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# Practical consolidation experience with Oracle and Linux on System z


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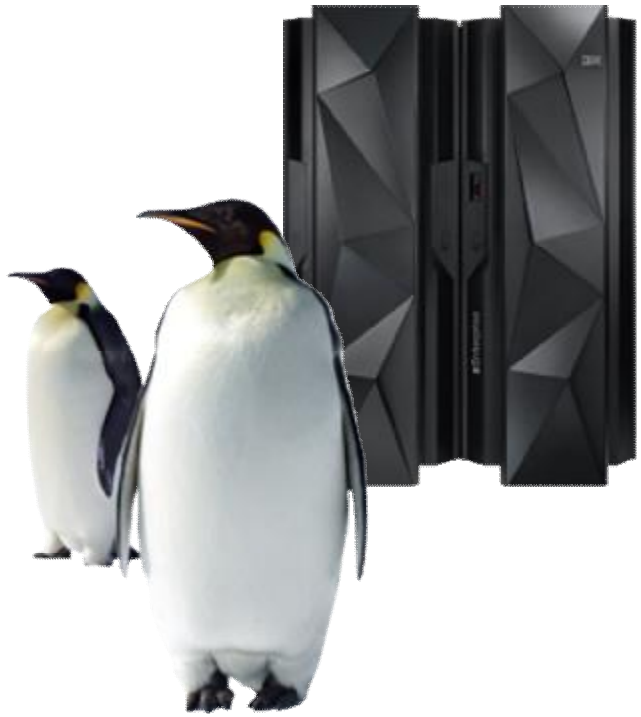
# Discussion Topics

## Consolidating Oracle database servers

### *A Real Customer Example*

- **Migration services**
- **Performance tuning results**
- **Best practices**



# A Real Customer Example

- Large Oracle database consolidation project
  - Oracle 10*g*R2 databases (including a few 11*g*R2 databases)
- Consolidation from x86 (HP ProLiant blade servers) to z196
  - 16 IFL
  - DS8800 with FICON attached ECKD
  - z/VM V6.1
  - RHEL 5.6
- Migration of individual databases over a longer time period
  - Utilizing IBM Migration Services (“Migration Factory”)

Problem statement:

- Customer reported application performance issues with 3 out of approx. 50 databases
  - Business analytics application ‘A’: not completing within expectation
  - Business analytics application ‘B’: not completing within expectation
  - Application ‘C’: increasing number of time-outs (transactions exceeding 1 minute)

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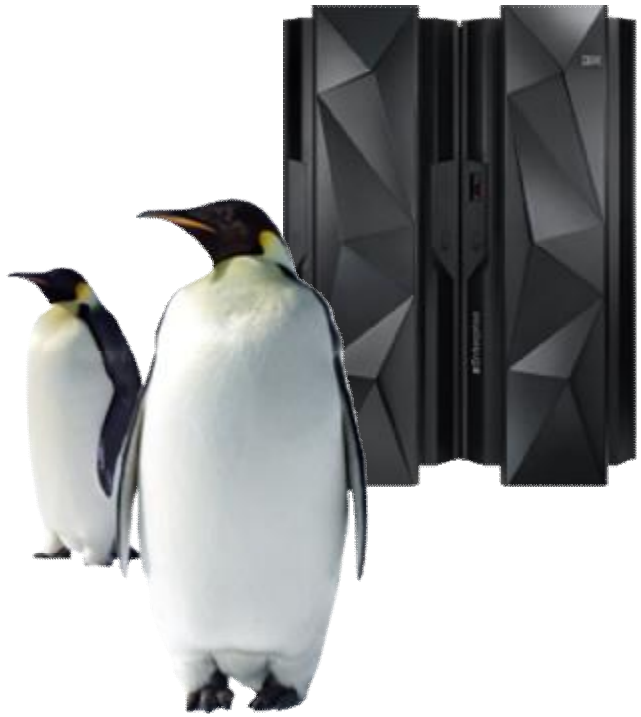
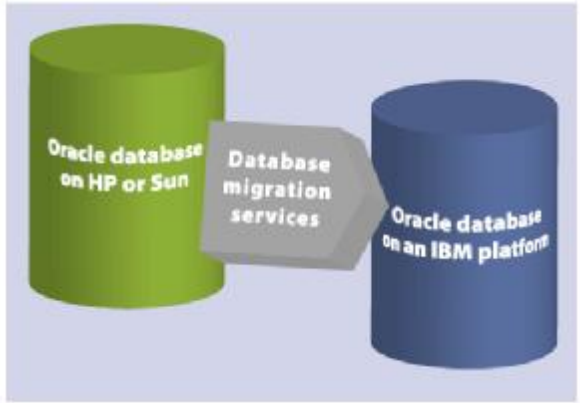


# Oracle Database Migration Services

## IBM Migration Factory (MF)

*How does it work?*

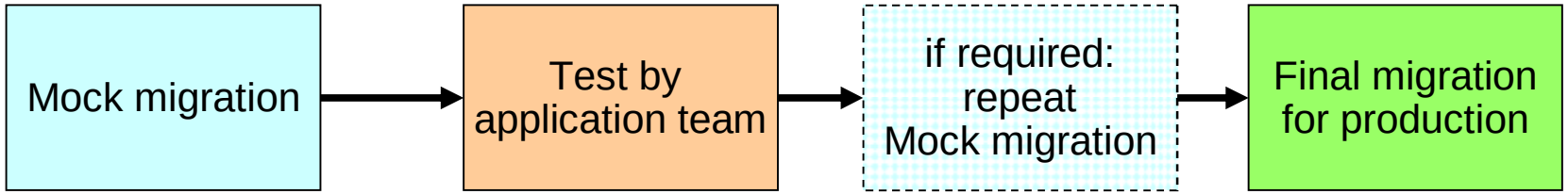
- Review your current database environment in a planning session with the MF team
- We tell you how long it will take and how much it would cost.
- We perform automated data collection to establish the metrics for the databases to be migrated.
- We work with you to establish testing requirements and a cutover strategy.
- We prepare a detailed project plan.
- We manage and perform the migration of the required databases according to the plan to help ensure that risk, schedule and cost are correctly managed.
- We confirm that the migrated databases meet your testing requirements.
- We support you during cutover into production.
- We provide basic skills transfer for an established number of your personnel on the migration tasks performed during these services.




