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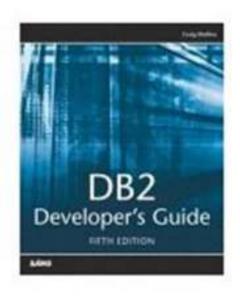
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The Tuning Progression

Problem Resolution



Application

- SQL
- Host Language Code

Database

- Indexing
- Database and Index Organization
- Database Design (normalization / denormalization)

DB2 Subsystem

ZPARMs, Pools, Locking, IRLM, DDF, etc.

Environment

- Network
- TP Monitor (CICS, IMS/TM)
- Operating System



Basic Tuning Rules

80% of the results of tuning come from 20% of the tuning effort -and-

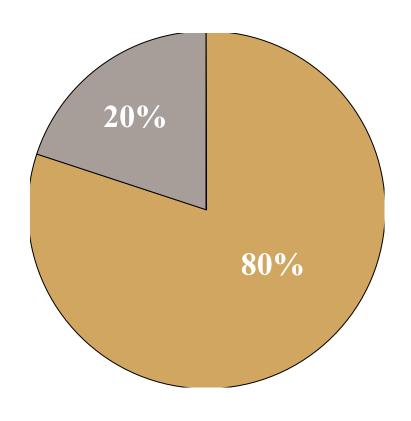
20% of your DB2 applications cause 80% of your problems

Tune one thing at a time

How else do you know whether the action helped or not?

All tuning optimizes:

CPU, I/O or concurrency

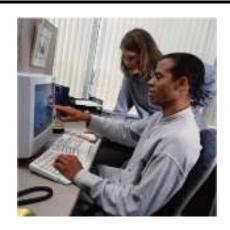


A Few General Performance Themes to Remember



Almost never say always or never.

There are rarely any "rules" that always apply.



Don't panic and remain humble.

Remaining calm and open to all solutions, even ones that recognize "you" as being the culprit, is important for building efficient DB2 databases and applications.

It is better to design for performance from the start.

The further into the development process you are, the more painful it becomes to makes changes.



The Cardinal Rule...

It depends!

Understand your circumstances and apply what makes sense.

NOTE

The cardinal rule of RDBMS development is "**It depends!**" Most DBAs and SQL experts resist giving a straight or simple answer to a general question because there is no simple and standard implementation that exists. Every situation is different, and every organization is unique in some way.

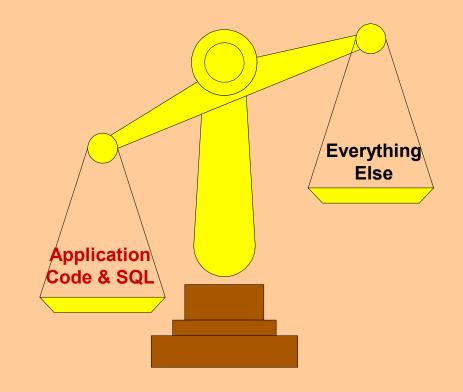
Don't be discouraged when you ask the local expert which statement will perform better, and the answer is "It depends." The expert is just doing his or her job. The secret to optimizing DB2 performance is being able to answer the follow-up question to "It depends"—and that is "What does it depend on?"

The key to effective SQL performance tuning is to document each SQL change along with the reason for the change. Follow up by monitoring the effectiveness of every change to your SQL statements before moving them into a production environment. Over time, trends will emerge that will help to clarify which types of SQL formulations perform best.

Application Code and SQL

Most relational tuning experts agree that the majority of performance problems with applications that access a relational database are caused by poorly coded programs or improperly coded SQL...

as high as 70% to 80%





Application Tuning: SQL

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Simpler is better, but complex SQL can be efficient In general, let SQL do the work, not the program Retrieve the absolute minimum # of rows required Retrieve only those columns required - never more

Always provide join predicates (i.e. no Cartesian products)

Favor Stage 1 and Indexable predicates

Host variable data type/length should match column

Avoid tablespace scans for large tables *(usually)*Avoid sorting when possible:

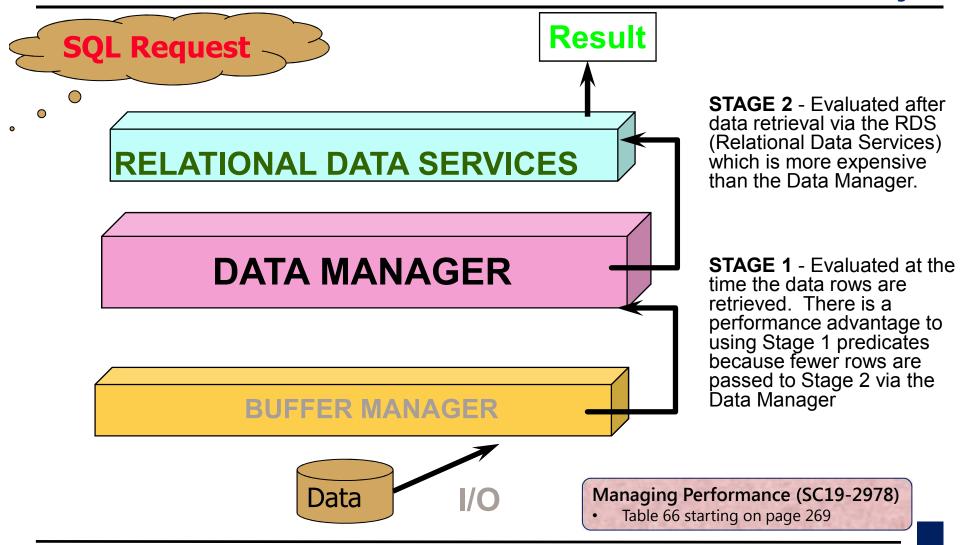
- indexes for ORDER BY and GROUP BY
- judicious use of DISTINCT
- UNION ALL versus UNION (if possible)





Application Tuning: Stage 1 and 2

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Not an actual Stage but

- It can be helpful to think of moving predicates from SQL into your programs as Stage 3
- Stage 1 better than Stage 2
- Stage 2 better than Stage 3



Ask Only for What You Absolutely Need



Retrieve the absolute minimum # of rows required

- Code appropriate WHERE clauses
- The only rows that should be returned to your program should be those that you need to process

Retrieve only those columns required: never more

- Don't ask for what you don't need
- Sometimes shortened to → Avoid SELECT *
 - This is a good idea for several reasons:
 - 1. Insulation of programs from change
 - Performance
 - But it is not enough...





What is Wrong with this SQL?

SELECT LAST_NAME, FIRST_NAME, JOB_CODE, DEPT **FROM EMP** WHERE **JOB_CODE** = 'A' **DEPT** = 'MIS'; AND Why are we asking for things we already know?



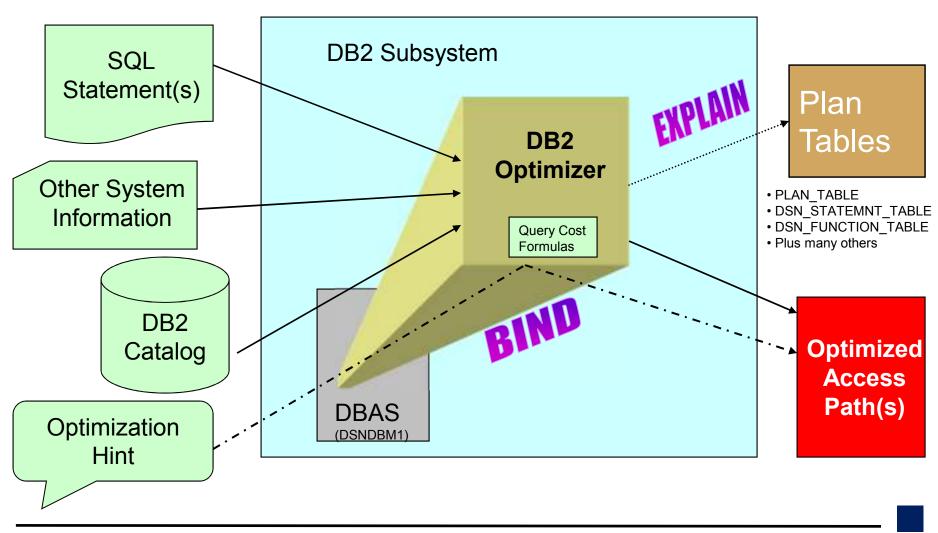
Avoid Black Boxes



Application Tuning: Optimization



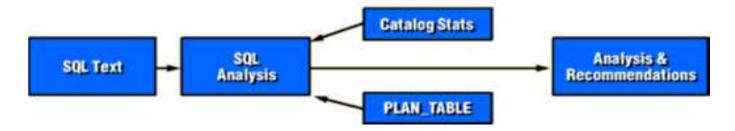
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101011 10101 101010 010110 1101110/1101 010 10110 100 001 0010' 1100 101 1110 1101

Application Tuning: EXPLAIN Analysis

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Hint used?

Index used?

Single, Multiple

Matching column(s)?

Index only?

TS scan (page range)

Type of Join?

- Nested Loop
- Merge Scan
- Hybrid

SQL Text

Table & Index Information

- DDL
- Stats

Cardinality

Other Stuff

- Triggers
- RI
- Constraints

Prefetch?

- Sequential
- List

Parallelism used?

- I/O, CPU, Sysplex
- Degree

Sort required?

