

IBM Advanced Technical Support

A Beginner's Guide to Measuring and Understanding z/VM Guest Performance

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NOTES:

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Agenda

Introduction

Resources that affect z/VM guest performance

- > What they are?
- How they can be controlled?
- ➢ How to recognize a problem using the Performance Toolkit for VM[™]

Summary



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Performance Problem Analysis Methodology

- Performance problems result from high usage and/or contention for key system resources
 - > CPU
 - Memory
 - > I/O
- Begin measuring system performance characteristics before problems occur
 - Use Performance Toolkit for VM or other monitoring tools
 - Understand performance characteristics of well running system
- When a problem occurs:
 - Identify the resource in contention
 - Reduce or eliminate the contention
 - Reduce the demand for the resource
 - Increase the resource capacity
 - Reallocate the usage of the existing resource



Performance Toolkit for VM



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z/VM Internals - Definitions

Virtual Machine Definition Block (VMDBK)

- > One per virtual processor defined to a virtual machine
- > CP representation of a virtual machine

Scheduler lists

- Dormant list
- Eligible list
- Dispatch list

Start Interpretive Execution (SIE)

Method z/VM uses to dispatch a virtual processor on a real processor

Working Set Size (WSS)

> Scheduler estimate of memory pages needed to run a virtual machine



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Resources That Affect Performance

•• • Processor

Memory

I/O

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Physical Processors

- z/VM5.3 supports up to 32 physical processors
- Physical processors may be dedicated to virtual machines
 - Not commonly done
 - > A dedicated processor not available to run other virtual machines
- z/VM dispatches virtual processors on physical processor using the Start Interpretive Execution (SIE) instruction
 - Intercept occurs based upon criteria specified by CP
 - SET SRM DSPSLICE amount of time a virtual CPU can use a physical CPU before being interrupted



Physical Processors...

• Utilization of physical processors is controlled by:

Number of VMDBKs in the dispatch list

Dispatch list is controlled by:

>SET DSPBUF command

> Other SRM commands such as STORBUF & LDUBUF

Utilization reported by:

>CP INDICATE LOAD command – reports utilization based upon a 4 minute smoothed average

Performance Toolkit for VM and CP Monitor also report utilization without the smoothed average



Physical Processors...

Utilization is divided into:

- Time where CP is running on real processor (%CP)
 - > Performing tasks not charged to a specific user (%SYS)
 - Scheduling, and other housekeeping
 - Performing tasks on behalf of virtual processors (%CP %SYS)
 - Instruction simulation, etc.
- >Time where a virtual processor is running on a real processor (%EMU)
 - Actual user instructions



Virtual Processors

- Virtual machines may be defined with up to 64 virtual processors
 - Recommend not defining more virtual processors than physical
- Virtual processors may be dedicated or shared
 - Recommend not mixing shared and dedicated processors
 - Dedicated processors are rarely necessary
- SHARE setting represents portion of system resources virtual processors should receive
 - Absolute or Relative
 - > Target minimum and maximum values
 - Maximum values can be either hard or soft limits
 - > SHARE value divided between virtual processors



1. CPU load and trans.



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21. User resource usage

FCX112	CPU	J 7060	SER :	10431	Interv	val 14	4:08:19	9 - 14:0	39:19	ZV	'MV 4R 40
•	<	CPU	Load	···>	<	vi	rtual	10/s	>		
Userid	%CPU	TCPU	VCPU	Ratio	Total	DASD	Avoid	Diag98	UB	Pa/s	User Status
>Sustem<	14.2	8.538	8.531	1.0	.0	.0	.0	.0	.0	.0	
WOJLINUX	99.7	59.79	59.75	1.0	.1	.1	.0	.0	.0	. 0	ESA, CL0, DIS
WOJLIN02	99.5	59.72	59.68	1.0	.0	.0	.0	.0	.0	.0	ESA, CL0, DIS
VMRTM	.03	.018	.005	3.6	.0	.0	.0	.0	.0	.0	ESA,, DOR
PERFSVM	.01	.005	.004	1.3	.1	.1	.0	.0	.0	. 0	ESA,,DOR
DISKACNT	0	0	0		0	0	0	0	0	0	ESA,,DOR
EREP	0	0	0		0	0	0	0	0	0	ESA,,DOR
FTPSERVE	.00	.000	.000		.0	. 0	.0	.0	. 0	. 0	XC,,DOR
MAINT	.00	.000	.000		.0	. 0	.0	.0	. 0	. 0	ESA,,DOR
OPERATOR	0	0	0		0	0	0	0	0	0	ESA,,DOR
OPERSYMP	0	0	0		0	0	0	0	0	0	ESA,,DOR
TCPIP	.00	.001	.000		.1	. 0	.0	.1	. 0	. 0	ESA,CL0,DIS
VMSERVR	0	0	0		0	0	0	0	0	0	ESA,,DOR
VMSERVS	0	0	0		0	0	0	0	0	0	XC,,DOR
VMSERVU	0	0	0		0	0	0	0	0	0	XC,,DOR

From this report we can see that the Linux guests were using all of the CPU – the z/VM scheduler distributed processor resource fairly evenly across the virtual machines that needed it (each has SHARE Relative 100).