**THE IBM MIGRATION FACTORY HELPS ANSWER KEY QUESTIONS**

- "Can it be done?"
- "How is it done?"
- "What will it cost?"
- "How long will it take?"
- "What are the risks?"

# DB Migration Approach – Supported by IBM Migration Factory




- Take flash copy of DB
- Uses special tool (Xenobridge)
- Highly parallel copy to IFL environment
- Mock migration process can take several hours depending on DB size
- Short outage of production server to take flash copy

- Verify correct application functionality
- Customer task

- Re-run of mock
- Uses special tool (Xenobridge)
- Highly parallel copy
- Mock migration process can take several hours depending on DB size
- Application not available during this time
- Switch IP address, assign source DB server IP (x86) to new server on IFL

§ Additional service offerings/tools available to minimize outage time during migration
  – Continuous data replication (“CDC”)
  – More complex set-up

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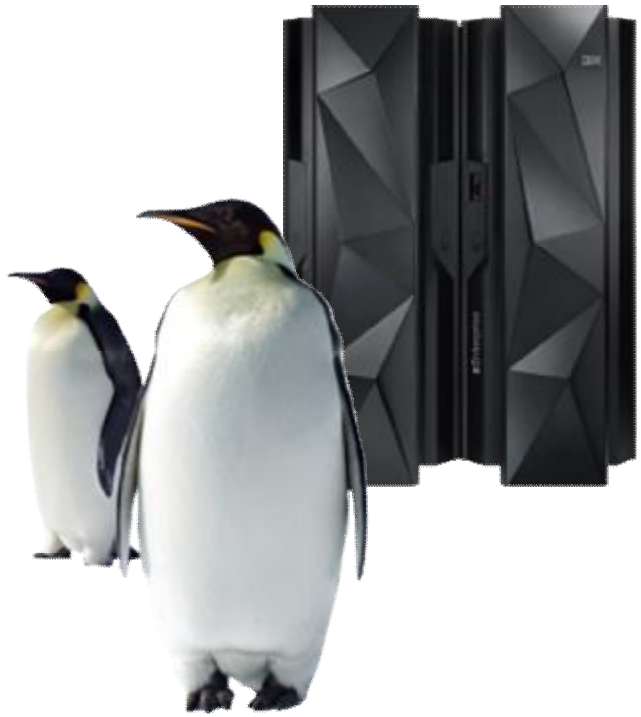


# Typical performance challenge

Customer reported performance issue:

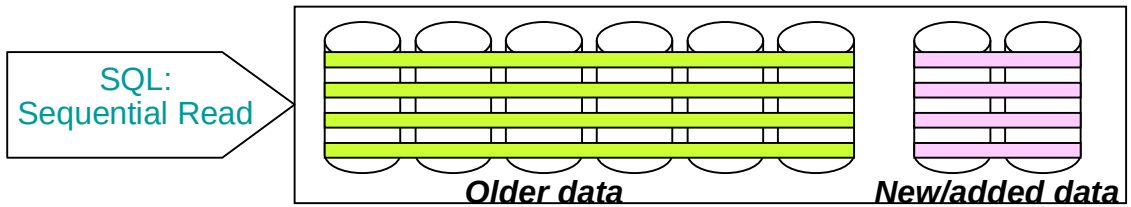
- Excessive run time for monthly business analysis run
- Application team states that no changes were made to the application

*However…..*

- Database size increased significantly
  - by about 12% in 3 month only
    April +45 GB, May +27 GB, June +32 GB
  - Added 2 Mod A disks (approx. 360 GB)

- Adding disk volumes has an impact on striping
  - New data striped over 2 volumes only (2 disks instead of 6)

ASM was not used!



