



e-business

VSE/ESA 2.6 and 2.7

Performance Considerations

Ingo Franzki

e-mail: ifranzki@de.ibm.com

VSE/ESA Development





e-business



Contents

- VSE/ESA 2.6
 - ▶ Performance Items (Overview)
 - ▶ Hardware Support
 - ▶ Virtual Tape Support
 - ▶ e-Business Enhancements
- VSE/ESA 2.7
 - ▶ Performance Items (Overview)
 - ▶ HiperSockets
 - ▶ Hardware Crypto support
- z800/z900 Remarks
- Dependencies for VSE/ESA Growth



e-business



VSE/ESA 2.6 Performance Items

- VSE/ESA 2.6 Base enhancements
 - ▶ Delete Label Function
 - ▶ LTA Offload for some AR commands
 - ▶ SVA-24 Phases moved above the line
 - ▶ Increased max number of SDL entries
 - ▶ SDL update from non-BG partitions
 - ▶ POWER Data file extension without reformat



e-business



VSE/ESA 2.6 Performance Items - continued

- VSE/ESA 2.6 Hardware Support
 - ▶ FICON Support (VSE/ESA 2.3 or higher)
 - ▶ New 2074 System Management Console
 - ▶ OSA Express Adapter (e.g. Gigabit Ethernet)
 - ▶ VSAM Support for large 3390-9 Disks (Shark)
 - ▶ Fastcopy Exploitation of ESS FlashCopy and RVA SnapShot



e-business



VSE/ESA 2.6 Performance Items - continued

- VSE/ESA 2.6 e-Business Enhancements
 - ▶ Updated Java-based connectors
 - ▶ VSAM SHROPT(4) avoidance for connectors
 - ▶ SSL for VSE/ESA exploitation
 - ▶ SSL enabled CICS Web Support
 - ▶ CICS External Call Interface
 - ▶ New VSAM Redirector
 - ▶ More samples (JConVSE, VSEPrint, etc.)
 - ▶ New JDBC Driver Layer for VSAM



e-business



Delete Label Function

- New function **DELLBL** in LABEL macro
- Must be explicitly exploited
 - ▶ Important for vendors with disk/tape management products
- Benefits
 - ▶ Saves recursive reads (GETNXGL) and write backs (ADDLBL/ADDNXL)
 - ▶ **Saves >90% of the SVCs** for this activity
- More Info
 - ▶ VSE Label Area -Layout and Capacity Consideratio
VSE/ESA Software Newsletter, 12/2000
 - ▶ <http://www-1.ibm.com/servers/eserver/zseries/os/vse/pdf/vsnew21/vseflab.pdf>



e-business



LTA Offload and SVA-24

- LTA Offload for some AR commands
 - ▶ Phases \$\$BATTNC and \$\$BATTNG are merged into \$\$BATTNA
 - ▶ Code of \$\$BATTNB is merged into IJBAR
 - ▶ Benefits
 - Less I/O by less FETCHes for LTA load
 - IGNORE, PAUSE, LOG, NOLOG, NEWVOL, START, BATCH
 - No LTA usage for MSG commands
- SVA-24 Phases moved above the line
 - ▶ \$IJBPRTY (6K)



e-business



SDL Entries

- Increased max. number of SDL entries
 - ▶ New IPL SVA parameter: $SDL=n$
 - Maximum value now **32765**
 - About 56 SDL entries per 4K page in shared space below
 - Theoretically would cost 2.28 MB
- SDL updates from non-BG partitions
 - ▶ SET SDL command can now be issued from any partition
 - ▶ Internal locking is done to assure correctness



e-business



Hardware Support

- New 2074 System Management Console
 - ▶ ESCON channel attached
 - ▶ Eliminates requirement for a non-SNA 3174 controller
- OSA Express Adapter Support
 - ▶ Available for G5 and above

	Gigabit Ethernet	Fast Ethernet 100 Mbps	ATM-LE 155 Mbps	Tokenring 4/16/100 Mbps
CHIPID TYPE=OSE (non-QDIO)	no	yes	yes	yes
CHIPID TYPE=OSD (QDIO)	yes	yes	yes	yes



e-business



Hardware Support

- Queued Direct I/O
 - ▶ Designed for very efficient exchange of data
 - ▶ Uses the QDIO Hardware Facility, without traditional S/390 I/O instructions
 - ▶ Without interrupts (in general)
 - ▶ Use of internal queues
 - ▶ With pre-defined buffers in memory for asynchronous use

- Exploitation by TCP/IP for VSE/ESA
 - ▶ see TCP/IP Performance Considerations



e-business



ESS Flashcopy

- The DASD Architecture of ESS allow copy of DASD's with the utility FlashCopy
 - ▶ The copy process takes a few seconds instead of hours
 - ▶ From Operating system view it is a real copy
 - ▶ From DASD controller view it is a virtual copy

- FlashCopy support is available for 3 VSE products
 - ▶ IXFP SNAP command
 - ▶ VSAM SYNONYM Backup
 - ▶ VSE/Fast Copy



e-business



ESS Flashcopy - continued

- Problems
 - ▶ Duplicate VOLIDs (DASD names) not allowed on a VSE system
 - ▶ Duplicate VSAM Catalog names not allowed on a VSE system
- FlashCopy of volume containing VSAM datasets would mean
 - ▶ duplicate VOLIDs
 - ▶ duplicate VSAM Catalog names
- Solution
 - ▶ IDCAMS SNAP command
 - Changes the VOLIDs of the copied volumes
 - ▶ IDCAMS SYNONYM BACKUP command
 - Uses a synonym list to access copied volumes



e-business



Virtual Tape Support

- Allows access to tape images residing in
 - ▶ A VSAM file (ESDS)
 - ▶ A remote file on a workstation or server
- Tape image is AWSTAPE format
 - ▶ Known from P/390, R/390 or FLEX/ES
- New VTAPE command
- Virtual Tape Simulator
 - ▶ Simulates channel program execution
- Virtual Tape Data Handler
 - ▶ Runs in a partition



e-business



Virtual Tape Support - continued

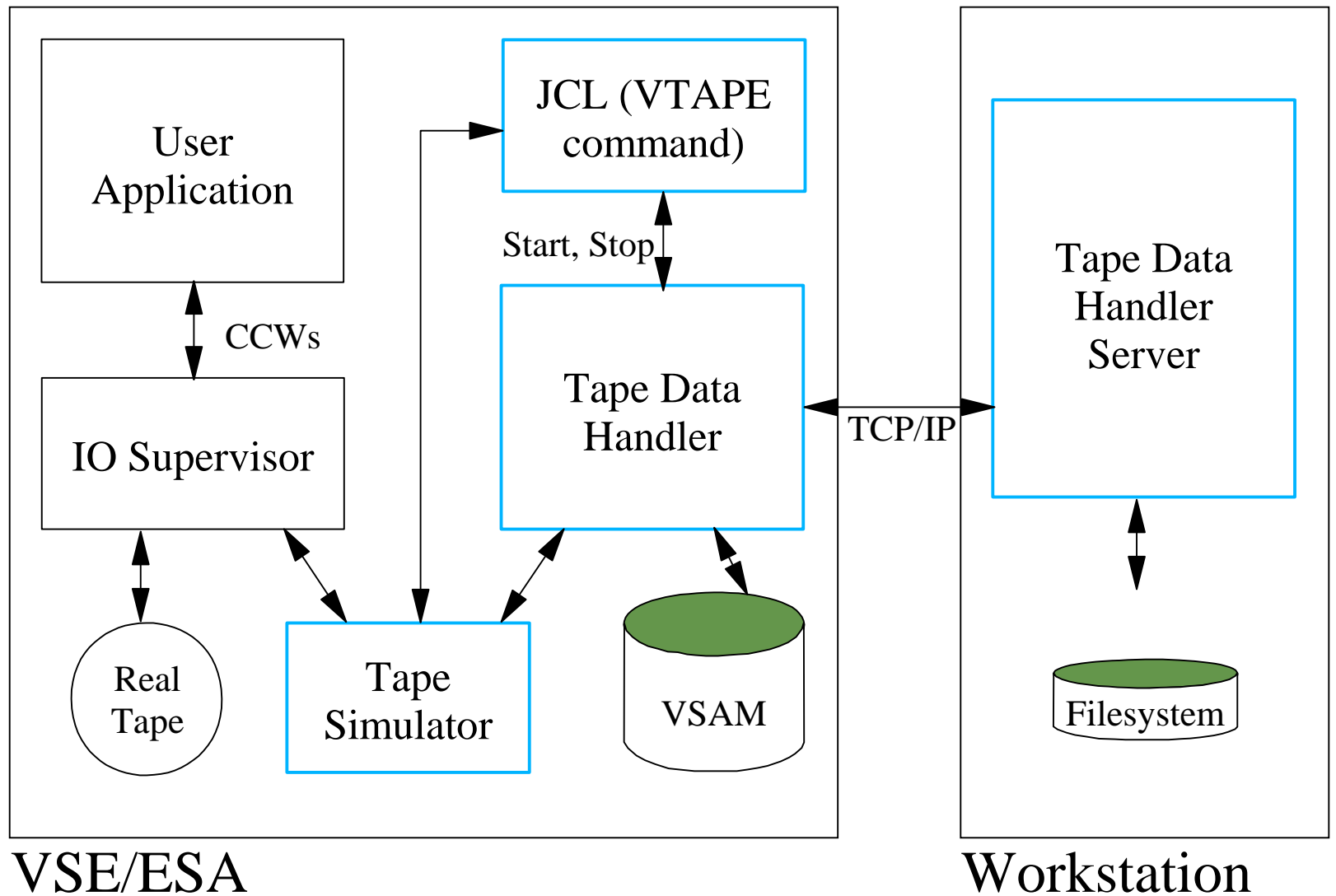
- Virtual Tape Server
 - ▶ Runs on a workstation or server (Java)
 - ▶ Allows to access a tape image remotely
 - ▶ Communicates via TCP/IP with Virtual Tape Data Handler
- Designed to allow e-Delivery and e-Service (future)
 - ▶ Download a tape image containing a product
 - ▶ Obtain a CD/DVD containing the tape image
 - ▶ Install the product via Virtual Tape directly from the workstation
- Also possible
 - ▶ Backup to a Virtual Tape + copy to CD
 - ▶ Restore directly from CD via Virtual Tape



e-business



Virtual Tape Support - continued





e-business



Updated Java-based Connector

- The Java-based connector has been updated to support the Java 2 platform (JDK 1.3)
- Introduced JDBC layer for VSAM access
 - ▶ Allows to issue SQL statements
- Adaptations for WebSphere 4.0
 - ▶ Enhanced connection pooling by support of JCA (Java Connector Architecture)
 - ▶ Connectors can be deployed as Resource Adapter and as (JDBC-) Data Source
- SSL enabled connections possible
 - ▶ Transparent use of secured connections



e-business



VSAM Share Options with Connectors

- SHROPT(4) Backgrounds
 - ▶ Using connectors to UPDATE a VSAM file already opened for output (e.g. by CICS) needs SHROPT(4)
 - ▶ SHROPT(4) has big overhead
- Performance implications
 - ▶ Bigger pathlength for processing of UPDATE requests due to VSAM internal locking
 - ▶ Each READ must be done from disk
 - ▶ Each WRITE must go to disk
 - ▶ Additional catalog I/Os for statistics
 - ▶ Influence on any application, not only connectors



e-business



VSAM SHROPT(4) Avoidance

- Connectors in VSE/ESA 2.5 require SHROPT(4) when updating VSAM files owned by CICS
- New VSAM-via-CICS Service avoids SHROPT(4) by routing the VSAM requests to CICS
- Communication between batch and CICS is XPCC
- New transactions related to VSAM-via-CICS:

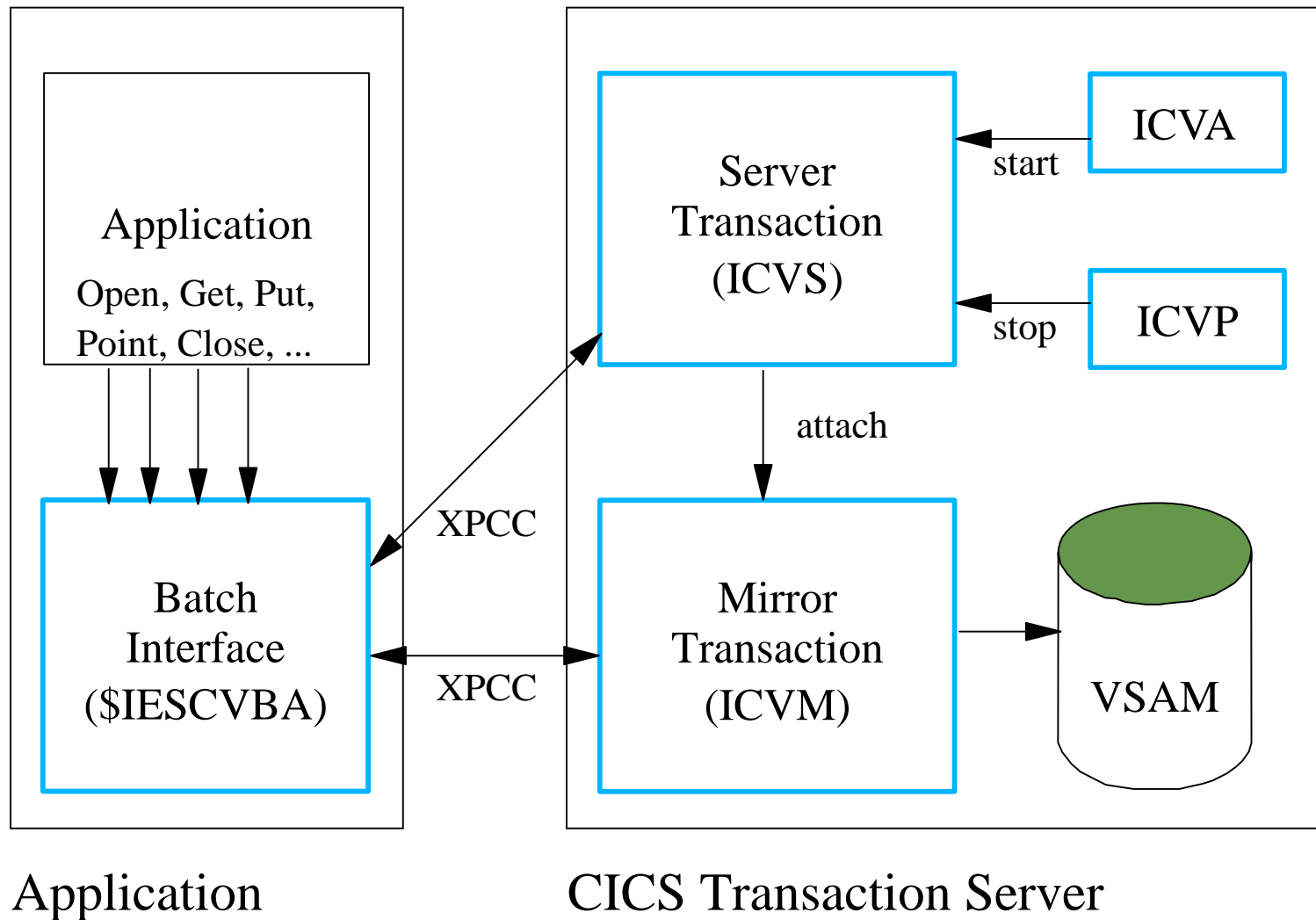
Transaction	Program	Description
ICVA	IESCVSTA	starts the service
ICVP	IESCVSTP	stops the service
ICVS	IESCVSRV	internal server task
ICVM	IESCVMIR	internal mirror task
none	IESCVSTI	internal start program



e-business



VSAM-via-CICS Service





e-business



VSAM-via-CICS Service - continued

- How VSAM-via-CICS works
 - ▶ Long running server transaction ICVS
 - ▶ Attaches a mirror transaction ICVM on request
 - ▶ Mirror transaction is attached for
 - "Open" from batch
 - Browse files from batch
 - ▶ Mirror transaction ends at "close" from batch
 - ▶ Service can run in multiple CICSes at the same time
 - ▶ Batch counterpart is implemented in phase \$IESCVBA



e-business



VSAM-via-CICS Service - continued

- Naming convention for "VSAM-via-CICS files"
 - ▶ Each CICS is treated as "virtual" catalog
 - ▶ Files defined in CICS (via CEDA DEFINE FILE) are visible within this catalog
 - ▶ "Virtual" catalog file id

#VSAM.#CICS.<applid>

indicates "virtual"
CICS catalog

APPLID of CICS region
owning the files within this
catalog

- "Virtual" cluster file id is the 7 character name known in CICS



e-business



VSAM-via-CICS Service - continued

- Example
 - ▶ Assume there is a CICS region DBDCCICS
 - ▶ CICS knows a file named MYFILE
 - ▶ Real VSAM files MY.VSAM.TEST.FILE resides in catalog MY.USER.CATALOG
 - ▶ "Batch only" name would be
 - **Catalog: MY.USER.CATALOG**
 - **Cluster: MY.VSAM.TEST.FILE**
 - ▶ "VSAM-via-CICS" name would be
 - **Catalog: #VSAM.#CICS.DBDCCICS**
 - **Cluster: MYFILE**



e-business



VSAM-via-CICS Service - continued

- VSAM-via-CICS files can only be accessed from the following applications
 - ▶ Java-based connector via VSE Java Beans
 - ▶ DB2-based connector via VSAM CLI (SQL)
 - ▶ REXX new VSAMIO function
- IDCAMS does NOT show these files
- "Virtual" names can NOT be specified in DLBLs
- No changes made in VSAM for this support
- No influence on "normal" VSAM processing
- But: Maps can be defined for a "virtual" file
 - ▶ Via Java-Based connectors
 - ▶ Via IDCAMS RECMAP function



e-business



VSAM Redirector

- New connector
 - ▶ VSE is client
 - ▶ PC / workstation is server
- Exploits VSAM exit IKQVEX01
- Allows to redirect one or more VSAM files to a PC workstation
- All VSAM requests of a particular file are redirected
 - ▶ Open / close
 - ▶ Get / put / point / delete / insert
- Transparent for applications
 - ▶ Usable from batch and CICS



e-business



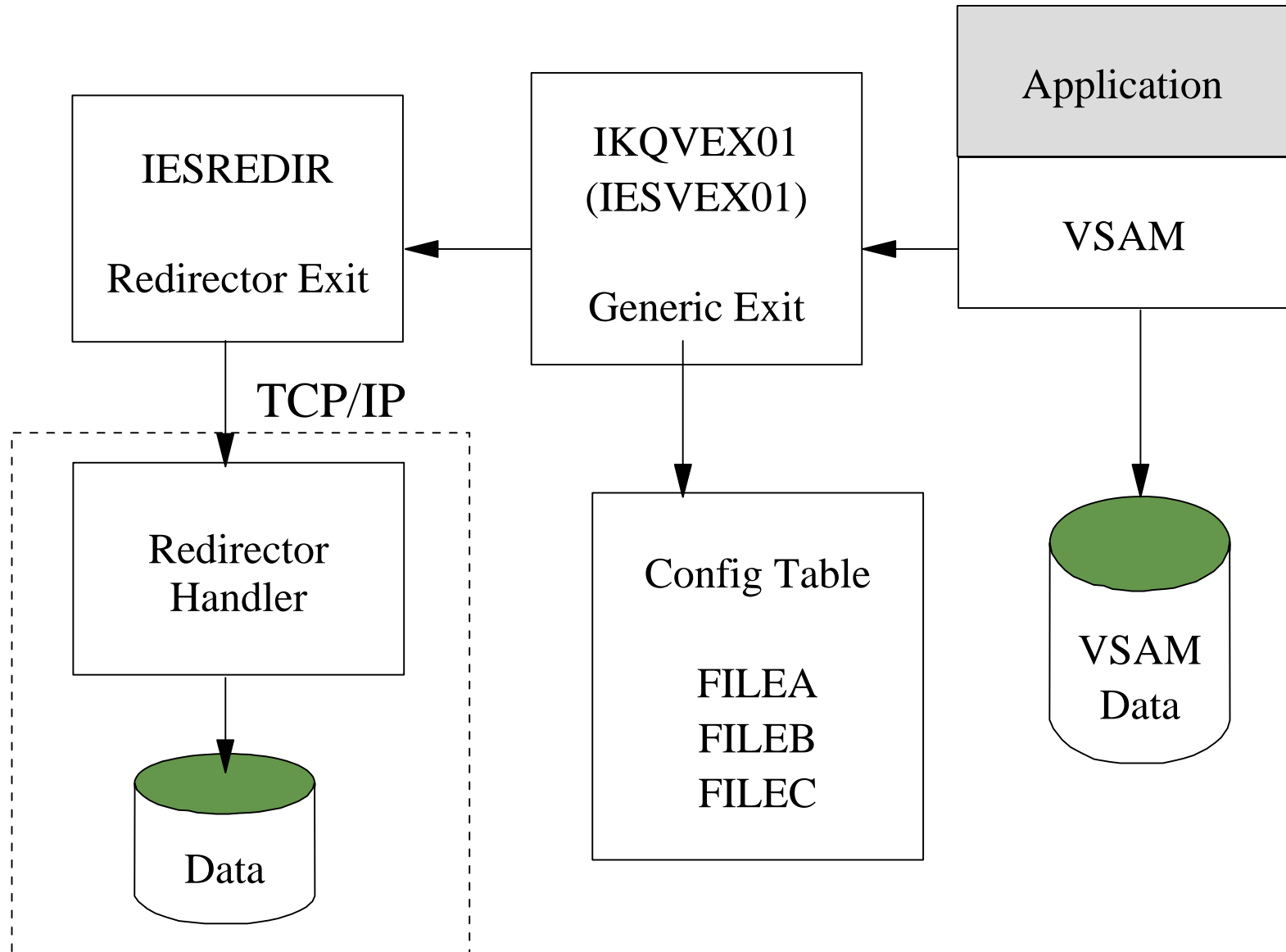
VSAM Redirector - continued

- Owner of data can be
 - ▶ VSAM
 - Requests are forwarded to workstation
 - VSAM still owns the data
 - VSAM executes the requests
 - Used for data replication/synchronisation
 - ▶ PC / workstations
 - VSAM does not execute the requests
 - Handler on workstation 'simulates' VSAM logic
 - A VSAM file with at least one dummy record is required (for open processing)



VSAM Redirector - continued

e-business





e-business



VSAM Redirector - continued

- Decision if a file is redirected or not is
 - ▶ Done at open time
 - ▶ Based on the config table (PHASE)
 - Catalog id and file id
 - ▶ Only a very small (open-)overhead for non-redirected files
 - ▶ No overhead for get/put/... if not redirected
- Generic Exit can also call a 'Vendor' exit instead of Redirector exit
 - ▶ Defined in the config table
 - ▶ Based on catalog id and file id



e-business



VSAM Redirector - Performance Implications

- Is the file redirected ?
 - ▶ No: only at OPEN time (very small overhead)
 - ▶ Yes: at each request
- Network overhead ?
 - ▶ Yes, if file is redirected
 - ▶ Depends on
 - Number of VSAM requests
 - Size of records
- Data ownership
 - ▶ OWNER=REDIR
 - no VSAM I/O



e-business



CICS TS Enhancements

- CICS Web Support
 - ▶ new: SSL enabled (https)