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3B. Proc. load & config (PROCSUM)

FCX239	CPU	7060	SER 1	9431	Inte	erval 1	3:30:19	14:25	:19	ZVMV4R4	40
	<	CPU)	>	>Spi	in Lock	Activi	.ty	>		
		KRat	io>		>1 -	·>	K S	Scheduler	>	<sie <="" td=""><td>/sec></td></sie>	/sec>
Interval	Pct	\frown	Cap-	On-	>ge	Pct	Locks	Average	Pct	In-	Inter-
End Time	Busy	T/W	ture	line	>ec	Spin	/sec	usec	Spin	struct	cept
>>Mean>>	99.9	1.00	.9979	2.0	>52	.000	.0	8.145	.000	424.3	21.8
14:04:19	99.9	1.00	.9979	2.0	469	.000	.0		0	422.6	20.5
14:05:19	99.9	1.00	9976	2.0		0	.0		0	423.5	21.3
14:06:19	99.9	1.00	9981	2.0		0	.0		0	423.1	20.7
14:07:19	99.9	1.00	9981	2.0		0	.0		0	422.0	19.9
14:08:19	99.9	1.00	9981	2.0		0	.0		0	422.7	20.5
14:09:19	99.9	1.00	.9975	2.0	026	.000	.0		0	422.0	20.2
14:10:19	99.9	1.00	.9981	2.0	483	.000	.0		0	427.8	24.5
14:11:19	99.9	1.00	.9981	2.0	211	.000	.0	3.063	.000	424.6	21.9
		\bigcirc								<u> </u>	1

T/V ratio shows low CP overhead (0%), further evidence of this is the ratio of SIE intercepts/second to SIE Instructions/second. Intercepts less than half the number of SIE Instructions is consistent with low CP overhead.

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1. CPU load and trans.



This example shows low CPU utilization, but fairly high I/O activity (virtual I/O rate is 205 per second). This results in a high number of privileged instructions per second and potential higher CP overhead (T/V ratio).

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21. User resource usage

FCX112	CPU	J 7060	SER	10431	Interv	val 1	5:19:19	9 - 15:3	20:19	ZV	/MV4R40
	<	- CPU K-Seco	Load onds->	; т/у	<	Vi	rtual :	IO/s	>		
Userid	%CPU	TCPU	VCPU	Ratio	Total	DASD	Avoid	Diag98	UR	Pg/s	User Status
>Sustem<	.36	.218	.115	1.9	14.6	14.4	.0	.1	.0	.0	,,
WOJLINUX	4.43	2.655	1.386	1.9	202	202	> .0	.0	.0	.0	ESA,CL0,DIS
WOJLINUZ	.50	.345	.204	1.7	. 2	. 2	.0	.0	. 0	.0	ESA,CL0,DIS
VMRTM	.03	.019	.005	3.8	.0	.0	. 0	.0	.0	.0	ESA,,DOR
MAINT	.02	.011	.003	3.7	. 6	.0	.0	.0	.0	.0	ESA,, DOR
PERFSVM	.02	.010	.008	1.3	. 2	. 2	. 0	.0	.0	.0	ESA,,DOR
TCPIP	.02	.013	.007	1.9	1.5	. 0	. 0	1.5	.0	.0	ESA, CL0, DIS
DISKACNT	0	Θ	Θ		0	0	Θ	0	0	0	ESA,, DOR
EREP	0	0	Θ		0	0	Θ	0	0	0	ESA,, DOR
FTPSERVE	.00	.000	.000		. 0	. 0	. 0	.0	.0	.0	XC,, DOR
OPERATOR	0	Θ	Θ		0	0	Θ	0	0	0	ESA,,DOR
OPERSYMP	0	0	Θ		0	0	Θ	0	0	0	ESA,, DOR
VMSERVR	0	0	Θ		0	0	Θ	0	0	0	ESA,, DOR
VMSERVS	0	0	0		0	0	0	0	0	0	XC,, DOR
VMSERVU	0	0	0		0	0	0	Θ	0	0	XC,,DOR
	FCX112 Userid >Sustem WOJLINUX WOJLINU2 VMRTM MAINT PERFSVM TCPIP DISKACNT EREP FTPSERVE OPERATOR OPERSYMP VMSERVR VMSERVS VMSERVU	FCX112 CPU Userid %CPU >Sustem< 35 WOJLINUX 4.43 WOJLINUZ 56 VMRTM .03 MAINT .02 PERFSVM .02 TCPIP .02 DISKACNT 0 EREP 0 FTPSERVE .00 OPERATOR 0 OPERSYMP 0 VMSERVE 0 VMSERVS 0 VMSERVU 0	FCX112 CPU 7060 . .	FCX112 CPU 7060 SER · · · · · · · · · · · · · · · · · · · · · · · · · · · · · · · · · · · · · · · · <t< td=""><td>FCX112 CPU 7060 SER 10431 </td><td>FCX112 CPU 7060 SER 10431 Intervision </td><td>FCX112 CPU 7060 SER 10431 Interval 19 </td><td>FCX112 CPU 7060 SER 10431 Interval 15:19:19 <td>FCX112 CPU 7060 SER 10431 Interval 15:19:19 - 15:3 </td><td>FCX112 CPU 7060 SER 10431 Interval 15:19:19 - 15:20:19 </td><td>FCX112 CPU 7060 SER 10431 Interval 15:19:19 - 15:20:19 ZV</td></td></t<>	FCX112 CPU 7060 SER 10431 	FCX112 CPU 7060 SER 10431 Intervision 	FCX112 CPU 7060 SER 10431 Interval 19 	FCX112 CPU 7060 SER 10431 Interval 15:19:19 <td>FCX112 CPU 7060 SER 10431 Interval 15:19:19 - 15:3 </td> <td>FCX112 CPU 7060 SER 10431 Interval 15:19:19 - 15:20:19 </td> <td>FCX112 CPU 7060 SER 10431 Interval 15:19:19 - 15:20:19 ZV</td>	FCX112 CPU 7060 SER 10431 Interval 15:19:19 - 15:3 	FCX112 CPU 7060 SER 10431 Interval 15:19:19 - 15:20:19 	FCX112 CPU 7060 SER 10431 Interval 15:19:19 - 15:20:19 ZV

The virtual machine doing most of the I/O (WOJLINUX) shows high CP overhead (T/V ratio \rightarrow 2.655 / 1.386 = 1.92 92% overhead). Lots of instruction simulation!