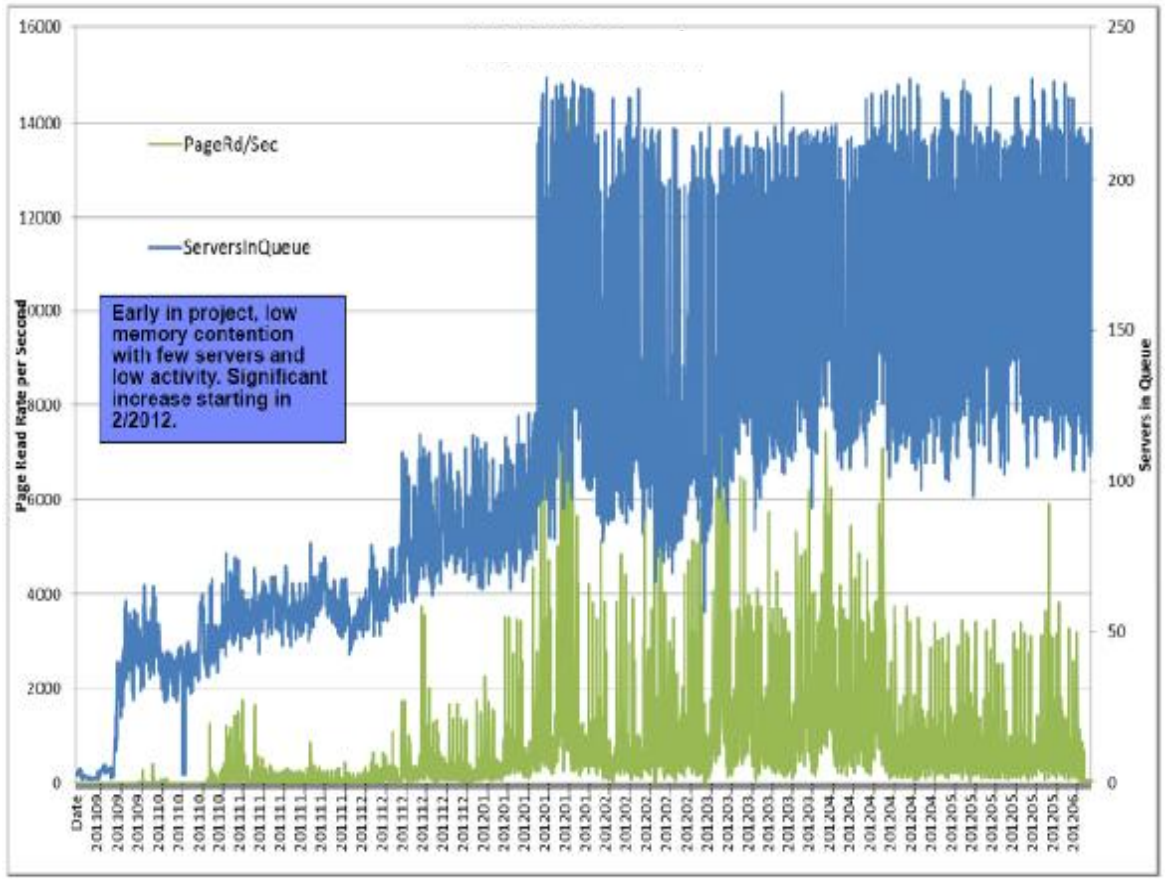
# Performance Degradation Over Time

## Problem:

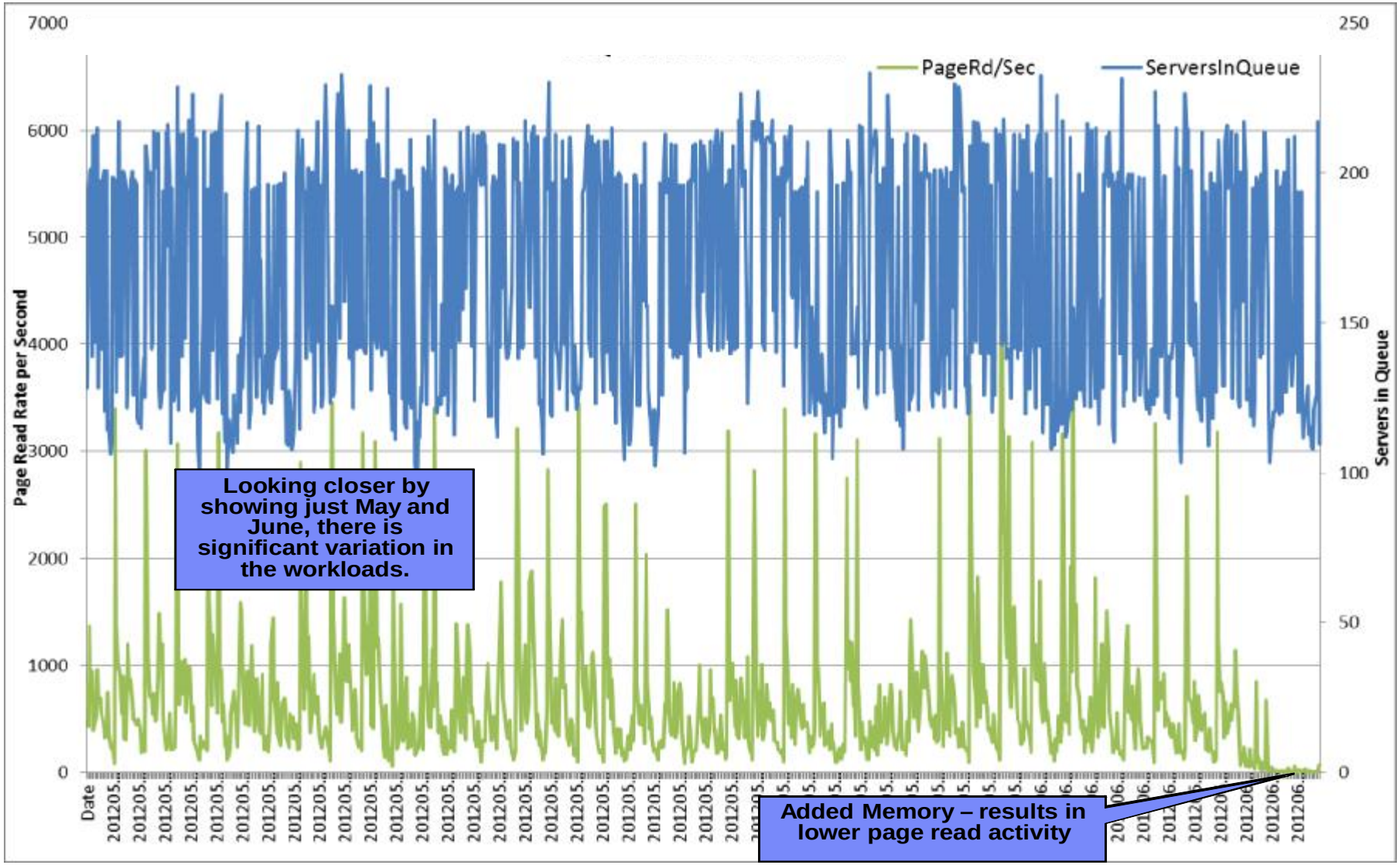
- The performance of selected servers/DB applications became worse over time with increased load on system

## Root cause:

- The add'l servers and increased activity led to increased memory contention
- Memory contention led to high paging rates to disks and internal systems management overhead (competing for memory between servers)