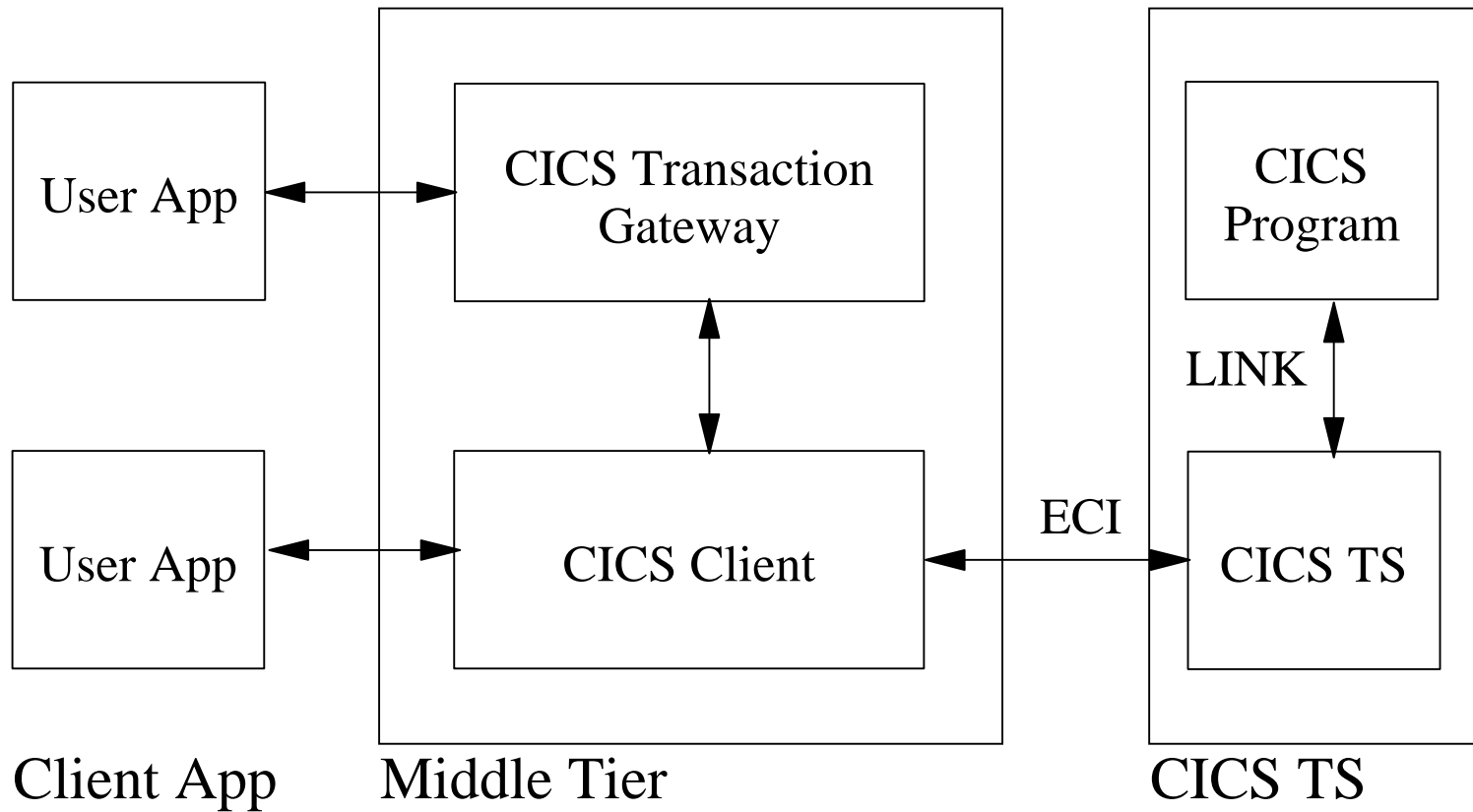
- External Call Interface (ECI)
 - ▶ Call a CICS program from a workstation
 - ▶ Prerequisites
 - CICS Client
 - CICS Transaction Gateway



e-business



External Call Interface (ECI)





e-business



General Performance Hints for Connectors

- Reduce amount of data transferred
 - ▶ Transfer only data that is needed
 - ▶ Issue only requests that are needed
- Use connection pooling
 - ▶ Reduce overhead of connection establishment
- Performance of connectors depends on
 - ▶ Network performance
 - ▶ Performance of "server"
 - ▶ Performance of "client" or middle tier



e-business



VSE/ESA 2.7 Performance Items

- VSE/ESA 2.7 hardware support
 - ▶ z800/z900, Multiprise 3000, G5/G6
 - ▶ HiperSockets
 - ▶ Hardware Crypto Support
 - ▶ 32760 cylinder 3390 support
 - ▶ 3590 buffered tape mark

- VSE/ESA 2.7 enhancements
 - ▶ New TCP/IP for VSE/ESA release 1.5
 - ▶ \$IJBLBR above the line
 - ▶ II User Status Record above the line
 - ▶ VTAPE: removed DVCDN/DVCUP
 - ▶ POWER: reallocate queue file during warm start



e-business



VSE/ESA 2.7 Hardware support

- VSE/ESA 2.7 runs on the following machines
 - ▶ z800 (2066)
 - ▶ z900 (2064)
 - ▶ 9672 Parallel Enterprise Server (G5/G6)
 - ▶ Multiprise 3000 (7060)
 - ▶ equivalent emulators (Flex-ES)
- VSE/ESA 2.7 is based on the hardware instruction set described in the manual 'ESA/390 Principles of Operation' (SA22-7201).
- With VSE/ESA 2.7 it is assumed that all the ESA/390 instructions and facilities described in that manual can be used.



e-business



z800/z900 Remarks

- Prior to z800/z900 there is one cache for data and instructions
- z800/z900 has split data and instruction cache
- Performance implications:
 - ▶ If program variables and code that updates these program variables are in the same cache line (256 byte)
 - Update of program variable invalidates instruction cache
 - Performance decrease if update is done in a loop
 - ▶ See APAR PQ66981 for FORTRAN compiler



e-business



32760 cylinder 3390 support

- With announcement 101-341 at 11/13/2001 IBM announced the new 32760 cylinder 3390 volumes of the IBM TotalStorage Enterprise Storage Server (ESS)
 - ▶ This enhancement of the ESS F models was made available 11/30/2001
- VSE/ESA 2.7 now supports these volumes
 - ▶ helps relieve address constraints
 - ▶ improves the disk resource utilization
 - ▶ can be used to consolidate multiple disk volumes into a single address
- VSAM can only address 10017 cylinders.



e-business



3590 Buffered Tape Mark support

- The 3590 control unit provides support for writing tape marks (TM) in buffered mode
- Writing TM's in "buffered" mode should enhance the performance
 - ▶ of all programs which write many TM's as part of their file creation process (e.g. POFFLOAD)
- All the TM's written during OPEN/CLOSE (label processing) will remain to be written "UNbuffered"
 - ▶ all the programs which write TM's mainly or only during OPEN/CLOSE will NOT benefit from this enhancement



e-business



\$IJBLBR phase moved above the line

- The \$IJBLBR.PHASE has been split into two phases
 - ▶ \$IJBLBR.PHASE
 - ▶ \$IJBLB31.PHASE
- \$IJBLBR.PHASE will continue to reside in SVA-24
- \$IJBLB31.PHASE will reside in SVA-ANY (high SVA)
 - ▶ This will free about 180KB in SVA-24



e-business



II User status record above the line

- During Logon each II user gets besides others two storage areas allocated
 - ▶ User_Status_Record USR (904 bytes)
 - ▶ Panel_Hierarchy_List PHL (1352 bytes)
 - ▶ originally located in the CICS DSA (below)
- With VSE/ESA 2.7 the USR and PHL has been moved to ESDSA (shared above)
 - ▶ frees 2.3 KB in DSA below per user
- ICCF TCTUALOC=ANY now supported
 - ▶ ICCF transaction programs has been changed to support a TCTUA (28 bytes) above the line



e-business



HiperSockets hardware elements (**'Network in a box'**)

- Synchronous data movement between LPARs and virtual servers within a zSeries server
 - ▶ Provides up to 4 "internal LANs" HiperSockets accessible by all LPARs and virtual servers
 - ▶ Up to 1024 devices across all 4 HiperSockets
 - ▶ Up to 4000 IP addresses
 - ▶ Similar to cross-address-space memory move using memory bus
- Extends OSA-Express QDIO support
 - ▶ LAN media and IP layer functionality (internal QDIO = iQDIO)
 - ▶ Enhanced Signal Adapter (SIGA) instruction
 - No use of System Assist Processor (SAP)



e-business



HiperSockets hardware elements (**'Network in a box'**) - continued

- HiperSockets hardware I/O configuration with new CHPID type = IQD
 - ▶ Controlled like regular CHPID
 - ▶ Each CHPID has configurable Maximum Frame Size
- Works with both standard and IFL CPs
- No physical media constraint, no physical cabling, no priority queuing
- Secure connections



e-business



Measurement Environment

- z800 (2066-004)
 - ▶ 4 processors
- VSE/ESA 2.7 GA Driver in an LPAR (native)
 - ▶ 1 CPU active (~2066-001)
 - ▶ TCPIP00 (F7): OSA Express Fast Ethernet
 - ▶ TCPIP01 (F8): HiperSockets
- Linux for zSeries in an LPAR (native)
 - ▶ 3 CPUs active (shared)
 - ▶ eth0: OSA Express Fast Ethernet
 - ▶ hsi10: HiperSockets



e-business



Latency (Round trip time) - results

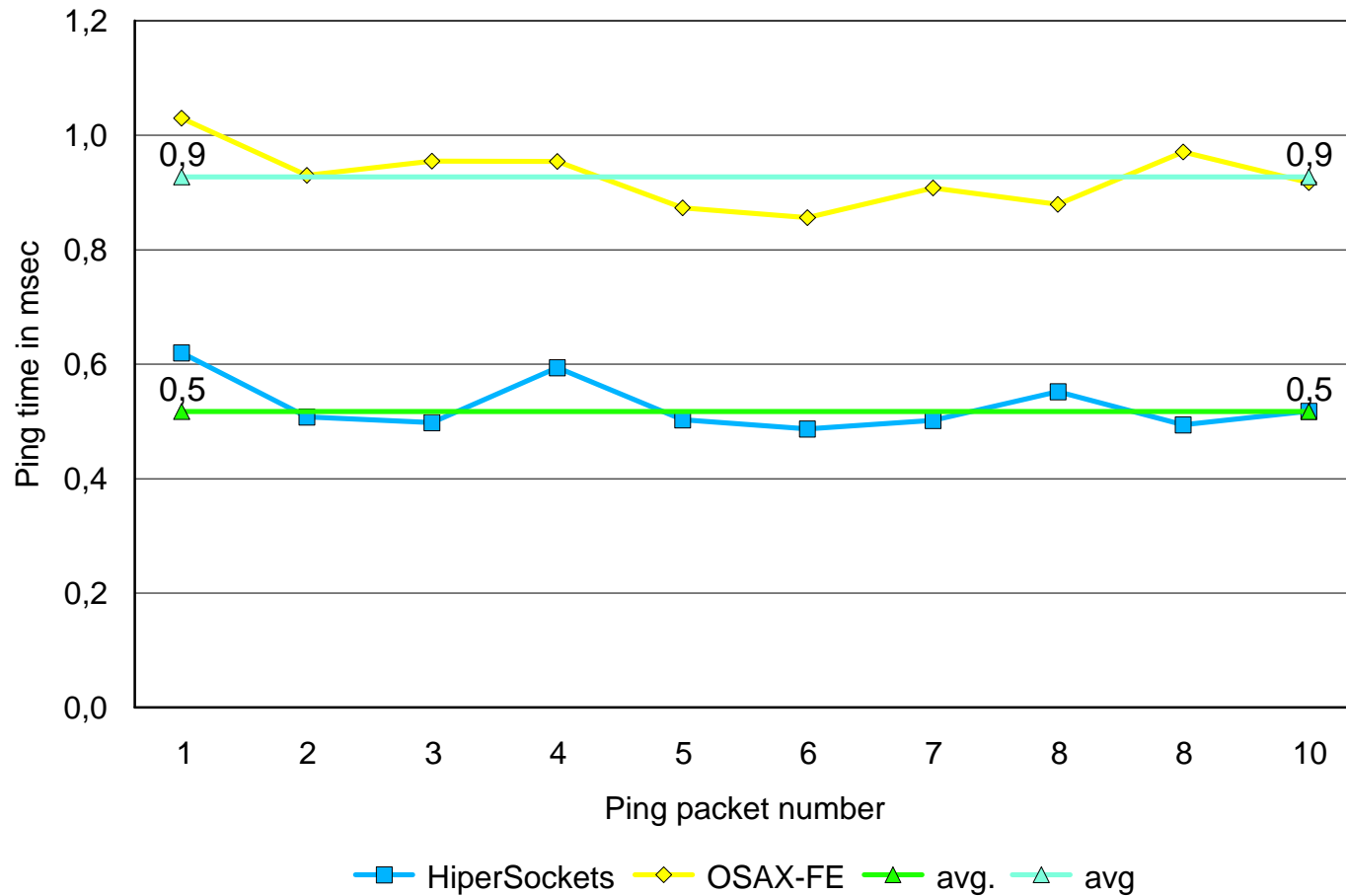
- Measurements has been done with PING command
 - ▶ Issued at Linux side
 - ▶ 10 Pings
 - ▶ PING sends a datagram to VSE
 - ▶ VSE sends an answer back to Linux
 - ▶ Time until answer arrives is measured
 - Round trip time



e-business



Latency (Round trip time) - results



HiperSockets is about 1.8 times faster in terms of latency



e-business



Throughput (MB/sec)

