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# VSE/ESA 2.6 and 2.7

## Performance Considerations

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VSE/ESA Development





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## Contents

- VSE/ESA 2.6
  - ▶ Performance Items (Overview)
  - ▶ Hardware Support
  - ▶ Virtual Tape Support
  - ▶ e-Business Enhancements
- VSE/ESA 2.7
  - ▶ Performance Items (Overview)
  - ▶ HiperSockets
  - ▶ Hardware Crypto support
- zSeries Remarks
- Turbo Dispatcher considerations
- z/VSE 3.1 Preview
- Dependencies for VSE/ESA Growth



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## VSE/ESA 2.6 Performance Items

- VSE/ESA 2.6 Base enhancements
  - ▶ Delete Label Function
  - ▶ LTA Offload for some AR commands
  - ▶ SVA-24 Phases moved above the line
  - ▶ Increased max number of SDL entries
  - ▶ SDL update from non-BG partitions
  - ▶ POWER Data file extension without reformat



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## VSE/ESA 2.6 Performance Items - continued

- VSE/ESA 2.6 Hardware Support
  - ▶ FICON Support (VSE/ESA 2.3 or higher)
  - ▶ New 2074 System Management Console
  - ▶ OSA Express Adapter (e.g. Gigabit Ethernet)
  - ▶ VSAM Support for large 3390-9 Disks (Shark)
  - ▶ Fastcopy Exploitation of ESS FlashCopy and RVA SnapShot



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## VSE/ESA 2.6 Performance Items - continued

- VSE/ESA 2.6 e-Business Enhancements
  - ▶ Updated Java-based connectors
  - ▶ VSAM SHROPT(4) avoidance for connectors
  - ▶ SSL for VSE/ESA exploitation
  - ▶ SSL enabled CICS Web Support
  - ▶ CICS External Call Interface
  - ▶ New VSAM Redirector
  - ▶ More samples (JConVSE, VSEPrint, etc.)
  - ▶ New JDBC Driver Layer for VSAM



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## Delete Label Function

- New function **DELLBL** in LABEL macro
- Must be explicitly exploited
  - ▶ Important for vendors with disk/tape management products
- Benefits
  - ▶ Saves recursive reads (GETNXGL) and write backs (ADDLBL/ADDNXL)
  - ▶ **Saves >90% of the SVCs** for this activity
- More Info
  - ▶ VSE Label Area -Layout and Capacity Consideratio  
VSE/ESA Software Newsletter, 12/2000
  - ▶ <http://www-1.ibm.com/servers/eserver/zseries/os/vse/pdf/vsenew21/vseflab.pdf>





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## LTA Offload and SVA-24

- LTA Offload for some AR commands
  - ▶ Phases \$\$BATTNC and \$\$BATTNG are merged into \$\$BATTNA
  - ▶ Code of \$\$BATTNB is merged into IJBAR
  - ▶ Benefits
    - Less I/O by less FETCHes for LTA load
      - IGNORE, PAUSE, LOG, NOLOG, NEWVOL, START, BATCH
    - No LTA usage for MSG commands
- SVA-24 Phases moved above the line
  - ▶ \$IJBPRTY (6K)





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## SDL Entries

- Increased max. number of SDL entries
  - ▶ New IPL SVA parameter:  $SDL=n$ 
    - Maximum value now **32765**
    - About 56 SDL entries per 4K page in shared space below
    - Theoretically would cost 2.28 MB
- SDL updates from non-BG partitions
  - ▶ SET SDL command can now be issued from any partition
  - ▶ Internal locking is done to assure correctness



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## Hardware Support

- New 2074 System Management Console
  - ▶ ESCON channel attached
  - ▶ Eliminates requirement for a non-SNA 3174 controller
- OSA Express Adapter Support
  - ▶ Available for G5 and above

	<b>Gigabit Ethernet</b>	<b>Fast Ethernet 100 Mbps</b>	<b>ATM-LE 155 Mbps</b>	<b>Tokenring 4/16/100 Mbps</b>
CHIPID TYPE=OSE (non-QDIO)	no	yes	yes	yes
CHIPID TYPE=OSD (QDIO)	yes	yes	yes	yes



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## Hardware Support

- Queued Direct I/O
  - ▶ Designed for very efficient exchange of data
  - ▶ Uses the QDIO Hardware Facility, without traditional S/390 I/O instructions
  - ▶ Without interrupts (in general)
  - ▶ Use of internal queues
  - ▶ With pre-defined buffers in memory for asynchronous use
  
- Exploitation by TCP/IP for VSE/ESA
  - ▶ see TCP/IP Performance Considerations



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## ESS Flashcopy

- The DASD Architecture of ESS allow copy of DASD's with the utility FlashCopy
  - ▶ The copy process takes a few seconds instead of hours
  - ▶ From Operating system view it is a real copy
  - ▶ From DASD controller view it is a virtual copy
  
- FlashCopy support is available for 3 VSE products
  - ▶ IXFP SNAP command
  - ▶ VSAM SYNONYM Backup
  - ▶ VSE/Fast Copy



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## ESS Flashcopy - continued

- Problems
  - ▶ Duplicate VOLIDs (DASD names) not allowed on a VSE system
  - ▶ Duplicate VSAM Catalog names not allowed on a VSE system
- FlashCopy of volume containing VSAM datasets would mean
  - ▶ duplicate VOLIDs
  - ▶ duplicate VSAM Catalog names
- Solution
  - ▶ IDCAMS SNAP command
    - Changes the VOLIDs of the copied volumes
  - ▶ IDCAMS SYNONYM BACKUP command
    - Uses a synonym list to access copied volumes



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## Virtual Tape Support

- Allows access to tape images residing in
  - ▶ A VSAM file (ESDS)
  - ▶ A remote file on a workstation or server
- Tape image is AWSTAPE format
  - ▶ Known from P/390, R/390 or FLEX/ES
- New VTAPE command
- Virtual Tape Simulator
  - ▶ Simulates channel program execution
- Virtual Tape Data Handler
  - ▶ Runs in a partition



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## Virtual Tape Support - continued

- Virtual Tape Server
  - ▶ Runs on a workstation or server (Java)
  - ▶ Allows to access a tape image remotely
  - ▶ Communicates via TCP/IP with Virtual Tape Data Handler
- Designed to allow e-Delivery and e-Service (future)
  - ▶ Download a tape image containing a product
  - ▶ Obtain a CD/DVD containing the tape image
  - ▶ Install the product via Virtual Tape directly from the workstation
- Also possible
  - ▶ Backup to a Virtual Tape + copy to CD
  - ▶ Restore directly from CD via Virtual Tape

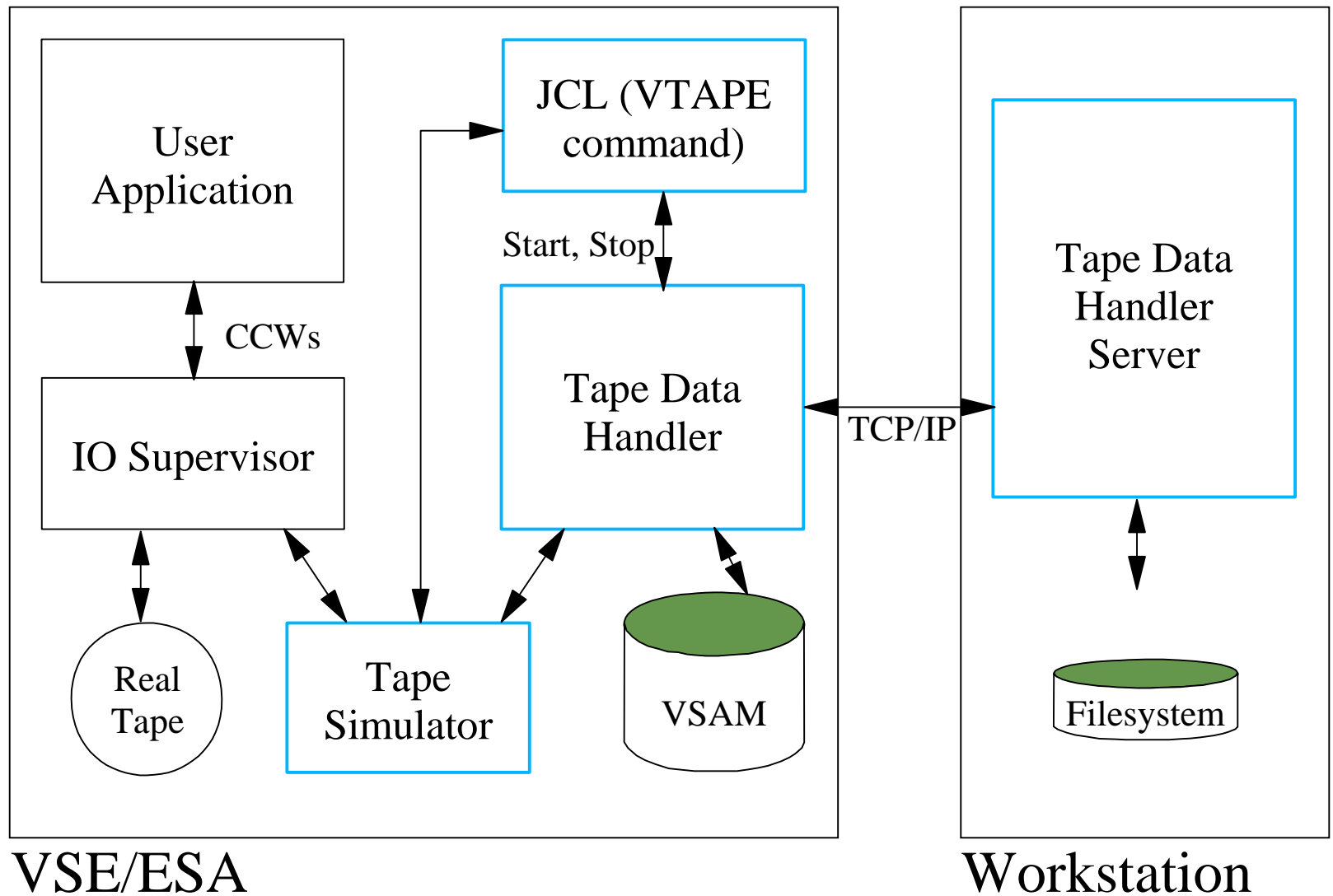




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## Virtual Tape Support - continued





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## Updated Java-based Connector

- The Java-based connector has been updated to support the Java 2 platform (JDK 1.3)
- Introduced JDBC layer for VSAM access
  - ▶ Allows to issue SQL statements
- Adaptations for WebSphere 4.0
  - ▶ Enhanced connection pooling by support of JCA (Java Connector Architecture)
  - ▶ Connectors can be deployed as Resource Adapter and as (JDBC-) Data Source
- SSL enabled connections possible
  - ▶ Transparent use of secured connections



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## VSAM Share Options with Connectors

- SHROPT(4) Backgrounds
  - ▶ Using connectors to UPDATE a VSAM file already opened for output (e.g. by CICS) needs SHROPT(4)
  - ▶ SHROPT(4) has big overhead
- Performance implications
  - ▶ Bigger pathlength for processing of UPDATE requests due to VSAM internal locking
  - ▶ Each READ must be done from disk
  - ▶ Each WRITE must go to disk
  - ▶ Additional catalog I/Os for statistics
  - ▶ Influence on any application, not only connectors



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## VSAM SHROPT(4) Avoidance

- Connectors in VSE/ESA 2.5 require SHROPT(4) when updating VSAM files owned by CICS
- New VSAM-via-CICS Service avoids SHROPT(4) by routing the VSAM requests to CICS
- Communication between batch and CICS is XPCC
- New transactions related to VSAM-via-CICS:

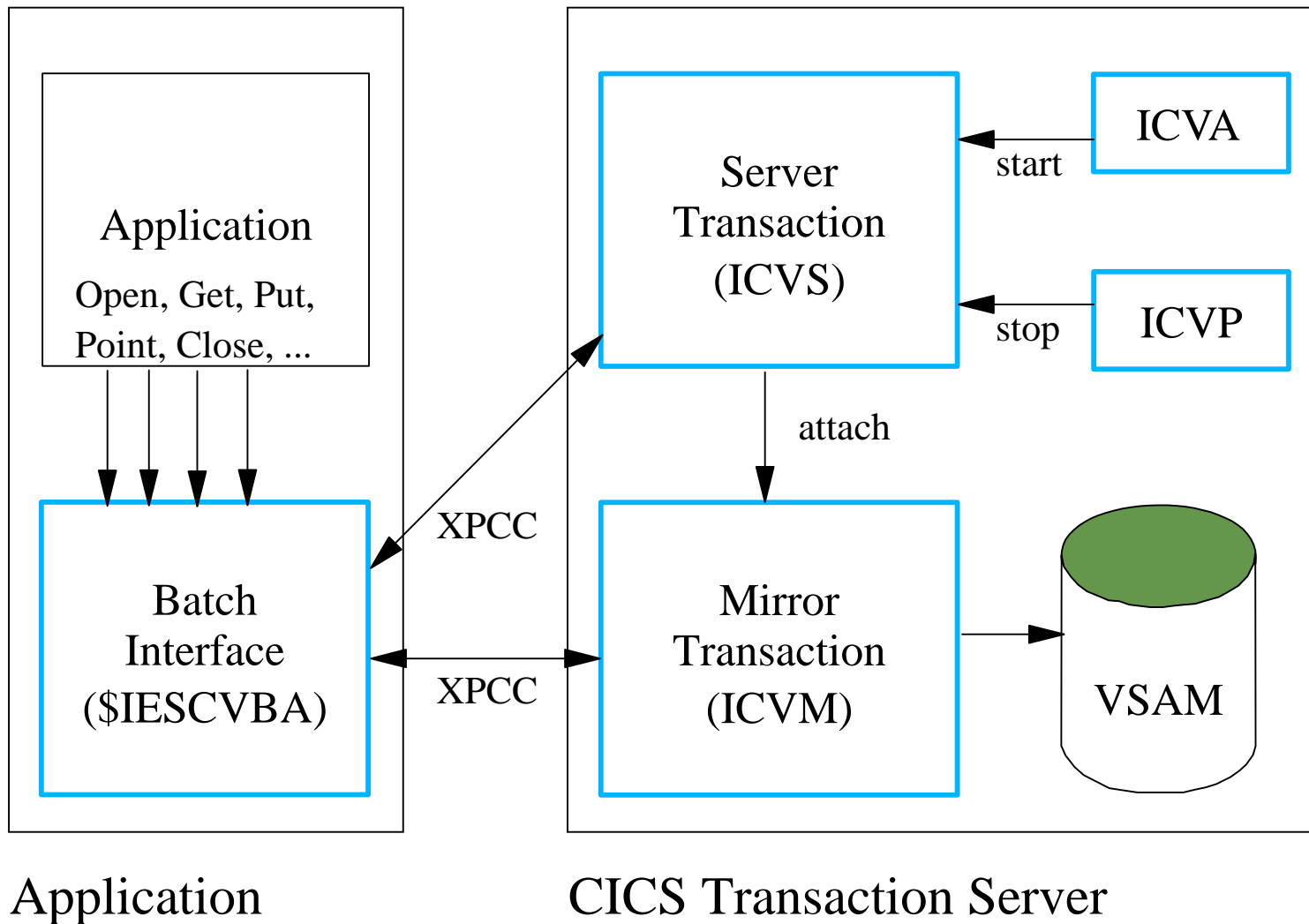
Transaction	Program	Description
ICVA	IESCVSTA	starts the service
ICVP	IESCVSTP	stops the service
ICVS	IESCVSRV	internal server task
ICVM	IESCVMIR	internal mirror task
none	IESCVSTI	internal start program