Join, Unique, Group By, Order By

Locking



Application Tuning: Locking

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Minimize deadlocks by coding updates in the same sequence regardless of program

Issue data modification SQL statements as close to the end of the UOW as possible

the later in the UOW the update occurs, the shorter the duration of the lock



- ISOLATION(CS) / CURRENTDATA(NO)
- Can be used only by read only cursors

Use LOCK TABLE judiciously

Consider ISOLATION(UR) to avoid locking







Avoid Bachelor Programming Syndrome



Fear of COMMITing

Plan and implement a COMMIT strategy

or experience TIMEOUTs and DEADLOCKs





Application Tuning: Program

Do not embed efficient SQL in inefficient program logic

Classic example: finely tuned, efficient SQL inside of a program loop that executes 3,510,627 times!

Let SQL do the work when possible

e.g.) do not code "program" joins

Sign of trouble: SQL followed by lots of IF...ELSE or CASE statements

If you are only going to retrieve one row, consider coding a singleton SELECT (usually)

Consider adopting multi-row FETCH

Multiple tests have shown that moving to multi-row FETCH can yield between a 50% to 60% improvement for 100 and 1000 FETCHes



Application Tuning:

Online vs. Batch



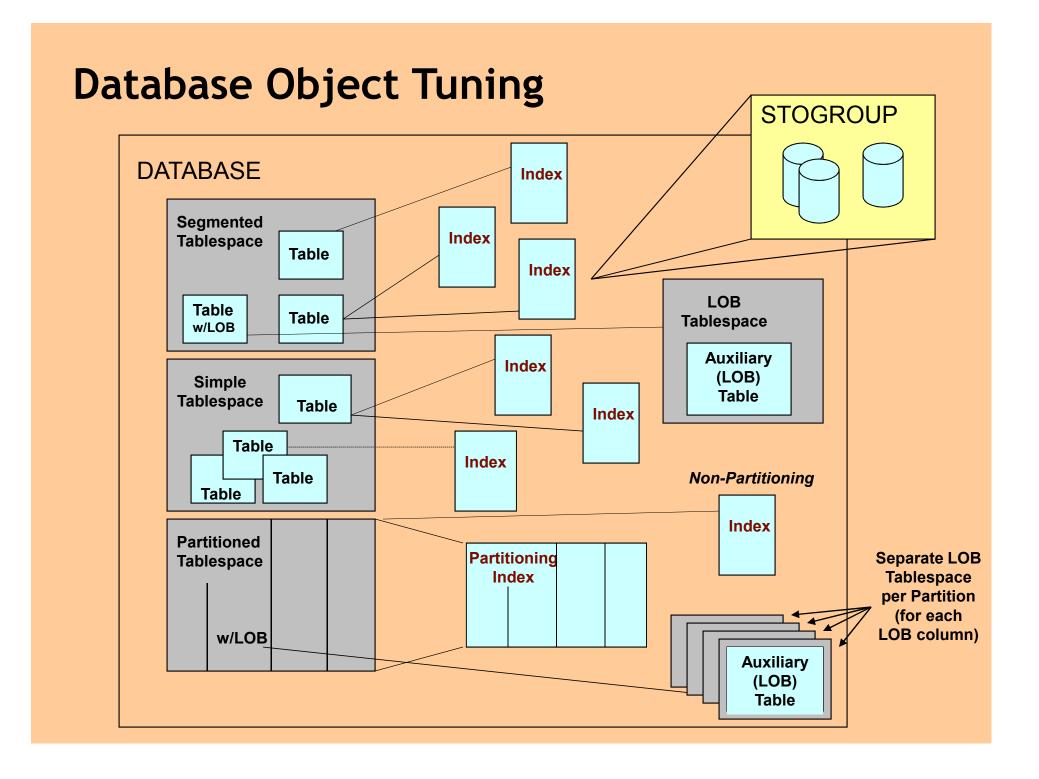
When designing online transactions, limit the amount of data to be retrieved to a reasonable amount

No one reads hundreds of pages/screens online!

Limit online sorting and joining (but be reasonable)

Consider OPTIMIZE FOR 1 ROW to disable list prefetch

- With list prefetch, DB2 acquires a list of RIDs from a matching index, sorts the RIDs, & accesses data by the RID list
- Can be very inefficient for a multiple page transaction







Combine the space management of segmented table spaces with the organization of partitioned table spaces.

Types of universal table space:

- Partition by range (or range partitioned)
- Partition by growth

Benefits of universal table space:

- Flexibility
- Growth
- Uses space map pages, like segmented
 - More information about free space space than just partitioned
 - Improved mass delete performance
 - Immediate reuse of segments after a table is dropped or mass deleted

Universal Table Spaces Partition By Growth (PBG)



Ideal when a table is expected to exceed 64 GB but there is no suitable partitioning key

- Begins as a single-partition table space
- Grows automatically: partitions are added as needed, as the data volume expands
- Can grow up to 128 TB
 - The maximum size is determined by the MAXPARTITIONS and DSSIZE values that you specified and the page size

Consider UTS PBG TS as replacement for simple TS



Universal Table Spaces Partition By Range (PBR)





Partition by range, or range-partitioned, universal table spaces are created by specifying both SEGSIZE and NUMPARTS on the CREATE TABLESPACE statement.

