

# Cost-Complexity Pruning Process

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Assuming a CART or QUEST tree has been grown successfully using a learning sample, this document describes the automatic cost-complexity pruning process for both CART and QUEST trees. Materials in this document are based on *Classification and Regression Trees* by Breiman et al (1984). Calculations of the risk estimates used throughout this document are given in “Assignment and Risk Estimation” (TREE-assignment-risk.pdf).

## Cost-Complexity Risk of a Tree $T$

Given a tree  $T$  and a real number  $\alpha$ , the cost-complexity risk of  $T$  with respect to  $\alpha$  is

$$R_\alpha(T) = R(T) + \alpha |\tilde{T}|,$$

where  $|\tilde{T}|$  is the number of terminal nodes and  $R(T)$  is the resubstitution risk estimate of  $T$ .

## Smallest Optimally Pruned Subtree

**Pruned subtree:** For any tree  $T$ ,  $T'$  is a pruned subtree of  $T$  if  $T'$  is a tree with the same root node as  $T$  and all nodes of  $T'$  are also nodes of  $T$ . Denote  $T' \preceq T$  if  $T'$  is a pruned subtree of  $T$ .

**Optimally pruned subtree:** Given  $\alpha$ , a pruned subtree  $T'$  of  $T$  is called an optimally pruned subtree of  $T$  with respect to  $\alpha$  if  $R_\alpha(T') = \min_{T'' \preceq T} R_\alpha(T'')$ . The optimally pruned subtree may not be unique.

**Smallest optimally pruned subtree:** If  $T' \preceq T''$  for any optimally pruned subtree  $T'' \preceq T_0$  such that  $R_\alpha(T') = R_\alpha(T'')$ , then  $T'$  is the smallest optimally pruned subtree of  $T_0$  with respect to  $\alpha$ , and is denoted by  $T_0(\alpha)$ .

## Cost-Complexity Pruning Process

Suppose that a tree  $T_0$  was grown. The cost-complexity pruning process consists of two steps:

1. Based on the **learning sample**, find a sequence of pruned subtrees  $\{T_k\}_{k=0}^K$  of  $T_0$  such that  $T_0 \succ T_1 \succ T_2 \succ \dots \succ T_K$ , where  $T_k$  has only the root node of  $T_0$ .
2. Find an “honest” risk estimate  $\hat{R}(T_k)$  of each subtree. Select a right sized tree from the sequence of pruned subtrees.

## Generate a sequence of smallest optimally pruned subtrees

To generate a sequence of pruned subtrees in step 1, the cost-complexity pruning technique developed by Breiman et. al. (1984) is used. In generating the sequence of subtrees, only the learning sample is used. Given any real value  $\alpha_{\min}$  ( $\alpha_{\min} = 0$  in any SPSS implementation) and an initial tree  $T_0$ , there exists a sequence of real values  $-\infty < \alpha_1 = \alpha_{\min} < \alpha_2 < \dots < \alpha_K < +\infty$  and a sequence of pruned subtrees  $T_0 \succ T_1 \succ \dots \succ T_K$ , such that the smallest optimally pruned subtree of  $T_0$  for a given  $\alpha$  is

$$T_0(\alpha) = \begin{cases} T_0 & \alpha < \alpha_1 \\ T_0(\alpha_k) = T_k & \alpha_k \leq \alpha < \alpha_{k+1} \quad 1 \leq k < K, \\ T_0(\alpha_K) = T_K & \alpha_K \leq \alpha \end{cases}$$

where

$$\alpha_{k+1} = \min_{t \in T_k} g_k(t), \quad T_{k+1} = \{t \in T_k : g_k(s) > \alpha_{k+1} \text{ for all ancestors of } t\},$$

$$g_k(t) = \begin{cases} \frac{R(t) - R(T_{k,t})}{|\tilde{T}_{k,t}| - 1} & t \in T_k - \tilde{T}_k, \\ +\infty & t \in \tilde{T}_k \end{cases},$$

$\tilde{T}_{k,t}$  is the branch of  $T_k$  stemming from node  $t$ , and  $R(t)$  is the resubstitution risk estimate of node  $t$  based on the learning sample.