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# VSAM-via-CICS Service





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## VSAM-via-CICS Service - continued

- How VSAM-via-CICS works
  - ▶ Long running server transaction ICVS
  - ▶ Attaches a mirror transaction ICVM on request
  - ▶ Mirror transaction is attached for
    - "Open" from batch
    - Browse files from batch
  - ▶ Mirror transaction ends at "close" from batch
  - ▶ Service can run in multiple CICSes at the same time
  - ▶ Batch counterpart is implemented in phase \$IESCVBA



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## VSAM-via-CICS Service - continued

- Naming convention for "VSAM-via-CICS files"
  - ▶ Each CICS is treated as "virtual" catalog
  - ▶ Files defined in CICS (via CEDA DEFINE FILE) are visible within this catalog
  - ▶ "Virtual" catalog file id

**#VSAM.#CICS.<applid>**

indicates "virtual"  
CICS catalog

APPLID of CICS region  
owning the files within this  
catalog

- "Virtual" cluster file id is the 7 character name known in CICS





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## VSAM-via-CICS Service - continued

- Example
  - ▶ Assume there is a CICS region DBDCCICS
  - ▶ CICS knows a file named MYFILE
  - ▶ Real VSAM files MY.VSAM.TEST.FILE resides in catalog MY.USER.CATALOG
  - ▶ "Batch only" name would be
    - **Catalog: MY.USER.CATALOG**
    - **Cluster: MY.VSAM.TEST.FILE**
  - ▶ "VSAM-via-CICS" name would be
    - **Catalog: #VSAM.#CICS.DBDCCICS**
    - **Cluster: MYFILE**



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## VSAM-via-CICS Service - continued

- VSAM-via-CICS files can only be accessed from the following applications
  - ▶ Java-based connector via VSE Java Beans
  - ▶ DB2-based connector via VSAM CLI (SQL)
  - ▶ REXX new VSAMIO function
- IDCAMS does NOT show these files
- "Virtual" names can NOT be specified in DLBLs
- No changes made in VSAM for this support
- No influence on "normal" VSAM processing
- But: Maps can be defined for a "virtual" file
  - ▶ Via Java-Based connectors
  - ▶ Via IDCAMS RECMAP function



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## VSAM Redirector

- New connector
  - ▶ VSE is client
  - ▶ PC / workstation is server
- Exploits VSAM exit IKQVEX01
- Allows to redirect one or more VSAM files to a PC workstation
- All VSAM requests of a particular file are redirected
  - ▶ Open / close
  - ▶ Get / put / point / delete / insert
- Transparent for applications
  - ▶ Usable from batch and CICS



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## VSAM Redirector - continued

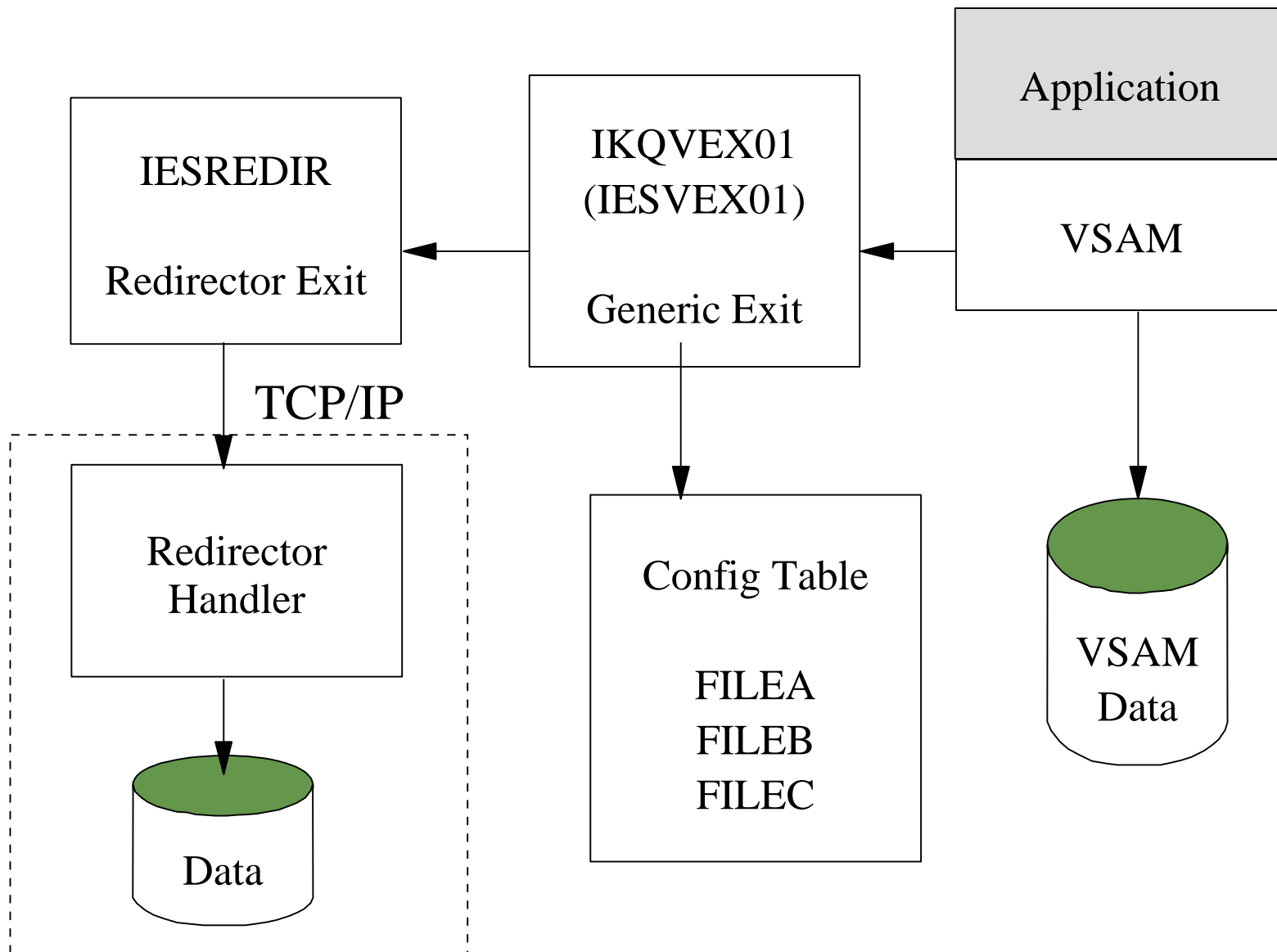
- Owner of data can be
  - ▶ VSAM
    - Requests are forwarded to workstation
    - VSAM still owns the data
    - VSAM executes the requests
    - Used for data replication/synchronisation
  - ▶ PC / workstations
    - VSAM does not execute the requests
    - Handler on workstation 'simulates' VSAM logic
    - A VSAM file with at least one dummy record is required (for open processing)



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## VSAM Redirector - continued





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## VSAM Redirector - continued

- Decision if a file is redirected or not is
  - ▶ Done at open time
  - ▶ Based on the config table (PHASE)
    - Catalog id and file id
  - ▶ Only a very small (open-)overhead for non-redirected files
  - ▶ No overhead for get/put/... if not redirected
- Generic Exit can also call a 'Vendor' exit instead of Redirector exit
  - ▶ Defined in the config table
  - ▶ Based on catalog id and file id



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## VSAM Redirector - Performance Implications

- Is the file redirected ?
  - ▶ No: only at OPEN time (very small overhead)
  - ▶ Yes: at each request
- Network overhead ?
  - ▶ Yes, if file is redirected
  - ▶ Depends on
    - Number of VSAM requests
    - Size of records
- Data ownership
  - ▶ OWNER=REDIR
    - no VSAM I/O





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## CICS TS Enhancements

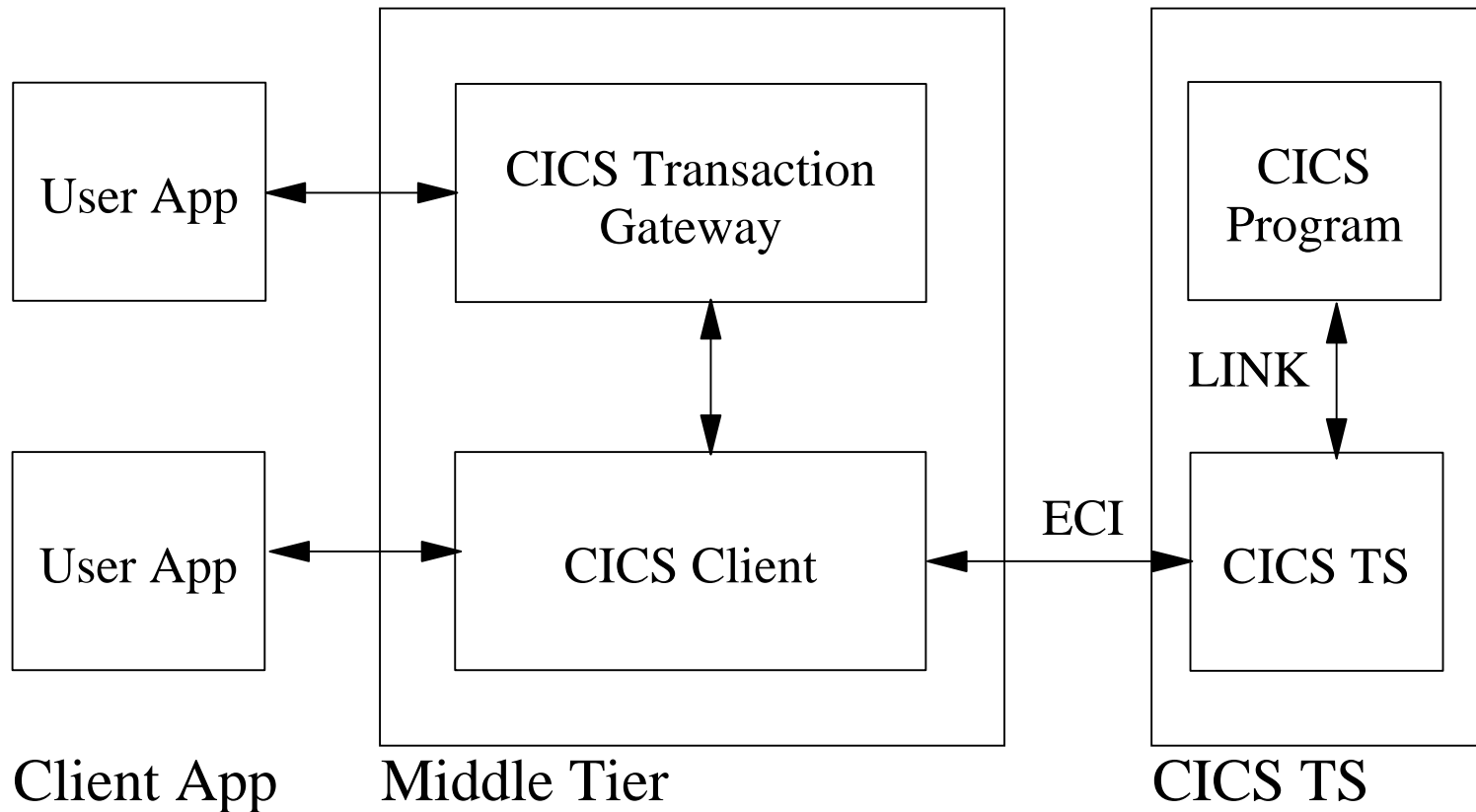
- CICS Web Support
  - ▶ new: SSL enabled (https)
  
- External Call Interface (ECI)
  - ▶ Call a CICS program from a workstation
  - ▶ Prerequisites
    - CICS Client
    - CICS Transaction Gateway



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# External Call Interface (ECI)





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## General Performance Hints for Connectors

- Reduce amount of data transferred
  - ▶ Transfer only data that is needed
  - ▶ Issue only requests that are needed
- Use connection pooling
  - ▶ Reduce overhead of connection establishment
- Performance of connectors depends on
  - ▶ Network performance
  - ▶ Performance of "server"
  - ▶ Performance of "client" or middle tier



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## VSE/ESA 2.7 Performance Items

- VSE/ESA 2.7 hardware support
  - ▶ z800/z900, Multiprise 3000, G5/G6
  - ▶ HiperSockets
  - ▶ Hardware Crypto Support
  - ▶ 32760 cylinder 3390 support
  - ▶ 3590 buffered tape mark
  
- VSE/ESA 2.7 enhancements
  - ▶ New TCP/IP for VSE/ESA release 1.5
  - ▶ \$IJBLBR above the line
  - ▶ II User Status Record above the line
  - ▶ VTAPE: removed DVCDN/DVCUP
  - ▶ POWER: reallocate queue file during warm start



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## VSE/ESA 2.7 Hardware support

- VSE/ESA 2.7 runs on the following machines
- zSeries: z800, z900, z990, z890
- 9672 Parallel Enterprise Server (G5/G6)
- Multiprise 3000 (7060)
- equivalent emulators (Flex-ES)
- VSE/ESA 2.7 is based on the hardware instruction set described in the manual 'ESA/390 Principles of Operation' (SA22-7201).
- With VSE/ESA 2.7 it is assumed that all the ESA/390 instructions and facilities described in that manual can be used.



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## zSeries Remarks

- Prior to zSeries there is one cache for data and instructions
- zSeries has splited data and instruction cache
- Performance implications:
  - ▶ If program variables and code that updates these program variables are in the same cache line (256 byte)
    - Update of program variable invalidates instruction cache
    - Performance decrease if update is done in a loop
  - ▶ See APAR PQ66981 for FORTRAN compiler



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## Supported VSE Releases

- **VSE/ESA 2.4/2.3:** already out of service
  - ▶ runs also on zSeries (z800, z900)
  - ▶ does not run on z990, z890 (Hardwait during IPL)
- **VSE/ESA 2.5:** end of service 12/31/2003
  - ▶ runs also on zSeries (z800, z900)
  - ▶ runs also on z990 with additional PTF
- **VSE/ESA 2.6**
  - ▶ runs also on zSeries (z800, z900)
  - ▶ runs also on z990, z890 with additional PTF
- **VSE/ESA 2.7**
  - ▶ runs on zSeries (z800, z900, z990, z890, G5/G6, MP3000)
- **OSA Express:** Supported with VSE/ESA 2.6 and 2.7
- **HiperSockets and PCICA (Crypto)**
  - ▶ Supported with VSE/ESA 2.7





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## 32760 cylinder 3390 support

- With announcement 101-341 at 11/13/2001 IBM announced the new 32760 cylinder 3390 volumes of the IBM TotalStorage Enterprise Storage Server (ESS)
  - ▶ This enhancement of the ESS F models was made available 11/30/2001
- VSE/ESA 2.7 now supports these volumes
  - ▶ helps relieve address constraints
  - ▶ improves the disk resource utilization
  - ▶ can be used to consolidate multiple disk volumes into a single address
- VSAM can only address 10017 cylinders.