All actions that are allowed on exclusively partitioned or exclusively segmented table spaces are allowed on range-partitioned universal table spaces.

Ranges for range-partitioned universal table space can be specified on subsequent CREATE TABLE (or CREATE INDEX statements).

General Table Space Recommendations



As of DB2 V9, favor universal table spaces over segmented or traditional partitioned table spaces

- UTS are the future of DB2 table spaces
- At some point, other table space types are likely to be deprecated (like simple already have been)

In most cases limit yourself to one table per table space

You can still use a segmented table space when you must have multi-table TS

DSSIZE < 4GB unless you definitely need large TS



Database Organization

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Be sure to run RUNSTATS

- as data volume changes, new data structures added
- followed by (RE)BIND with /EXPLAIN(YES)

Review statistics (or RTS) to determine when to REORG

- NEARINDREF and FARINDREF
- LEAFDIST, PERCDROP
- For clustering indexes
 - ♦ NEAROFFPOSF and FAROFFPOSF
 - ◆ CLUSTERRATIOF
- Migrate to Real Time Statistics!
- Analyze access patterns before reorganizing
 - ◆ Random vs. sequential
 - Consider automation





Table Design Basics

As normalized as possible, but performance before aesthetics; normalization optimizes "update" at the expense of "retrieval"

Don't let data modelers dictate "physical" design

Do not create base table views

Avoid the defaults - they are usually wrong

Determine amount of free space

- PCTFREE amount of each page to remain free during REORG
- FREEPAGE after this many pages of data, keep an empty page
- Based on volatility
- Don't just let everything default (for example, to 10).



Database Design: Rows & Columns

Avoid wasted space (page size?)

- Row length > 4056 requires larger page size
- Row length 2029 4056 = one row per page
 - Ex) 2500 bytes: page size?
- Row length < 15 wastes space (max 255 rows/page)</p>

Sequence columns based on logging

- Infrequently updated non-variable columns first
- Static (infrequently updated) variable columns
- Frequently updated columns last
- Frequently modified together, place next to each other

New DB2 9 Format

Reordered Row Format



Database Design: Data Types

Use NULL sparingly

Use appropriate DB2 data types

- Use DATE instead of CHAR or numeric for dates
- Store numeric data using a numeric data type
 - INTEGER, SMALLINT, DECIMAL, etc.
- INTEGER versus DECIMAL(x,0)
 - Control over domain vs. storage requirements
- "DATE and TIME" versus TIMESTAMP
 - Ease of use/storage vs. precision/arithmetic

Compression versus VARCHAR

- Compression = less overhead (no 2 byte prefix)
- Compression requires no programmatic handling

Be Aware Of New DB2 9 Data Types

BIGINT DECFLOAT VARBINARY XML



Database Design: Integrity

Use DB2 declarative RI instead of program RI (usually)

- performance and ease of use
- ensure integrity for planned and ad hoc database modification

Do not use RI for lookup tables (overkill)

consider CHECK constraints vs. lookup tables



Use triggers only when declarative RI is not workable

- Triggers are less efficient (usually) than RI
 - but usually better than enforcing in application programs

Specify indexes on foreign keys



Database Design: Indexes

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A proper indexing strategy can be the #1 factor to ensure optimal performance

First take care of unique & PK constraints

Then for foreign keys (usually)

Heavily used queries - predicates

Overloading of columns for IXO

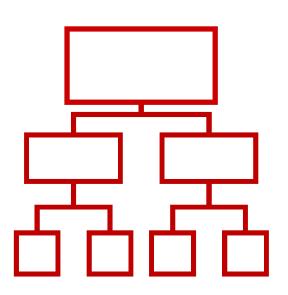
Index to avoid sorting

ORDER BY, GROUP BY

Consider INS / UPD / DEL implications

Consider how to index variable cols - [PADDED | NOT PADDED]

Index not used? SYSIBM.SYSINDEXSPACESTATS.LASTUSED



Database Design: Buffer Pools, Pt. 1



Bufferpool allocations - do not default everything to BPO

Explicitly specify a buffer pool for every table space and index

Ideas:

- isolate the catalog in BP0
- separate indexes from table spaces
- isolate heavily hit data
- isolate sort work area
- optimize BP strategy for your data & app processing mix: sequential vs. random
- there is no "silver bullet" approach
 - more on bufferpool tuning coming up!



REORG/RUNSTATS/REBIND

Straddling the line between application and database performance we have The Five R's!





BIND and REBIND

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The BIND and REBIND commands determine the access paths that the SQL in your programs will use to get to the data

■ **BIND** - accepts a DBRM (Data Base Request Module) and binds it into a package (or plan, until V10, at which point DBRMs will not be able to be bound directly to plans)



- SQL can change
- REBIND take a package (or plan) that has already been bound and re-assesses the access paths based on the latest statistics



SQL cannot change



Why Rebind?