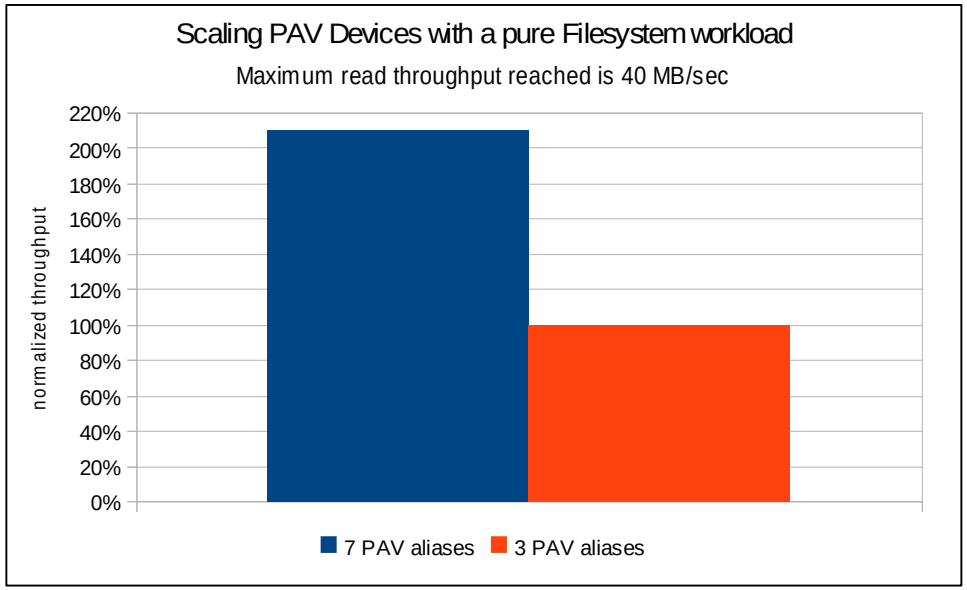


# Memory Over-commitment Changes



# PAV – Parallel Access Facility

- DASD and PAV devices are directly attached to the guests
  - For disk I/O intensive database workloads is this is the recommended setup
  - It is a requirement for using HyperPAV in Linux
- In case of Minidisk usage
  - Virtual PAV devices and a multipath setup for the Linux guest is required **and**
  - Physical PAV or HyperPAV devices in z/VM are required
- The amount of PAV devices is a critical parameter for disk throughput
- With 7 PAV devices the system can drive 2x more I/Os than with 3 PAV devices
- **Measurements showed that disk access is not a bottleneck with 7 PAV devices**
- *Measurement results are random I/O access pattern (not sequential I/O)*




*Notes:*

- *HyperPAV is not supported with RHEL 5.6 (supported with RHEL 5.9 & 6 and SLES 11)*
- *HyperPAV substantially reduces disk management (PAV-aliases do not need to be considered)*

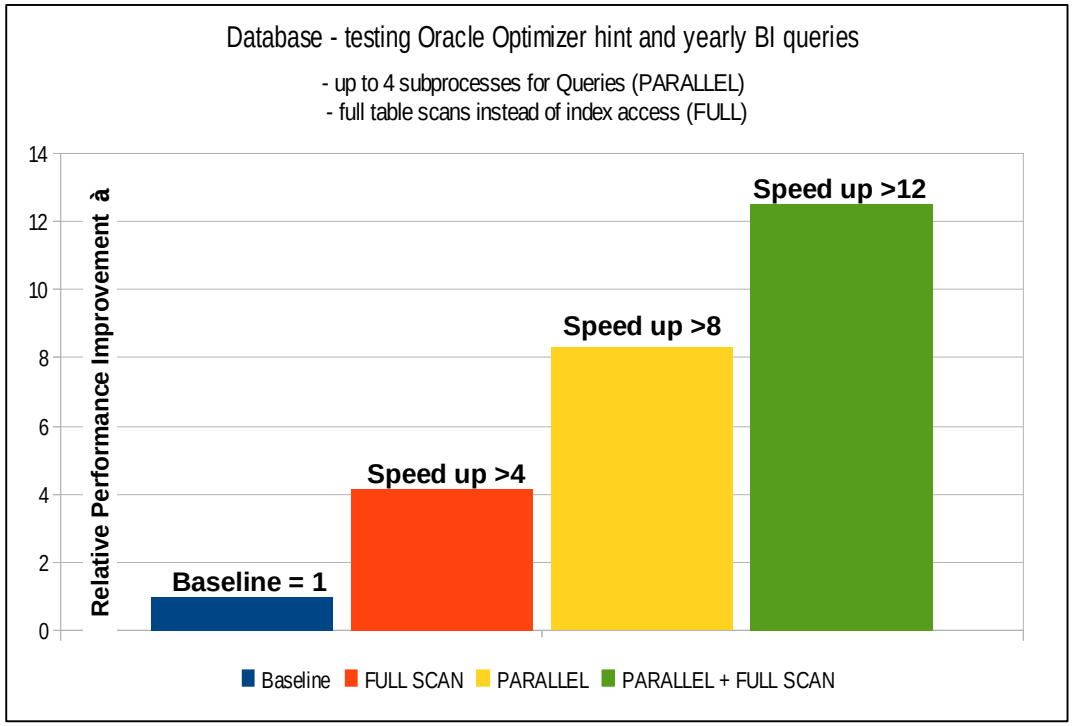
# Oracle DB Tuning Activities – Business Analytics Application 'A'

- **Actions taken – results:**
  - DB and application copied to a "sandbox" environment
    - Recreation of problem successful
    - Test runs with historical data from 2011
  - Used FIO (flexible I/O) tool to emulate a database like disk load and stress the disk devices (test achievable disk subsystem bandwidth)
    - Number of PAV devices (data striping – parallel access) increased from 3 to 7 per disk volume
    - Bandwidth increased from 4 MB/s to 8 MB/s
    - *rr_min_io* changed from 1000 to 1 (Linux default = 1000)
    - Bandwidth increased from 8 MB/s to 20 MB/s (in test)
  - Ø **Significant throughput increase for queries in monthly/yearly run**

