



# VM Performance Update

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4/14/2008

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## Introduction

- **Some Post z/VM 5.2.0 News**
  - PAV
  - OMEGAMON XE
- **z/VM 5.3.0 Performance**
  - Line Items that have an impact
- **APARs of Interest**
- **z10 Performance**
- **See Performance Report on web for details**
  - <http://www.vm.ibm.com/perf/reports/>

## PAV Exploitation for VM Minidisks

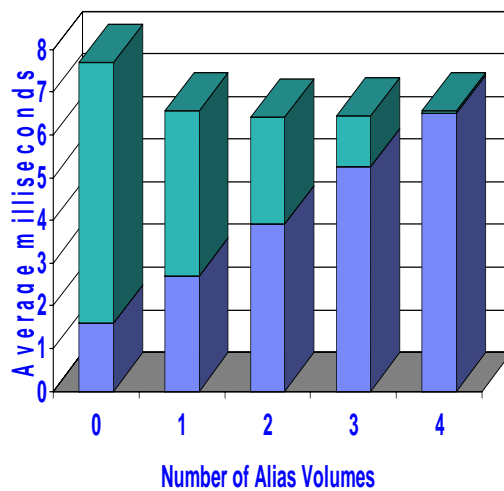
- **Previously only supported as dedicated disks.**
- **VM63855 (for z/VM 5.2.0, available May 2006):**
  - CP uses PAV to potentially decrease response time on minidisk I/O
  - We tightened the rules about ATTACHing or DEDICATEing PAV devices
- **VM63855 – virtualizes PAV for minidisks.**
- **Useful for environments where queuing on I/O occurs for minidisk I/O.**
- **Sometimes referred to as SYSTEM-owned PAV volumes**
- **PAV Base and Alias volumes defined on the Storage CU**
- **Summary of Results**
  - Varies depending on DASD CU Model
  - Varies depending on read-write mix
  - Helpful when I/O queuing occurs
  - Law of diminishing return; that is, defining more Alias than needed can lower performance

## PAV – Rules of Thumb

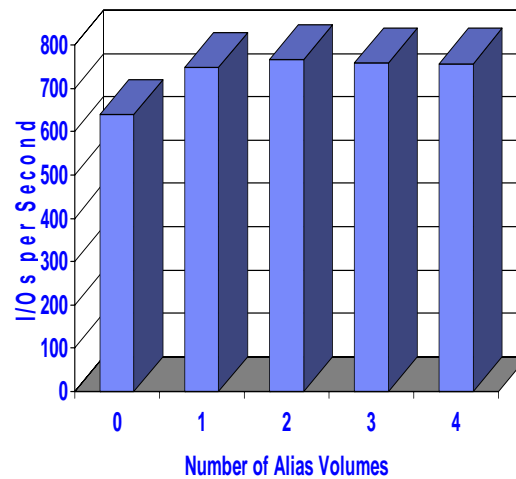
- **Symptom:**
  - I/O wait queue forming at real volume where minidisks are
  - See Performance Toolkit FCX168 reports (or equivalent)
- **Remedy:**
  - Configure a PAV alias device in the storage controller
  - Make sure the alias device is varied online
  - Make sure the alias device is ATTACHED to SYSTEM
- **Measure:**
  - Re-run your workload
  - Look again at those disk performance reports
- **Success criterion:**
  - Response time equals service time (no wait queue)