Data volumes have changed

- More data
- Less data

Data skew has changed

Data distribution has changed

DB2 has changed

- PTF
- New version

Environment has changed





Scheduling Rebinds...

REBIND is critical for application performance It is a wise course of action to plan your REBIND strategy There are several common approaches:

- Regular maintenance: REBIND after RUNSTATS
 - Perhaps not every day, but REBIND are done after RUNSTATS
- Global REBIND after migration to new DB2 version
- Global REBIND after installing new PTFs
 - Above two mean access paths only change when DB2 changes
- REBIND after x days / weeks / months ...
- Let it Ride! ("If it ain't broke, don't fix it.")



Let It Ride

Programs once bound, are (almost) never rebound.

- Reason:
 - Fear of access path degradation
- Result:
 - No improvement to access paths
 - No CPU savings from new DB2 efficiencies
 - Sub-optimal performance
 - Every DB2 program potentially suffers for fear that one or two SQL statements will become inefficient





Regular REBIND

Better Approach: Regular REBINDing

- The Three R's (next slide)
- Why is this better?
 - Access paths are more up-to-date based on the current state of the data.
- Result:
 - Generally, improved access paths
 - CPU savings from new DB2 efficiencies
 - Optimal performance
- Of course, you can still get those "problem" access paths.



Monthlylip

The Three R's

- REORG
- RUNSTATS
- REBIND

Good, old advice but...

Still a couple of R's short.

The 3 R's: Rules for Running RUNSTATS



How accurate is the RUNSTATS utility? Does RUNSTATS use estimates derived from data sampling or does it actually access each row to collect and accumulate full measurement statistics? Also, what are some "rules of thumb" to use for scheduling RUNSTATS?



Statistics are collected by the RUNSTATS utility using both of the methods that you describe. When RUNSTATS INDEX is executed, exact statistics are collected. When RUNSTATS TABLESPACE is executed, the statistics for COLCARD are estimated using a technique called collective sample counting. However, the estimates are very accurate and reliable.

Some "rules of thumb" moveming the execution of RUNSTATS follow:

- Cansider running RUNSTATS whenever 10% or more of the data in a table has been modified. This includes INSURTs, UPDATEs, DELETEs, and LOADs.
- Collect column statistics only for those columns used in SQL predicates. The
 collection of column statistics can be sets expensive and should be performed
 only when it can impact access paths.
- Keep a lisitory of each application's statistics. After running RUNSTATS, select the statistics from the DB2 Catalog and insent them into a table or tables with a timestating or each row. These tables can be analyzed to show data growth mends.
- Punlose statistics reports using either the REPORT YES option of RUNSTATS or an SQL query against the DB2 Catalog. The SQL query will produce a more modelsfor report, but the REPORT YES option is easier to implement.
- Du not blindly REBIND every puckage and plan after exacuting RUNSTATS.
 REBIND only if the state changes significantly or if performance is suffering.
- Optimally, statistics should reflect the status of the data during the period of highest data access, if providing subsedute RUNSTATS to achieve this.
- Analyze RUNSTATS data to dotermine when REORG is necessary. Always run RUNSTATS after a REORG.

Originally published February 1993 for DB2#V2R3







Problems With the Three R's

They pose a lot of questions...

- When should you REORGanize?
 - To properly determine requires statistics (available in RTS).
 - So perhaps it should be RTS, REORG, RUNSTATS, REBIND?
- When should you run RUNSTATS?
 - To properly determine you need to know the make-up, usage, and volatility of your data.
- When should you REBIND?
 - When statistics have changed significantly enough to change access paths.
 - Knowing when this happens can be tricky.

A Better Approach: The 3 5 R's



RTS (or RUNSTATS)

REORG

RUNSTATS

REBIND

Recheck

- In other words, what did the REBIND do?
 - Did any access paths change?
 - Are they better or worse?
 - Does anything need attention?

Access Path Degradation Correction



Problem access paths can still occur. If so:

- Absolutely determine you are not at fault by re-rechecking
 - Statistics up-to-date?
 - Correct statistics run?
 - Package rebound with latest statistics?
- If problems persist, one of the following approaches could work for you:
 - Plan Stability (V9)
 - Tweak SQL (ways to "coerce" the optimizer)
 - Code and Use an Access Path Hint
 - Manipulate statistics (caution)



System & DB2 Subsystem Tuning

Allied **Agent Address Spaces** • CICS

- IMS
- TSO

SPAS/WLM

Stored Procedure Address Space

SSAS

System Services Address Space

Logging Recovery Access to TP **DBAS**

Database Services Address Space

RDS - Optimizer Data Manager **Buffer Manager**

IRLM

Internal Resource Lock Manager

Locking

VSAM Media Manager **Data** **DDF**

Distributed Data Facility

Relational database systems "love" memory

Performance improves if important information and run-time details are cached in memory instead of being read from disk every time they are needed.