- **Tests with Oracle optimizer show dramatic further speed-up**

# Oracle DB Tuning Activities – Business Analytics Application ‘A’

§ **Oracle optimizer hints** are specific for the SQL statement where specified
- 'FULL' force table scans vs index access
- 'PARALLEL' forces breaking up the statement into parts which can be executed in parallel in the same time
- 'PARALLEL' and 'FULL'




§ Risks
- Forcing a table scan can result in a severe performance degradation, when index access is the appropriate access method
- There might be reasons that a certain statements can not be executed parallel, then the behavior will not change

# Oracle DB Tuning Activities – Business Analytics Application ‘B’

- Multi-part workflow for data analysis
  - DB copied to a “sandbox” environment, directed the original workload against the “sandbox” system
  - Workload consist of
    - 3 steps (S, R, and D) with different workflows
    - only the last two steps (R and D ) are performance critical

- **Baseline**: **13 hours** run time for analysis with full year data
  - Initial migrated setup

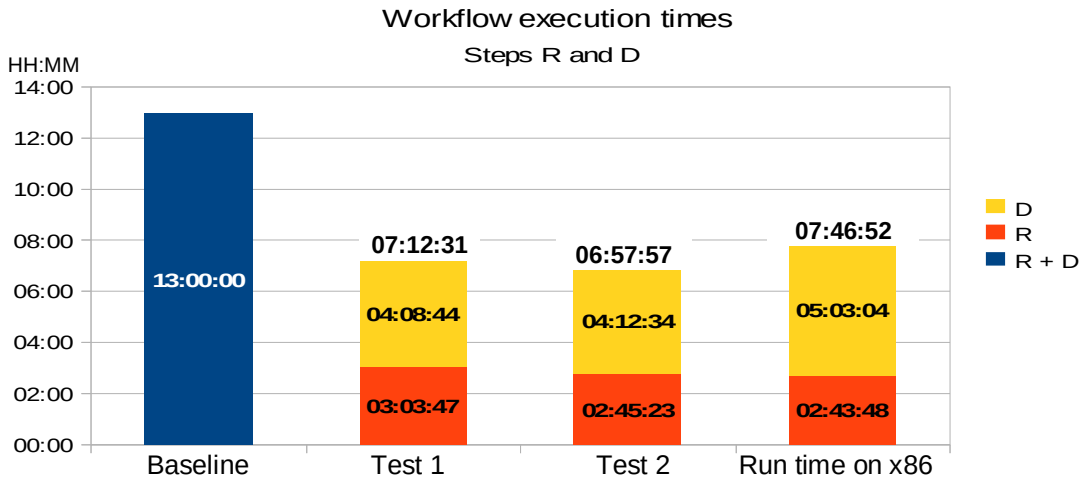
- **Test 1** (run time 07:12:31)
  - Environment related tuning (memory, disk setup, etc.)
  - Nearly factor 2x improvement

- **Test 2** (run time 06:57:57)
  - All tuning changes from Test 1 and
  - Database specific tuning (Oracle parameters)
  - Both tuning steps together provide an improvement of slightly more than factor 2x against the baseline

# Oracle DB Tuning Activities – Business Analytics Application 'B'

## *Parameter changes:*

§ Test 1 (run time 07:12:31)
- Added memory to LPAR
- Enabled 7 PAV devices per DASD device, directly attached to the guest,
- Multipath setup: round robin with *rr_min_io=1*




§ Test 2 (run time 06:57:57)
- Ensure that huge pages are really used → caused a SGA reduction from 8192MB to 7600MB *(better solution would have been to increase the amount of configured huge pages)*
- Profile parameter changes:

  *db_writer_processes=2* (prior 8),

  *filesystemio_options=setall* (prior asynch),

  *parallel_degree_policy=auto* (prior manual),

  *pga_aggregate_target=3700M* (prior 3,221,225,472)
- Added parameters:

  *log_buffers=104,857,600*
- Removed parameters:

  *disk_asynch_io,*

  *log_checkpoint_timeout,*

  *optimizer_index_caching,*

  *optimizer_index_cost_adj,*

  *shared_pool_size*

# Oracle DB Tuning Activities – Application 'C'

Oracle back-end for Windows application server - transaction workload

§ **Critical limit:**
- Requests should finish within 60 seconds
- Only 30 time-outs (>60 sec) are acceptable within 24 hour window

| | Known as Good case | Problem Case | After tuning action part 1 | After tuning action part 2 |
|---|---|---|---|---|
| Measurement Duration | 24 h | 23 h | 17.25 h | 48 h |
| Less than 3 Sec | 91,79% | 88,37% | 88,31% | 99,97% |
| 3 to 5 Sec | 2,74% | 3,35% | 3,69% | 0,02% |
| 5 to 10 Sec | 2,74% | 3,50% | 3,20% | 0,01% |
| 10 to 60 Sec | 2,58% | 4,48% | 4,27% | **0,00%** |
| More than 60 Sec | 0,16% | 0,30% | 0,53% | **0,00%** |
| More than 60 Sec | **13 requests** | **29 requests** | **24 requests** | **0 requests** |

§ **Tuning actions part 1:**
- Increased PAV devices from 3 to 7
- *rr_min_io = 1*
- Shut down inactive servers (reducing memory pressure)
- Further analysis showed a correlation with swapping activities - increased virtual memory size of Linux guest by 2 GB and activate direct I/O

  Environment monitoring showed good results, still getting time-outs

§ **Tuning actions part 2:**
- Increased number of vCPUs from 2 to 4, increased SGA by 2 GB