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3B. Proc. load & config (PROCSUM)

FCX239	CPU	7060	SER 1	0431	Inte	erval 1	3:30:19	9 - 16:27	:19	ZVMV4R4	10
	<	CPU		>	>Spi	in Lock	Activ:	ity	>		
		KRat	io>		>1 -	·>	K 9	Scheduler	>	<sie <="" td=""><td>'sec></td></sie>	'sec>
Interval	Pct	\frown	Cap-	On-	>ge	Pct	Locks	Average	Pct	In-	Inter-
End Time	Busy	T/W	ture	line	>ec	Spin	/sec	usec	Spin	struct	cept
>>Mean>>	37.0	1.01	.9924	2.0	>56	.000	1.8	2.324	.000	530.5	373.9
15:16:19	2.6	2.65	7482	2.0	636	.001	9.6	1.895	.001	1525	1512
15:17:19	1.8	2.71	.7232	2.0	459	.001	7.4	1.599	.001	1042	1032
15:18:19	1.9	2.36	.7428	2.0	418	.001	7.0	1.604	.001	977.7	967.6
15:19:19	1.7	2.92	.6992	2.0	583	.001	7.1	1.952	.001	955.3	941.8
15:20:19	3.3	2.45	7723	2.0	551	.001	10.0	1.695	.001	1870	1851
15:21:19	3.8	2.57	.7645	2.0	572	.001	12.5	1.701	.001	2170	2152
15:22:19	2.3	2.68	.7401	2.0	382	.001	6.9	1.573	.001	1334	1323
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T/V ratio confirmed to be high. SIE Intercept/s to SIE Instruct/s ratio is high also, which is consistent with high I/O instead of high CPU.



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Resources That Affect Performance

Processor



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Central Storage

z/VM 5.3 supports up to 256GB of real memory as central storage for:

- > Virtual machine page frames
- z/VM nucleus
- Control blocks and other CP storage requirements
- z/VM virtual disk blocks
- Minidisk Cache

Prior to z/VM 5.2:

- Page frames used for I/O buffers and other similar uses were required to reside below 2GB
 - Movement of page frames from above 2GB to below 2GB is recorded in monitor records



Expanded Storage

- Used for:
 - High speed paging
 - Minidisk cache
 - Virtual disks
- Page frames must be moved from expanded to central storage before addressing the contents
- May be dedicated to a virtual machine
 - Linux can use expanded storage as a swap device
 - Dedicated XSTORE not available for use by CP
- Even with large amounts of central storage, it's best to define some expanded storage
 - Creates paging hierarchy page to xstore first, then migrate old pages to DASD
 - Good starting point is to define expanded storage equal to about 25% of central storage size (1GB-4GB)



Virtual Machine Storage

Defined on USER statement in USER DIRECT file (default & maximum)

> z/Architecture virtual machines may have greater than 2 GB storage

Virtual machine storage can be "locked" or "reserved"

- CP LOCK USERID username firstpagenum lastpagenum
 - Least flexible method to prevent paging
 - Requires knowledge of what specific pages should stay in central storage
- CP SET RESERVED userid nnn
 - Best choice to reduce paging for a virtual machine
 - Specify number of pages to maintain in central storage for virtual machine, but not the exact pages
 - CP will allow nnn pages to remain resident for the specified virtual machine at all times
- For Linux guests, keep virtual machine size as small as possible to help reduce allocation of I/O buffer and file system cache



Storage Planning

- Amount of central storage allocated does not need to be equal to sum of all logged on virtual machine sizes
 - Central storage size is a function of virtual machine working set sizes, and page turn over rate
 - > Hard to estimate for a completely new system
- Amount of DASD paging space allocated needs to be greater than the sum of all logged on virtual machine sizes
 - > DASD paging area utilization should not exceed 50%
 - > Higher utilizations reduce z/VM paging efficiency



Storage Planning...

Monitor effectiveness of minidisk cache:

- If little benefit and storage is constrained, set maximum cap for MDC or turn off minidisk cache so that pages can be used for virtual machine paging
- > CP SET MDCACHE STORAGE 0M 128M set cap for central storage
- CP SET MDCACHE XSTORE 0M 0M eliminates use of expanded storage
- > CP SET MDCACHE SYSTEM OFF turns off MDC
- Reduce size of CP TRACE table if storage is constrained
 - > CP SET TRACEFRAMES MASTER 100



Storage Control

- Scheduler can control storage and paging device utilization
- Virtual machines in dispatch list will have pages resident in central storage
 - Access to dispatch list limited by scheduler when central storage is constrained
 - Scheduler estimates working set size for each virtual machine placed in dispatch list
 - > Virtual machines exceeding storage thresholds placed in eligible list
 - > No access to physical processors from eligible list
 - Scheduler moves virtual machines from eligible list to dispatch list as central storage becomes available

CP SET SRM STORBUF and CP SET SRM LDUBUF commands influence scheduler behavior



Storage Control...

STORBUF

- Changes scheduler view of the amount of storage available for virtual machine pages
- > Typically need to modify default for Linux guest environment
- Modification over commits central storage, so a robust paging subsystem is necessary
- > CP SET SRM STORBUF 300 250 200

LDUBUF

- Changes scheduler view of what the paging subsystem can handle
- Applies to virtual machines classified as "loading users" (i.e. more than 5 page faults in a minor time slice)
- Leave at default or set all 3 values to 100, depending on number and size of paging devices available

> CP SET SRM LDUBUF 100 100 100



Storage Control...

QUICKDSP

- > Virtual machine attribute
 - > CP SET QUICKDSP userid ON
 - > OPTION QUICKDSP in USER DIRECT
- Instructs scheduler to NEVER place the virtual machine in an eligible list, even if central storage is constrained
- Only use this attribute for critical service machines that other virtual machines depend on
 - > TCP/IP virtual machine
 - Linux virtual machines acting as routers, database servers, etc.
- Specifying this attribute for all virtual machines inhibits scheduler's ability to manage central storage



Paging Device Guidelines

- Keep DASD page space utilization less than 50%
- Monitor blocks read per paging request
 - Greater than 10 is good
- Allocate page space over multiple volumes and multiple paths for best performance
- Do not mix paging areas with minidisk or other DASD usage
 - > CP uses never ending channel programs for paging devices
- For storage subsystems, make sure cache is enabled

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Monitoring Memory - Examples 1. CPU load and trans.