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## 3590 Buffered Tape Mark support

- The 3590 control unit provides support for writing tape marks (TM) in buffered mode
- Writing TM's in "buffered" mode should enhance the performance
  - ▶ of all programs which write many TM's as part of their file creation process (e.g. POFFLOAD)
- All the TM's written during OPEN/CLOSE (label processing) will remain to be written "UNbuffered"
  - ▶ all the programs which write TM's mainly or only during OPEN/CLOSE will NOT benefit from this enhancement



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## **\$IJBLBR phase moved above the line**

- The \$IJBLBR.PHASE has been split into two phases
  - ▶ \$IJBLBR.PHASE
  - ▶ \$IJBLB31.PHASE
- \$IJBLBR.PHASE will continue to reside in SVA-24
- \$IJBLB31.PHASE will reside in SVA-ANY (high SVA)
  - ▶ This will free about 180KB in SVA-24



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## II User status record above the line

- During Logon each II user gets besides others two storage areas allocated
  - ▶ User\_Status\_Record USR (904 bytes)
  - ▶ Panel\_Hierarchy\_List PHL (1352 bytes)
  - ▶ originally located in the CICS DSA (below)
- With VSE/ESA 2.7 the USR and PHL has been moved to ESDSA (shared above)
  - ▶ frees 2.3 KB in DSA below per user
- ICCF TCTUALOC=ANY now supported
  - ▶ ICCF transaction programs has been changed to support a TCTUA (28 bytes) above the line



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## HiperSockets hardware elements (**'Network in a box'**)

- Synchronous data movement between LPARs and virtual servers within a zSeries server
  - ▶ Provides up to 4 "internal LANs" HiperSockets accessible by all LPARs and virtual servers
  - ▶ Up to 1024 devices across all 4 HiperSockets
  - ▶ Up to 4000 IP addresses
  - ▶ Similar to cross-address-space memory move using memory bus
- Extends OSA-Express QDIO support
  - ▶ LAN media and IP layer functionality (internal QDIO = iQDIO)
  - ▶ Enhanced Signal Adapter (SIGA) instruction
    - No use of System Assist Processor (SAP)



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## HiperSockets hardware elements (**'Network in a box'**) - continued

- HiperSockets hardware I/O configuration with new CHPID type = IQD
  - ▶ Controlled like regular CHPID
  - ▶ Each CHPID has configurable Maximum Frame Size
- Works with both standard and IFL CPs
- No physical media constraint, no physical cabling, no priority queuing
- Secure connections



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## Measurement Environment

- z800 (2066-004)
  - ▶ 4 processors
- VSE/ESA 2.7 GA Driver in an LPAR (native)
  - ▶ 1 CPU active (~2066-001)
  - ▶ TCPIP00 (F7): OSA Express Fast Ethernet
  - ▶ TCPIP01 (F8): HiperSockets
- Linux for zSeries in an LPAR (native)
  - ▶ 3 CPUs active (shared)
  - ▶ eth0: OSA Express Fast Ethernet
  - ▶ hsi10: HiperSockets





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## Latency (Round trip time) - results

- Measurements has been done with PING command
  - ▶ Issued at Linux side
  - ▶ 10 Pings
  - ▶ PING sends a datagram to VSE
  - ▶ VSE sends an answer back to Linux
  - ▶ Time until answer arrives is measured
    - Round trip time



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## Latency (Round trip time) - results



HiperSockets is about 1.8 times faster in terms of latency



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## Throughput (MB/sec)

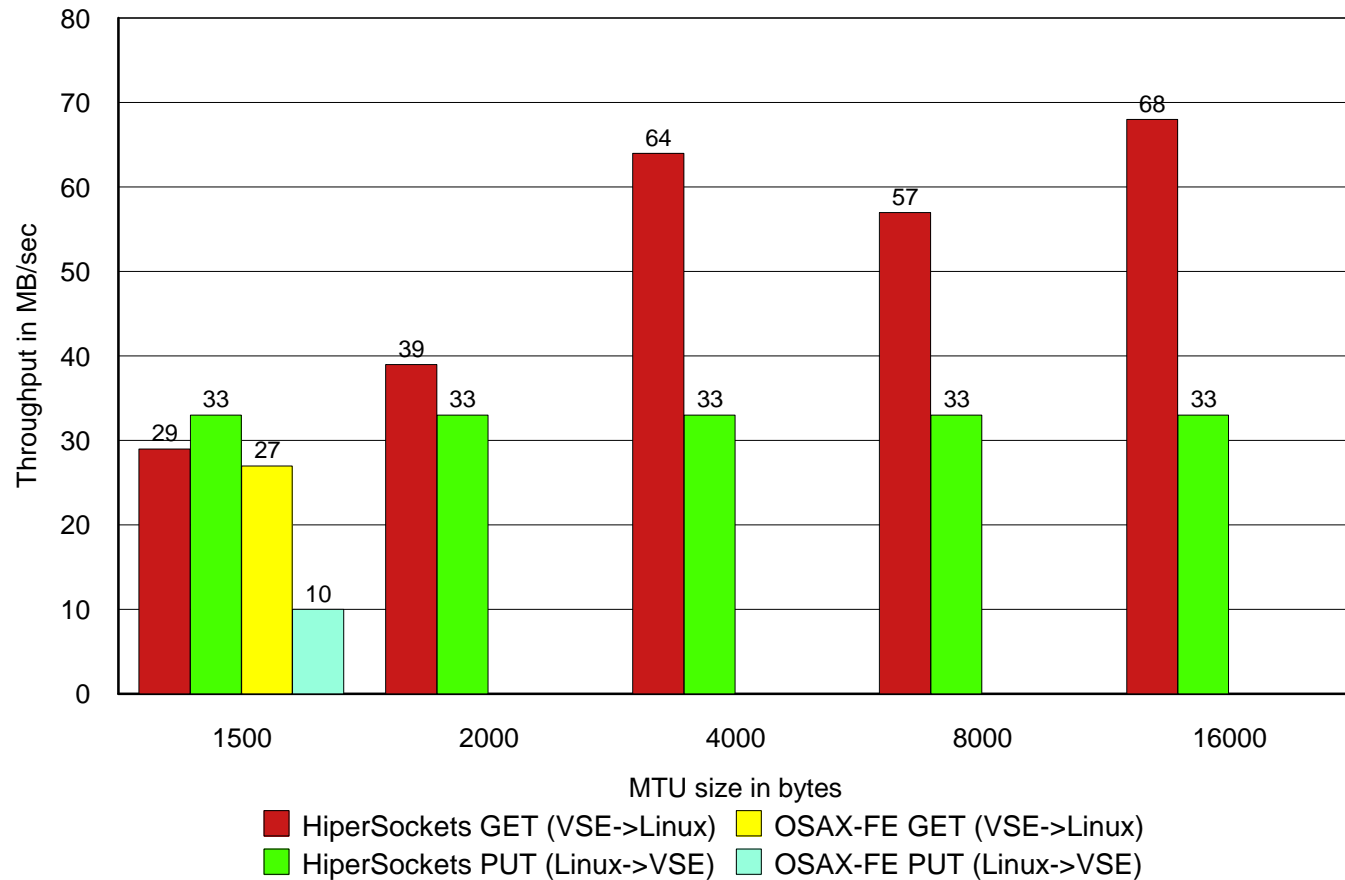
- Measurements has been done with FTP
  - ▶ Initiated at the Linux side
  - ▶ Transferring 1GB (1000MB)
    - without translation (binary)
    - 1 to 5 parallel streams
  - ▶ PUT: send data to VSE
    - VSE inbound
    - sending a 1GB file to \$NULL file (in memory file)
    - No file I/O is done by VSE/Linux
  - ▶ GET: receive data from VSE
    - VSE outbound
    - receiving \$NULL file (in memory file) into /dev/null
    - No file I/O is done by VSE/Linux



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## Throughput (MB/sec) - results



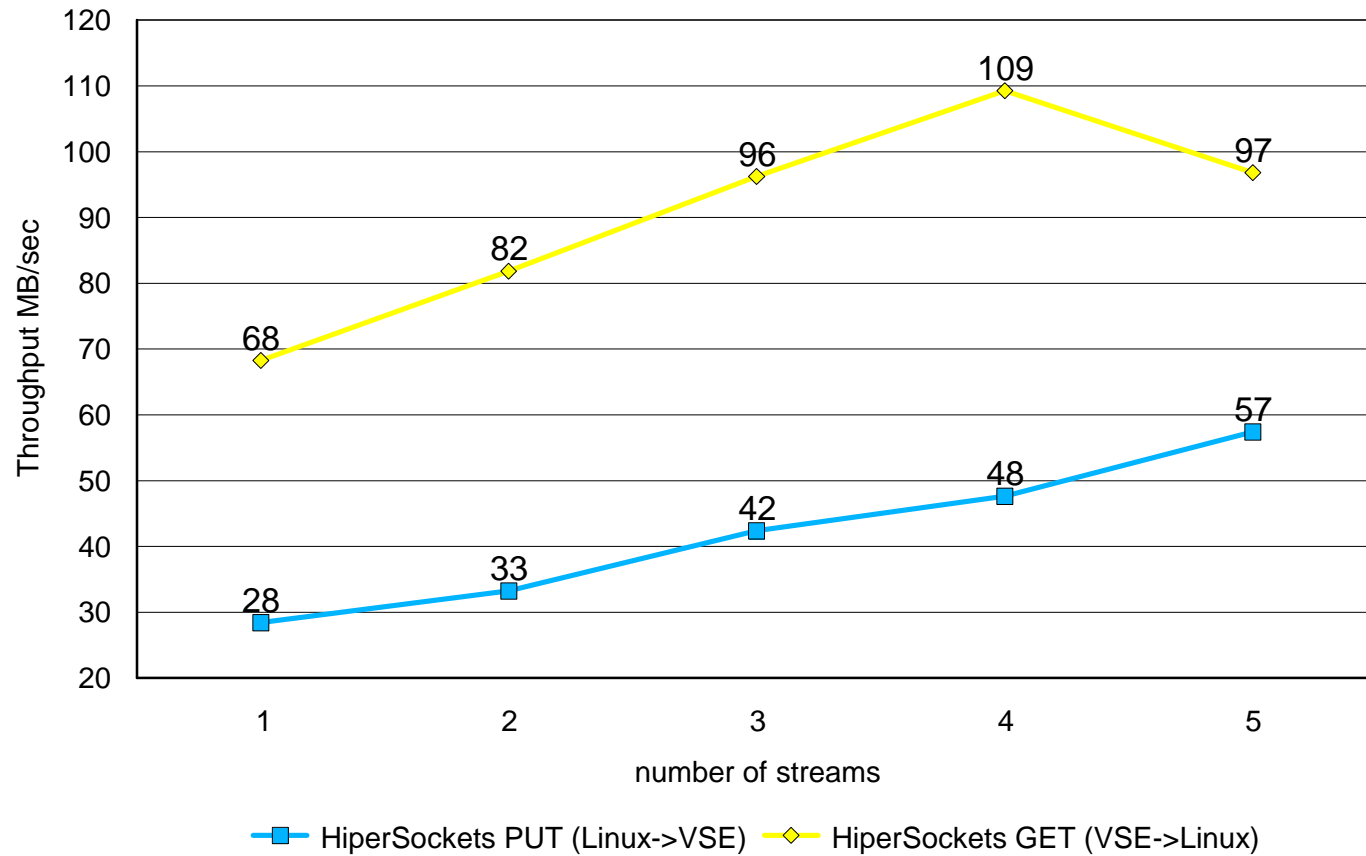
HiperSockets throughput is between 30-80 MB/sec



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## Throughput (MB/sec) - results (2)



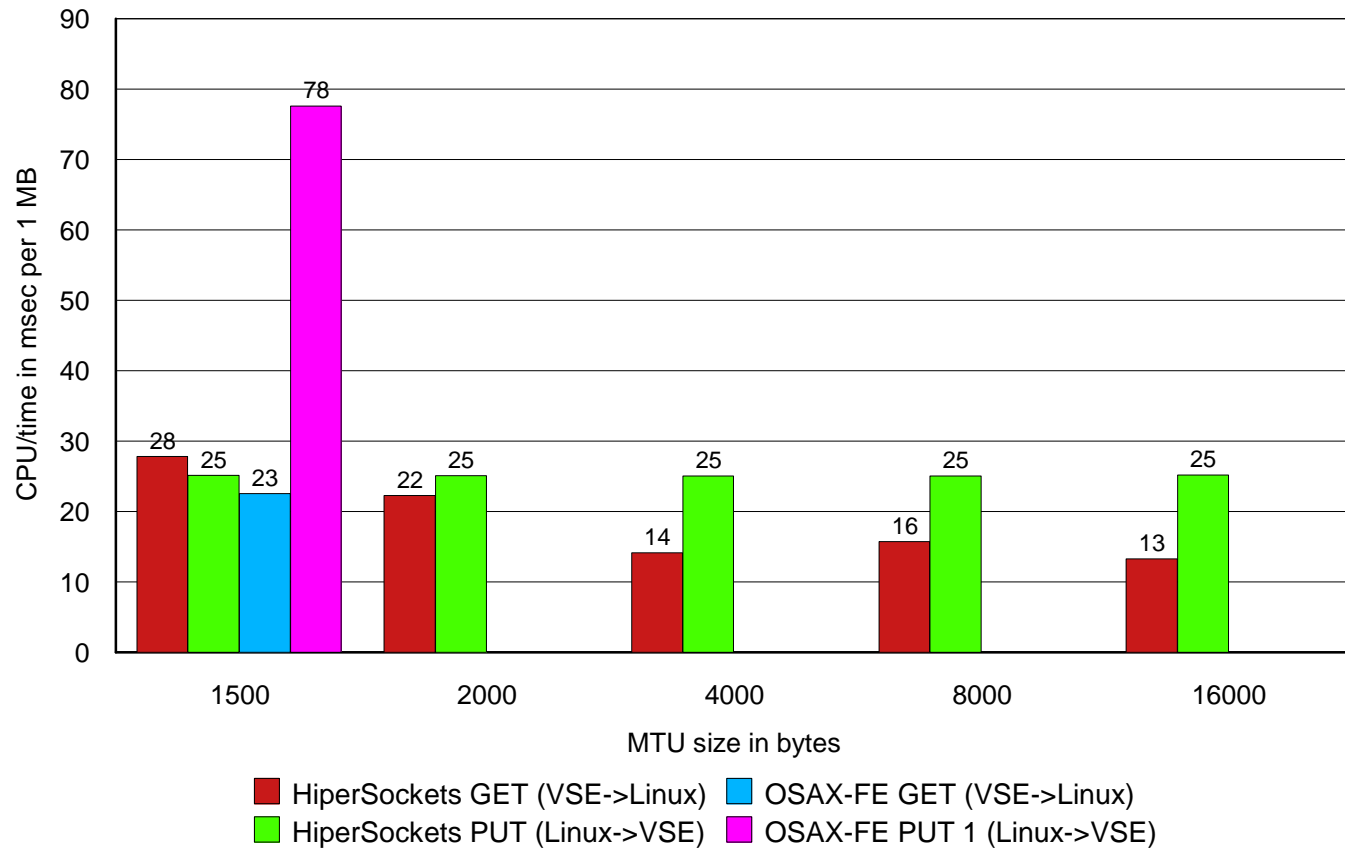
Maximum HiperSockets throughput of 109 MB/sec at 4 concurrent connections