## System Owned PAV Results – DS8100 – 100% Writes

### Device Response Time

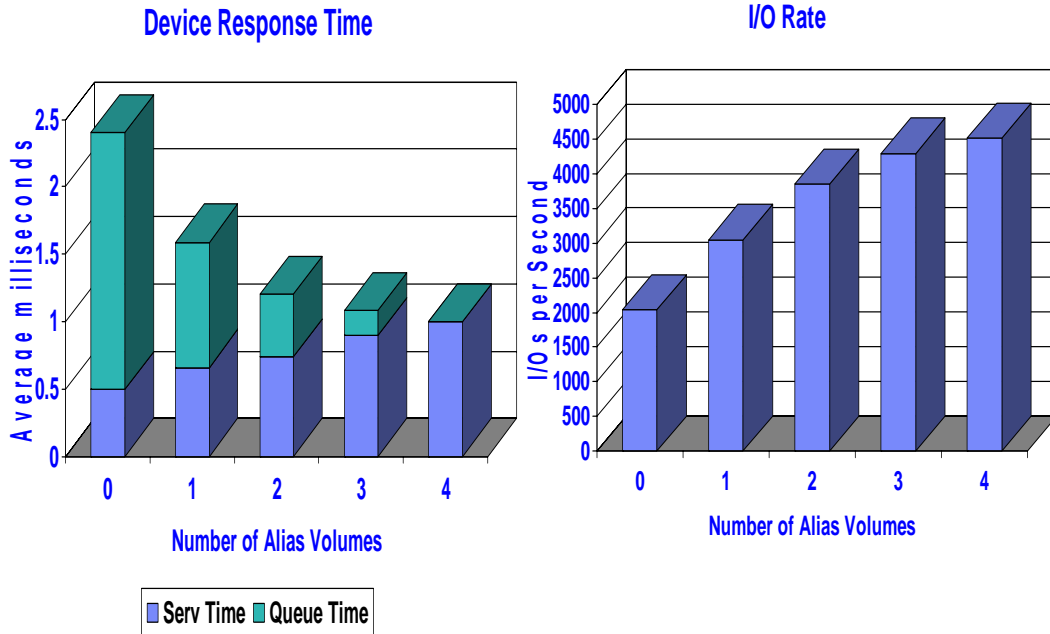


### I/O Rate

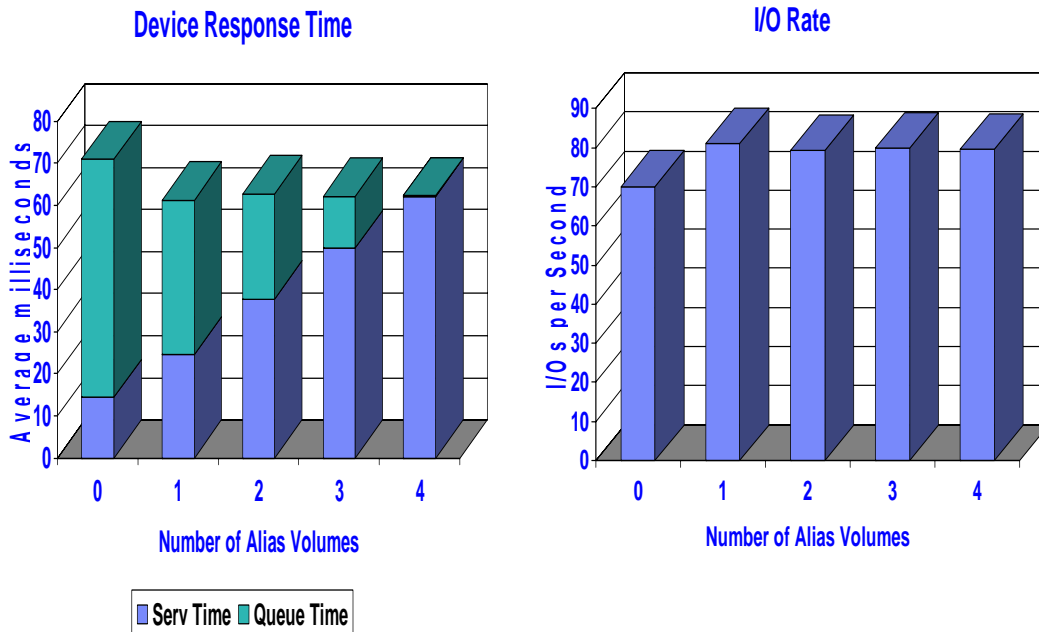


■ Serv Time ■ Queue Time

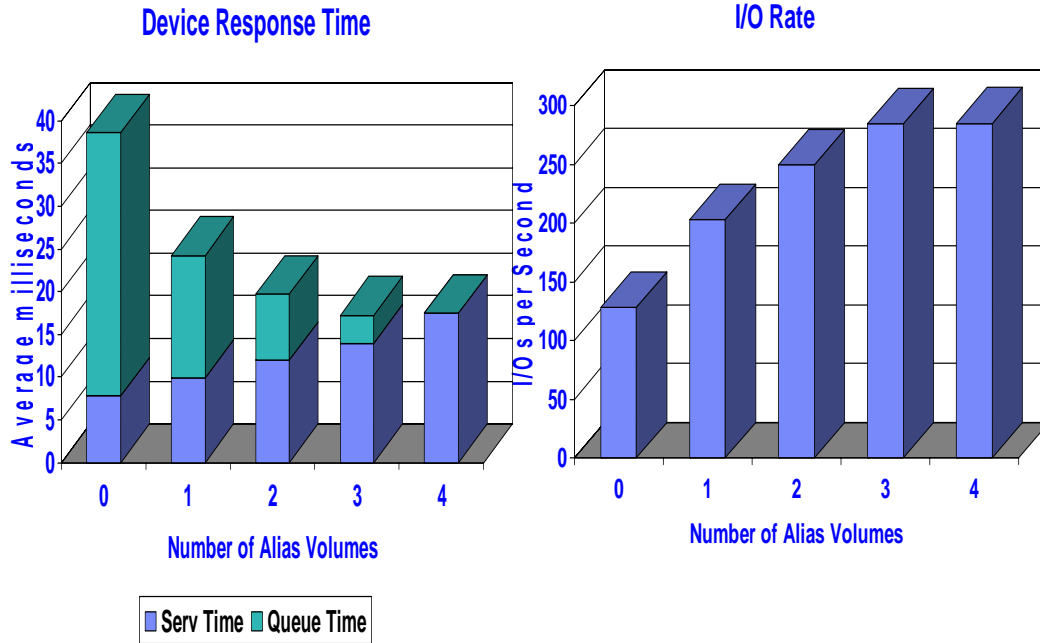
## System Owned PAV Results – DS8100 – 100% Reads



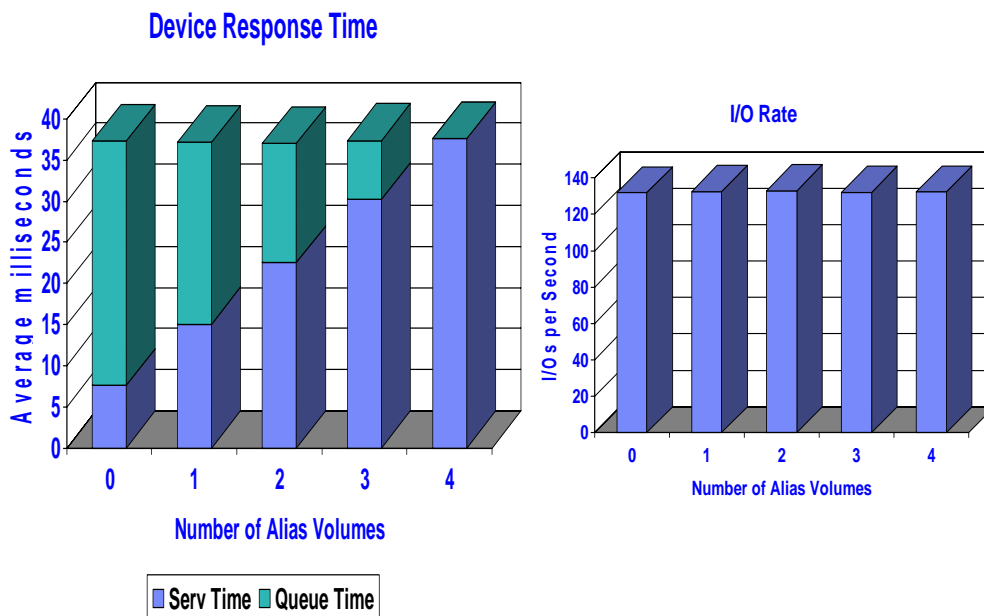
## System Owned PAV Results – DS6800 – 100% Writes