FCX100	Data for	2003/12/02	Interva	l 16:45:3	37 - 16:46:37	Monitor S	ican
CPU Load PROC %CPU P00 21 P01 28 P02 28	XCP XEMU 2 20 1 27 1 26	%WT %SYS 79 1 72 0 72 1	XSP XSIC 0 88 0 82 0 81	%LOGLD 21 28 28	Vector Facil XVTOT XVEMU 0 0 0 0 0 0	ity Statu REST ded. .0 .0	is or User
Total SSCH/F Virtual I/O Total rel. S	RSCH rate SHARE	89/s 4a 8/s 88 8867 To	ge rate TORE pagi t. abs SH	140.7/ ng 548.9/ ARE 0	Priv. i Diagnos	nstruct. 4 e instr. 1	165/s .92/s
Queue Statis VMDBKs in qu VMDBKs loadi Eligible VMD El. VMDBKs l Tot. WS (pag Expansion fa 85% elapsed	stics: Jeue DBKs Jes) 30 Sctor time 1	Q0 Q1 4 0 0 0 0011 0 .472 .184	Q2 0 0 0 0 1.472	Q3 24 0 1279k 8.832	User Status: # of logged # of dialled # of active # of in-queu % in-Q users % elig. (res	on users Lusers Lusers Le users Lin PGWAIT Lin IOWAIT Cource wait)	42 34 28 97 0
Transactions Average user Trans. per s Av. time (se UP trans. ti MP trans. ti System ITR (Emul. ITR (s Q- sec. ec) 1 ime ime trans.pe	Disp tri .3 .260 er sec. tot	vial no .0 .053 .053 .000 . CPU) . CPU)	n-tru .0 .000 .000 .000 .5	User Extreme Max. CPU % Max. VECT % Max. IO/sec Max. PGS/s Max. RESPG Max. MDCIO Max. MDCIO	S: ZSITL001 MONWRITE LPWWPE02 LPWWPE02 12 ZDEVL014 16	36.1 2.1 60.4 0517

System has low CPU utilization, but is experiencing paging to both XSTORE and DASD (migration of pages out of XSTORE).

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Monitoring Memory - Examples

2. Storage Utilization

	FCX103	Data for	~ 2003/12	/02	Interval	16:43:37 -	16:44:37	Monitor	Scan
	Main storage Total real 3 Total avail3 Offline stor SYSGEN stor3 CP resident Shared stor3 FREE storage FREE storage	e utiliz: storage able rage fran age size nucleus age e pages subpools	ation: mes	6'14 6'14 6'14 6'19 2'18 2'18 7'18 1'94	4MB 4MB 0KB 4MB 2KB 2KB 2KB 2KB	XSTORE uti Total avai Att. to vin Size of CP CP XSTORE (Low thresho XSTORE all Average age	lization: lable rt. machines partition utilization old for migr cation rate e of XSTORE e at migratio	2 2 . 2 blks on	'048MB 0kB '048MB 99% '720kB 200/s 4405s s
	Total DPA s Locked pages Trace table Pageable Storage uti Tasks waitin V=R area:	r. utili; ize s lization ng for a ng for a	frame page	2'02: 2'09: 1'00: 1'93: 10:	3% 1MB 2kB 0kB 0MB 9% 0 0/s	MDCACHE ut. Min. size Max. size Ideal size Act. size Bias for XS Min. size Max. size Ideal size	ilization: in XSTORE in XSTORE in XSTORE STORE in main stor in main stor in main stor	2 301 166 . 6 c. 559	0KB '048MB '604KB '664KB 1.00 '144MB '144MB '916KB '916KB
<	FREE storage V=R recover V=R user Paging / spo Page moves	y area in coling ac <208 for age-in r	tivity:		2/s	MDCACHE lin Users with MDISK cache MDISK cache MDISK cache MDISK cache	mit / user mit / user MDCACHE ins e read rate e read hit r e read hit r	erts ate atio	1.00 1.044kB 2 12/s /s 7/s 59%
	Long path pa Long path pa Page read ra Page write n Page write n Page write n Migrate-out Paging SSCH SPOOL read n SPOOL write	age-in r; age-out ; ate locking blocking blocking rate rate rate rate	ate rate factor factor g factor	19	5/s 5/s 7/s 0/s 6 • • 5/s 0/s	VDISKs: System lim User limit Main store Expanded s Pages on Di	it (blocks) (blocks) page frames tor. pages ASD	Un Un	lim. lim. 0 36

Check to see if there is a below 2GB constraint.