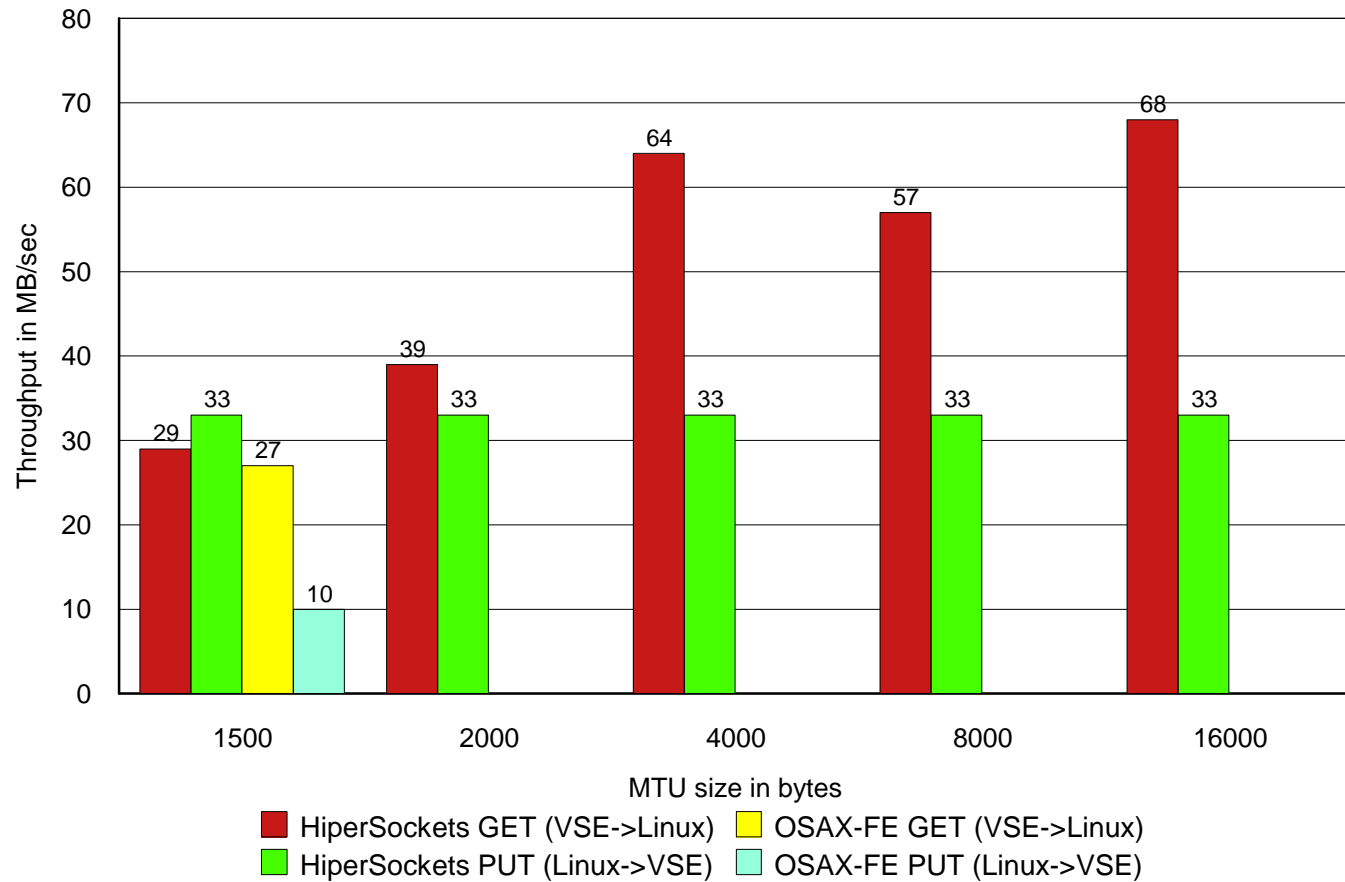
- Measurements has been done with FTP
 - ▶ Initiated at the Linux side
 - ▶ Transferring 1GB (1000MB)
 - without translation (binary)
 - 1 to 5 parallel streams
 - ▶ PUT: send data to VSE
 - VSE inbound
 - sending a 1GB file to \$NULL file (in memory file)
 - No file I/O is done by VSE/Linux
 - ▶ GET: receive data from VSE
 - VSE outbound
 - receiving \$NULL file (in memory file) into /dev/null
 - No file I/O is done by VSE/Linux



e-business



Throughput (MB/sec) - results



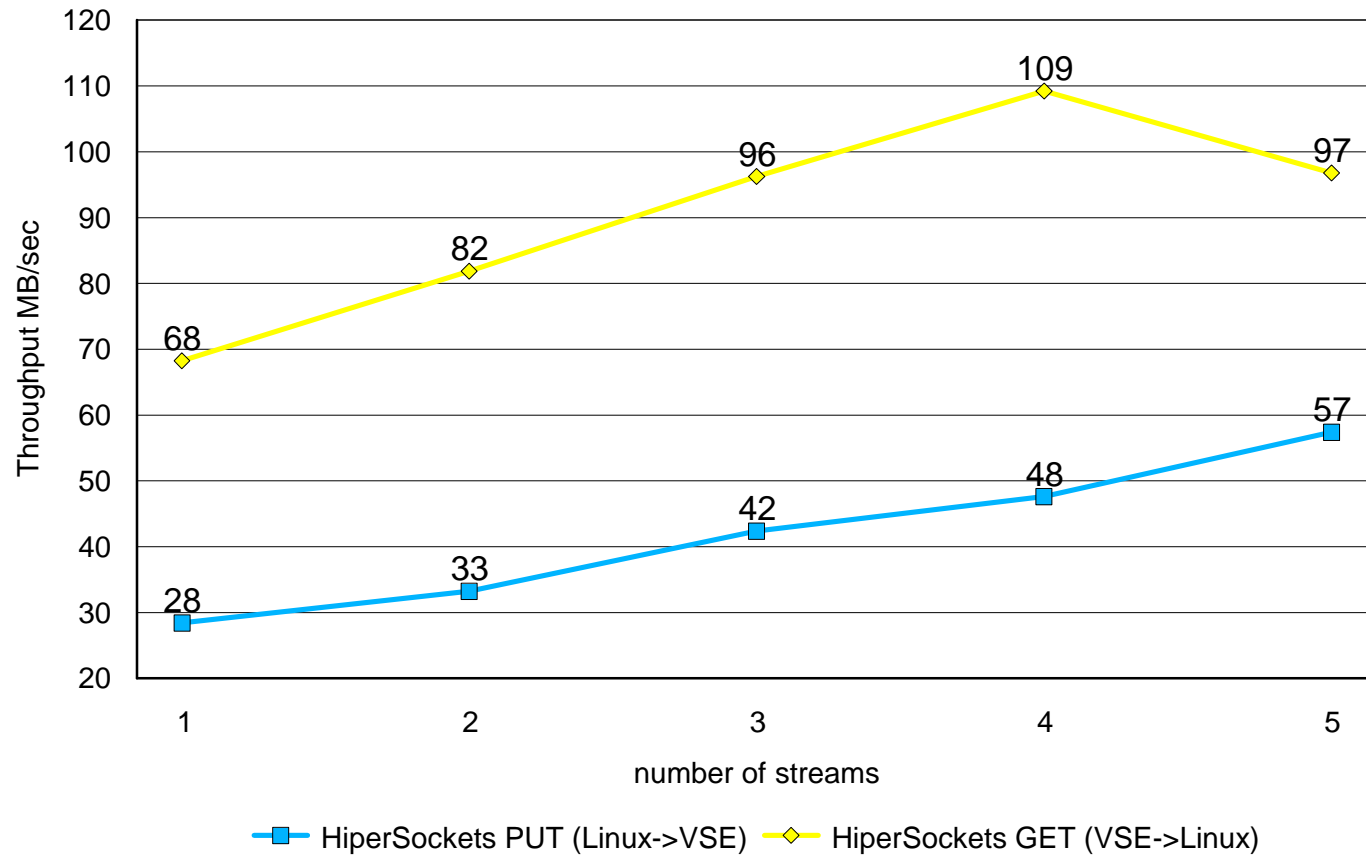
HiperSockets throughput is between 30-80 MB/sec



e-business



Throughput (MB/sec) - results (2)



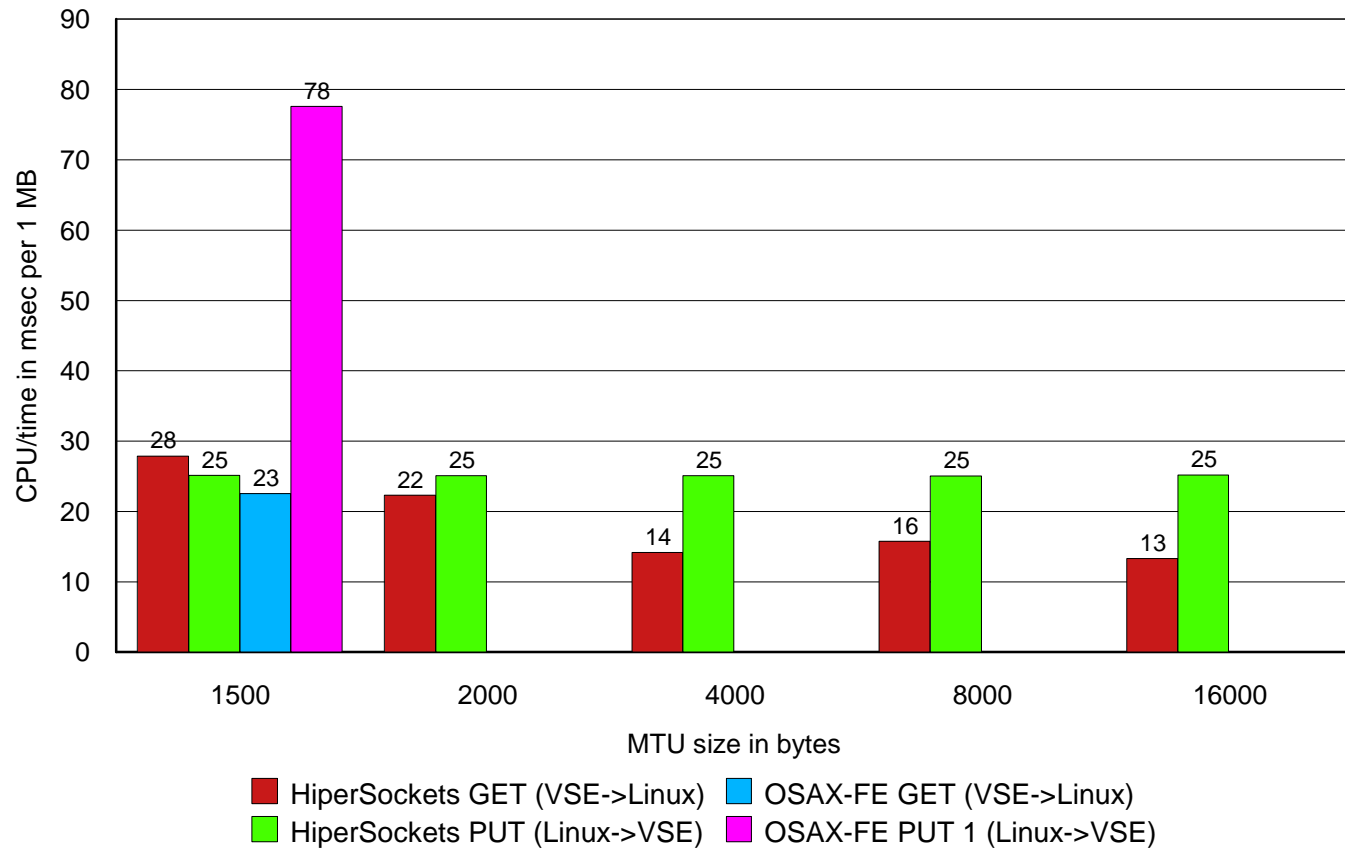
Maximum HiperSockets throughput of 109 MB/sec at 4 concurrent connections



e-business



CPU time per MB - results



About 15-30 msec CPU time per MB for HiperSockets
(on a z800 2066-001)