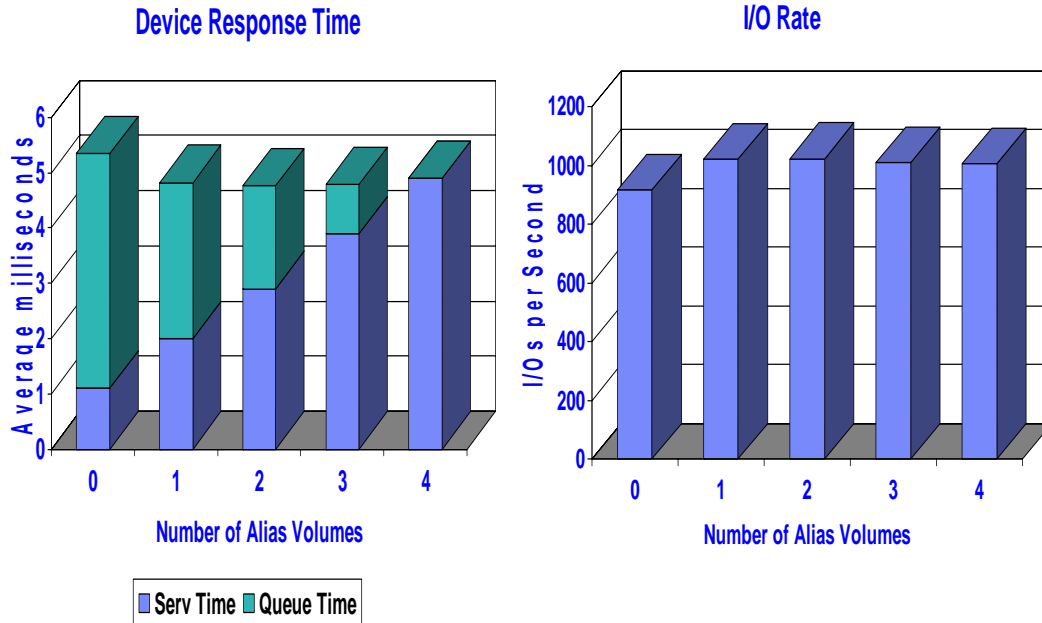
## System Owned PAV Results – DS6800 – 100% Reads



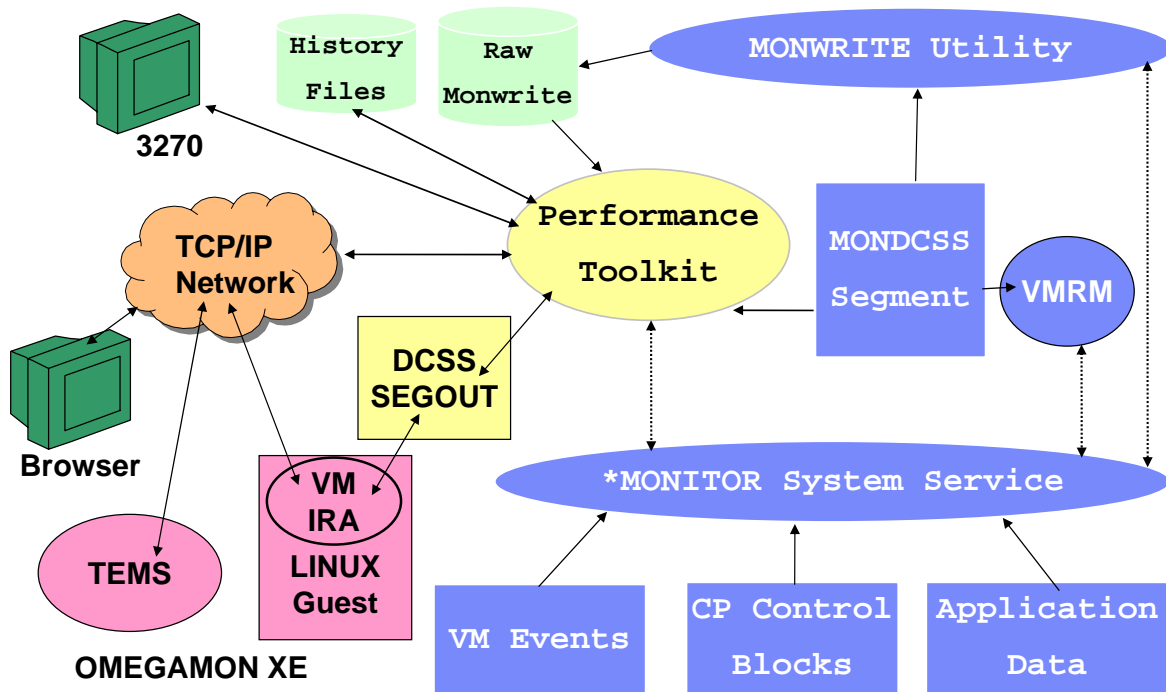
## System Owned PAV Results – ESS F20 – 100% Writes