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## CPU time per MB - results



About 15-30 msec CPU time per MB for HiperSockets  
(on a z800 2066-001)



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## Transaction per second

- Measurements has been done with an ECHO server
  - ▶ Client on Linux sends 100 bytes to server
  - ▶ Server on VSE echoes 100 bytes
  - ▶ Per TCP connection 10000 transactions are driven
  - ▶ Variations: Number of TCP connections
    - 1,2,3,4,5
    - 10,50,100
  - ▶ Measurements
    - Transactions per second
    - CPU time per transaction

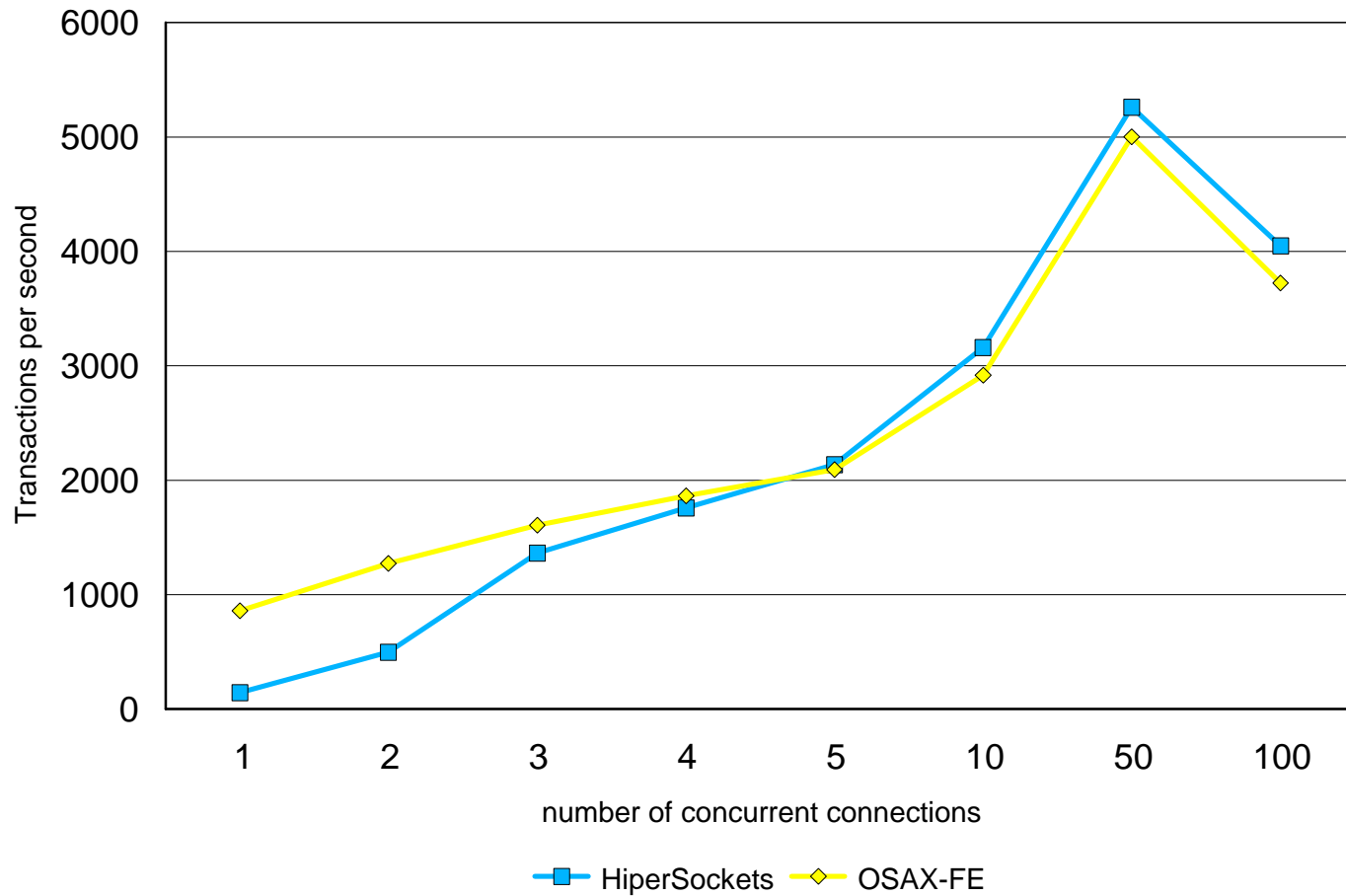




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## Transactions per second - results



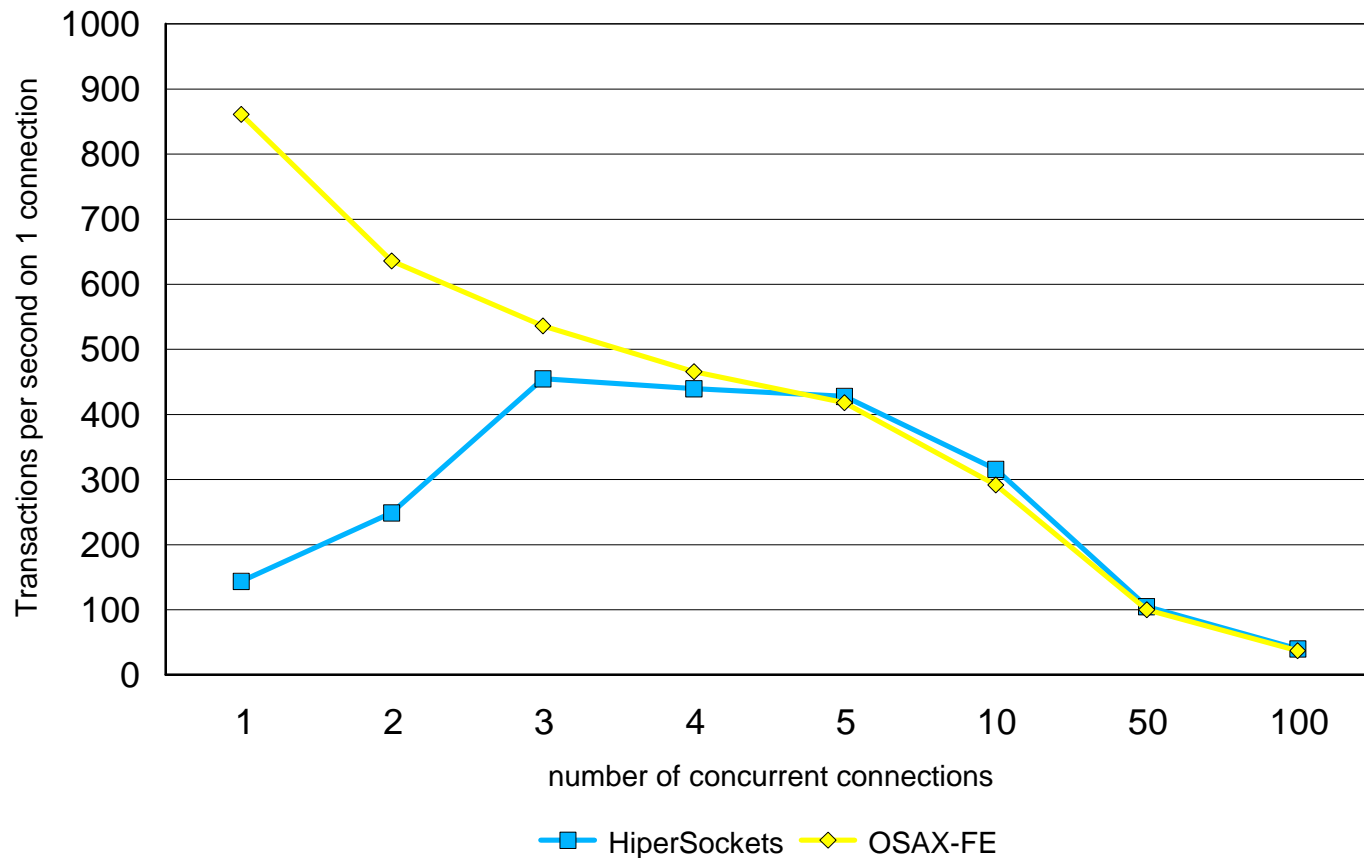
Maximum of 5200 transactions per second at 50 concurrent connections



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# Transactions per second on 1 connection - results



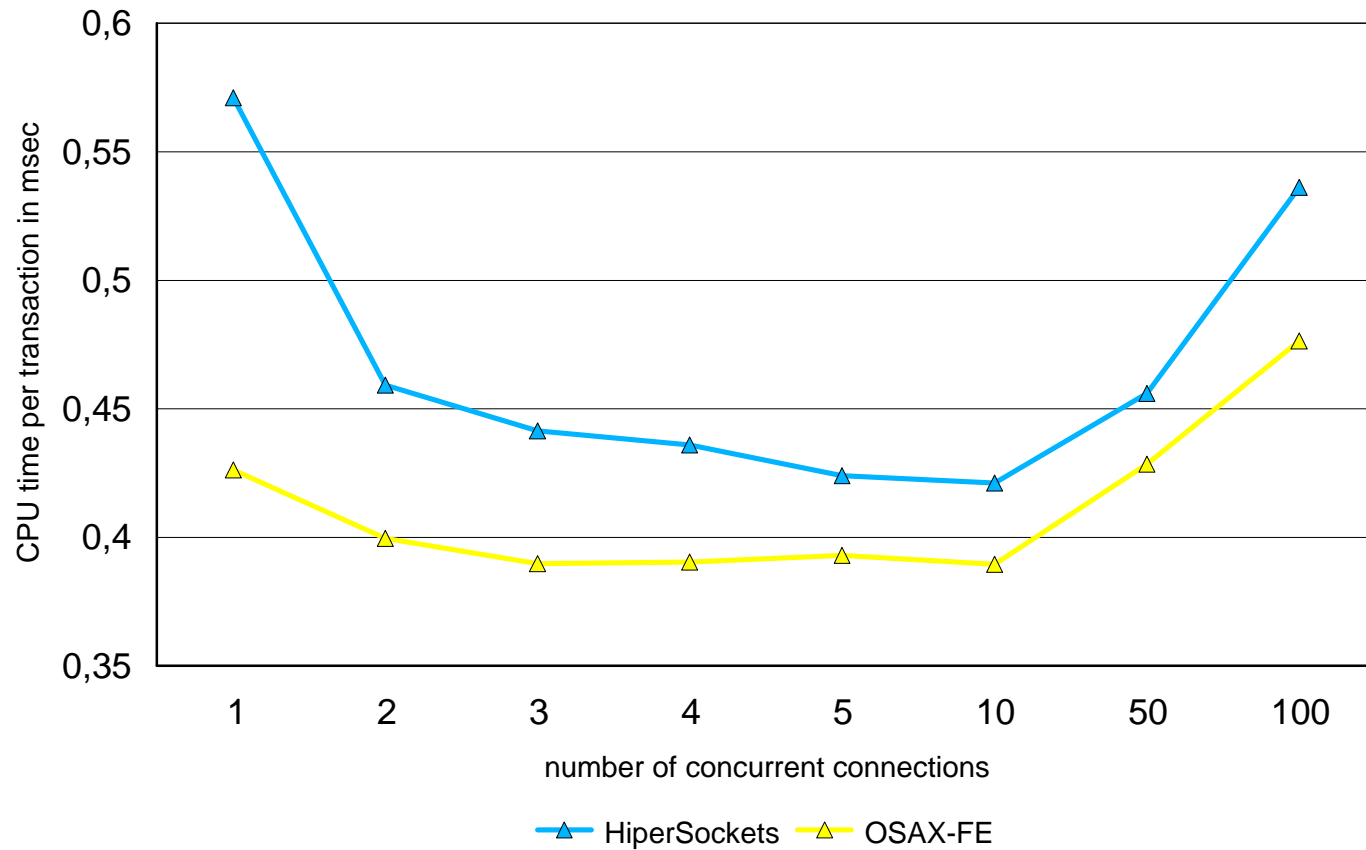
HiperSockets: Maximum of about 450 transactions per second on 1 connection (= about 2 msec response time)



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# CPU time per transaction



HiperSockets: About 0.45 msec CPU time per transaction  
for 2-50 connections



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## Measurement Results - conclusion

- HiperSockets
  - ▶ Throughput
    - Between 30-80 MB/sec
    - Maximum throughput of 109 MB at 4 concurrent connections
    - About 15-30 msec CPU time per MB
  - ▶ Transactions per second
    - Maximum of 5200 Transactions per second at 50 concurrent connections
    - About 0.4-0.45 msec CPU time per transaction



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## Hardware Crypto Overview

- Requires VSE/ESA 2.7 and TCP/IP for VSE/ESA 1.5
- Supported crypto cards
  - ▶ PCI Cryptographic Accelerator (PCICA)
    - Feature code 0862
    - Available for zSeries (z800, z900)
- Only RSA (asymmetric) is supported
  - ▶ Of benefit for Session initiation (SSL-Handshake)
- Also supported with
  - ▶ z/VM 4.2 + APAR VM62905
  - ▶ z/VM 4.3