e-business



Transaction per second

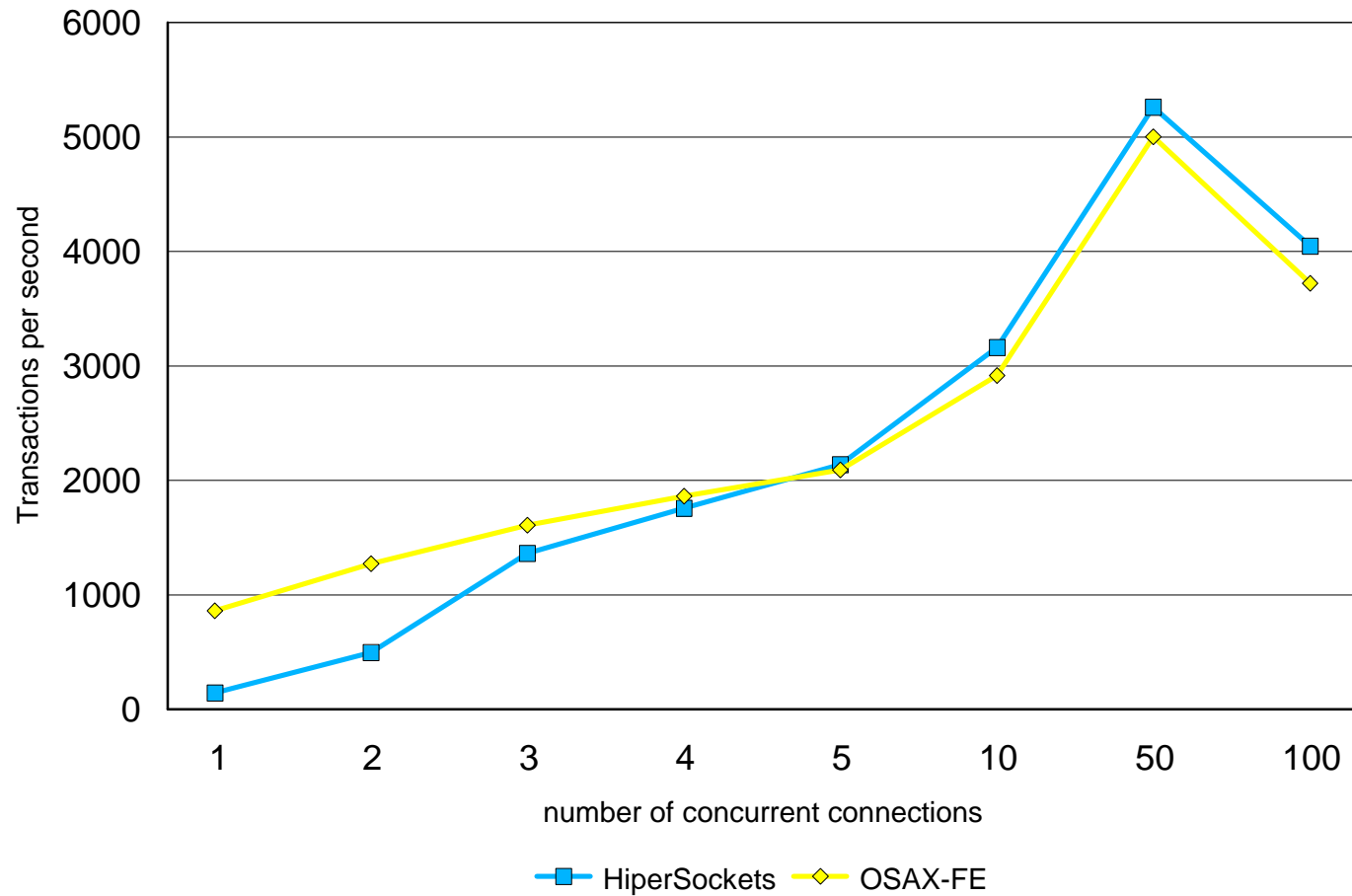
- Measurements has been done with an ECHO server
 - ▶ Client on Linux sends 100 bytes to server
 - ▶ Server on VSE echoes 100 bytes
 - ▶ Per TCP connection 10000 transactions are driven
 - ▶ Variations: Number of TCP connections
 - 1,2,3,4,5
 - 10,50,100
 - ▶ Measurements
 - Transactions per second
 - CPU time per transaction



e-business



Transactions per second - results



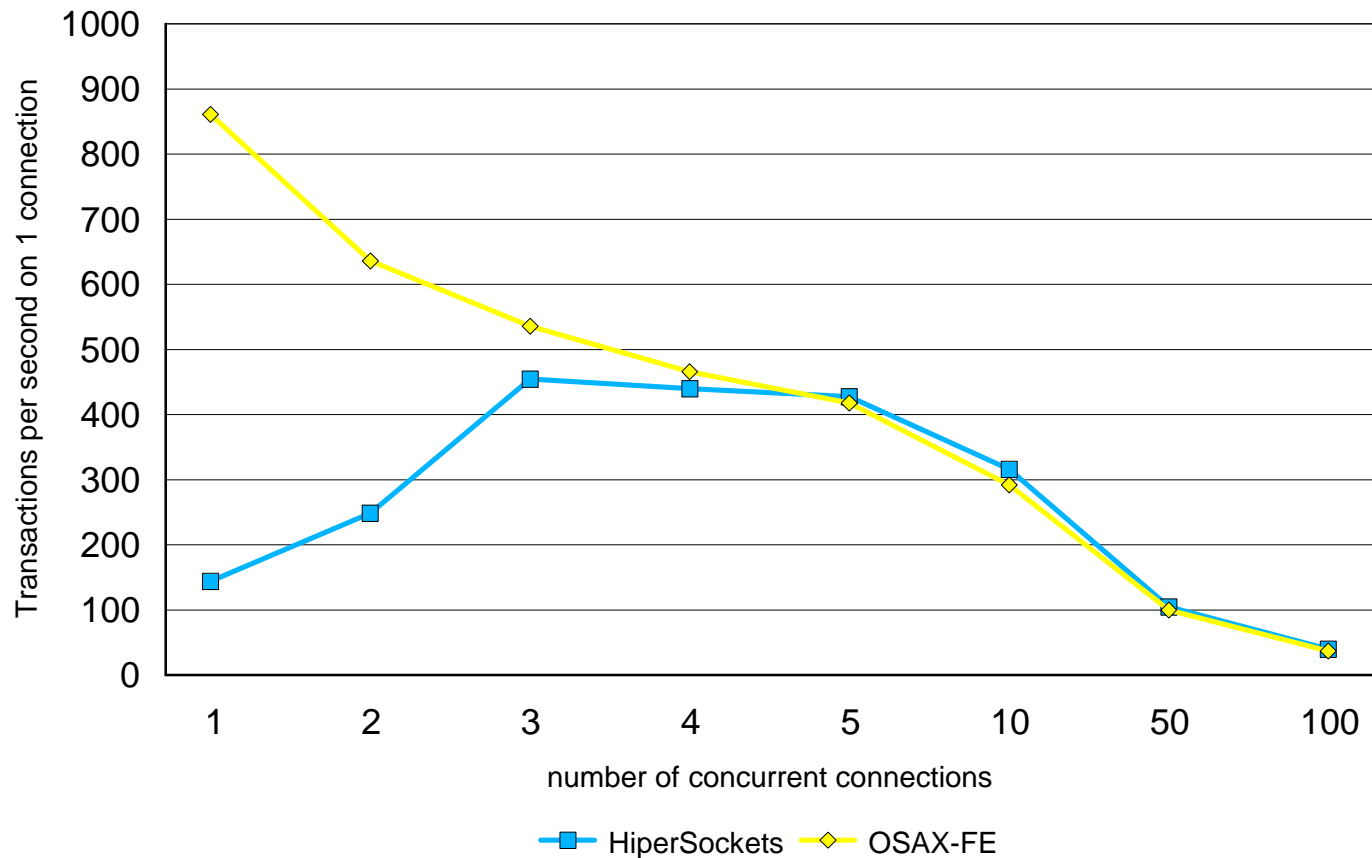
Maximum of 5200 transactions per second at 50 concurrent connections



e-business



Transactions per second on 1 connection - results



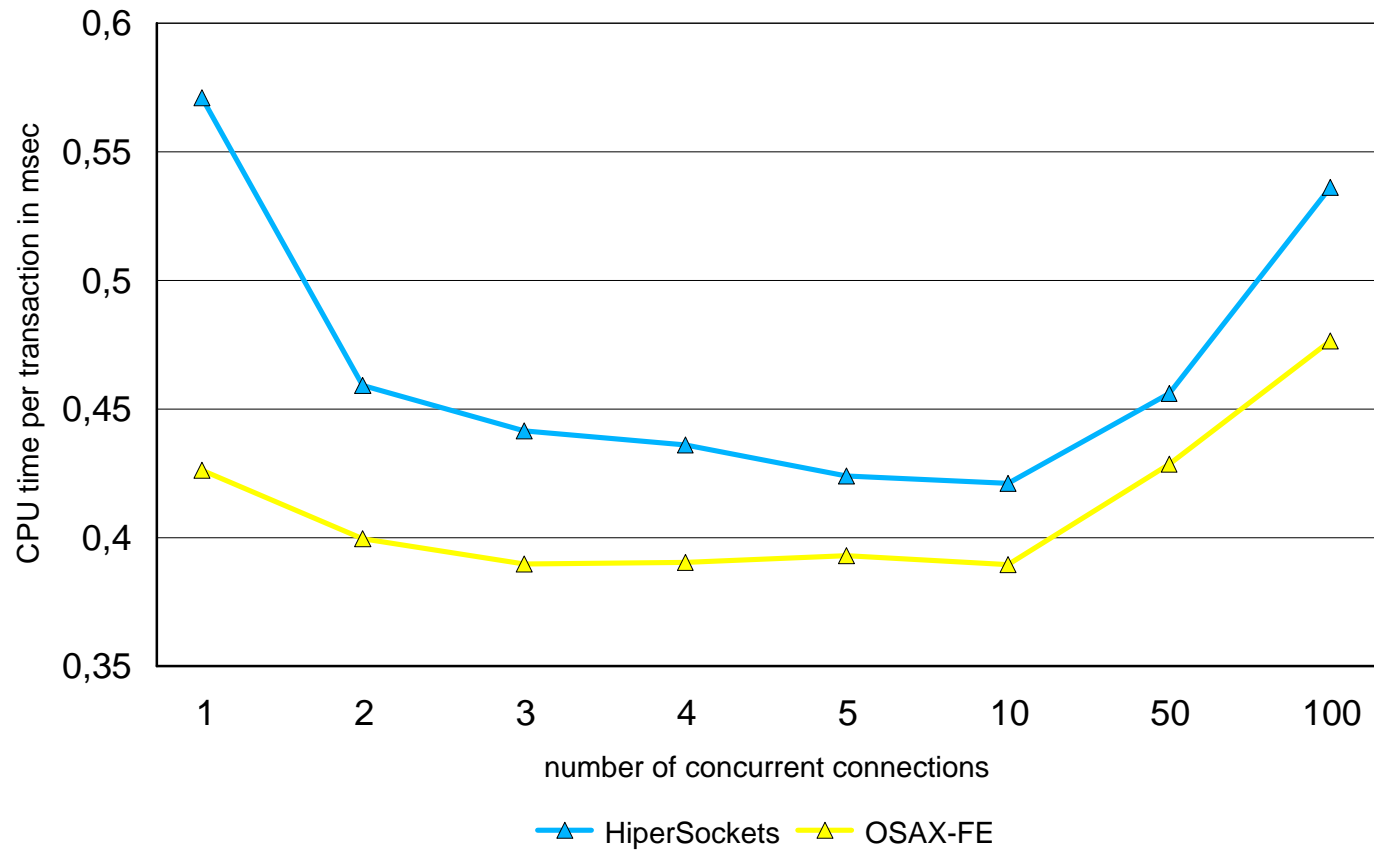
HiperSockets: Maximum of about 450 transactions per second on 1 connection (= about 2 msec response time)



e-business



CPU time per transaction



HiperSockets: About 0.45 msec CPU time per transaction for 2-50 connections



e-business



Measurement Results - conclusion

- HiperSockets
 - ▶ Throughput
 - Between 30-80 MB/sec
 - Maximum throughput of 109 MB at 4 concurrent connections
 - About 15-30 msec CPU time per MB
 - ▶ Transactions per second
 - Maximum of 5200 Transactions per second at 50 concurrent connections
 - About 0.4-0.45 msec CPU time per transaction