Monitoring Memory - Examples

3A. Paging activity

ECV142	Data	for 2	003/12	202 T		1 16:30	• 27 -	17:02:	37 Ma	niton (Roan
FUNING	Data	FOF 2	003/12/	02 1	nuerta	C 10.30	. 57 -	17.02.	or ne	nii core a	scan
	<		Expand	ded St	orage		>	<	Real	Storad	ae
			Fast-		7 \	Est.	Page	DPA	Non-	Resid	Mean
Interval	Paging	PGIN	Path	PGOUT	Total	Life	Miğr	Pgable	pgable	Shared	Avail
End Time	Blocks	/s	%	/s	/s	sec	7s	Frames	Pages	Pages	List
>>Mean>>	484804	93.2	97.4	114.5	207.8	4233	20.5	494092	23353	497	485
16:20:37	483790	85.4	99.7	112.6	198.0	4298	.0	495354	22091	528	2029
16:21:37	483790	39.8	99.9	47.9	86.8	10301	.0	495398	22047	530	1725
16:22:37	483790	63.5	99.9	11.9	135.0	6768		495224	22221	529	872
16:23:37	483790	34.5	99.7	42.4	76.9	11419	. 0	495865	21580	529	806
16:24:37	403730	40 7	99.9	42.4	62.0	11219	2.5	490004	22031	523	F00
16:25:37	483790	36.3	99.8	- 23.8	80.2	11037	6.1	494505	22940	528	530
16:27:37	483790	487.8	99.3	483.4	971.7	1000	1.0	493371	24074	527	360
16:28:37	483790	123.2	85.3	148.6	271.8	3256	28.2	492728	24717	529	328
16:29:37	482112	29.8	97.3	10.2	40.0	47189	23.0	492809	24636	528	855
16:30:37	479523	98.3	98.4	71.8	170.1	6680	12.4	493690	23755	525	1647
16:31:37	487253	63.7	96.0	86.4	150.1	5637	.0	492884	24561	528	1249
16:32:37	487249	413.8	99.2	414.6	828.4	1175	. 0	494534	22911	547	274
16:33:37	487249	50.6	94.7	88.3	138.9	5518	.0	494619	22826	527	306
16:34:37	487249	78.3	85.3	135.0	213.3	3608	.0	493812	23633	550	251
16:35:37	487249	48.9	98.1	61.6	110.4	7916		494649	22796	530	285
16:36:37	484434	32.6	98.0	120 6	39.6	59536 240c	32.1	493696	23749	551	999
16:37:37	484434	100 2	100.0	171 0	200.1	3436	51.0	435274	22171	530	660
16.30.37	404434	42.2	49.0	51 4	94 6	9425	00.1	493470	22967	570	575
16:40:37	484434	55.8	99.7	65.6	121.4	7383	6.2	493478	23967	531	258
16:41:37	484434	137.2	99.1	223.7	660.6	2166	91.4	493489	23956	534	270

Total paging into and out of XSTORE by time, shows spikes up to 971/s.

IBM Advanced	Technical	Sup	oort
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Monitoring	g Mer	nory	' - Exa	ampl	es		3 <i>A</i>	A. Pag	jing a	ctivity
FCX143	Data f	or 2003.	/12/02 I	nterval	16:30:3	37 - 17	:02:37	Мо	nitor	Scan
<pre>Interval End Time >>Mean>> 16:20:37 16:22:37 16:22:37 16:24:37 16:24:37 16:24:37 16:26:37 16:26:37 16:29:37 16:29:37 16:30:37 16:31:37 16:34:37 16:35:37 16:40:37 16:40:37 16:40:37 16:40:37 16:40:37 16:40:37</pre>	<pre>> Real > Non- > pgables > 233531 4 220947 4 222221 5 21580 4 220947 4 222940 4 222940 1 24035 5 22940 1 24074 8 244717 9 24636 0 23755 4 229511 9 22826 9 223749 4 222911 9 22826 9 223749 4 222917 9 223967 9 223967 8 223967 8 223958 8 239958 8 23958 8 23958 8 223023 6 225189 4 22051</pre>	Storid Resides Shages Shages S209955229 522955228 522955228 522955228 522955228 522955228 522955228 522955228 52200 55000 55000 55000 55000 55000 55000 55000000	e Hean Es Mean Es Avait Pa 1,485 364 1,485 1057 8760 115 8760 115 8760 115 8760 115 5300 107 5300 107 5300 107 8601 15 5300 107 8601 15 5300 107 8601 15 5300 107 8600 107 855 148 12494 11 3066 200 9999 135 6604 20 5758 685 2855 868 2769 275 8697 251 2697 251 269 11 368 269 11 269 126 269 11 269 126 269 11 269 126 269 126 269 11 269 126 269 11 269 126 269 11 269 11 269 126 269 11 269 1	-> < t. Rea/s fel 26.9 fel 26.9 21.3 59.3 21.4 50.3 21.4 50.3 21.4 50.3 21.4 50.3 21.4 50.3 21.4 50.3 2017 21.4 50.3 2017 21.4 50.3 2017 20	Write 19.7 .0 .0 .0 2.2 6.1 1.3 28.3 23.6 12.4 .1 .0 .0 .0 .0 .0 .0 .0 .0 .0 .0 .0 .0 .0	Pagin Total 46.99 91.34 77.48 10.77 531.32 10.64 329.06 322.17 10.77 531.32 10.77 531.32 10.60 322.11 140.17 140.14 17894.33 1994.33 1994.34 1994.33 1994.34 1994.3	g to D Sh 4.645543487682810519604880133 33334.6345519604880133 4.5556344333335463	g SDists 1.353322224274488430524436753 2.2224274488430524436753 1.353322224274488430524436753		> ds> Total 1.07 .35 .3322228 2.282744 1.88430524436753 1.3 5.9 .24436753

Total paging to DASD per second by time – shows spikes also. Mean Avail List is quite small at times (274*4096~1MB). Storage is constrained.

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Monitoring Memory - Examples 14. CP owned disks

FCX109 Data fo	or 2003/12/02 In	terval	16:45:3	37 - 10	5:46:31	7 Mo	onitor	Scan
Page / SPOOL Alloca PAGE slots availabl PAGE slot utilizati T-Disk cylinders av T-Disk space utiliz	ation Summary le 4831200 ion 59% vail zation%		SPOOL SPOOL DUMP S DUMP S	slots slot u slots a slot uf	avail: utiliz: availat tilizat	able ation ole tion	63	2880 15% 0
<pre></pre>	Area Area Type Extent PAGE 1-3338 PAGE 1-3338 SPOOL 1-3338 SPOOL 79- 256 PAGE 257- 390 PAGE 802880 PAGE 1-3338 PAGE 1-3338	Used 78 79 12 83 73 100 24 79 77	<pre>> User > User > Inter > feres 1 1 0 1 1 1 1 1 1</pre>	Queue Lngth 0 0 0 0 0 0 0 0 0 0	Serv Time /Page 1.7 1.2 .6 5.1 1.0 .7	MLOAD Resp Time 1.7 .9 1.2 .6 5.1 1.0 .7	Brock Page Size 3 4 9 4 5	Alloc 62 90 71 57

Since the system is paging, check paging efficiency. Block page size is small – the result of over allocated paging extents. A block page size greater than 10 is good. Add paging extents to correct this problem.