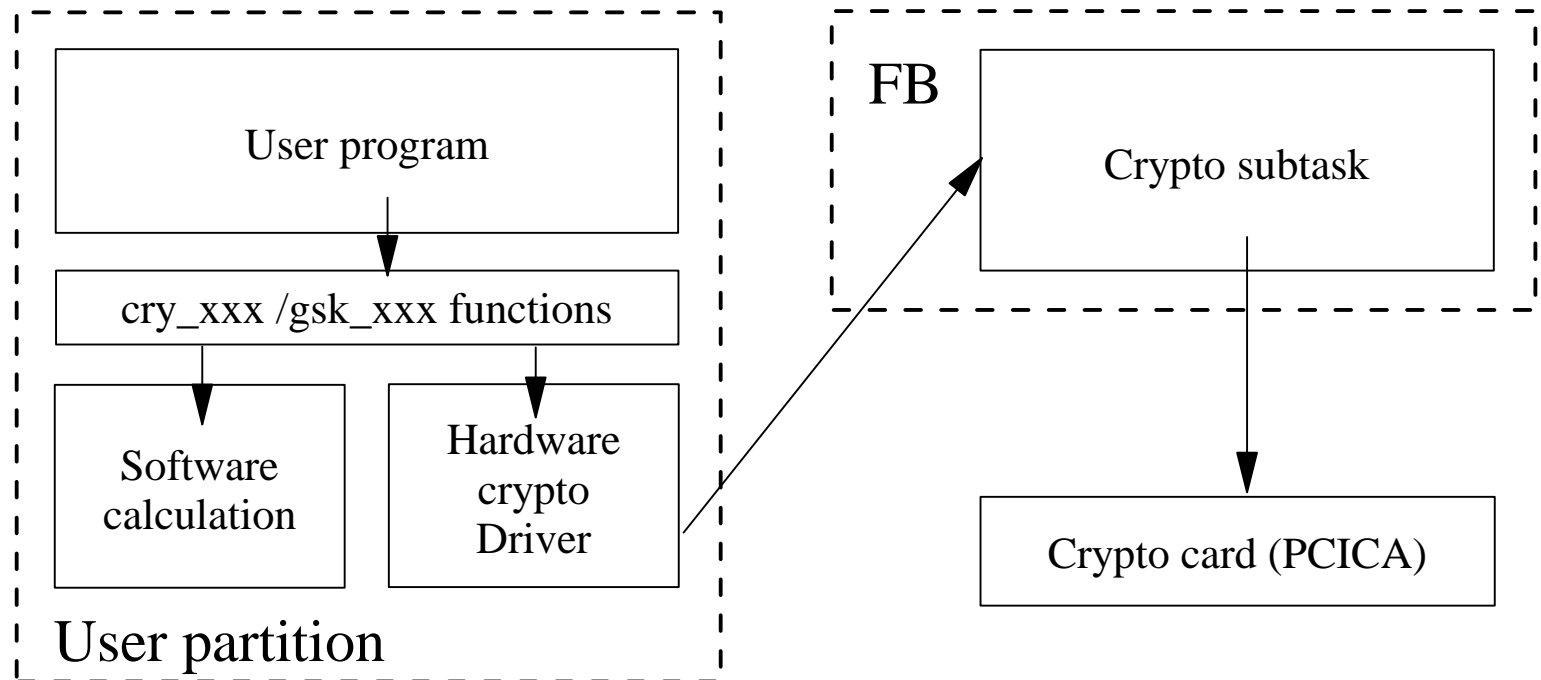


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## Hardware Crypto Overview - continued

- New crypto subtask in Security Server (SECSECV) running in FB
  - ▶ Or as separate job if no SECSECV is running
  - ▶ Crypto card is polled by crypto task





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## Measurement Environment

- VSE/ESA 2.7 running on a z900 (2064-109)
  - ▶ on 1 processor (~2064-101)
  - ▶ with a PCI Cryptographic Accelerator
- Testcase programs on VSE
  - ▶ Crypto operations measurements
    - calling cry\_xxx functions (RSA, DES, SHA, MD5)
    - each crypto operation is performed 10000 times
  - ▶ Secured data transfer (SSL)
    - performs SSL handshake
    - performs encrypted data transfer
    - counterpart program running on Windows (SSL-client)
- All RSA operations are measured
  - ▶ with Hardware Crypto support
  - ▶ with Software Crypto
    - support already available with TCP/IP 1.4/1.5 as shipped in VSE/ESA 2.6





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## Measurement Environment - continued

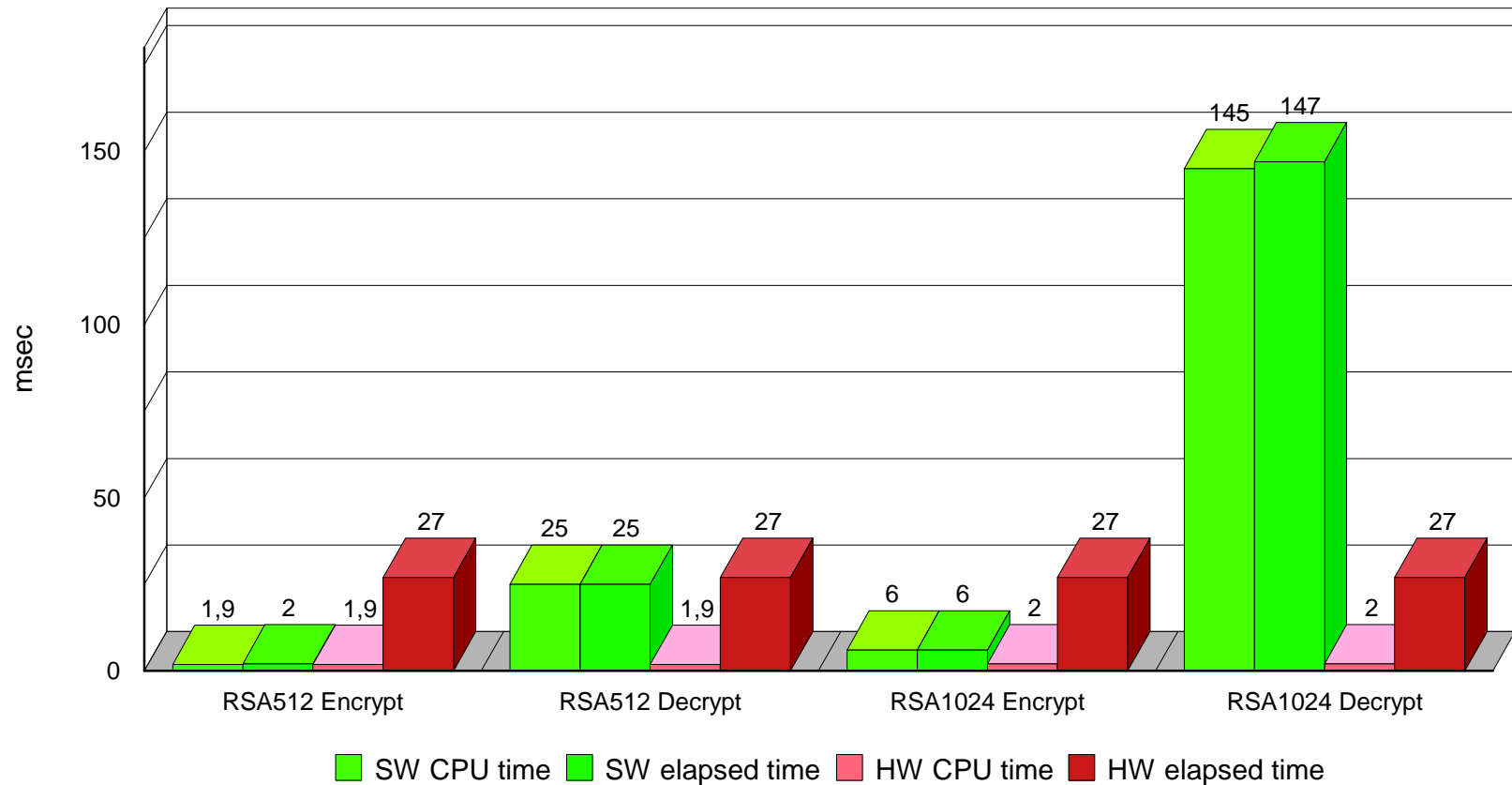
- Variations
  - ▶ RSA encrypt/decrypt
    - 512 / 1024 bit key
  - ▶ DES, DES CBC, 3DES CBC encrypt/decrypt
    - software crypto only
    - message length (128, 256, 512 bytes)
  - ▶ SHA Hash, MD5 Hash, SHA HMAC, MD5 HMAC
    - software crypto only
    - message length (128, 256, 512, 1K, 2K bytes)
  - ▶ SSL handshake/data transfer
    - 01 RSA512\_NULL\_MD5
    - 02 RSA512\_NULL\_SHA
    - 08 RSA512\_DES40CBC\_SHA
    - 09 RSA1024\_DES\_CBC\_SHA
    - 0A RSA1024\_3DES\_EDE\_CBC\_SHA



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## Measurements Results - RSA



### HW Crypto:

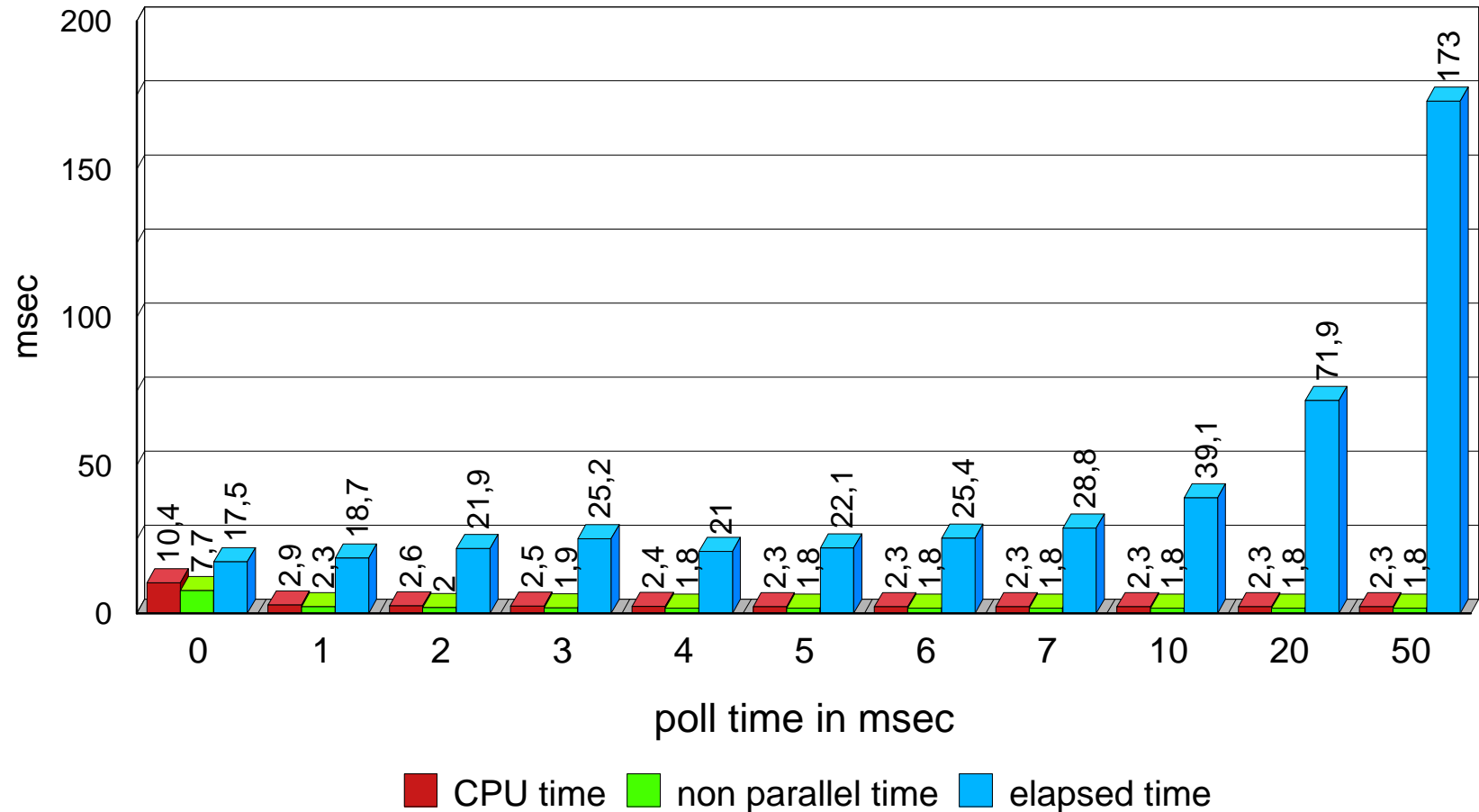
- CPU time and elapsed time is independent of operation / key length
- RSA operation takes about 2 msec CPU time and 28 msec elapsed time
- CPU time is always less than software crypto



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## Measurements Results - RSA polltime



Per default a polltime of 7 msec is used.

Can be changed with: `MSG FB,DATA=WAITTIME=nn`

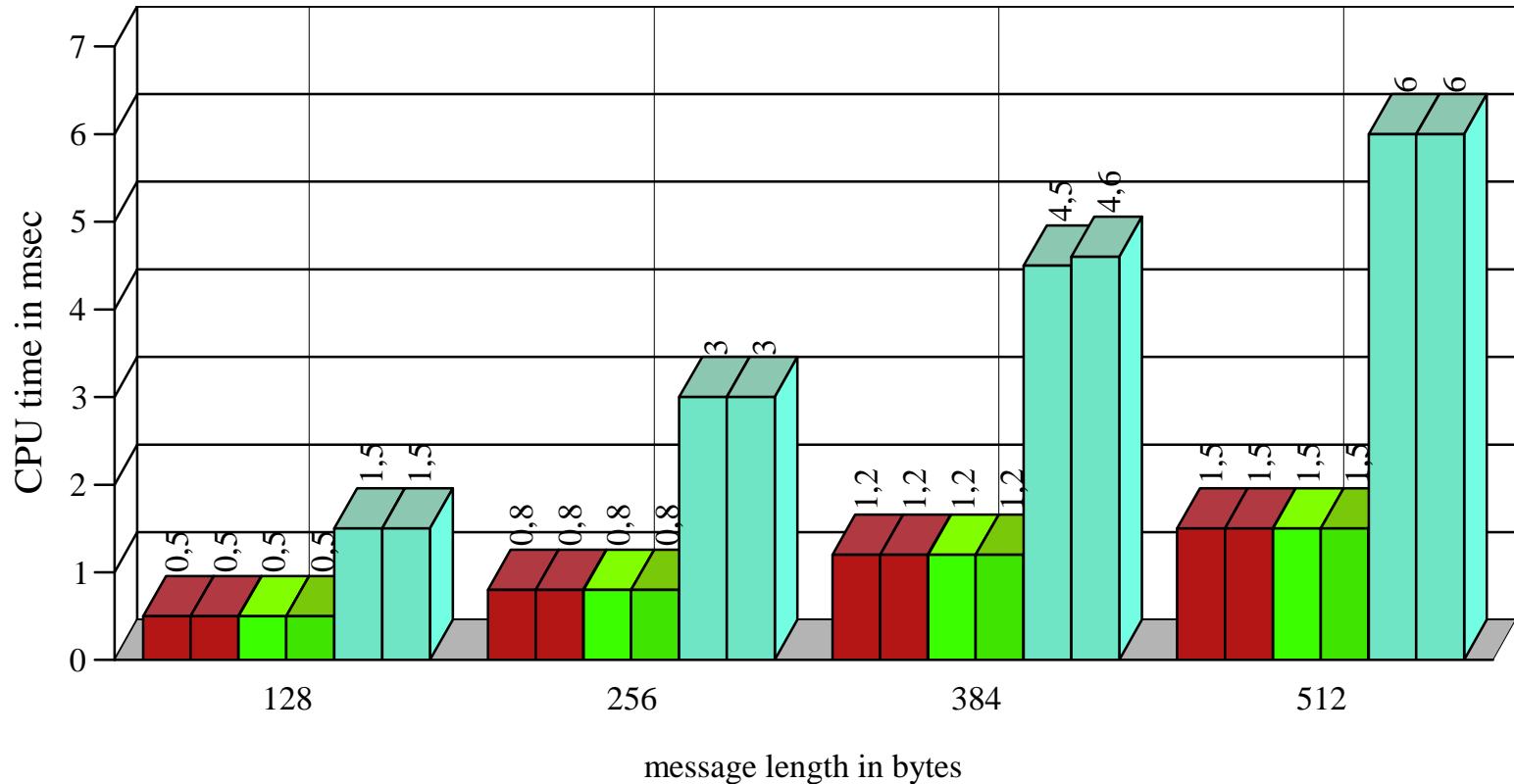
Smaller values increase CPU time, higher values increase elapsed time



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## Measurements Results - DES, DES CBC, 3DES CBC (symmetric)



■ DES Encrypt    ■ DES CBC Encrypt    ■ 3DES CBC Encrypt  
■ DES Decrypt    ■ DES CBC Decrypt    ■ 3DES CBC Decrypt

Software Crypto only!