e-business



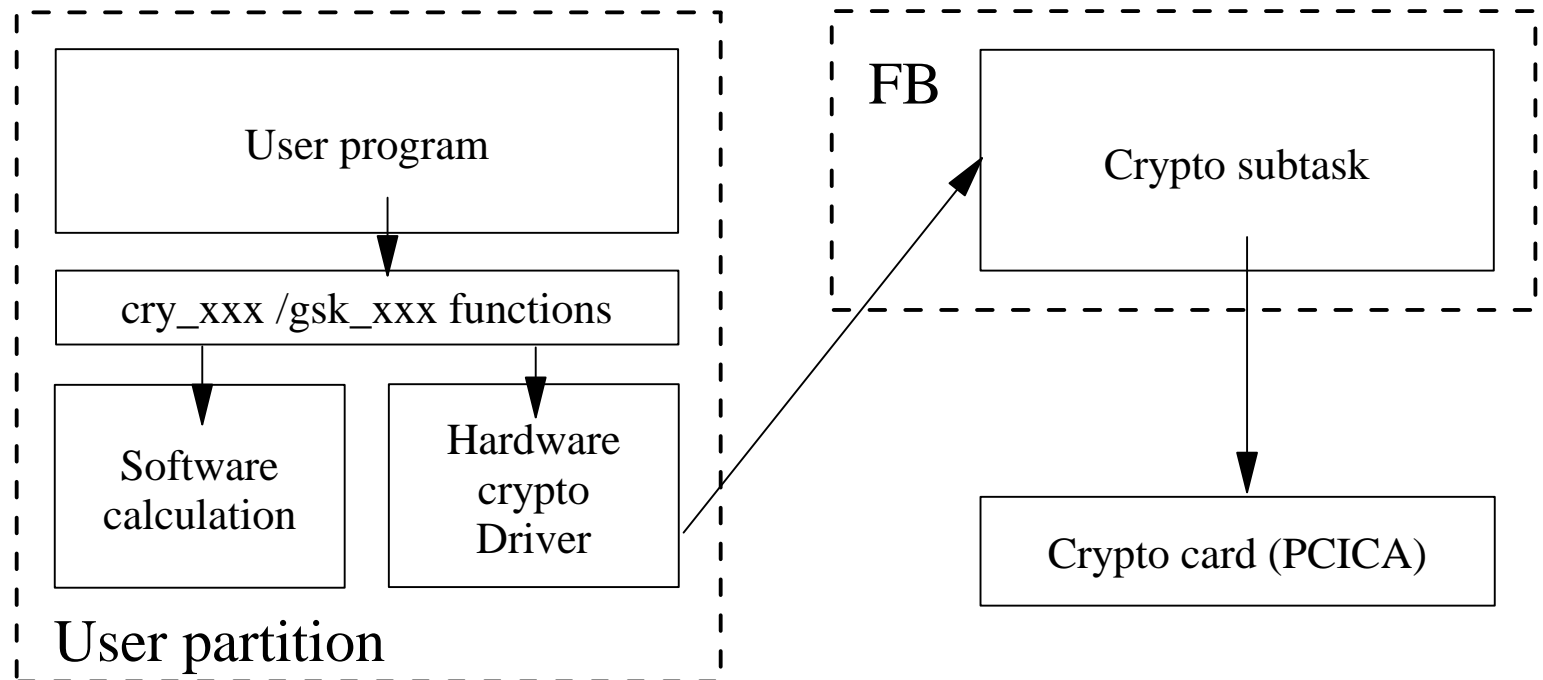
Hardware Crypto Overview

- Requires VSE/ESA 2.7 and TCP/IP for VSE/ESA 1.5
- Supported crypto cards
 - ▶ PCI Cryptographic Accelerator (PCICA)
 - Feature code 0862
 - Available for zSeries (z800, z900)
- Only RSA (asymmetric) is supported
 - ▶ Of benefit for Session initiation (SSL-Handshake)
- Also supported with
 - ▶ z/VM 4.2 + APAR VM62905
 - ▶ z/VM 4.3



Hardware Crypto Overview - continued

- New crypto subtask in Security Server (SECSECV) running in FB
 - ▶ Or as separate job if no SECSECV is running
 - ▶ Crypto card is polled by crypto task





e-business



Measurement Environment

- VSE/ESA 2.7 running on a z900 (2064-109)
 - ▶ on 1 processor (~2064-101)
 - ▶ with a PCI Cryptographic Accelerator
- Testcase programs on VSE
 - ▶ Crypto operations measurements
 - calling cry_xxx functions (RSA, DES, SHA, MD5)
 - each crypto operation is performed 10000 times
 - ▶ Secured data transfer (SSL)
 - performs SSL handshake
 - performs encrypted data transfer
 - counterpart program running on Windows (SSL-client)
- All RSA operations are measured
 - ▶ with Hardware Crypto support
 - ▶ with Software Crypto
 - support already available with TCP/IP 1.4/1.5 as shipped in VSE/ESA 2.6



e-business



Measurement Environment - continued

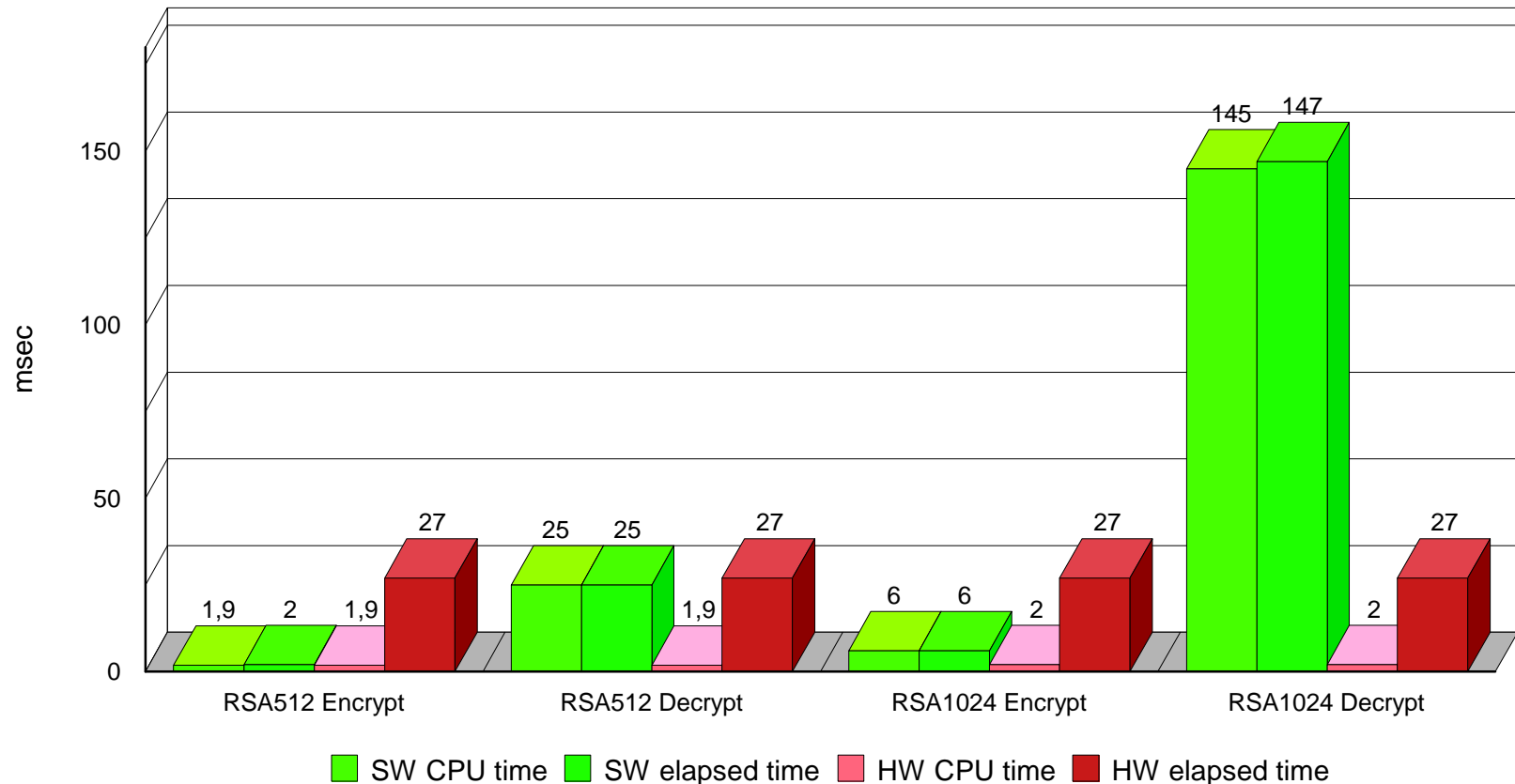
- Variations
 - ▶ RSA encrypt/decrypt
 - 512 / 1024 bit key
 - ▶ DES, DES CBC, 3DES CBC encrypt/decrypt
 - software crypto only
 - message length (128, 256, 512 bytes)
 - ▶ SHA Hash, MD5 Hash, SHA HMAC, MD5 HMAC
 - software crypto only
 - message length (128, 256, 512, 1K, 2K bytes)
 - ▶ SSL handshake/data transfer
 - 01 RSA512_NULL_MD5
 - 02 RSA512_NULL_SHA
 - 08 RSA512_DES40CBC_SHA
 - 09 RSA1024_DES_CBC_SHA
 - 0A RSA1024_3DES_EDE_CBC_SHA



e-business



Measurements Results - RSA



HW Crypto:

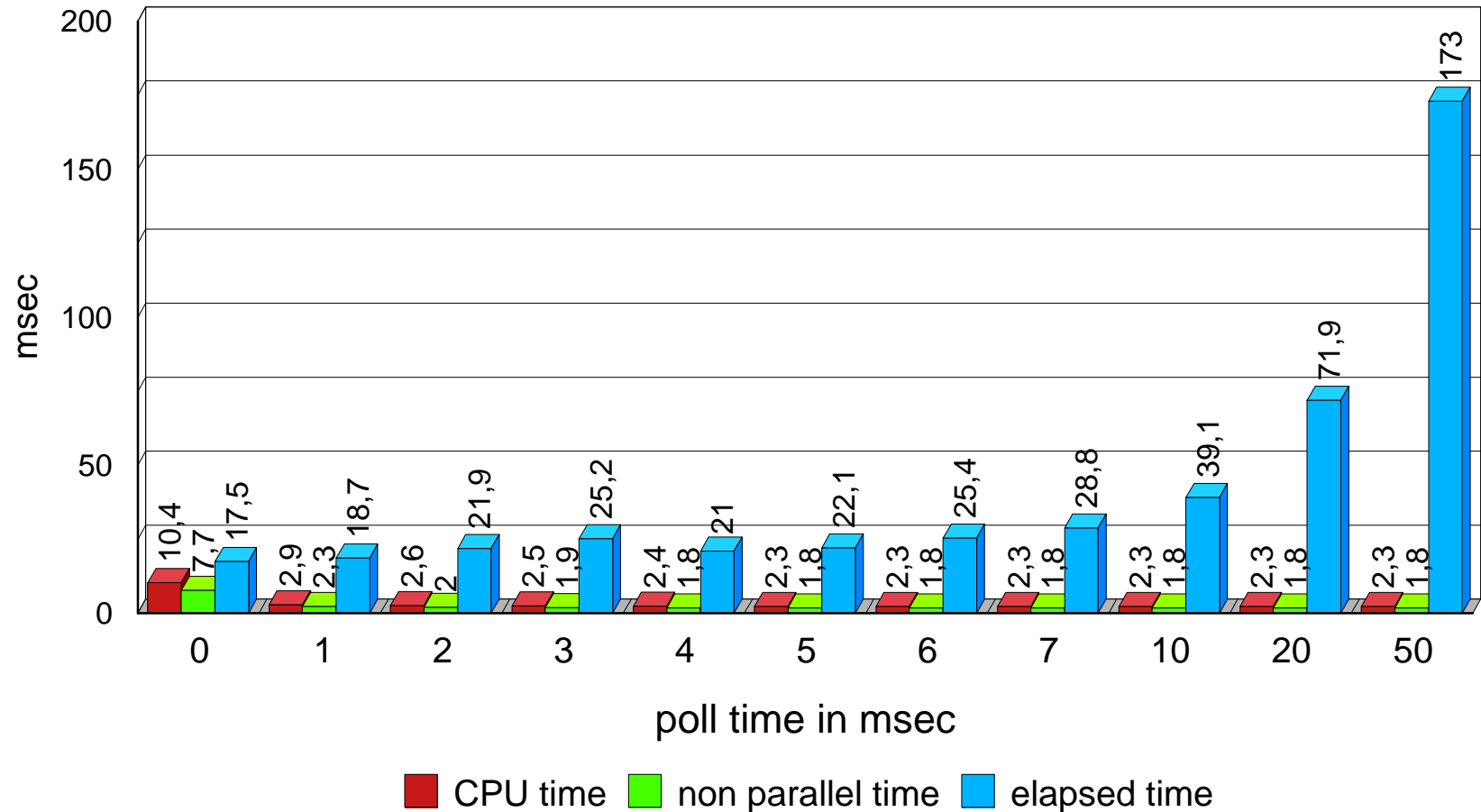
- CPU time and elapsed time is independent of operation / key length
- RSA operation takes about 2 msec CPU time and 28 msec elapsed time
- CPU time is always less than software crypto