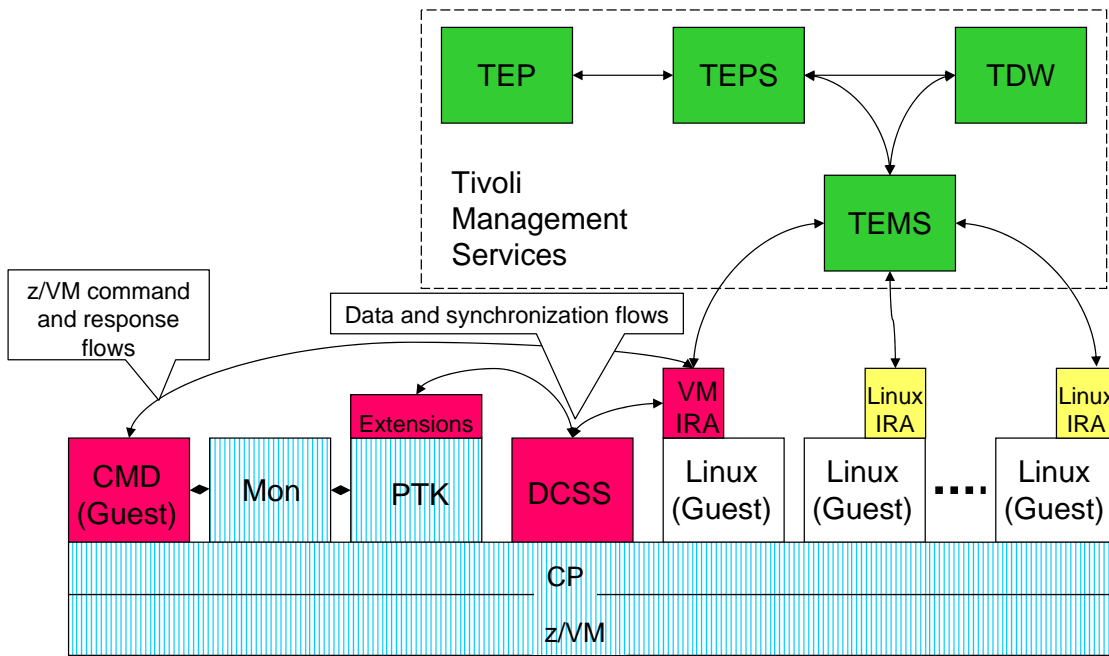
## System Owned PAV Results – ESS F20 – 100% Reads