  **Dramatic improvement – no time-outs**

  **Results confirmed by longer term monitoring**

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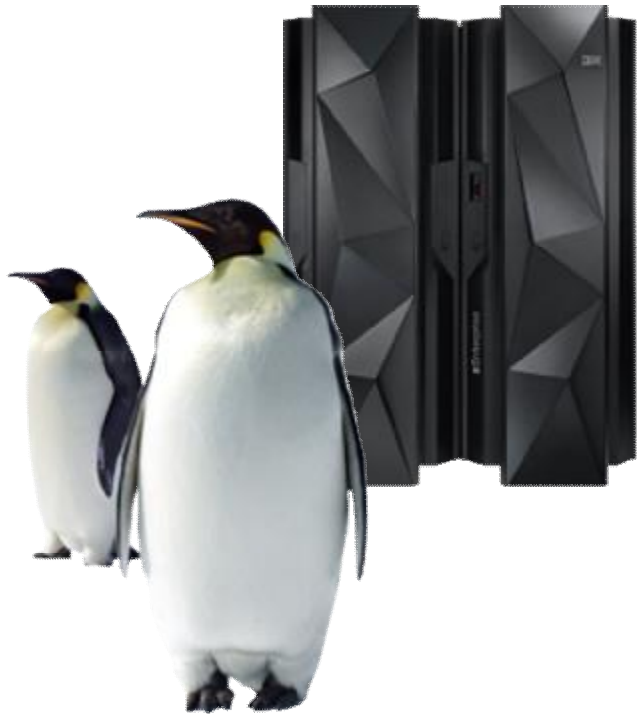


# General Recommendations – Monitoring

**Establish permanent monitoring**

§ z/VM Performance Toolkit

§ Linux sadc/sar

§ Tivoli OMEGAMON® XE on z/VM® and Linux
  – Tivoli Composite Application Manager (ITCAM) for Applications – Oracle Agent

- **Pro-active systems management**
  - **Detect potential problems/bottlenecks before users complain**
- **Capacity planning**
- **Accounting – charge back**

# General Recommendations – z/VM

- **z/VM Performance Toolkit**
- Ensure the virtual to real memory ratio stays in an appropriate range for the workloads
  - Indicators of impact:

    z/VM Paging activity

    Report 'User Paging Activity and Storage Utilization' (UPAGE, FCX113)

    Columns: 'X>DS' paging to DASD, critical: Reads paging from DASD

    z/VM Guest Waits

    Report 'Wait State Analysis by User' (USTAT,FCX114)

    Especially columns %PGW, %PGA, and %CFW

    z/VM CPU load

    Report 'System Performance Summary by Time' (SYSSUMLG, FCX225)

    Report 'General CPU Load and User Transactions' (CPU, FCX100)

- Disable Page reorder for guests larger than 8 GB
  - Find more information at http://www.vm.ibm.com/perf/tips/reorder.html

# General Recommendations - Linux

**Two possible disk devices for System z:**

§ Fixed (512-byte) blocks SCSI, connected with Fiber Channel Protocol (FCP) connection technology

- SCSI storage can be faster because it supports multiple parallel I/Os to a storage device
- FCP requires that you manually install FCP and configure multipath

§ DASD Disk I/O (FICON attached ECKD disks)

- Required: sufficient PAV devices (minimum 7 per disk) or HyperPAV (20 per LCU)
- In case of MDISKs use virtual PAV devices in Linux and physical PAV devices in z/VM. Use of HyperPAV would be the preferred method (supported in RHEL 6 and SLES 11).
- Multipath setup: *set rr_min_io* parameter to *1* (used for BI workloads)

  The rr_min_io value is storage dependent
  - For DS8K rr_min_io=100 provided good results for transaction processing
  - XIV recommends rr_min_io=15
- ECKD uses less CPU per transaction (utilizes SAP processors)

# General Recommendations - Linux

**Memory requirements:**

- Don’t over-configure Linux memory because -
  - Excess memory allocated to the Linux guest is used by Linux for I/O buffer and File system cache
  - In a virtualized environment under z/VM, oversized guests place unnecessary stress on the VM paging subsystem
  - Real memory is a shared resource, caching pages in a Linux guest reduces memory available to other Linux guests.
  - Larger virtual memory requires more kernel memory for address space management.

- Consider setting *vm.swapiness* to 0 (sysctl.conf) for all systems which are running primarily databases using page cache I/O
  - Defines a preference to reuse page cache pages instead of swap application pages

# General Recommendations – Linux
# Huge Pages

- If huge pages are configured, this amount of memory is no longer available for applications using 4K pages
  - Oracle 11*g* can use huge pages automatically
    If the SGA can not be allocated as a whole in huge pages, the fall back is to allocated the whole SGA in 4KB pages, which can produce a heavy memory pressure.
  - Ensure to have enough huge pages defined that the full SGA from **all** Oracle 11*g* databases in that system server fits into

- Check /proc/meminfo
  - HugePages_Total: configured huge pages,
    e.g via *vm.nr_hugepages*
  - HugePages_Free: unused part from HugePages_Total,
    but might be, not all are allocate-able due to memory fragmentation
  - HugePages_Rsvd: these are huge pages in any case available
  - pre-allocate huge pages on the kernel boot command line by specifying the "*hugepages=N*" parameter, where 'N' = the number of huge pages requested.
    This is the most reliable method for pre-allocating huge pages as memory has not yet become fragmented!

- To verify usage of Hugepages
  - Monitor value of HugePages_Free: When starting Oracle 11*g* the amount value of HugePages_Free must be lower (reduced by the SGA size)

# General Recommendations – Oracle parameters

- **Highly recommended:** parameter *filesystemio_options=setall*
  - In combination with this, remove definitions of parameter disk_asynch_io
- When defining **SGA_TARGET**, Oracle Database 10*g* automatically sizes the most commonly configured components, including:
  - The shared pool (for SQL and PL/SQL execution)
  - The Java pool (for Java execution state)
  - The large pool (for large allocations such as RMAN backup buffers)
  - The buffer cache
  - The Streams pool
  - Consider removing the existing definitions (if not sure) and let Oracle handle the sizing