### Explicit algorithm

The algorithm can be used to generate a sequence of subtrees of  $T_0$  for a given initial value  $\alpha = \alpha_{\min}$ , and an initial tree  $T_0 = \{1, \dots, \#T_0\}$  where  $\#T_0$  is the number of nodes in  $T_0$ . For node  $t$ , let

$$lt(t) = \begin{cases} 0 & t \text{ is terminal} \\ \text{left child of } t & \text{otherwise} \end{cases}, \quad rt(t) = \begin{cases} 0 & t \text{ is terminal} \\ \text{right child of } t & \text{otherwise} \end{cases},$$

$$pa(t) = \begin{cases} 0 & t \text{ is root node} \\ \text{parent of } t & \text{otherwise} \end{cases}$$

$$\tilde{N}(t) = \begin{cases} 1 & t \text{ is terminal} \\ |\tilde{T}_{k,t}| & \text{otherwise} \end{cases}, \quad S(t) = \begin{cases} R(t) & t \text{ is terminal} \\ R(T_{k,t}) & \text{otherwise} \end{cases},$$

$$G(t) = \min_{s \in T_{k,t}} g_k(s).$$

The explicit algorithm is shown below.

1. Set  $k = 1$ ,  $\alpha = \alpha_{\min}$ .

For  $t = \#T_0$  to 1 {

if  $t$  is a terminal node, set

$$\tilde{N}(t) = 1, \quad S(t) = R(t), \quad g(t) = +\infty, \quad G(t) = +\infty,$$

else (i.e., if  $t$  is not a terminal node), set

$$\tilde{N}(t) = \tilde{N}(lt(t)) + \tilde{N}(rt(t))$$

$$S(t) = S(lt(t)) + S(rt(t))$$

$$g(t) = (R(t) - S(t)) / (\tilde{N}(t) - 1)$$

$$G(t) = \min\{g(t), G(lt(t)), G(rt(t))\}$$

}

2. If  $G(1) > \alpha$ ,

$$\alpha_k = \alpha \text{ and } T_k = \{t \in T_{k-1} : g(s) > \alpha_k \text{ for all ancestor } s \text{ of } t\}.$$

$$\alpha = G(1), \quad k = k + 1.$$

Else

if  $\tilde{N}(1) = 1$ , terminate this process.

3. Set  $t = 1$ .

$$\text{While } G(t) < g(t), \quad t = \begin{cases} lt(t) & G(t) = G(lt(t)) \\ rt(t) & \text{otherwise} \end{cases}.$$

4. Make current node  $t$  terminal by setting

$$\tilde{N}(t) = 1, \quad S(t) = R(t), \quad g(t) = +\infty, \quad G(t) = +\infty.$$

5. Update ancestor's information of current node  $t$ .

While  $t > 1$  (i.e.  $t$  is not the root node) {

$$t = pa(t)$$

$$\tilde{N}(t) = \tilde{N}(lt(t)) + \tilde{N}(rt(t))$$

$$S(t) = S(lt(t)) + S(rt(t))$$

$$g(t) = (R(t) - S(t)) / (\tilde{N}(t) - 1)$$

$$G(t) = \min\{g(t), G(lt(t)), G(rt(t))\}$$

}

6. Then repeat steps 2 to 6 until the termination condition  $\tilde{N}(1) = 1$  in Step 2 is satisfied.

## Selecting the Right Sized Subtree

To select the right sized pruned subtree from the sequence of pruned subtrees  $\{T_k\}_{k=0}^K$  of  $T_0$ , an “honest” method is used to estimate the risk  $\hat{R}(T_k)$  and its standard error  $se(\hat{R}(T_k))$  of each subtree  $T_k$ . Two methods can be used: the resubstitution estimation method and the test sample estimation method. Resubstitution estimation is used if there is no test sample. Test sample estimation is used if there is a testing sample. Select the subtree  $T_{k^*}$  as the right sized subtree of  $T_0$  based on one of the following rules.

### Simple rule

The right sized tree is selected as the  $k^* \in \{0, 1, 2, \dots, K\}$  such that

$$\hat{R}(T_{k^*}) = \min_k \hat{R}(T_k).$$

### The b-SE rule

For any nonnegative real value  $b$  (default  $b = 1$ ), the right sized tree is selected as the largest  $k^{**} \in \{0, 1, 2, \dots, K\}$  such that

$$\hat{R}(T_{k^{**}}) \leq \hat{R}(T_{k^*}) + b \cdot se(\hat{R}(T_{k^*})).$$

## References

Breiman, L., Friedman, J.H., Olshen, R., and Stone, C.J., 1984. *Classification and Regression Trees* Wadsworth & Brooks/Cole Advanced Books & Software, Pacific California.