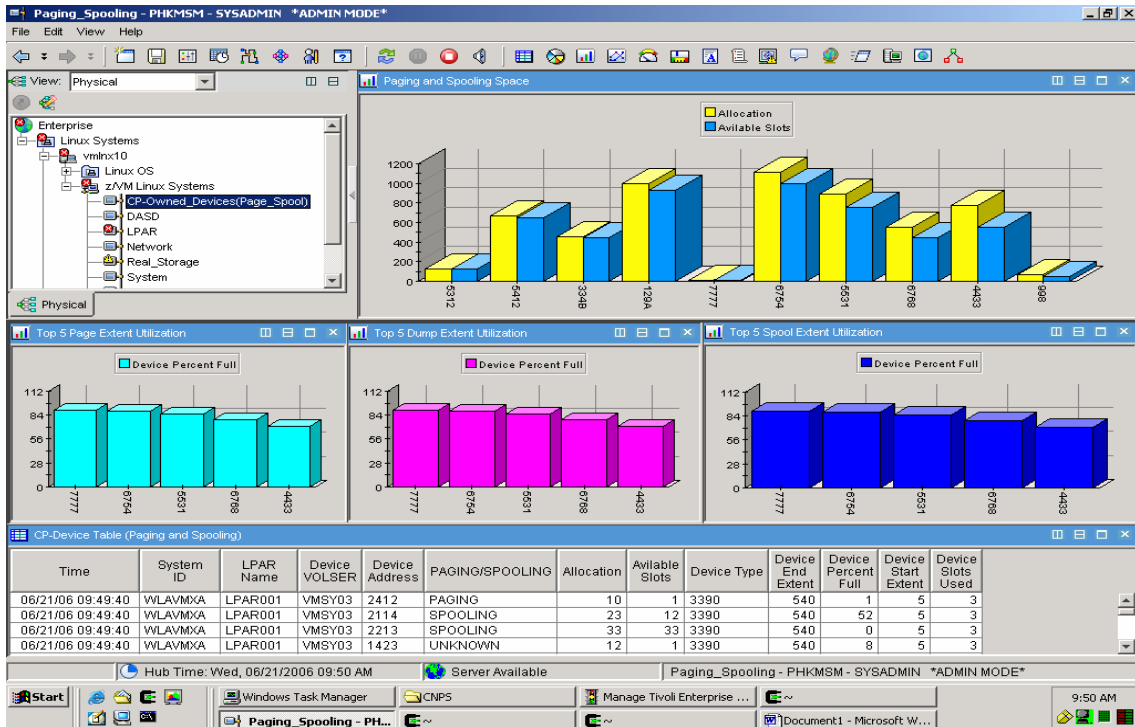
## 5,000 Foot View



# Basic Architecture



## PAGING and SPOOLING Utilization

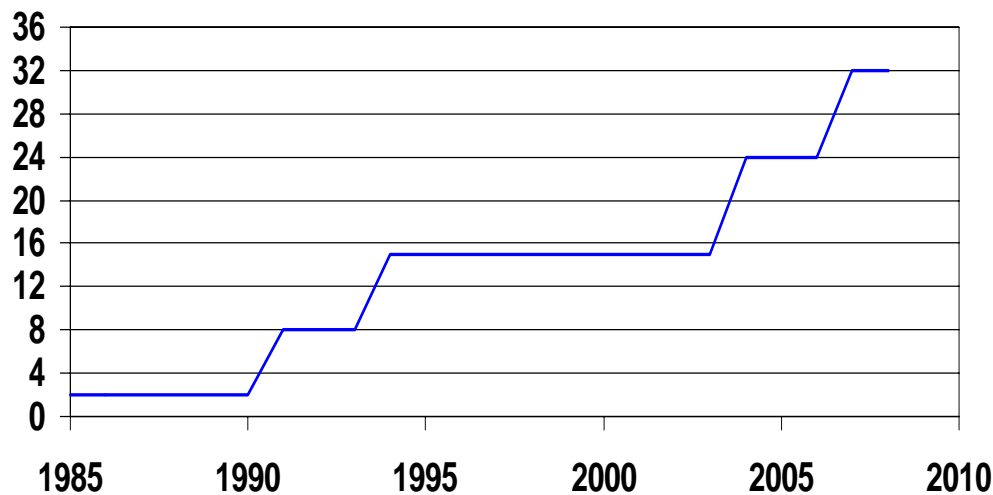


## z/VM 5.3.0

- **GA June 29, 2007**
  - Many more details available in the z/VM Performance Report:
    - <http://www.vm.ibm.com/perf/reports/>
- **Scalability and capability extended in several directions**
  - Processors, Memory, I/O, Network
  - What were the old limits?
  - What are the new limits?
- **Other Performance Enhancements**

## Processor Scaling

### Number of Supported Processors

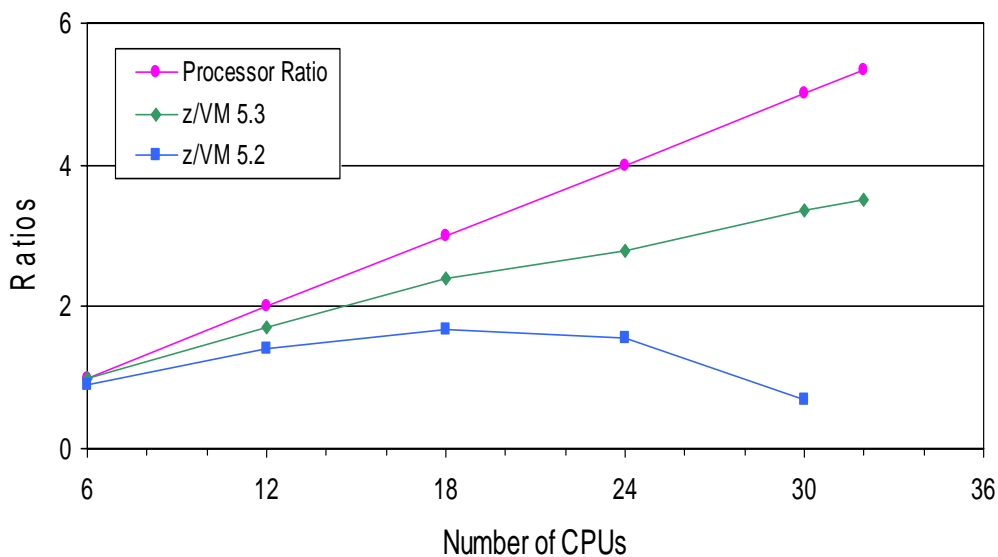




## Greater than 24 CPU Support

- **While z/VM 5.2.0 would run on up to 31 processors, it only supported 24 due to performance limitations**
- **z/VM 5.3.0 supports 32 processors**
- **Serialization Changes**
  - General support for exclusive and shared formal spin locks
  - First to exploit is the Scheduler Lock (SRMSLOCK)
  - New lock associated with each Processor Local Dispatch Vector (PLDV) for dispatching (DSVLOCK)
- **Performance is Workload Dependent**
  - Watch for Master Processor Limitations
    - Tend to be more traditional workloads, not Linux environments
  - Single non-MP virtual machine limits
    - Example: DB2 for z/VM & VSE can only use 1 processor