e-business



Measurements Results - RSA polltime



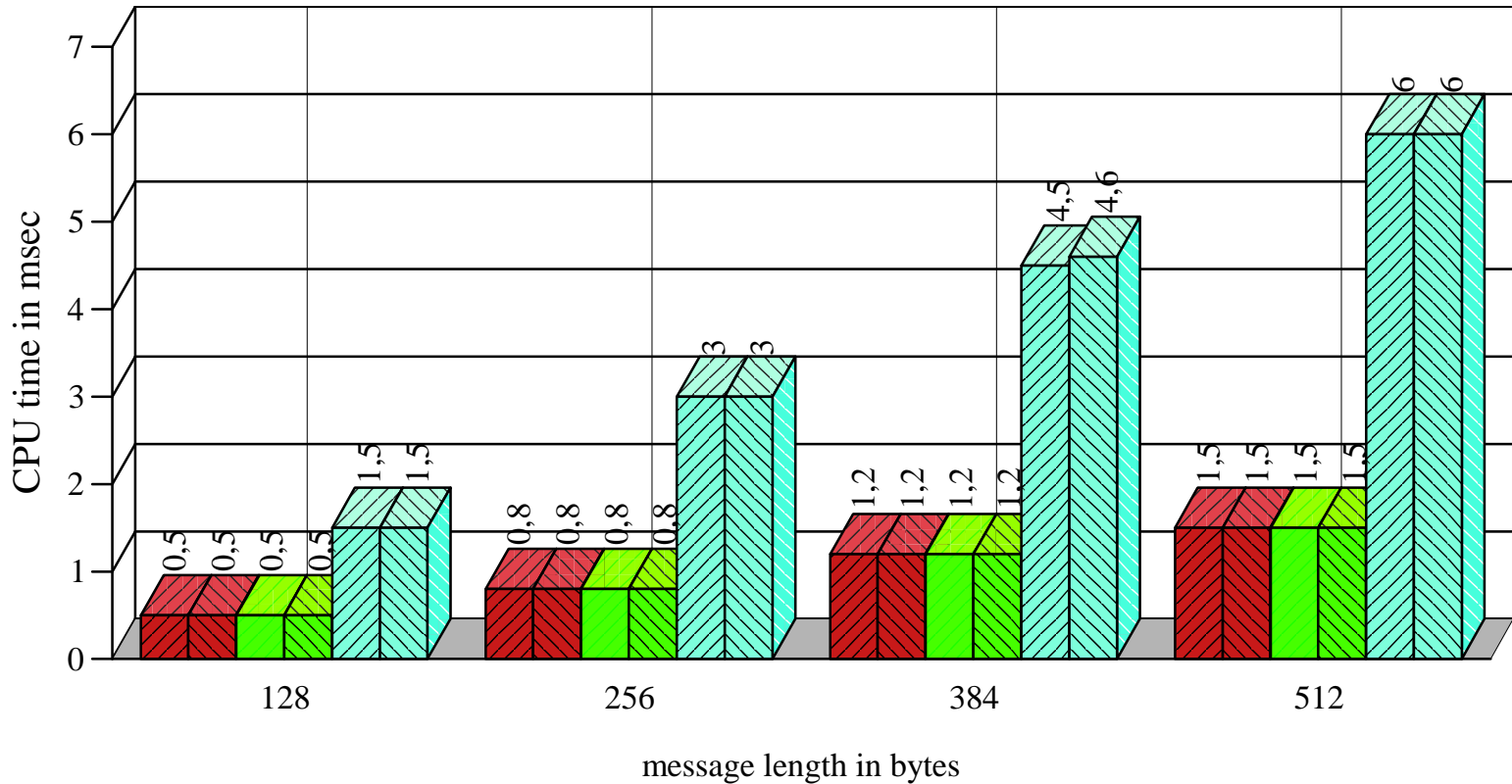
Per default a polltime of 7 msec is used.

Can be changed with: `MSG FB,DATA=WAITTIME=nn`

Smaller values increase CPU time, higher values increase elapsed time



Measurements Results - DES, DES CBC, 3DES CBC (symmetric)



■ DES Encrypt ■ DES CBC Encrypt ■ 3DES CBC Encrypt
■ DES Decrypt ■ DES CBC Decrypt ■ 3DES CBC Decrypt

Software Crypto only!

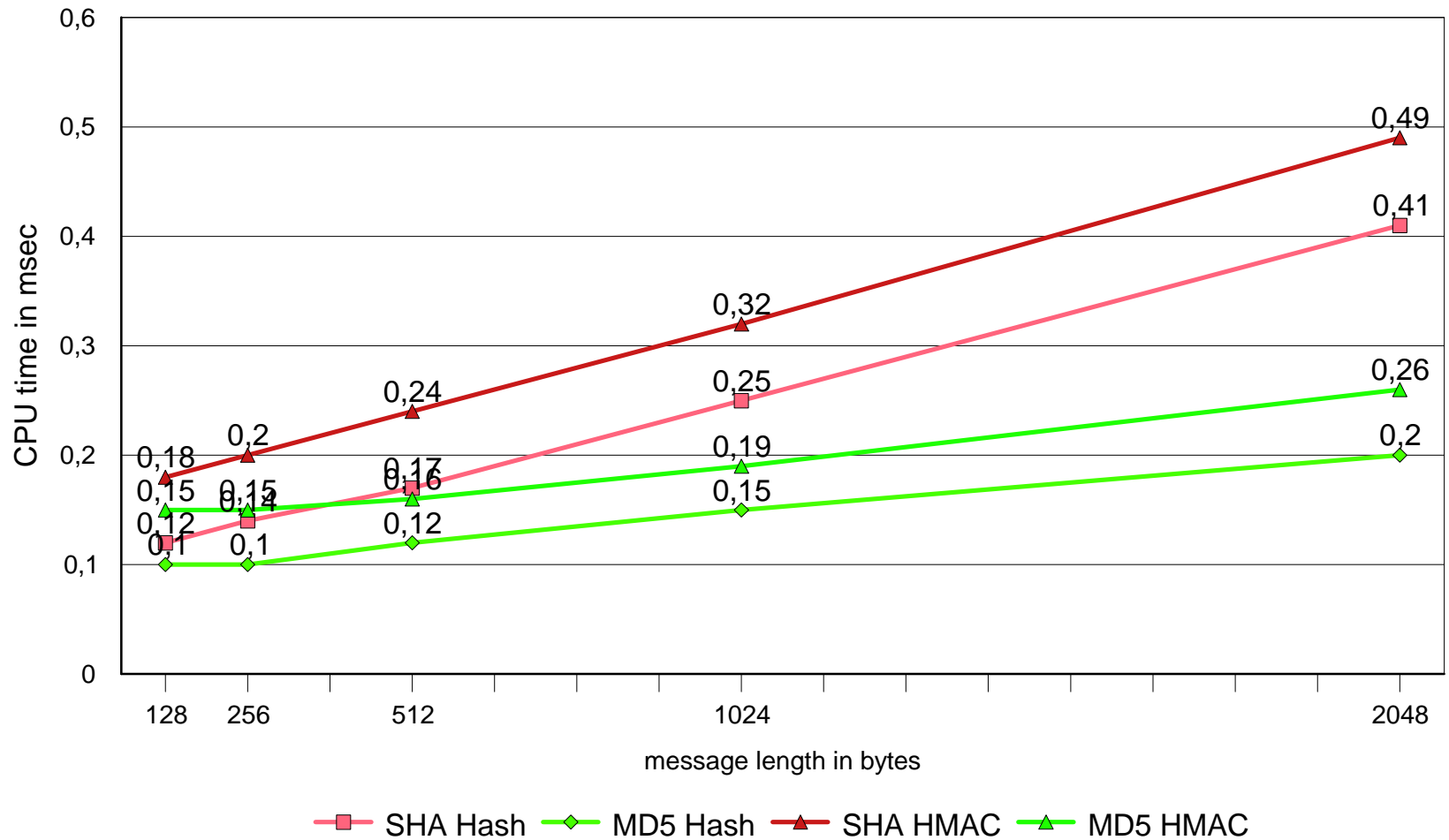
DES and DES CBC takes similar CPU times, 3DES CBC about 3.8 times



e-business



Measurements Results - SHA, MD5



SHA takes about 1.8 times more CPU time compared to MD5

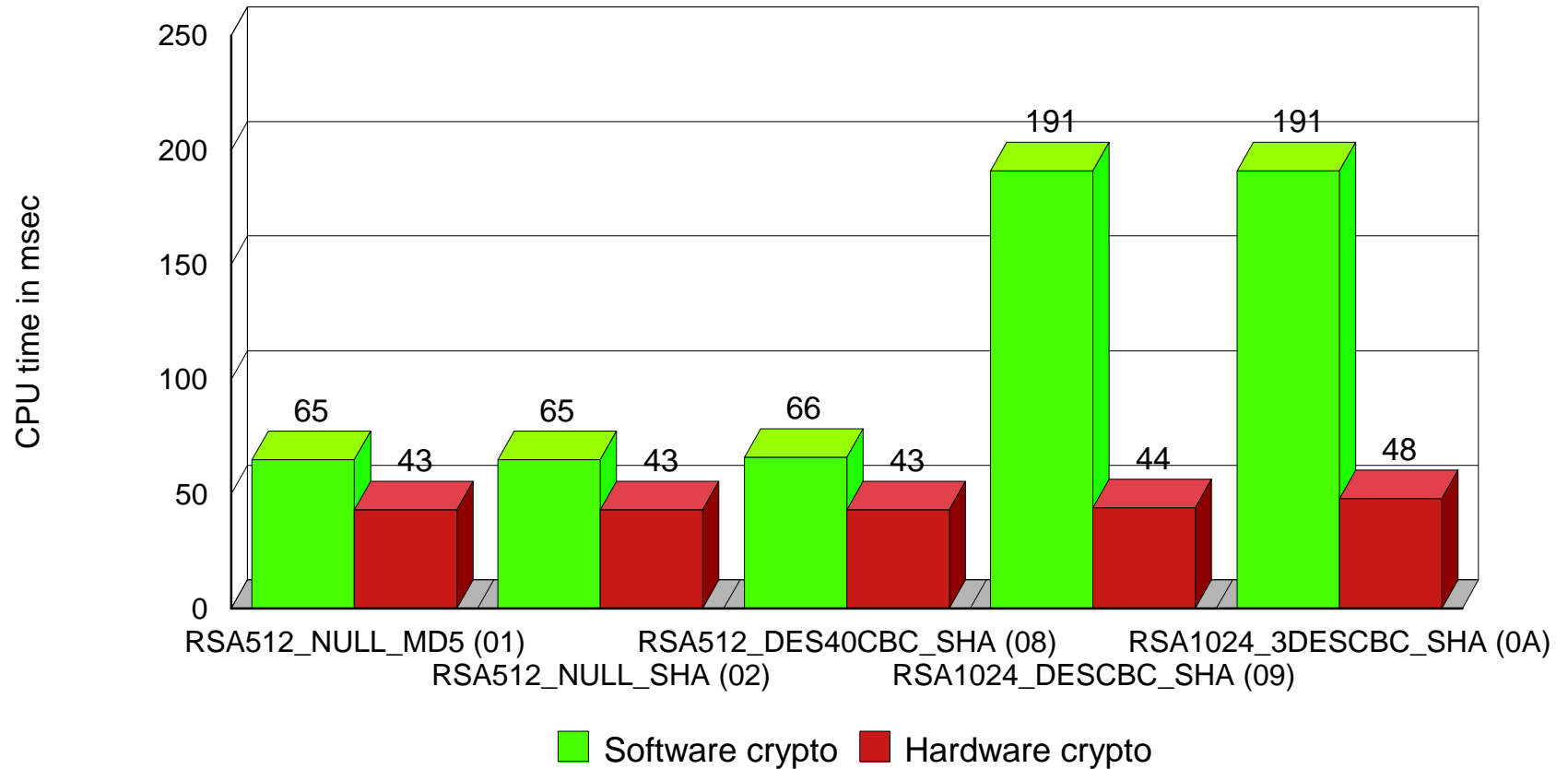
Software Crypto only!