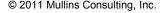




Swimming in DB2's "Pools"

DB2 uses four types of "pools" - or memory structures to cache data and information to avoid costly disk I/O

- Buffer Pools used to cache data in memory when it is read from disk.
- RID Pool used to sort RIDs (record identifiers) for List Prefetch, Multiple Index Access, and Hybrid Joins.
- EDM Pool used to cache program details (access paths, dynamic PREPARE, authorization) and database structural information (DBD).
- Sort Pool used when DB2 must sort data.





Subsystem: Buffer Pools, Pt. 2

DB2 provides up to 80 buffer pools - USE THEM!

- 4K: BP0 thru BP49 - 8K: BP8K0 thru BP8K9

- 16K: BP16K0 thru BP16K9 - 32K: BP32K thru BP32K9

Consider reserving a bufferpool for tuning

Move problem objects there to isolate for tuning

DB2 V8 significantly increased buffer pool storage

- 1TB new limit for buffer pools
- No more hiperpools; no more buffer pools in data spaces

Monitor hit ratio: % times a page is found in the buffer pool

The higher the ratio the better

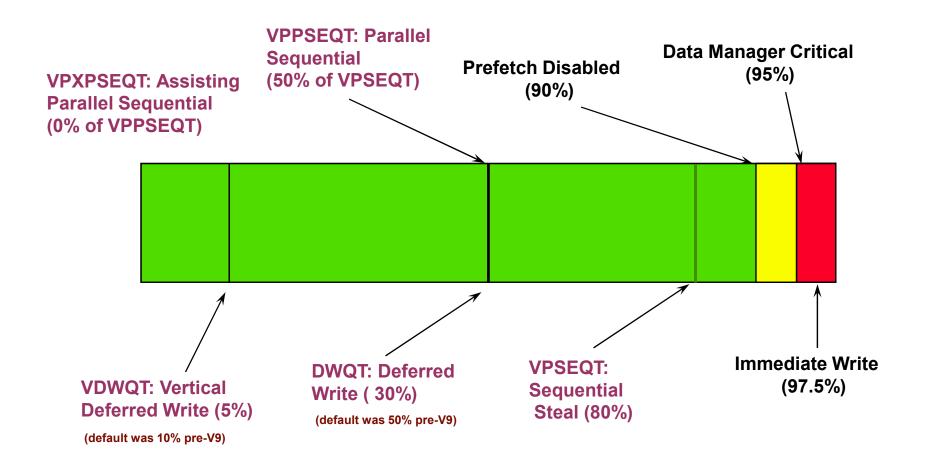
(GETPAGES – PAGES READ) / GETPAGES

SYNC I/O + ASYNC I/O



Buffer Pools: Tune Thresholds

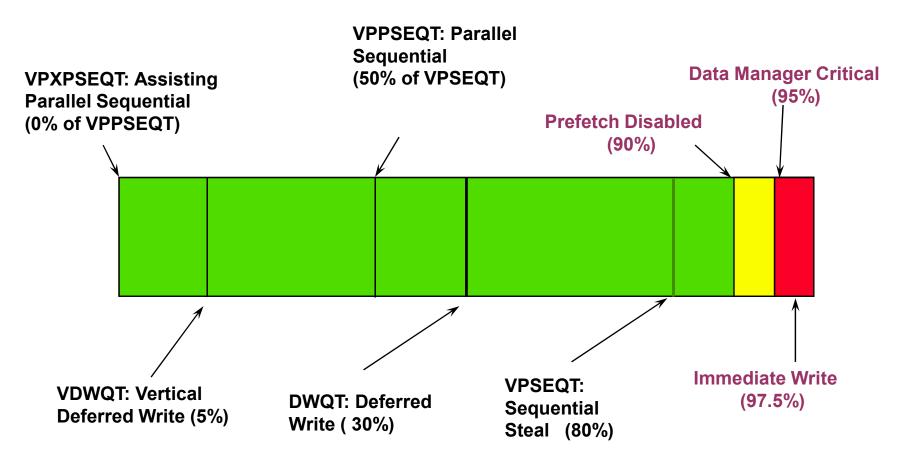
Variable Thresholds





Buffer Pools: Monitor Thresholds

Fixed Thresholds





Subsystem: RID Pool

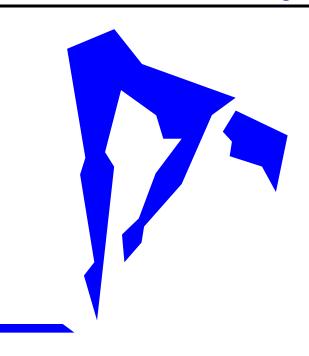
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One RID pool for all processing.

The default RID pool size is 4 MB. Can be changed.