## Large N-Way Effects on ITR Ratio



Linux Guests with Apache Webserving

## Metrics for Formal Spin Locks

FCX265 CPU 2094 SER 19B9E Interval 02:31:51 - 12:34:01 GDLVM7

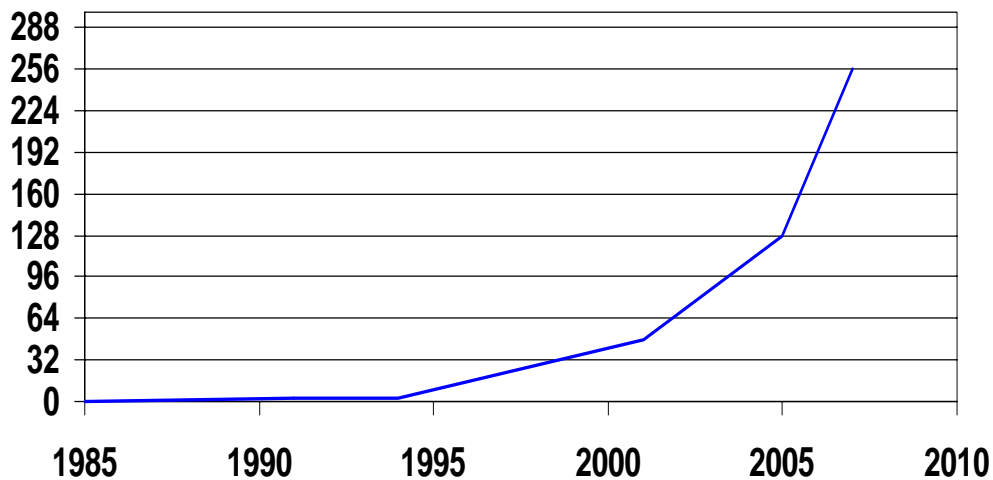
```

<----- Spin Lock Activity ----->
<----- Total -----> <--- Exclusive ---> <----- Shared ----->
Interval                Locks Average  Pct  Locks Average  Pct  Locks Average  Pct
End Time LockName      /sec   usec  Spin  /sec   usec  Spin  /sec   usec  Spin
>>Mean>> SRMATDLK      1.9    .539  .000   1.9    .539  .000    .0    .000  .000
>>Mean>> RSAAVCLK       .0    2.015  .000    .0    2.015  .000    .0    .000  .000
>>Mean>> FSDVMLK        .0   24.97  .000    .0   24.97  .000    .0    .000  .000
>>Mean>> SRMALOCK        .0     .000  .000    .0     .000  .000    .0    .000  .000
>>Mean>> HCPTRQLK       4.1     .195  .000    4.1     .195  .000    .0    .000  .000
>>Mean>> SRMSLOCK      34.0    1.096  .001   32.7    1.037  .001    1.3    .001  .000

```

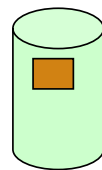
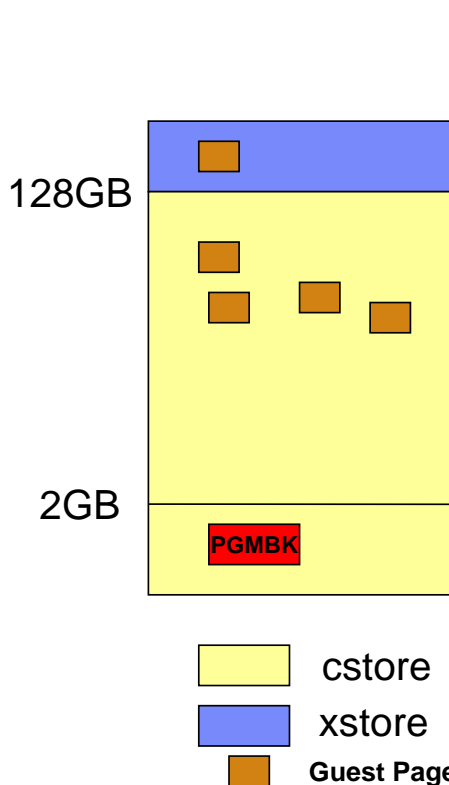
## Memory Scaling

### Effective Real Memory Use Limits



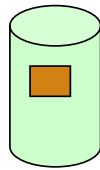
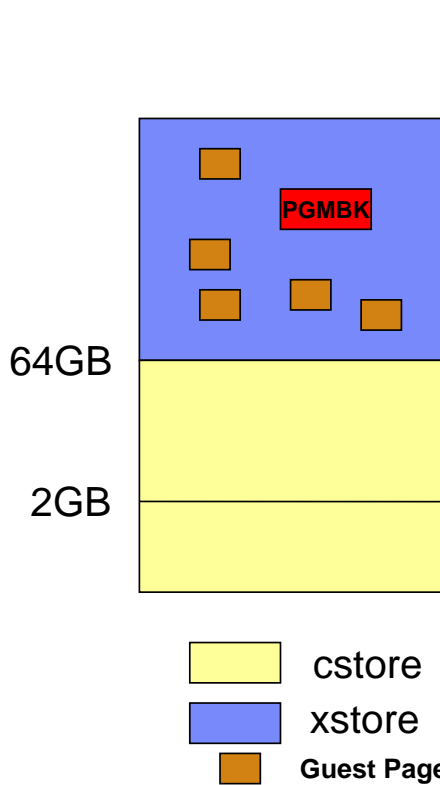
## Greater than 128GB Memory Support (256GB)

- **z/VM 5.3 Improvements:**
  - PGMBKs allowed to be allocated above 2GB
    - Each PGMBK is 8KB (2 contiguous frames)
  - Enhanced contiguous frame management
- **Also seeing improvements to smaller configurations that are memory constrained**
- **Be careful with memory terminology**
  - Try to define various terms when you use them or hear them
  - Examples:
    - Defined
    - Resident
    - Backed
    - Active
    - Actively Referenced
    - Addressable



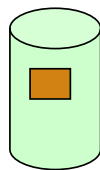
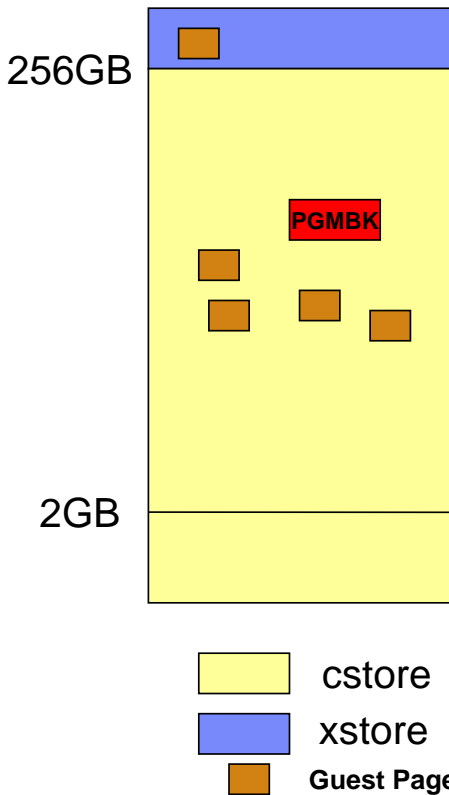
### z/VM 5.2.0 Traditional Xstore:Cstore Ratio

- PGMBK is required for any guest page that is acted on
- PGMBK is 8KB (2 contiguous frames)
- Resident PGMBK located below 2GB
- PGMBKs are pageable
- PGMBK resident if any guest pages it represents are resident
- Perfect world: ~256GB of virtual memory actively being referenced
- Realistic world: ~128GB of virtual memory actively being referenced



### z/VM 5.2.0 Larger Xstore:Cstore Ratio

- More expanded storage can increase the likely hood that guest pages are moved out of central storage.
- This allows PGMBKs to also be moved.