e-business



Measurements Results - SSL Handshake



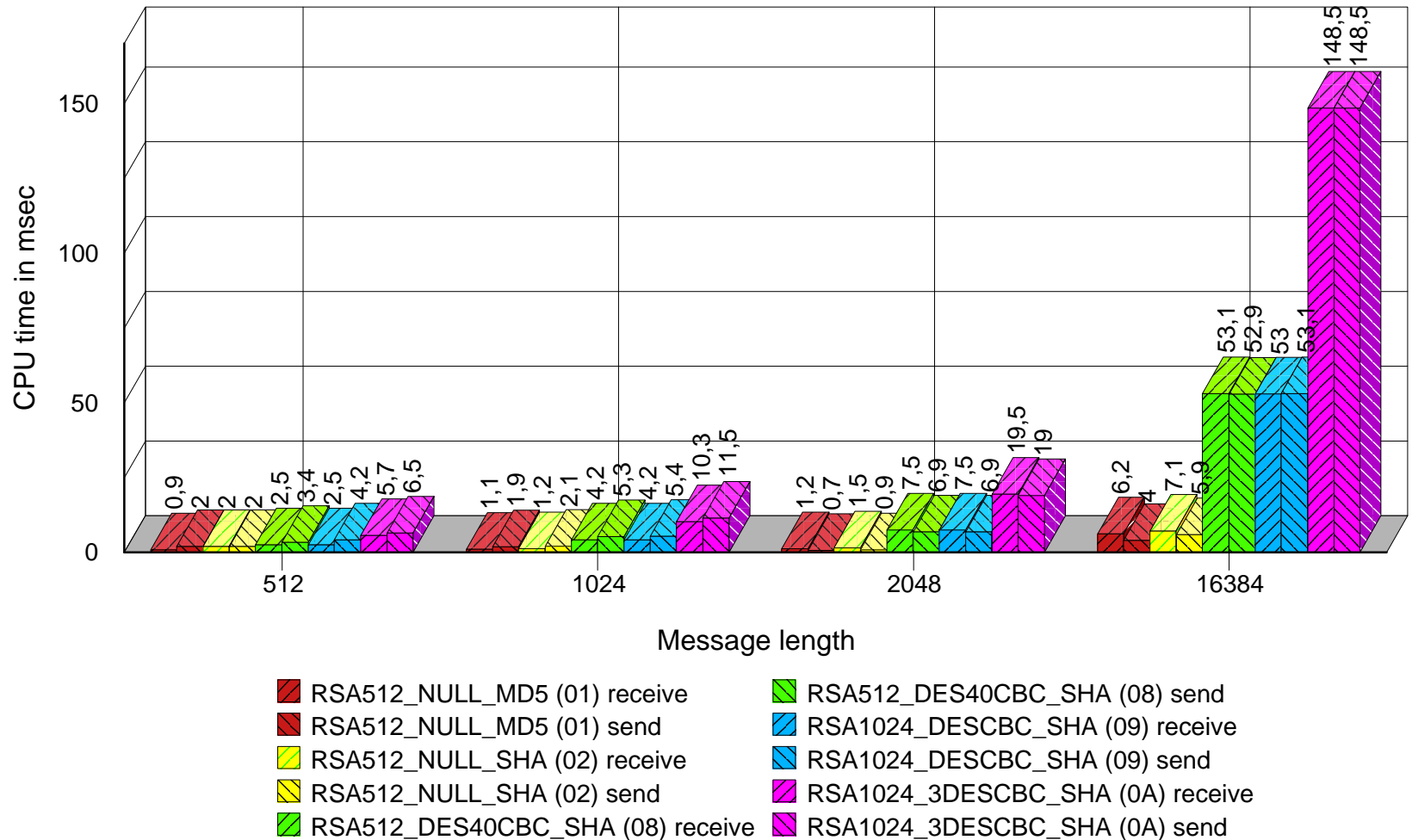
HW Crypto:

- CPU time and elapsed time is independent of cipher suite used
- SSL handshake takes about 43-48 msec CPU time (connection establishment)



Measurements Results - SSL data transfer

e-business



CPU time depends on used hashing (SHA/MD5) and encryption algorithm (DES/3DES)

Software Crypto only!



e-business

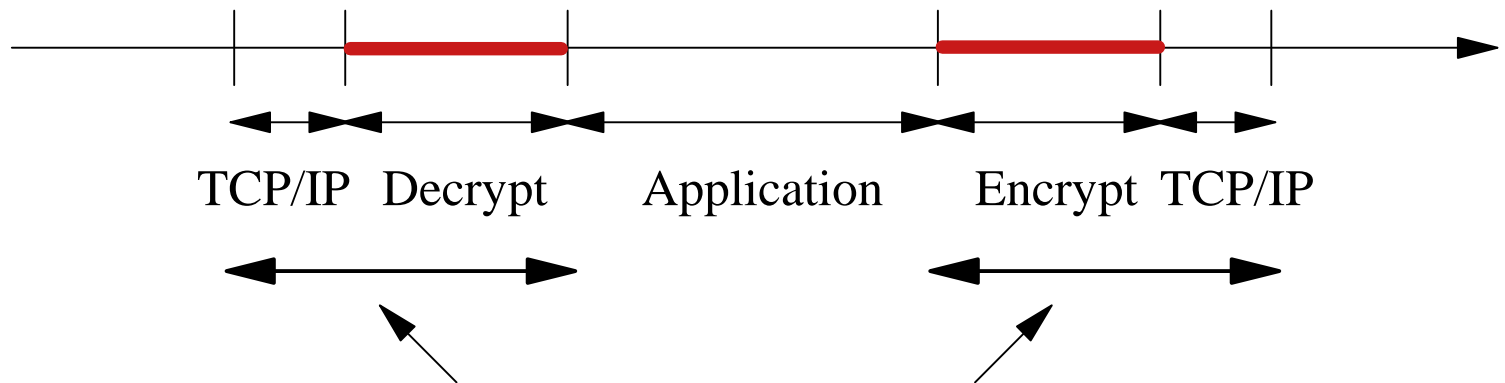


SSL data transfer overhead

Non SSL



SSL



this has been
measured



e-business



Measurements Results - conclusion

- HW Crypto
 - ▶ Supports RSA operations only (e.g. used by SSL handshake)
 - ▶ CPU time/elapsed time is independent of operation and key length
 - ▶ Software RSA encryption is faster in terms of elapsed time (on large processors)
 - but hardware crypto saves CPU time
- SW Crypto
 - ▶ CPUtime /elapsed time is very dependent on CPU speed and utilization



e-business



SSL Performance Recommendations

- Use SSL only if there is a need for
 - ▶ If at least one of the following is required
 - Keeping secrets
 - Proving identity
 - Verifying information
- Cipher Suites 01 and 02 has less CPU-time consumption, but NO data encryption
 - ▶ RSA512_NULL_MD5, RSA512_NULL_SHA
- If data encryption is required
 - ▶ Use cipher suites 08, 09 or 0A
 - ▶ 08 uses 512 bit keys, others 1024
 - ▶ 1024 bit RSA keylength is recommended (from a security point of view)



e-business



Dependencies for VSE/ESA Growth

- System dependencies
 - ▶ Many control-blocks etc.. still below the line
 - ▶ VTAM IOBUF areas in System GETVIS-24
 - ▶ Non-Parallel-Share limits n-way support
 - ▶ Number of tasks
 - Up to 255, 32 per partition, 208 subtasks in total
- Application dependencies
 - ▶ Integrated system concepts/functions
 - ▶ Functions/Applications dependencies
 - ▶ Number of users per TCP/IP partition



e-business



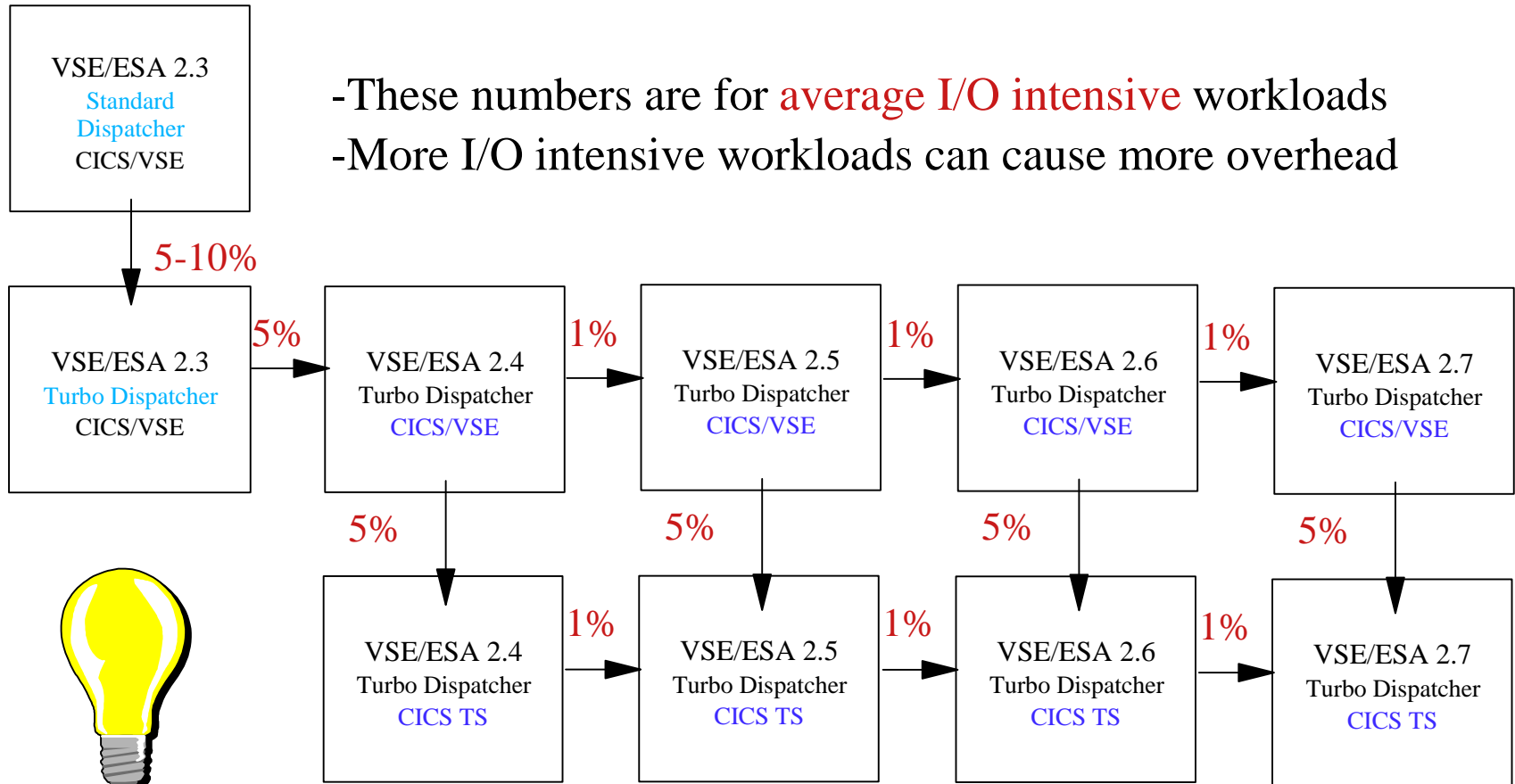
Dependencies for VSE/ESA Growth - continued

- Not being considered to be a limit
 - ▶ Number of partitions
 - 12 static + 150-200 dyn. partitions
 - ▶ Real storage (max. 2 GB)
 - ▶ Total virtual storage (max. 90 GB)
 - ▶ Total number of devices (3 digit CUU)
 - Max. 1024 devices (and 16 channels)
 - ▶ Total number of logical units
 - 255 per partition and $12 \times 255 = 3060$ in total
 - ▶ Label area
 - Max. about 9000 in total, and 712 in sub areas



Overhead Deltas for VSE Releases

e-business



Remember that you get a lot of **new functions** that in most cases helps you to **increase VSE system performance and throughput**:

Partition Balancing, PRTY SHARE (Turbo Dispatcher), FlashCopy (ESS), Buffer Hasing, Shared data Tables (CICS TS)



e-business



Further Information

- VSE Homepage:
<http://www.ibm.com/servers/eserver/zseries/os/vse/>
- VSE Performance Homepage:
<http://www.ibm.com/servers/eserver/zseries/os/vse/library/vseperf.htm>
- Performance Documents from W. Kraemer
 - ▶ available on the Performance Homepage
- VSE/ESA e-business Connectors User's Guide
<http://www.ibm.com/servers/eserver/zseries/os/vse/pdf/ieswue20.pdf>