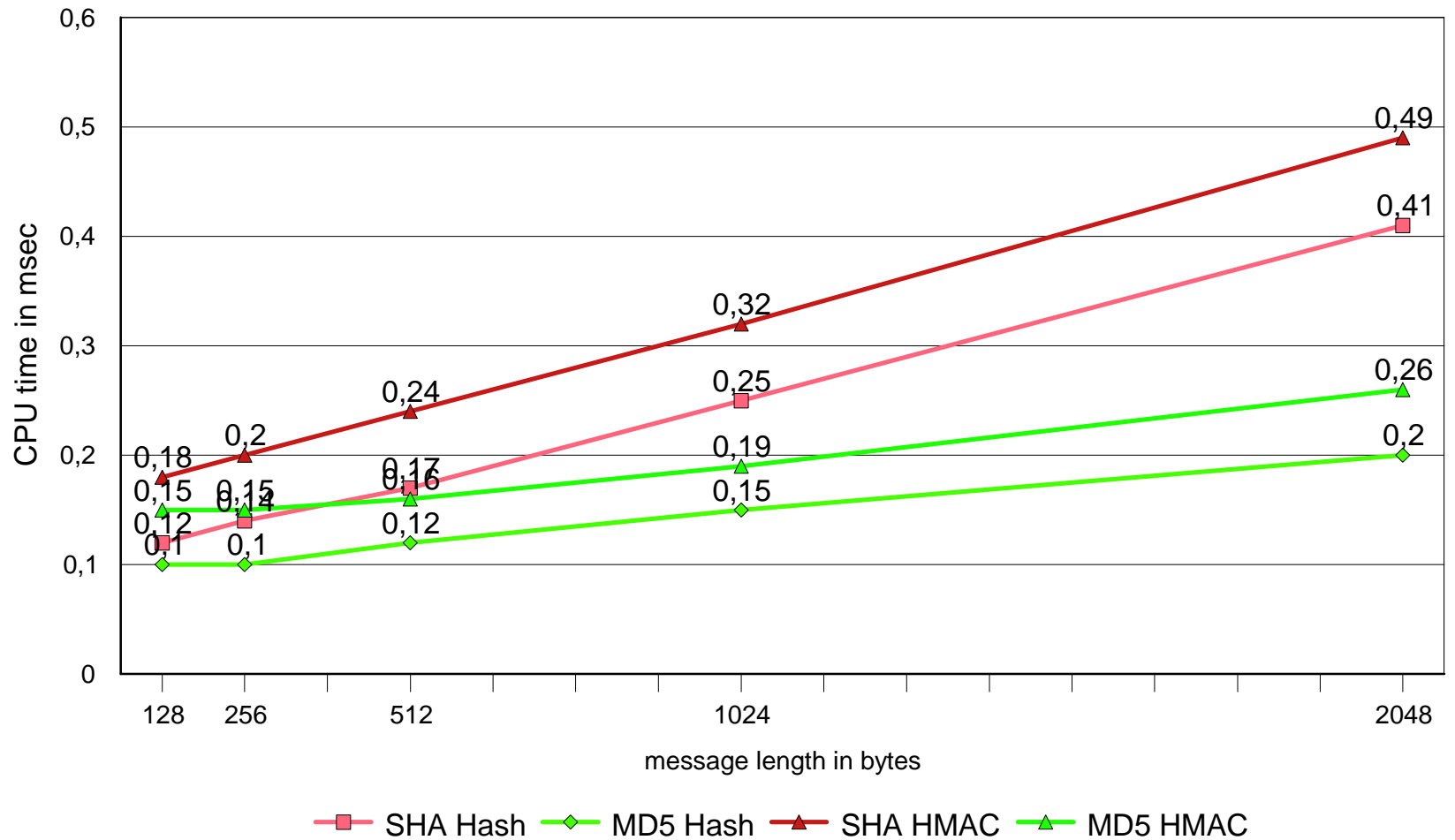
DES and DES CBC takes similar CPU times, 3DES CBC about 3.8 times



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## Measurements Results - SHA, MD5



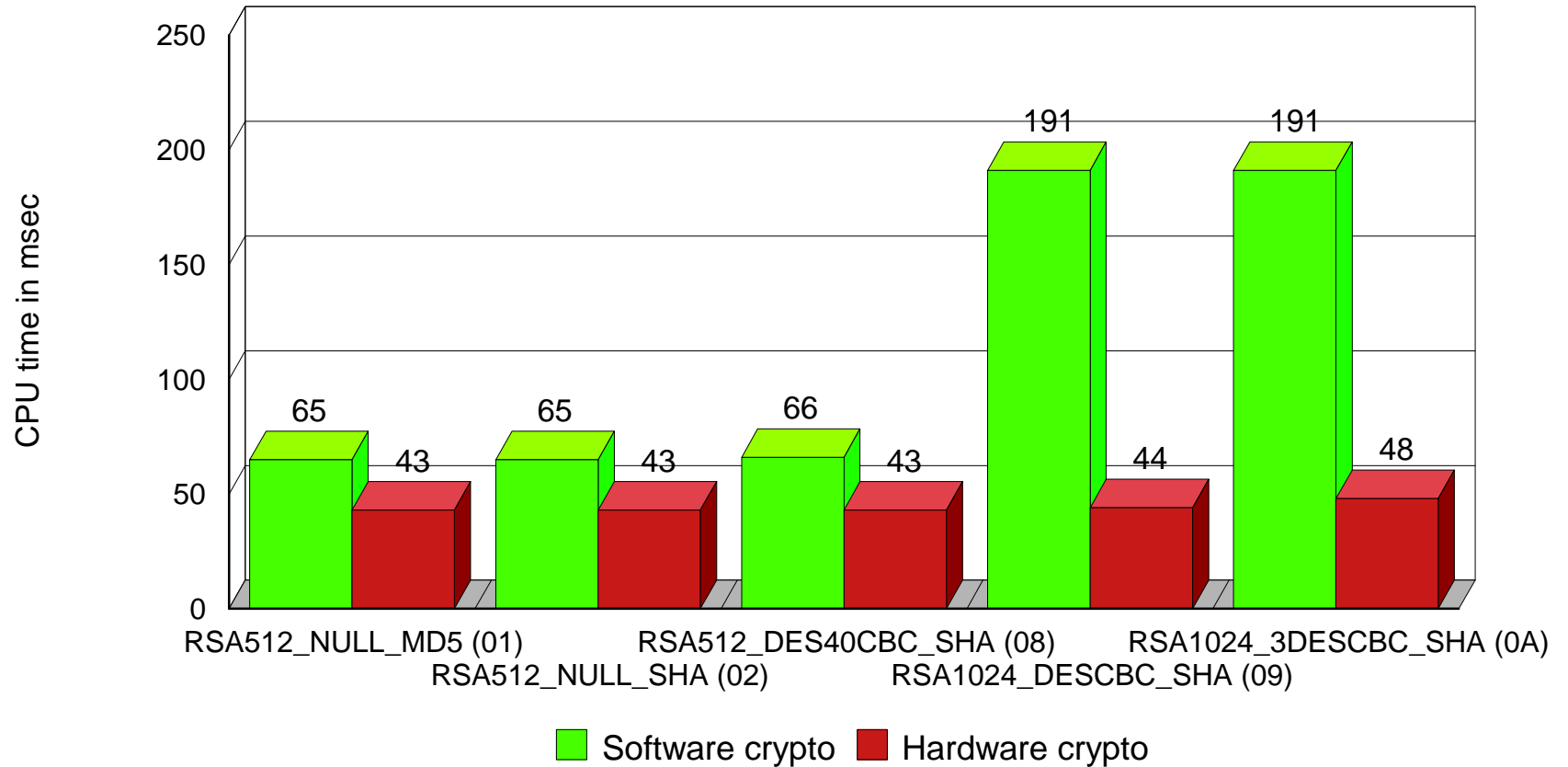
SHA takes about 1.8 times more CPU time compared to MD5 Software Crypto only!



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# Measurements Results - SSL Handshake



HW Crypto:

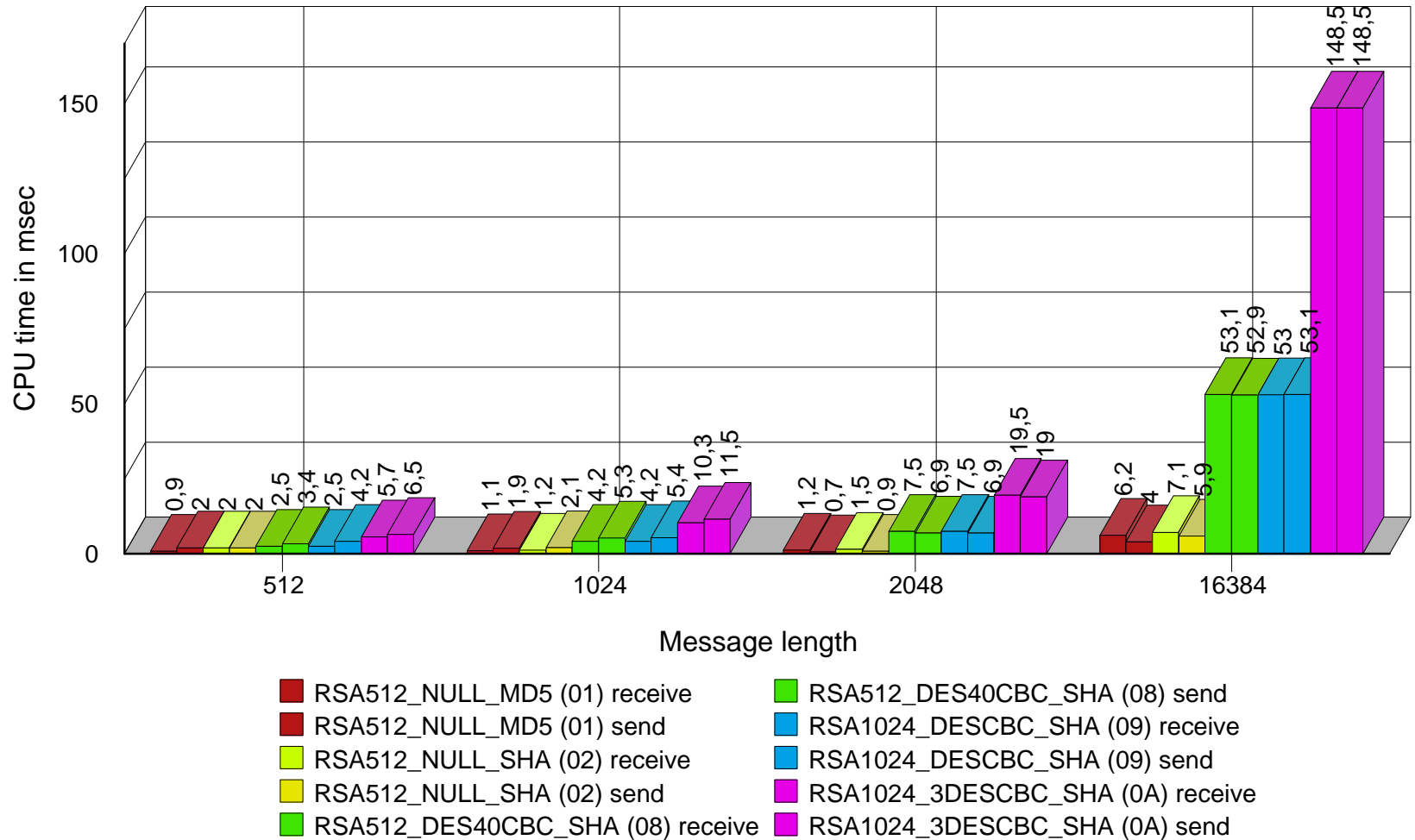
- CPU time and elapsed time is independent of cipher suite used
- SSL handshake takes about 43-48 msec CPU time (connection establishment)



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# Measurements Results - SSL data transfer



CPU time depends on used hashing (SHA/MD5) and encryption algorithm (DES/3DE  
Software Crypto only!





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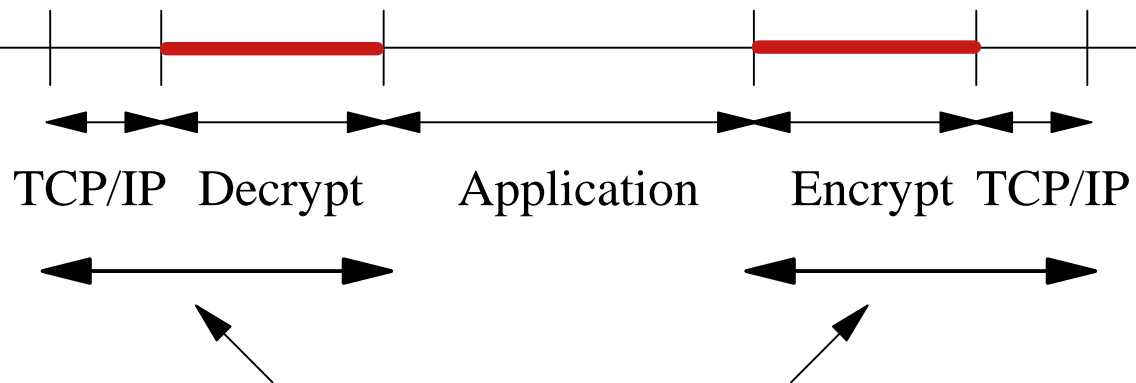


## SSL data transfer overhead

Non SSL



SSL



this has been measured



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## Measurements Results - conclusion

- HW Crypto
  - ▶ Supports RSA operations only (e.g. used by SSL handshake)
  - ▶ CPU time/elapsed time is independent of operation and key length
  - ▶ Software RSA encryption is faster in terms of elapsed time (on large processors)
    - but hardware crypto saves CPU time
- SW Crypto
  - ▶ CPUtime /elapsed time is very dependent on CPU speed and utilization



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## SSL Performance Recommendations

- Use SSL only if there is a need for
  - ▶ If at least one of the following is required
    - Keeping secrets
    - Proving identity
    - Verifying information
- Cipher Suites 01 and 02 has less CPU-time consumption, but NO data encryption
  - ▶ RSA512\_NULL\_MD5, RSA512\_NULL\_SHA
- If data encryption is required
  - ▶ Use cipher suites 08, 09 or 0A
  - ▶ 08 uses 512 bit keys, others 1024
  - ▶ 1024 bit RSA keylength is recommended (from a security point of view)



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## Turbo Dispatcher - Overview

- Turbo Dispatcher
  - ▶ available since 1995
  - ▶ VSE/ESA 2.1-2.3 Standard and Turbo Dispatcher
  - ▶ since VSE/ESA 2.4 only Turbo Dispatcher
  - ▶ last changes:
    - VSE/ESA 2.6.2 (APAR DY45869)
    - VSE/ESA 2.7.0 (APAR DY45926)
  - ▶ Supports basic (native), LPAR and VM mode
  - ▶ Runs on Uni- and n-Way-procercssors
    - CPUs have "equal" rights
    - more than 3 CPUs are not recommended



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## Turbo Dispatcher - Overview (2)

- IPL is done on 1 CPU only
  - ▶ after IPL other CPUs can be started
  - ▶ CPUs can be started or stopped without re-IPL
  - ▶ at least 1 CPU (IPL CPU) must always be active

```
SYSDEF TD,START=n|ALL
```

```
SYSDEF TD,STOP=n|ALL
```

```
SYSDEF TD,STOPQ=n|ALL
```

```
QUERY TD
```



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## Turbo Dispatcher - Quiesced CPUs

- SYSDEF TD,STOPQ=n to set a CPU in quiesced mode
  - ▶ Implemented for z/VM guest systems
    - **Not started guest CPUs stop IOASSIST**
    - STOPQ remains IOASSIST active, and avoids TD Overhead, (CPU will no longer participate in work unit selection)
    - quiesced CPUs will not process any workunits
    - quiesced CPUs will not handle any interrupt
    - quiesced CPUs can be started with SYSDEF TD,START



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## Turbo Dispatcher - Design

- **TD dynamically assigns partitions to CPUs**
  - ▶ **Work unit** = from assignement to one CPU until next interrupt/SVC
  - ▶ If one task (subtask) of a partition is active, no other task of the same partition will be selected
  - ▶ TD dispatches on partition-basis, not on task-basis
  - ▶ A job running in a partition is processed in several work units.