### z/VM 5.3.0 PGMBKs above 2GB Bar

- PGMBKs can reside anywhere but still must be better in the hierarchy than the guest pages they represent
- Next Limit: amount of PTRM space
  - Each space 4GB in size mapping 500GB of memory
  - Limit of 16 Spaces
  - Totals: 8TB of virtual machine memory

## Performance Toolkit DSPACESH Report

```

FCX134      CPU 2094  SER 19B9E  Interval 13:04:01 - 13:09:01  GDLVM7
              <-----Number of Pages----->
Owning              <--Resid--> <-Locked--> <-Aliases-->
Userid   Data Space Name   Total Resid R<2GB   Lock L<2GB   Count   Lockd   XSTOR   DASD
SYSTEM   PTRM0000           1049k 35602  1104      0      0      0      0      980 7502

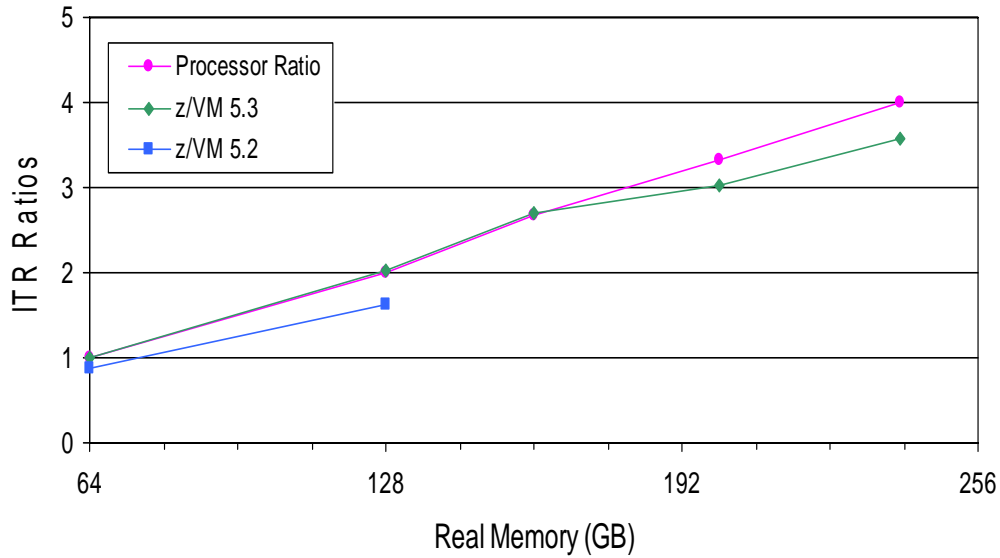
```

- Slightly edited FCX134 report
- PGMBKs live in PTRM0000, PTRM0001, ... PTRM000F
- Most systems will just have a PTRM0000

## Limitations for Memory

- **Memory limitations dependent on workload & configuration**
  - 256GB Real memory
  - 8TB of 'addressable' virtual machine memory – Limit of Page Tables
  - Paging Space (optimal when <50% full)
    - 11.2TB for ECKD
    - 15.9TB for Emulated FBA on FCP SCSI
  - Virtual Machine Size (HW Dependent)
    - 1TB on z9

## Scaling Memory Results - Apache Webserving



## Improvement in Memory Constrained Environments

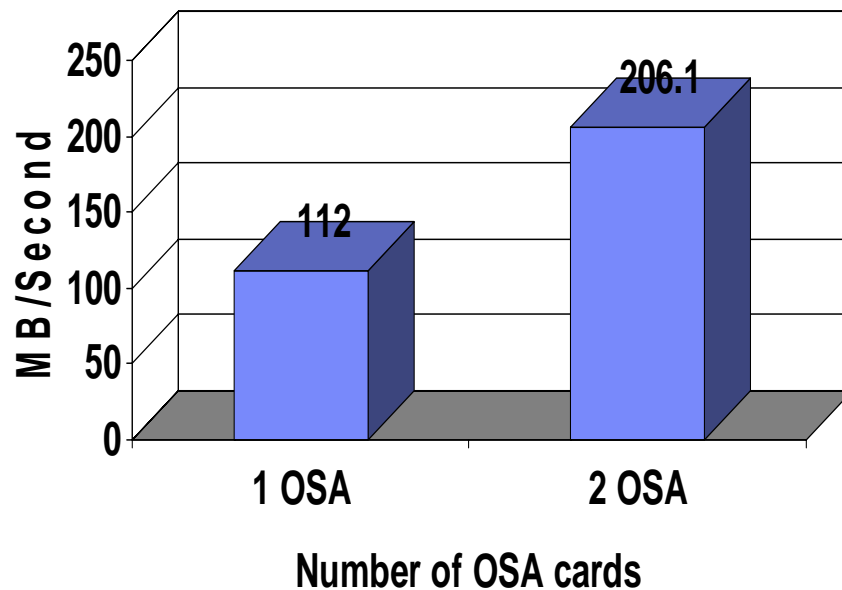
Scenario	Contention (Page Reqs per CPU-second)	Delta Thruput	Delta Total CPU/Tx
<b>3G/4G 2084 3-way</b>	<b>2159</b>	<b>+10.3%</b>	<b>-9.5%</b>
<b>64G/2G 2094 3-way</b>	<b>0</b>	<b>+1.0%</b>	<b>-0.9%</b>
<b>64G/2G 2094 3-way</b>	<b>352</b>	<b>+15.5%</b>	<b>-12.7%</b>
<b>128G/2G 2094 6-way</b>	<b>291</b>	<b>+21.6%</b>	<b>-19.4%</b>