IBM Advanced	Technical	Support
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22. User paging load

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Monitoring Memory - Examples FCX113 Data for 2003/12/02 Interval 16:45:37 -16:46:37 Monitor Scan -- Paging Activity/s Data Numbe (Page Rate) <--Page Migration--> >2GB> X>MS MS>X X>DS Page Spaces. Reads Write Steals Userid Owned MS>X X>DS WSS Lockd Resrvd 423 Suste 6 7.4 . 1 5.6 7.4 1.8 37745 Θ LPWWPE02 2.0 2546 Θ 60.4 . 0 . 0 . 2 4.1 .0 129591 Θ . 1 LPWWPE04 .0 .0 . 2 3.8 .0 126167 2883 Θ 0 • Z . 0 .0 Θ ZDEVL007 Θ .6 . 0 1.0 66.6 .0 113718 9 . 0 . 0 . 0 2.8 19.4 LUWWPE01 Θ. .0 . 0 111966 9 Θ .0 Θ LPWWPE03 Θ. .0 .0 .0 .0 . 6 .0 111367 2274 . 0 . 3 34.3 14.8 Θ ZSITL001 Θ 3.1 14.7 . 0 102406 9 7.5 . 0 . 0 . 0 1181 ZDEVL984 -O .0 99803 Θ 0 . 0 50.1 50.1 L DWWAI01 Θ. .0 . 0 . 4 3.3 98736 9 Θ 5.8 ZDEVL017 .0 . 0 . 0 .0 . 1 . 0 95041 9 Θ υ . 0 138 31.0 LUWWPE02 Θ. . 0 .0 94253 9 Θ 8.0 .0 2.7 ZDEVL014 Θ .0 . 0 .0 .0 18.4 89763 9 Θ .0 . 2 LPWWPE01 Θ 1.0 .0 .0 . 2 . 2 .0 88841 3058 Θ . 0 . 0 . 0 ZDEVL006 Θ .0 . 0 . 0 . 1 83700 12 Θ . 0 . 0 Θ .0 . 0 .0 . 0 56792 9 Θ ZSITL002 . 0 LNXSYS01 Θ .0 .0 .0 .0 .0 . 0 . 0 29374 9 Θ . 1 ZSYSL002 Θ . 2 . 0 . 0 .0 . 0 . 0 24078 12 Θ . 0 . 0 . 0 Θ .0 . 0 . 0 . 0 21886 9 Θ ZDEVL015 . 0 . 0 ZDEVL008 Θ .0 . 0 . 0 . 0 . 0 20749 9 Θ . 6 . 0 Θ 2.2 . 0 .0 8.6 19961 9 Θ LIWWAI01 . 1

Since system is paging, check to see which guests are paging.

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23. User wait states

Monitoring Memory - Examples

FCX114 for 2003/12/02 Interval 16:45:37 -16:46:37 Monitor Scan Data <-SVM and-> %CPU Userid XACT ×RUN %LDG (PGW κιοω XSIM %CFW %EL %IOA %PGA ×ΤIW. XTI. %DM 67 Θ Θ >Sustem< Ð. Θ Θ Θ Θ Θ Θ Θ 96 Θ LPWWPE02 100 Θ Θ Θ 12 Θ Θ Θ Θ Θ Θ Θ Θ 88 DATAMOVE Θ . . . DIRMAINT Θ . . . EREP Θ Θ FTPSERVE LDSLDI01 100 Θ 0 0 Θ Θ Θ Θ Θ Θ. Θ. Θ 100 Θ Θ Θ Θ Θ Θ LDWWAI01 100 Θ Θ Θ Θ Θ Θ Θ 100 Θ Θ Θ 100 Θ 100 Θ Θ Θ Θ Θ Θ Θ Θ LIWWAI01 Θ Θ Θ Θ Θ Θ Θ Θ LNXSYS01 100 Θ Θ Θ. Θ 100 Θ Θ LPSLDI01 100 Θ Θ Θ Θ Θ Θ Θ Θ Θ 100 Θ LPSLDI02 100 Θ Θ Θ Θ Θ Θ Θ Θ Θ Θ. Θ 100 Θ Θ Θ Θ Θ Θ Θ Θ Θ Θ Θ Θ. Θ LPWWAE01 100 100 93 4 Θ Θ Θ Θ Θ Θ Θ Θ Θ 96 Θ LPWWPE01 0 Θ Θ Θ Θ LPWWPE03 100 Θ Θ Θ Θ Θ Θ. Θ 100 Θ 2 Θ Θ Θ Θ Θ Θ Θ Θ Θ Θ Θ LPWWPE04 100 98 ø 2 Θ Θ Θ Θ Θ LUSLDI01 100 Θ Θ Θ. 0 98 Θ Θ Θ Θ Θ Θ Θ Θ Θ Θ LUWWPE01 100 Θ Θ. Θ 100 LUWWPE02 100 7 Θ Θ Θ Θ Θ Θ Θ Θ Θ Θ 93 Θ MAINT Θ MONWRITE 7 Θ Θ Θ Θ 50 50 Θ Θ Θ Θ Θ Θ • OPERATOR Θ OPERSYMP 0 з 50 Θ Θ Θ Θ 97 TCP IP Θ 0 Θ Θ Θ Θ Θ

Since guests are paging and the block page size is small, check to see if guests are frequently in page wait.