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## Turbo Dispatcher - Design (2)

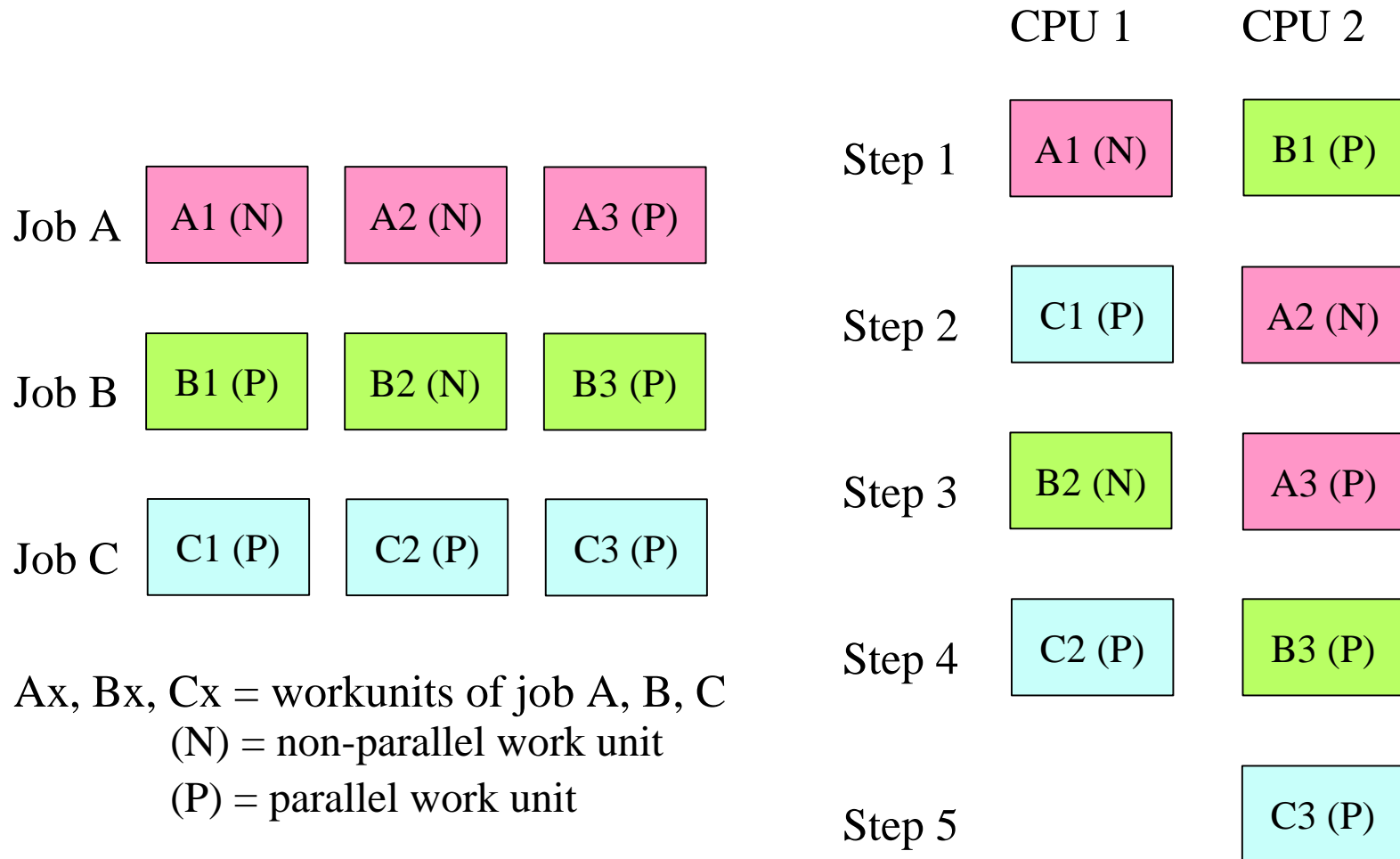
- **parallel work units**
  - ▶ Application code (CICS, Batch)
  - ▶ may run on any CPU concurrently with other parallel or non-parallel work units.
- **non-parallel work units**
  - ▶ System code (Services, VTAM, Vendor code)
  - ▶ As long as one non-parallel work unit is active on one CPU, no other non-parallel work unit can execute on any other CPU.



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# Turbo Dispatcher - Design - Example 1





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## Turbo Dispatcher - Design - Example 2

CPU 1

select A

↓ A (P)  
SVC

↓ A (N) - SVC Code

↓ Dispatcher

↓ A (P)

↓ Interrupt

↓ (N)

↓ Dispatcher

↓ B (P)

CPU 2

select B

↓ B (P)  
SVC

↓ wait for (N) = spin or delay

↓ (Dispatcher)

↓ B (N) - SVC Code

↓ Dispatcher

↓ A (P)



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## Turbo Dispatcher - Exploitation

- Uni-Processor
  - ▶ new Partition Balancing Concept
    - Helps to set priorities of partitions
  - ▶ Determination of non-parallel share, to find out if a 2. or 3. CPU would be of use
- n-Way Processors (2-3 CPUs)
  - ▶ System tuning required for exploitation
  - ▶ Increased Capacity (dependent on workload)
    - Exploitation increases by reduction of non-parallel work units



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# Turbo Dispatcher - CPU time measurement

- CPU time measurement (overall system)
  - ▶ SYSDEF TD,RESETCNT
  - ▶ Workload (e.g. run a job)
  - ▶ QUERY TD (QUERY TD,INTERNAL)

CPU	STATUS	SPIN_TIME	NP_TIME	TOTAL_TIME	NP/TOT
00	ACTIVE	0	237100	416698	0.568
01	ACTIVE	0	157556	415229	0.379
02	QUIESCED	0	0	0	*.***
03	INACTIVE				
-----					
TOTAL		0	394656	831927	0.474

NP/TOT: 0.474

SPIN/(SPIN+TOT): 0.000

OVERALL UTILIZATION: 179%

NP UTILIZATION: 85%

ELAPSED TIME SINCE LAST RESET:

463433

NP/TOT = non-parallel share (NPS)

SPIN\_TIME = CPU time waiting for NP



# Display System Activity Dialog

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```

Session C - [32 x 80]
File Edit View Communication Actions Window Help
IESADMDA          DISPLAY SYSTEM ACTIVITY          15 Seconds  13:55:26
*----- SYSTEM (CPUs:  1 /  0 ) -----* *----- CICS : DBDCCICS -----*
| CPU      :      0%  I/O/Sec:    1  | | No. Tasks:  7,018  Per Second :    *  |
| Pages In :      0  Per Sec:    *  | | Dispatchable:    0  Suspended  :    3  |
| Pages Out:      0  Per Sec:    *  | | Peak Active  :    7  MXT reached:    0  |
*-----* *-----*
Priority: Z,Y,S,R,P,C,BG,FA,F9,F8,F6,F5,F4,F2,F7,FB,F3,F1

