESC
 - But must sort for
 - ORDER BY DATE ASC, TIME ASC or ORDER BY DATE DESC, TIME DESC



INSERT within SELECT statement

- Benefits:
 - Enables user to immediately determine values inserted in tables by DB2 (identity, sequence, defaults, etc.)and before triggers
 - Cuts down on network cost in application programs
 - Cuts down on procedural logic in stored procedures
- What is it?
 - INSERT statement is now allowed in the FROM clause of a:
 - Select statement that is a subselect
 - SELECT INTO statement
 - Users can automatically retrieve column values created by DB2 INSERT in single SELECT statement
 - Identity columns, sequence values
 - User-defined defaults, expressions
 - Columns modified by BEFORE INSERT triggers
 - ROWIDs



Current Package Path Special Register

Benefits

- Reduce network traffic and improve CPU/elapsed time for application using DRDA from a z/OS requester
- Allows nested procedure, user-defined function to be implemented without concern for invoker's runtime environment and allows multiple collections to be specified
- Easier to switch to/from JDBC and SQLJ
- What is it?
 - Current Package Path Special Register:
 - □Used for package collection resolution
 - Means for application to specify a list of package collections to DB server (similar to PKLIST on BIND PLAN)
 - DB server (rather than application requester) can search through list and find first
 - package that exists with specified package name
 - Control for applications that do not run under a DB2 plan



Package Resolution today

- Given that multiple collections can be used for packages, how is package resolution managed today?...
 - CURRENT PACKAGESET special register
 - Set to single collection id to indicate any package to be invoked belongs to that collection
 - Application must issue SET CURRENT PACKAGESET before each package is invoked if collection for the package is different from previous package
 - BIND PLAN PKLIST
 - Ability to specify list of collection ids for packages for local OS/390 applications that use plans at execution time



On the Border

- Identity Column Enhancements
- Sequence Objects
- Common Table Expressions
- Materialized Query Tables
- UNICODE SQL, Multiple CCSIDs
- XML Publishing



Where

Data Converge

Business &

Where Business & Data Converge

Summary

- A wealth of SQL changes
- More functionality
- Improved performance



Where Business & Data Converge

New & Cool SQL: Version 8 *Session: G10*

William Favero

North American Lab Services DB2 for z/OS IBM Software Group wfavero@attglobal.net



Slide 36 of 36