IBM Advanced Technical Support

Resources That Affect Performance

Processor

Memory



4/25/2008

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Physical Devices

- Physical devices can be exploited by CP or supported for guest use, for example:
 - > CP supports allocation of system areas on 3390 DASD
 - > CP supports OSA devices for guest use only
- Devices supported for guest use only must be dedicated to a virtual machine
 - Devices exploited by CP may also be dedicated or attached to a service virtual machine
 - Dedicated devices may only be used by the virtual machine controlling the device



Virtual Devices

Virtualized devices

- > CP manages the underlying physical device
- > CP provides appearance that each virtual machine has the device
- > The underlying physical device must be present on the system
- > E.g. minidisks, crypto cards

Simulated devices

- > Complete representation of a physical device
- No physical device present
- E.g. 2540 card punch, Guest LAN, virtual disk
- Virtual devices are defined in the virtual machine directory entry, or dynamically created using CP commands



Controlling I/O

CP does not manage general I/O as a tunable resource

Limited tuning for real devices provided by:

- > CP SET THROTTLE command
- Slow I/O's per second from guest operating system to a particular real device
- Indirect control through SHARE setting, or other scheduler controls such as STORBUF, DSPBUF, LDUBUF
 - Virtual machines not in the dispatch list, cannot execute I/O instructions

I/O Priority Queuing

- > CP can effect queue placement for DASD devices
- > HW can effect priority in channel usage

Queued I/O Assist

Interpretive execution assist for QDIO devices

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Additional I/O Considerations

- Dedicated I/O is not eligible for Minidisk Cache (MDC)
- MDC read performance is as good as virtual disk performance
- Both virtual disks and MDC require sufficient real memory



Monitoring I/O

Privileged CP commands

- INDICATE I/O shows users in I/O wait state, and real device number to which most recent virtual I/O operation was mapped
- INDICATE USER userid EXP displays total number of virtual I/O operations started since logon, repeated displays provide a rough idea of I/O activity for particular virtual machine

CP Monitor

- Detailed I/O information available within monitor for each real device, including seeks information
- Total I/O information for individual virtual machines

Performance Toolkit for VM

> Detailed device I/O information and virtual machine I/O information

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1. CPU load and trans.

	FCX100	CPU 7	060 SER	10431 In	terval	11:52:19	- 11	:53:19 ZVN	1V4R40
	CPU Load PROC %CPU P00 13 P01 12	XCP XI 4 4	EMU %WT 9 87 9 88	XSYS XSP 1 0 1 0	×SIC × 98 98	LOGLD 13 13	Vecto %VTOT not not	r Facility %VEMU REST installed installed	Status or ded. User Master Alternate
<	Total SSCH/F Virtual I/O Total rel. S	RSCH rate HARE	809/s 808/s 3200	Page r XSTORE Tot. al	ate paging bs SHAR	.0/s .0/s E 0%	<	<u>Priv. instruc</u> Diagnose inst	st. 246 <u>1/s</u> tr. 25/s
	Queue Statis VMDBKs in qu VMDBKs loadi Eligible VMD El. VMDBKs l Tot. WS (pag Expansion fa 85% elapsed	tics: Jeue Ing DBKs .oadin Jes) actor time	9 12493 6.864	Q1 0 0 0 1 .858 6.1	Q2 0 0 0 0 0 864 41	Q3 0 0 0 1 .18	User # of # of # of % in- % eli	Status: logged on use dialled users active users in-queue user Q users in PC Q users in IC g. (resource	ers 14 5 0 7 5 3WAIT 0 3WAIT 35 Wait) 0
	Transactions Average user Trans. per s Av. time (se	s sec. sc)	Q-Disp .4 .4 1.261	trivial .0 .2 .002	non-	tru .0 .0	User Max. Max. Max.	Extremes: CPU % WOJL1 VECT % IO/sec WOJL1	(NUX 22.6

System shows I/O activity (virtual I/O almost equal to SSCH/RSCH rate). Also note that the privileged instruction rate is high. This is consistent with virtual I/O.

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13. I/O device load

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<pre></pre>	Qued	
→ AUL DHSD XX3 .0 .0 8.4 2.6 11.0 11.0 .0	.00	
A00C 3390-3 WOJ002 2 1 .2 .0 .0 1.0 .1 1.1 1.1 .0	.00	
A005 3390-3 440W02 CP 8 1 .2 .0 .0 .9 .1 1.0 1.0 .0	.00	
0124 3380 OS39H7 0 1 .0 .0	.00	
0125 3380 HFSUS1 0 1 .0 .0	.00	
0A80 3390-3 0839R7 0 1 .0 .0	.00	
0A82 3390-2 0S3R7A 0 1 .0 .0	.00	
UA85 3390-1 M3KPLX U 1 .0 .0	.00	
	.00	
$4341 3390^{-3} = 1001 0 0 2 0 0 0 0 0 0 0 0 0 0 0 0 0 0 $.00	
4343 3390-3 VMLTN2 0 2 10 10 11 10 13 14 14 16		
4344 3390-3 DEBIAN 0 2 .0 .0 .1 .0 .3 .4 .4 .6		
A003 3390-3 440RES CP 132 1 .0 .0 .1 .1 .0 .2 .2 .0	.00	
A004 3390-3 440W01 CP 47 1 .0 .0 .0 .0 .0 .0 .0 .0 .0	.00	
A009 3390-3 PAG001 CP 0 1 .0 .0 .1 .0 .0 .1 .1 .0	.00	
A00A 3390-3 BASRES 0 1 .0 .0 .0 .1 .0 .1 .1 .0	.00	
A00D 3390-3 440U01 3 1 .0 .0 .2 .1 .0 .3 .3 .6	.00	
A00E 3390-3 BASW01 0 1 .0 .0 .0 .0 .0 .0 .0 .0	.00	
A00F 3390-3 M2K353 0 1 .0 .0 .1 .1 .0 .2 .2 .6	.00	
H010 3390-3 BHSW0Z 0 1 .0 .0 .0 .0 .0 .0 .0 .0 .0	.00	
	.00	

Since I/O is occurring, look to see which devices are affected and whether response time is poor. Note the device address and number of paths for next screen.