  ID S JOB NAME      PHASE NAME      ELAPSED        CPU TIME      OVERHEAD      %CPU          I/O
  F1 1 POWSTART      IPWPOWER       29:23:33        1.23          .37           6,000
  F3 3 VTAMSTRT      ISTINCVT       29:23:28       18.13         5.65          304,230
  FB 8 SECSESV       BSTPSTS        29:23:33         .03           .01           213
 *F7 7 TCPIP00      IPNET          29:23:28         1.61          .77           814
  F2 2 CICSICCF      DFHSIP         29:23:28       597.71       169.82         8,718
  F4 4 <=WAITING FOR WORK=>
  F5 5 <=WAITING FOR WORK=>
  F6 6 <=WAITING FOR WORK=>
  F8 8 <=WAITING FOR WORK=>
  F9 9 <=WAITING FOR WORK=>
  FA A <=WAITING FOR WORK=>
  BG 0 <=WAITING FOR WORK=>
PF1=HELP      2=PART.BAL.  3=END          4=RETURN      5=DYN.PART    6=CPU
  
```

MA c 01/001

Connected to remote server/host boevmct1 using port 23 Print to Disk - Append





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## Migration aspects

- Consider hard-/software requirements:
  - ▶ Does my largest partition still fit into a single CPU of the target processor?
    - Note: a partition can only run on 1 CPU at a time
  - ▶ Is the processor capacity and speed still sufficient to run the workload?
  - ▶ Does multiprocessing help to run the workload?
    - What about non-parallel share (on 1-Way)?
    - Are there many parallel batch jobs?
      - A large CICS partition does not benefit of a 2. CPU





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## Migration overhead

- Uni-Processor
  - ▶ increased overhead because of
    - Release migration (VSE/ESA 2.6 vs. 2.7)
    - TD overhead (Standard Dispatcher vs. TD)
    - CICS/VSE vs. CICS TS
- N-Way Processor
  - ▶ CPU time increases when migrating from uni to n-Way Processor (for the same workload)
    - For PACEX Workload: Factor 1.4 (2 CPUs)
    - TD overhead for multiprocessor exploitation
    - z/VM Overhead



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# Migration

VSE/ESA 2.3  
Standard Dispatcher  
CICS/VSE 2.3

VSE/ESA 2.3  
**Turbo Dispatcher**  
CICS/VSE 2.3

**VSE/ESA 2.7**  
(Turbo Dispatcher)  
CICS/VSE 2.3

VSE/ESA 2.7  
(Turbo Dispatcher)  
**CICS TS 1.1**

Change only  
one thing at a time!

Allows you to see which step  
has introduced a problem.



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## Performance Tips

- A partition can only exploit **1 CPU** at a time
  - ▶ 2 CPUs do not have any benefit for a CICS partition
  - ▶ Use as many partitions as required for selected n-way
- Use/define only as many CPUs as really needed
  - ▶ additional CPUs create more overhead, but no benefit
- Partitions setup
  - ▶ Set up more batch and/or (independent) CICS partitions
  - ▶ Split CICS production partitions into multiple partitions



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## Performance Tips (2)

- **1 CPU** must be able to handle **all non-parallel workload**
- Non-parallel code limits the n-Way exploitation
  - ▶ QUERY TD:  $NP/TOT = NPS$
  - ▶ Measure NPS before migration
  - ▶ **max CPUs =  $0.9 / NPS$**

<b>NPS</b>	<b>#CPUs</b>	<b>NPS</b>	<b>#CPUs</b>
0.20	4.5 (4)	0.40	2.2 (2)
0.25	3.6 (3)	0.45	2.0 (2)
0.30	3.0 (3)	0.50	1.8 (1)
0.35	2.6 (2)	0.55	1.6 (1)



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## Performance Tips (3)

- Non-parallel code limits the maximum MP exploitation
- System code (Key 0) increases non-parallel share
  - ▶ Vendor code can have significant impact
- Overhead increases when NP code limits throughput
- Data In Memory (DIM) reduces non-parallel code
  - ▶ less system calls (I/Os)
  - ▶ may increase throughput
- In general **ONE faster CPU** is better than multiple slower ones
  - ▶ Even if sum of slower CPUs is higher than one faster CPU



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## CICS Implications

- Single CICS
  - ▶ Can consume processing power of one CPU only
  - ▶ parallel batch jobs may exploit 2. CPU
- Multiple CICS partitiones
  - ▶ Number of CPUs depends on non-parallelen share (NPS)
  - ▶ Function shipping and Transaction routing
    - AOR, TOR, FOR



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## Partition Balancing

- Balanced Group is defined with PRTY:
  - ▶ PRTY BG, **C=F5=F8**, F2, F3, F1
  - ▶ Each partition/class of the group has a default-SHARE (100)
  - ▶ Dynamic partitions gets the SHARE of its class
- To set a SHARE (1-1999)
  - ▶ PRTY SHARE, **F5=50**
  - ▶ SHARE = 0 means the lowest priority within the group

PRTY

AR 0015 PRTY BG, C=F5=F8, F2, F3, F1

AR 0015

AR 0015 SHARE F5= 50, F8= 100, C= 100

MSECS

AR 0015 MSECS 976 <---- influences task selection

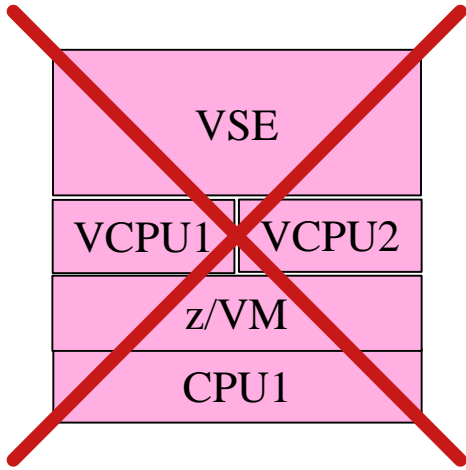




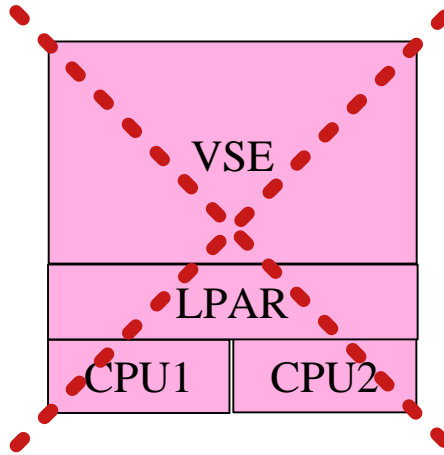
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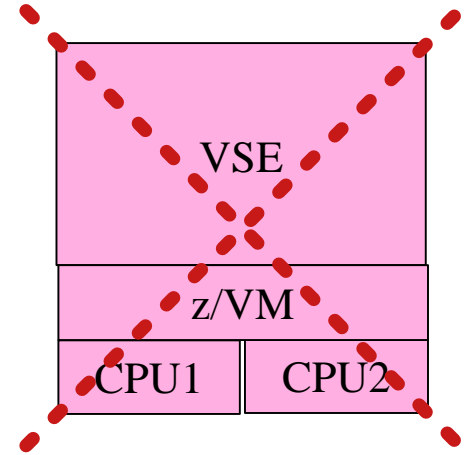
# Do's and Don't Do's



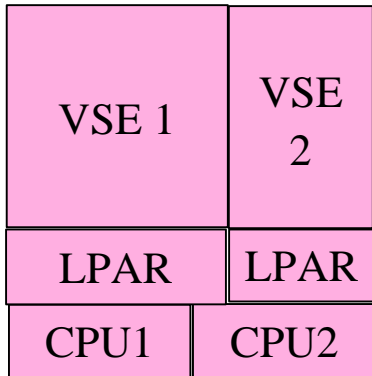
no virtual CPUs!  
(creates overhead)



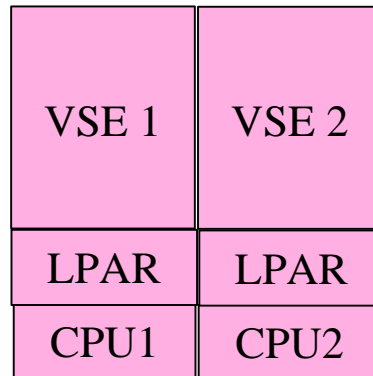
only if NPS < 4.5



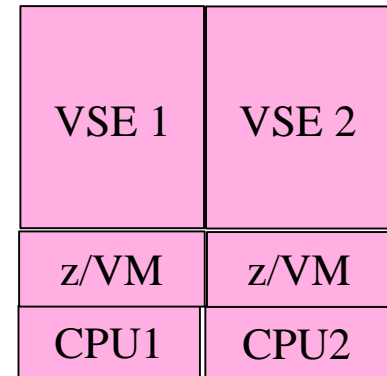
only if NPS < 4.5



VSE 1 = Production  
VSE 2 = Test



dedicated CPU  
per VSE



dedicated CPU  
per VSE



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## Do's and Don't Do's (2)

The fastest  
uni-processor  
is (almost always)  
the best processor !



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## New: z/VSE 3.1 preview announcement

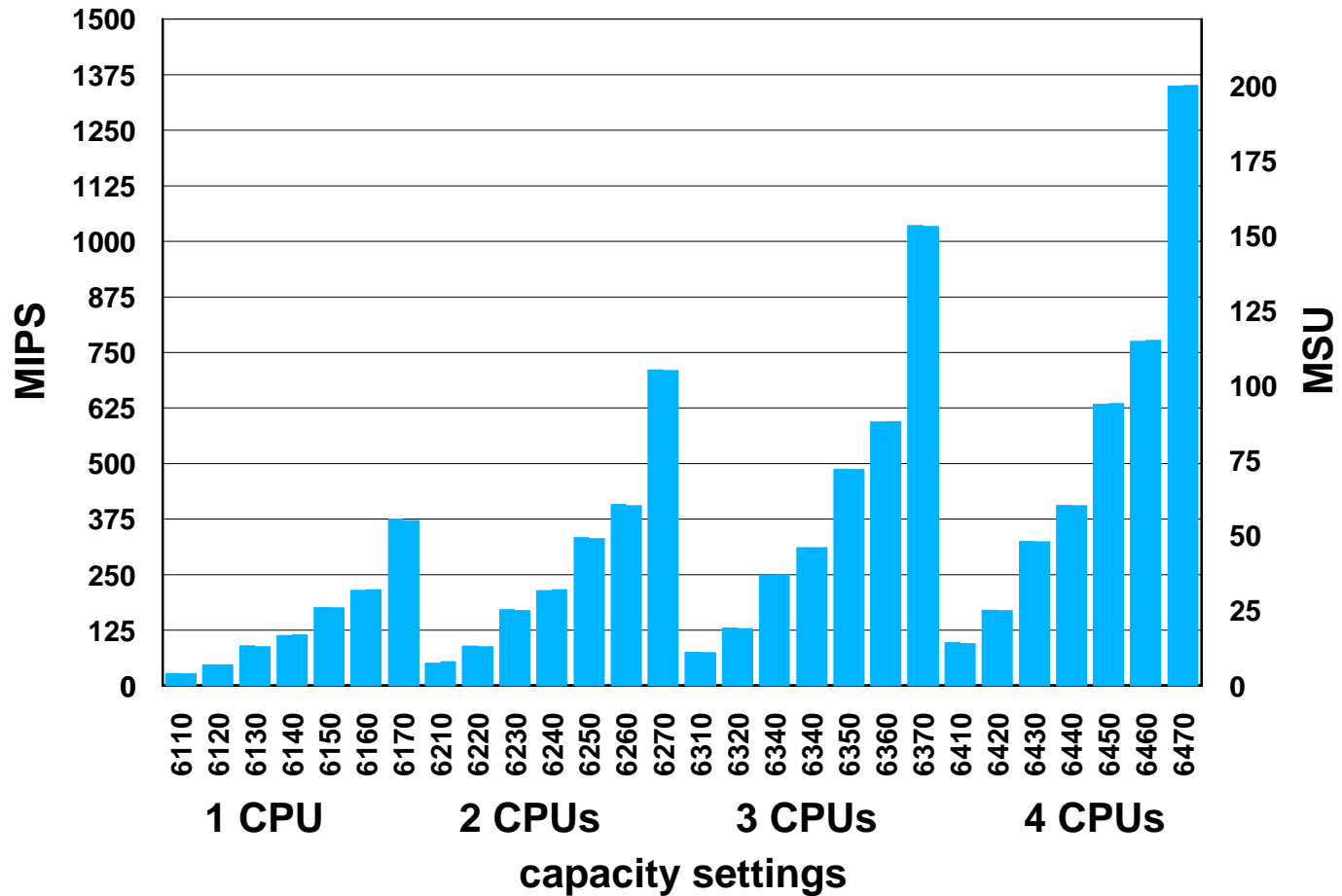
- z/VSE 3.1 is planned to be able to support:
  - ▶ z890, z800, z900, z990
  - ▶ Multiprise 3000, G5 and G6
  - ▶ Fibre Channel Protocol for SCSI — FCP channels
- IBM plans to continue to ship CICS/VSE V2.3 together with CICS TS for VSE/ESA
  - ▶ at no additional charge.
- z/VSE plans to offer simplified packaging
  - ▶ LE will become a component of VSE Central Functions
- Fast Service Upgrade possible from VSE 2.7 and 2.6
  - ▶ using equivalent ECKD disks
  - ▶ NOT: from ECKD to SCSI-FCP disks.



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## New: IBM eServer zSeries 890



z890 consists of one Model (A04) and 28 capacity settings



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## **New: Enterprise Storage Server Model 750 (Baby-Shark)**

- The ESS Model 750 is based on the same architecture as the ESS Model 800 to support functionality, stability, and reliability
- up to 64 disk drives
- 4.6 terabytes (TB) of physical capacity
- A two-way processor
- 8 GB of cache
- 2 GB of Non Volatile Storage (NVS)
- up to 6 Fibre Channel/FICON or ESCON host adapters
- Support for 72.8 GB and 145.6 GB 10,000 rpm drives
- configured as RAID 5, RAID 10, or a combination of both



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## Dependencies for VSE/ESA Growth

- System dependencies
  - ▶ Many control-blocks etc.. still below the line
  - ▶ VTAM IOBUF areas in System GETVIS-24
  - ▶ Non-Parallel-Share limits n-way support
  - ▶ Number of tasks
    - Up to 255, 32 per partition, 208 subtasks in total
- Application dependencies
  - ▶ Integrated system concepts/functions
  - ▶ Functions/Applications dependencies
  - ▶ Number of users per TCP/IP partition



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## Dependencies for VSE/ESA Growth - continued

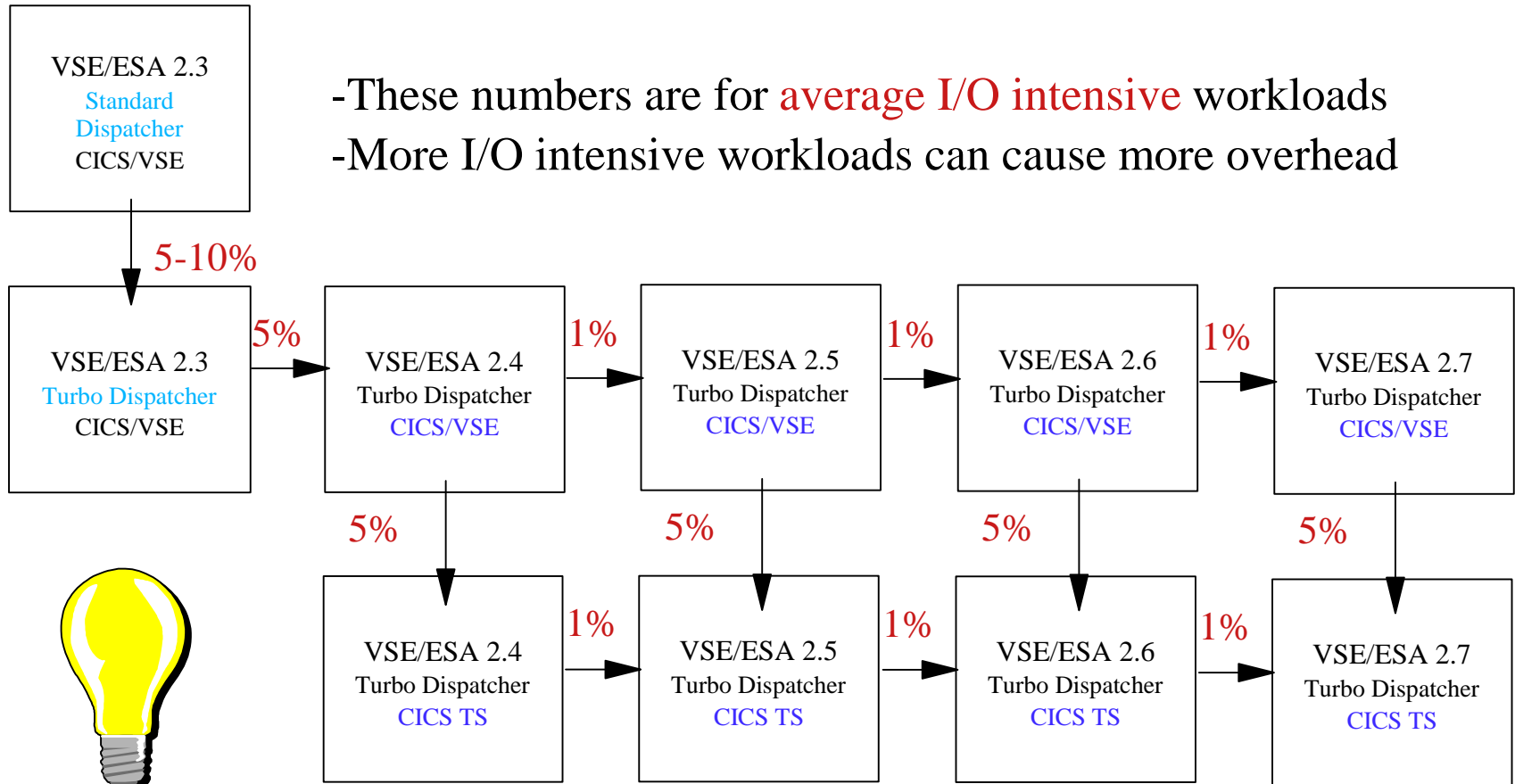
- Not being considered to be a limit
  - ▶ Number of partitions
    - 12 static + 150-200 dyn. partitions
  - ▶ Real storage (max. 2 GB)
  - ▶ Total virtual storage (max. 90 GB)
  - ▶ Total number of devices (3 digit CUU)
    - Max. 1024 devices (and 16 channels)
  - ▶ Total number of logical units
    - 255 per partition and  $12 \times 255 = 3060$  in total
  - ▶ Label area
    - Max. about 9000 in total, and 712 in sub areas





# Overhead Deltas for VSE Releases

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Remember that you get a lot of **new functions** that in most cases helps you to **increase VSE system performance and throughput**:

**Partition Balancing, PRTY SHARE (Turbo Dispatcher), FlashCopy (ESS), Buffer Hasing, Shared data Tables (CICS TS)**



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## VSE Health Check

- Goals
  - ▶ Recognize actual/upcoming problems
  - ▶ Optimize the system for new/current workload
- A-B-C analysis
  - ▶ A - concentrate on the essentials
    - 20 % work for 80 % results
  - ▶ B - more detailed analysis
    - 30 % work for 15 % results
  - ▶ C - analyze all details
    - 50 % work for 5 % results
- A-B analysis takes about 2 days
- C analysis takes about 1 week
- Should be done about once a year



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## VSE Health Check - continued

- What should be checked?
  - ▶ Processor (utilization, dispatching, z/VM, ...)
  - ▶ DASD, Tapes (I/O rate, cache, ...)
  - ▶ Network (network load, missrouted packets, ...)
  - ▶ System software
    - Turbo Dispatcher (PRTY, PRTY SHARE, ...)
    - VSAM (CA/CI sizes, shareoptions, buffers, ...)
    - CICS (MXT, DSA/EDSA sizes, SOS, ...)
    - Storage Layout (GETVIS 24, SVA, partitions, DSPACE, ...)
    - VTAM (bufferpool)
    - POWER (DBLK, DBLKGP, ...)
    - LE runtime options (Heap size, ...)
  - ▶ Application software



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## Hints and Tips for Performance

- Try to exploit Turbo Dispatcher functions
  - ▶ Priority settings
  - ▶ Partition balancing
  - ▶ Partition balancing groups
- Use as much data in memory (DIM) as possible
  - ▶ CICS Shared Data Tables
  - ▶ Large/many VSAM Buffers (with buffer hashing)
  - ▶ Virtual Disks
- Switch tracing/DEBUG off for production



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## Hints and Tips for Connector- and TCP/IP-Performance

- Reduce amount of data transferred
  - ▶ Transfer only data that is needed
  - ▶ Issue only requests that are needed
- Use connection pooling
  - ▶ Reduce overhead of connection establishment
- Performance of connectors depends on
  - ▶ Network performance
  - ▶ Performance of "server"
  - ▶ Performance of "client" or middle tier
- Reduce misrouted packets
- Use a packet filter
  - ▶ Unwanted packets increases TCP/IP and CPU load



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## Further Information

- VSE Homepage:  
<http://www.ibm.com/servers/eserver/zseries/os/vse/>
- VSE Performance Homepage:  
<http://www.ibm.com/servers/eserver/zseries/os/vse/library/vseperf.htm>
- Performance Documents from W. Kraemer
  - ▶ available on the Performance Homepage
- VSE/ESA e-business Connectors User's Guide  
<http://www.ibm.com/servers/eserver/zseries/os/vse/pdf/ieswue20.pdf>