    It defines lower limits and reduces the range Oracle can manage the buffers dynamically
- **Remove parameter *.log_checkpoint_timeout=0.**
  - It is not recommended to set this parameter unless FAST_START_MTTR_TARGET is set.
  - It is known as a potential cause for performance issues.
- **Define log_buffer = 104857600** or larger
- Be careful with specifying optimizer parameters (optimizer_...) as global parameters, because it might be an advantage only for some workloads.
  - Optimizer hints in the SQL statements are probably better because given for specific select statements

# General Recommendations – Oracle parameters

## § Log Setup

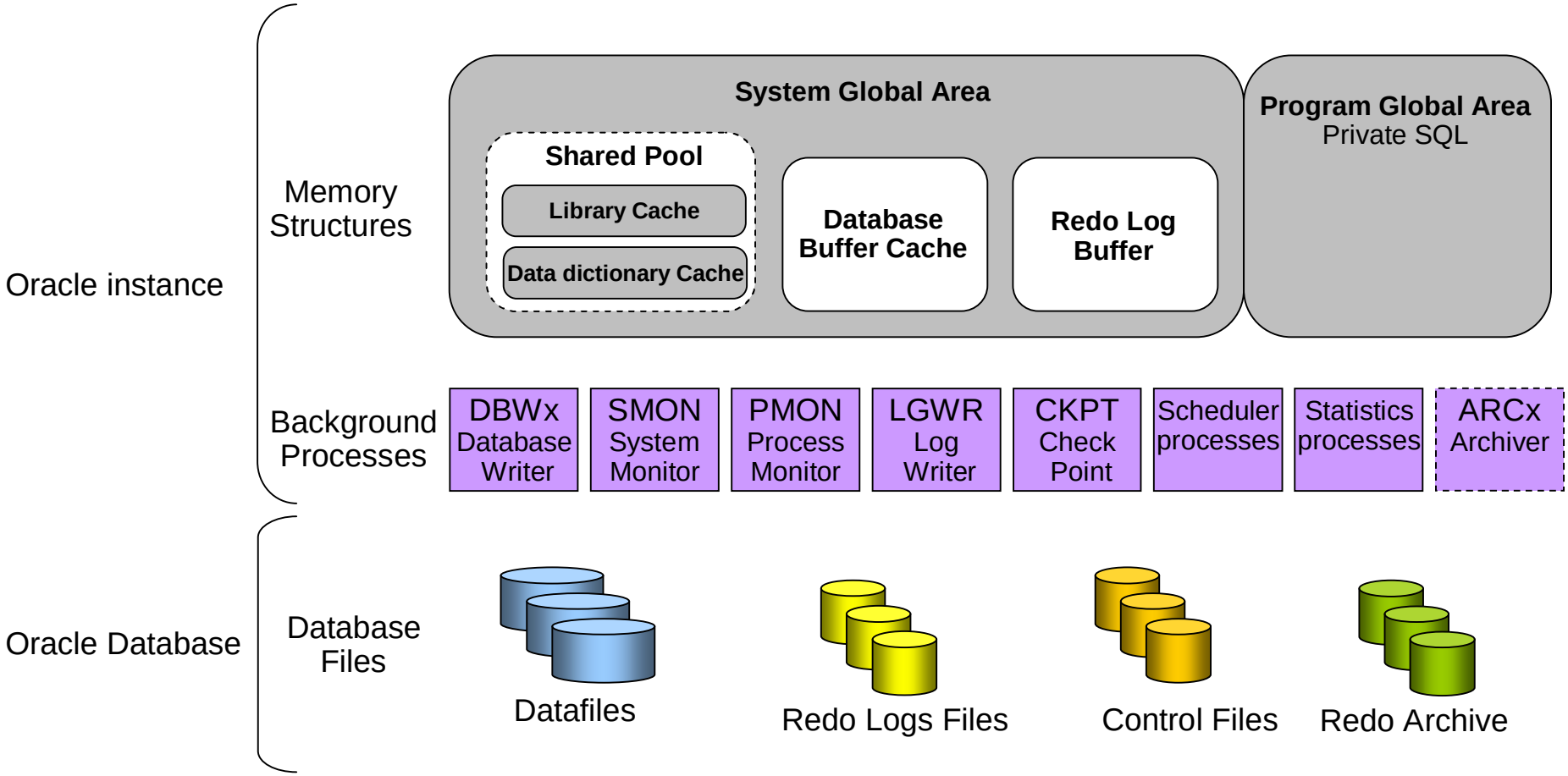
- Place redo logs on separate disks
  - Single disks are sufficient, striped LVM not needed
  - Ensure to have no other activity on these disks
- Recommendation: Usage of larger log files
  - e.g. 4x 1 – 1.5 GB to reduce the frequency of log switches

## § Review existing optimizer hints!

§ Customer workload specific experience with Oracle optimizer hints:

- Got very good improvements with the hints FULL(<table name>) and PARALLEL(<table name>, <number of CPUs>) for BI queries
- Suggest to review existing optimizer hints. Examples:
  - Combination of full(t) and parallel_index(t, 12) seems to be contradictory because usage of full table scan or index are mutually exclusive
  - Degree of parallelism specified with 12 seems to be much too high for a system with 4 vCPUs. A typical level for parallelism is <amount of vCPUs> or <amount of vCPUs + 1>, the upper limit is no more than 2X the number of cpus/virtual cpu
- For Oracle 11g consider to specify parallel_degree_policy=AUTO instead of explicit optimizer hints to let Oracle decide about parallelism

# Oracle server architecture



# Example of memory sizing

- Standard Memory estimation = sum of:
  - Memory required for Linux Kernel: 512 MB
  - Memory required for Oracle SGA: As per DBA estimation
  - Memory required for Oracle PGA: As per DBA estimation
  - Memory required for Oracle ASM: 256 MB to 512 MB (If ASM is used)
  - Memory required for additional agents like OEM, Tivoli etc., as needed by the application
  - Linux Overhead requirements: 5 % of the total memory