RID Pool is used for:

- enforcing unique keys while updating multiple rows
- sorting RIDs during the following operations:
 - List prefetch, including single index list prefetch
 - Access via multiple indexes
 - Hybrid Joins







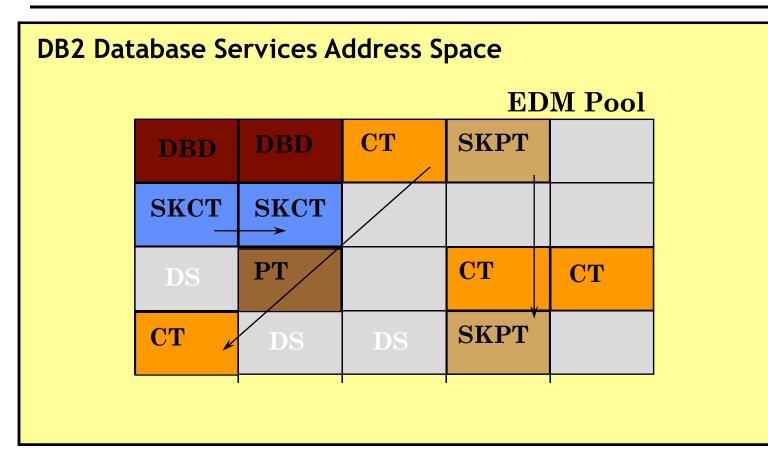
RID Pool Problems?

- 1. RID Pool Overflow (no storage)
 - Requests exceed ZPARM or DBM1 address space size
 - Limit: More than 16 million RID entries used or a single SQL statement consumes more than 50% of the RID Pool
 - The SQL causing this condition receives a -904.
 - Should not occur frequently.
- 2. DB2 anticipates RID pool access is not economical (RDS Limit)
 - Processing is suspended and access degrades to a table space scan
 - DB2 10 changes this behavior; write RIDs to disk instead of degrading to TS scan
 - You can disable the access paths listed above by specifying a RID pool size of 0.
 - Don't forget to REBIND to change access paths requiring RID pool



Subsystem: EDM Pool

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What's in
EDM Pool
DBDs
SKCTs
CTs
SKPTs
PTs
Auth Cache
Dyn SQL Prep
Free pages

V8 breaks each out into separate "pools"

Further break out in V9

General ROT: shoot for 80% efficiency; (1 in 5 DBD/SKPT/SKCT needs to be loaded)





The EDM Pool and V8, V9

V8: EDM Pool split into three specific pools:

- Below the 2GB Bar
 - EDMPOOL: EDM Pool stores only CTs, PTs, SKCTs, SKPTs
 - Should be able to reduce the size of this EDM pool
 - Provide some VSCR for below the 2GB Bar storage
- Above the 2GB Bar
 - EDMDBDC: DBDs
 - EDMSTMTC: Cached Dynamic Statements

V9: Introduces additional changes

- Above the 2GB Bar: EDM_SKELETON_POOL
 - All SKCTs and SKPTs
- A portion of the CT and PT is moved above the bar, too



Subsystem: Sort Pool

Sort Pool value is the maximum size of the sort work area allocated for each concurrent sort user.

The default Sort Pool size is 2 MB.

It can be changed on install panel DSNTIPC.

In general, estimate the required storage for a sort pool using the following formula:

```
32000 * (12 + sort key length + sort data length + 4))
```

Sorts that don't fit in SORTPOOL overflow to workfile

DSNDB07 for non-Data Sharing systems

Sort Performance

- In general, the larger the sort pool, the more efficient the sort is.
 - If the data fits into the sort pool, workfile database will not be required
- Allocate additional physical work files in excess of the defaults, and put those work files in their own buffer pool (e.g. BP7).
 - At least 5, sized the same, with no secondary
- The better sorted the data is originally, the more efficient the sort will be.
- Minimize the amount of data that needs to be sorted! →



Minimize Amount of Data to Sort

DB2 uses a tournament sort (next page) unless...

If Sorted Record > 4075, uses a <u>tag sort</u> (less efficient) because the data no longer fits on a 4K page:

- Data to sort put directly into 32K workfile database
 - For this reason be sure to always allocate at least one 32K workfile in DSNDB07
- Keys + RID are sorted
- Data retrieved from the sort using the RID





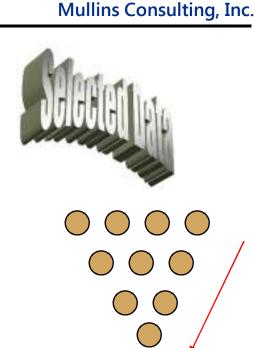
Tournament Sort

DB2 uses a Tournament Sort

Built into the hardware microcode and very efficient

How Does it Work?