Results of various Linux Apache measurements comparing z/VM 5.3 to z/VM 5.2

## Virtual Switch Link Aggregation

- **Ability to attach multiple OSAs to a single virtual Switch**
  - Aggregate bandwidth
  - Failover
- **Requires:**
  - z9 OSA-Express2 Support
  - Running in Layer 2 Mode
- **Dynamic Load Balancing**
  - Influenced by distribution of MAC addresses
  - Influenced by Physical Switch for inbound traffic
  - Cannot balance a single connection.
    - Example: a single data streaming connection will not get split across OSAs.

## Streaming Throughput Results



## I/O Improvements

- **PAV support for minidisks was provided in z/VM 5.2 via APAR VM63855 (May 2006)**
  - Should also apply VM64199 if using Minidisk Cache
- **z/VM 5.3 adds support for the HyperPAV feature of IBM System Storage DS8000**
  - Requires VM64248 (UM32072)
  - Allows for the creation of pools of Alias devices which the control unit will associate with different Bases as needed.
  - Performance characteristics similar to previous PAV support.

## Improved SCSI Disk Performance

- **Exploitation of SCSI write-same function of 2105 & 2107 improves CMS FORMAT of minidisks on SCSI volumes**
- **Additional pathlength reductions**
- **CP Paging to SCSI volumes now bypasses the FBA emulation, reducing processor resource requirements**



## Monitoring Enhancements

- **Lots of new fields in Monitor for new function:**
  - specialty engines
  - Scheduler changes
  - HyperPAV support
  - Memory management
- **New monitor Domain for Virtual Network Devices**
- **Additional flexibility in MONWRITE utility for starting/stopping**
- **Various changes in Performance Toolkit for VM**

## Service Must Haves: R530 APARs

- **VM64297 - PAGING SPIKE DEMAND SCAN SHUTS DOWN SYSTEM HANG**
  - PTF UM32197 Available and on RSU3
- **VM64269 - EXCESSIVE PAGING ACTIVITY DURING DEMAND SCAN**
  - PTF UM32133 Available and on RSU2
- **VM64287 - SLOW PING RESPONSE TIMES OVER QDIO WITH MORE THAN ONE VIRTUAL PROCESSOR**
  - PTF UM32158 available 10/18/07

## Must Haves: R530 APARs

- **VM64249 - PEVM63853 LPAR CHECK STOP DURING EDM H/W UPGRADE**
  - PTF UM32104 Available and on RSU2
  - HW Field Alert updated to include PTF numbers and additional information on z/VM 530 symptoms
  - Linux-only logical partitions when an Enhanced Driver Maintenance (EDM) occurs
- **VM64323 - PEVM63853 NO LPAR MONITOR RECORDS AFTER EDM UPGRADE**
  - PTF UM32196 Available and on RSU3

## z10 Performance

- **IT DEPENDS!!!**
- **Processor cycle time greatly improved over z9**
  - ~2.6 times faster (4.4 GHz)
  - Comparable to other platforms
- **Laws of Physics Must be Obeyed**
- **Tradeoffs made in order to achieve above**
  - Memory Differences
  - Key Ops
- **ITR Ratios (examples see LSPR for most current numbers)**
  - z/OS: z10 EC 701 up to **1.62 times** that of the z9 EC 701
  - LSPR z/VM Measurements: 1.30 to 1.60
  - z/VM Endicott Lab measurements: 1.23 to 2.05

## It Depends On....

- **Number of Processors**
  - Fewer processors, better ITRR
- **Storage References**
  - Smaller memory footprints, better ITRR
- **Data Movement**
  - Less data movement, better ITRR
- **Virtual I/O to Real Devices**
  - Less virtual I/O, better ITRR
- **Storage Overcommitment**
  - Less over commitment, better ITRR
- **Amount of memory involved in long searches**
  - Shorter & less frequent searches, better ITRR
- **Exploitation of New Features**
  - More exploitation of features, better ITRR

## Setting the proper expectations

- **z10 is a great machine, with a number of excellent attributes.**
- **Care must be taken when sizing migrations from z9 to z10.**
- **Additional Information:**
  - LSPR Q & A (complete)
    - Discuss range and factors affecting
    - Pointer to z/VM Web Page
  - z/VM Web Page
    - <http://www.vm.ibm.com/perf/z10.html>
  - “To MIPS or Not to MIPS, That is the Question!” by Gary King
    - [http://shareew.prod.web.sba.com/proceedingmod/abstract.cfm?abstract\\_id=17583](http://shareew.prod.web.sba.com/proceedingmod/abstract.cfm?abstract_id=17583)

## Summary

- **z/VM 5.2 Improvements via Service**
- **z/VM 5.3 significantly extends the capacity of:**
  - Processor
  - Memory
  - I/O
- **See z/VM Performance Report for more details**
  - <http://www.vm.ibm.com/perf/reports/>
- **Learn more about z/VM**
  - <http://www.vm.ibm.com/events/>