**Starting size = SGA + PGA + 0.5GB for Linux + ASM (if used)**

- Memory over-commitment (relationship of virtual to real memory)
  - Limit/avoid memory over-commitment for critical production databases
  - Test/development guests can benefit from z/VM memory over-commitment capability

# ASM

- Oracle ASM is an Oracle instance with a smaller SGA than regular database
- Oracle ASM is Oracle's methodology for striping database files across as many disk devices as possible.
- Oracle ASM is a form of software striping to raw or block devices
- When configuring ASM make sure that Disk/LUNs are assigned with the same size, type, and speed.
- Oracle ASM for Oracle 11g utilizes a 1 MB stripe size to stripe the database files across all the disk devices assigned to a particular disk group.
- Oracle REDO logs are also striped across the disk devices in the disk group, but are internally striped with a 128 KB stripe size.
- Oracle recommends the SAME approach for ASM files as well, by having one or two disk groups (if utilizing a Flash Recovery Area) and not separating the data and index data files into different disk groups.

# ASM or LVM

- LVM – Logical Volume Manager in Linux
- ASM – Automated Storage Management provided by Oracle
  - Oracle RAC One and Oracle RAC will require ASM

| | LVM | ASM |
|---|---|---|
| **Pro** | § Direct control on setting and layout<br>§ Can choose file system | § Automated, out of the box environment<br>§ Very good integration with Oracle |
| **Con** | § Complex setup | § RMAN required for backup |

- Overall recommendation: **ASM**

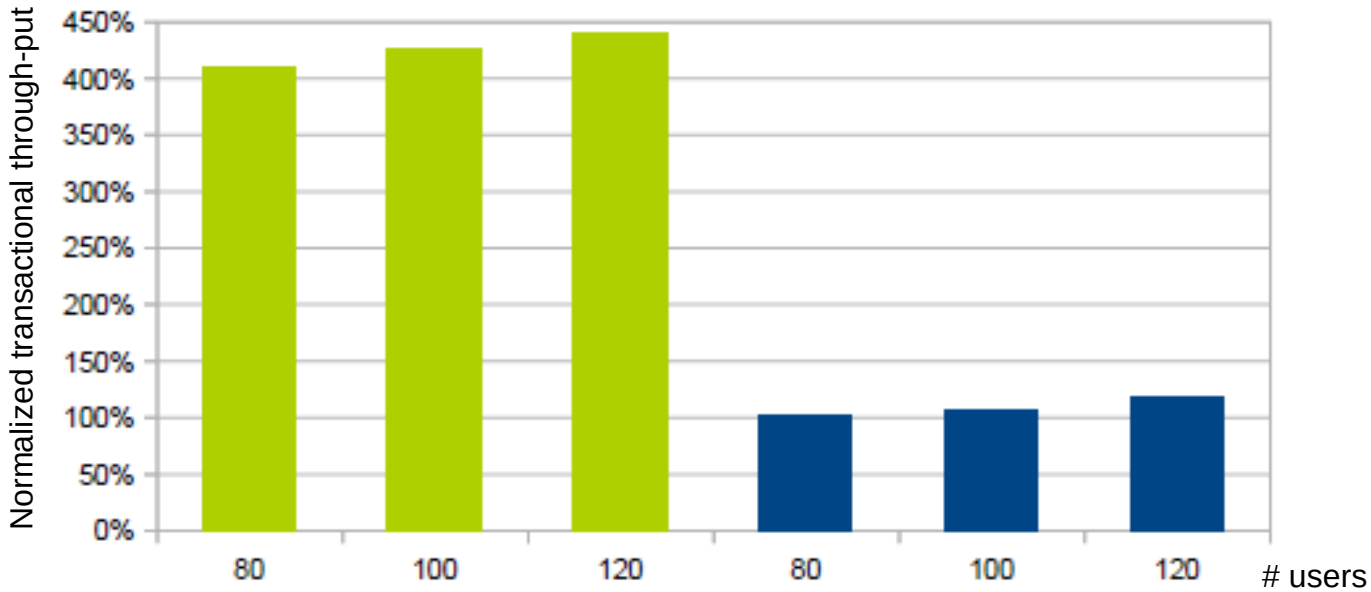
# Best practices – Oracle and Linux on System z

- Big database servers (SGA >100 GB) should be run in LPAR rather than as z/VM guest
- As z/VM guest use as few virtual processors as possible
  - The number of guest processors (virtual CPU) should be less or equal to the number of processors of z/VM LPAR
- Busy Linux database servers as z/VM guest should be given enough guest memory so that paging for this guest can be minimized
- There should be at least 2 GB of Expanded Storage defined for z/VM
- Size a Linux database server as z/VM guest that it just does not swap
- Use direct I/O for database files
  - Right-sizing the buffer pool is more beneficial than having additional Linux page cache
- Separate database disks and disks for logging/archive log
- Define sufficient I/O bandwidth for database disks
  - For SCSI discs, define multipathing and failover (understand & consider disk architecture)
  - For ECKD disks, use HyperPAV (SLES 11, RHEL 6) or define PAV aliases (more is better)
- Use data striping
  - ASM is Oracle's methodology for striping database files across as many disk devices as possible
  - XIV disk storage system has its own internal striping

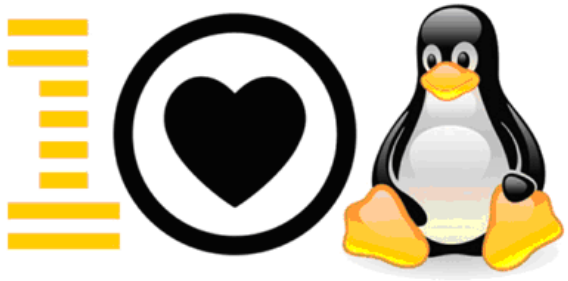
# Oracle 11*g* OLTP improvements

**Comparison: Oracle 10*g* versus 11*g* database**

User scaling – transactional through-put




Recommendation: upgrade to 11*g*R2 if not already done



# Questions?

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