- Input data to be sorted passes through a tree structure
- At each level in the tree the data is compared to data already there
- The 'winner' (lowest value for an ASC) moves up the tree
- At the top of the tree, the sorted entries are placed into <u>runs</u>
- Winning entries are removed from the tree and the next value inserted
- If there is more than one run, the runs must be merged







Subsystem: Logging

DB2 will only run as fast as the log Log Configuration

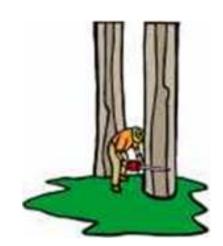
- Dual Active Logging is the preferred configuration
- Each log defined to separate devices and on separate channels

Output Buffer Size

- As BIG as possible please
- Waits occur if OUTBUFF is too small
- Max is 400000K

DB2 rollbacks from log data on DASD

Consider keeping archive logs on DASD*



 st and then migrate archive logs to tape after a specified period of time (HSM)





Subsystem: System Checkpoint

Periodically DB2 takes a checkpoint, containing:

- currently open unit of recoveries (UR) within DB2, all open page sets, a list of page sets with exception states, and a list of page sets updated by any currently open UR
- Dirty pages are written out at checkpoint and processing stops until they are written - so make sure DWQT is sized correctly!

Specified in the CHKFREQ* parameter in DSNZPARM

- Number of log records written
- Or, as of V7, number of minutes

Can be changed dynamically using:

- SET LOG or (temporary)
- SET SYSPARM (as of V7) (permanent)

5 minute intervals for checkpoints during peak processing times.

*CHKFREQ replaced LOGLOAD in DB2 V7





Subsystem Tuning: IRLM

MAXCSA

■ 250×(LOCKS PER USER)×(MAX USERS)

250 bytes of storage for each lock.

ITRACE=NO

- Do not use ITRACE;
- Instead, if needed, use DB2 lock traces.

DEADLOK

- 1. The number of seconds between two successive scans for a local deadlock
- 2. The number of local scans that occur before a scan for global deadlock starts



Environment

Operating System

version, memory, JCL, RACF, etc.

TP Monitors

CICS, IMS/TM, C/S GUI, web, etc.

Networking

TCP/IP, SNA, DRDA, stored procedures, etc.

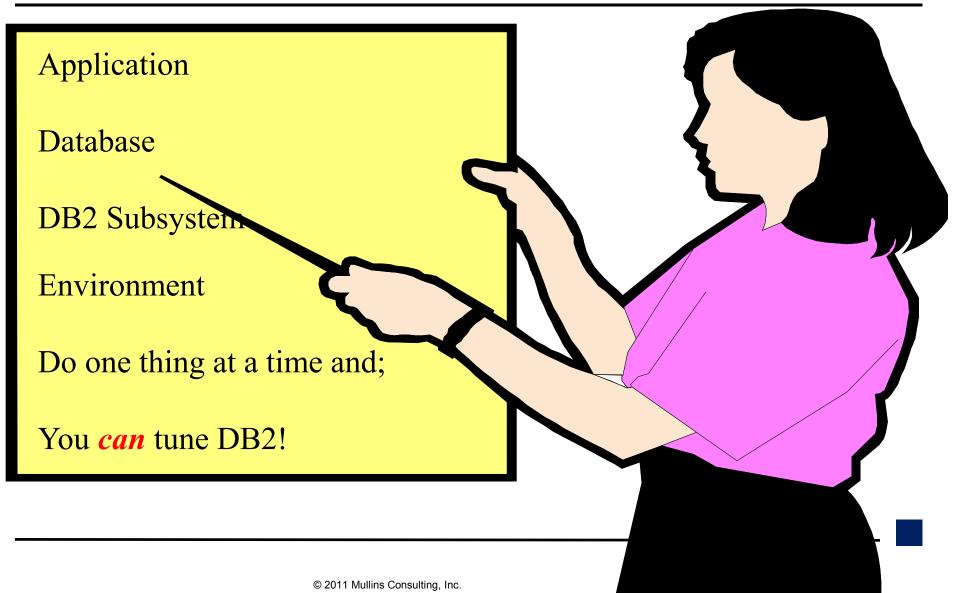
Hardware

storage, ESS/Shark, zIIP/zAAP, etc.



Summary

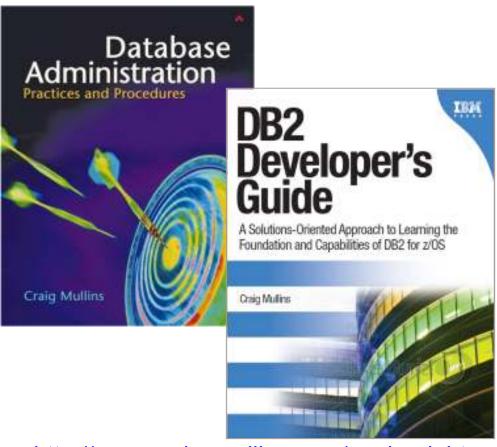
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