4/25/2008

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19. I/O configuration

FCX131	CPU 7060	SER 10431	Status 11:53:19	EVMV4R40
<pre>< Ran Device-No 000E-000F 0124-0125 0700-0702 0900-0901 0A80 0A82 0A85 0C20-0C23 0C24-0C27 0C24-0C27 0C28-0C2B 0C2C-0C2F 0E60 0E61 0F40-0F41 1200-1207 1300-1307 1400-1407 1500-1507 4340-4344</pre>	ges> SubchID 0188-0189 018C-0180 01C6-01C8 01CB-01CC 01CF 01D2 01D7-01DA 01D8-01DE 01DF-01E2 01E7 01E8 0207-0208 0208-0212 0213-021A 0223-022A 0238-0242 0283-0287	Device Type 1403 3380 (E) 3270-2 3270-2 3390-3 (E) 3390-2 (E) 3390-1 (E) 3490 3490 3490 3490 3490 3490 3490 3490	- Channel Path Ids 1 2 3 4 5 6 7 C C C C C C C	-> Control 8 Unit Status 2821-01 Online 3880-23 Online 3274-1D Online 3274-1D Online 3990 Online 3990 Online 3990 Online 3490-20 Offline 3490-20 Offline 3490-20 Offline 3490-20 Offline 3490-20 Offline 3490-20 Offline 3490-20 Offline 3490-20 Offline 3490-20 Online 3490-20 Offline 3490-20 Online 3088 Online 3088 Online 3088 Online
A009-A017 A018	02C9-02D7 02D8	3390-3 (E) 3390-1 (E)		. 3990 Online . 3990 Online

Look to see what chpids are associated with device A00B.



11. Channel load

FCX107	CPU 7060	SER 10431	Interval	11:29:1	9 - 11:	53:19	ZVMV4R40	
CHPID Chan (Hex) Desci 04 ESCOI 05 ESCOI 06 ESCOI 07 ESCOI 08 ESCOI 08 ESCOI 08 ESCOI 08 ESCOI 08 ESCOI 08 ESCOI 00 ESCOI 00 ESCOI 00 ESCOI 10 ESCOI 10 ESCOI 10 ESCOI FC FICOI FD ESCOI	- Group - Qual - Qual - Qual - 00 - 0	<%Busy> Cur Ave 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0	<pre>< CH 0-10 11-2 100 100 100 100 100 100 100 100 100 10</pre>	annel %B 0 21-30 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0	usy Dis 31-40 4 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0	tributio 1-50 51- 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0	n 11:29:19-1 60 61-70 71- 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0	1:5000000000000000000000000000000000000
Chpid FD channels o	shows all on 7060).	'', which	i means u	tilization	ı data is	s not av	ailable (inte	rna

IEM					
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13. I/O device load (select device)

FCX110 CPU 7	060 SER 10431	Interval 1	1:52:03 - 11:	54:30 ZVMV4F	840
Detailed Analysis Device type : 33 VOLSER : WO Nr. of LINKs: Last SEEK : SSCH rate/s : Avoided/s : Status: MDCACHE U	for Device A00 90-3 Funct J001 Disco 2 Conne 1224 Servi 7.9 Respo .0 CU qu SED	B (SYSTEM ion pend.: nnected : cted : ce time : nse time : eue time :) 9.5ms 3.0ms 12.5ms 12.5ms 12.5ms .0ms	Device busy I/O contention Reserved SENSE SSCH Recovery SSCH Throttle del/s	9% 0% 0% 0
Path(s) to device Channel path stat	A00B: FD us : ON				
Device DIR ADDR VOLSER	Overall CU-Cac IO/S %READ %RD	he Performa HIT %WRHIT	nce ICL/S BYP/S	Split IO/S %READ %RD	ніт
MDISK Extent	Userid Addr	IO/s VSEEK	Status LIN	IK VIOZS %MDC N	1DIO/s
C <u>1 - 50</u> C 51 - 3338	WOJLINUX 0191 WOJLINUX 0200	7.9 1173	WR WR	<u>1 .0</u> 1 7.9 0	.0 C 7.9 C

Detail device screen shows virtual machine doing I/O to a minidisk (0200) on physical device A00B.

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23. User wait states

FCX114	CPU	7060	SER :	10431	Inte	enval	11:53	3:19 -	11:5	54:1	9	ZVM	/4R40	
1 - C	1.1	1.1					1.1	1.1	1.1	<- \$1	VM år	hd-X	1.1	
Userid	XACT	%RUN	%CPU	%LDG	%PGW	×10W	%SIM	$\times T I W$	%CFW	XTI	%EL	%DM	%I0A	%PGA
>System<	24	5	2	0	0	0	•	45		- 0	- 0	18	28	0
PERFSVM	63	0	0	0	0	3	0	0	0	- 0	- 0	97	0	0
DISKACNT	0													1.1.1.1
EREP										1.1.1.1	1.1.1.1	1.1.1		
FTPSERVE														1.1.1.1.1
MAINT	10	0	0	0	0	0	0	100	0	0	0	0	0	0
OPERATOR	0													
OPERSYMP														
TCPIP	67	0	0	0	0	0	0	0	0	0	0	0	100	0
VMRTM	0		1.1.1.1								1.1.1	1.1.1		1.1.1.1
VMSERVR														
VMSERVS	0													
VMSERVU							1							
WOJLINUX	100	18	3	0	0	0	. 0	47		0	0	0	30	0
WOJLIN02	100	0	5	0	0	0	0	95	0	0	0	0	0	0

Check to see if the virtual machine performing the I/O, is consistently waiting on I/O.

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Summary

Primary resources that can affect z/VM guest performance are:

- Processor
- Memory
- Paging I/O

z/VM controls for managed resources include:

- SET SRM STORBUF memory
- SET SRM LDUBUF paging I/O
- SET SHARE processor

Monitoring z/VM performance

- CP commands watch out for smoothed averages
- Performance Toolkit for VM
- > Other performance software
- Start monitoring now!